## Short Notes on Capacity Measures and Additivity

Siqi Yao

SCHOOL OF DATA SCIENCE

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# Classical Channel Capacity

- The basic mathematical model for a communication system is shown in Figure 1, where:
  - X/Y are the channel input/output symbol taken in an alphabet  $\mathcal{X}/\mathcal{Y}$ .
  - $p(y \mid x) = p_{Y|X}(y \mid x)$  is the conditional probability distribution.
- Model X, Y as random variables. Since  $p_{Y|X}(y \mid x)$  is a fixed property of the channel, the choice of the marginal distribution  $p_X(x)$  completely determines the joint distribution  $p_{X,Y}(x,y)$ .
- Thus, we can define the **classical channel capacity** as

$$C = \sup_{p_X(x)} I(X;Y) \tag{1}$$

where  $I(X;Y) = \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}} P_{X,Y}(x,y) \log \frac{P_{X,Y}(x,y)}{P_X(x)P_Y(y)}$  is the **mutual information**.

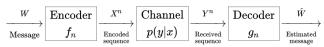


Figure 1: Communication Model. Source: [Wikipedia contributors, 2025]

## Additivity of Classical Channel Capacity

- Classical channel capacity is **additive** over independent **classical** channels: using two independent channels in parallel provides the same capacity as the sum of each.
- Let  $p_1, p_2$  be two independent channels. Define the product (parallel) channel  $p_1 \times p_2$  as:

$$(p_1 \times p_2) ((y_1, y_2) \mid (x_1, x_2)) = p_1(y_1 \mid x_1) \cdot p_2(y_2 \mid x_2).$$
 (2)

Then the additivity states that:

$$C(p_1 \times p_2) = C(p_1) + C(p_2).$$
 (3)

## Classical Channel Capacity of Quantum Channels

- However, classical channel capacity of quantum channels (classical input and output, transmission through quantum channels) does not satisfy additivity.
- The classical capacity of a quantum channel  $\mathcal N$  is given by the HSW theorem:

$$C(\mathcal{N}) = \lim_{n \to \infty} \frac{1}{n} \chi^*(\mathcal{N}^{\otimes n}) \tag{4}$$

where  $\chi^*(\mathcal{N})$  is the maximum Holevo information of  $\mathcal{N}$ :

$$\chi^*(\mathcal{N}) = \max_{\{p_x, \rho_x\}} \left[ S\left(\sum_x p_x \mathcal{N}(\rho_x)\right) - \sum_x p_x S(\mathcal{N}(\rho_x)) \right]$$
 (5)

• Since  $\chi^*(\mathcal{N})$  does not satisfy additivity, same is for  $C(\mathcal{N})$ .

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## Super-additivity of Holevo Information

- [Shor, 2004] showed that the following four conjectures are either all true or all false:
  - Additivity of the minimum entropy output of a quantum channel;
  - Additivity of the maximum Holevo information of a quantum channel;
  - Additivity of the entanglement of formation;
  - Strong super-additivity of the entanglement of formation.
- [Hastings, 2009] showed that all of these conjectures are false.
- In fact,  $\chi^*(\mathcal{N})$  satisfies **super-additivity**, i.e., for n independent quantum channels  $\{\Phi_i\}_{i=1}^n$ , when combined in parallel as a new channel  $\mathcal{N} := \Phi_1 \otimes \cdots \otimes \Phi_n$ , we have the inequality

$$\chi^*(\mathcal{N}) \ge \sum_{i=1}^n \chi^*(\Phi_i) \tag{6}$$

which can be strict.

#### Coherent Information and Quantum Capacity

• For a quantum channel  $\Phi$ , the **coherent information** is defined as:

$$Q^{(1)}(\Phi) := \max_{\rho} I_c(\rho, \Phi) = \max_{\rho} [S(\Phi(\rho)) - S(\Phi^c(\rho))]$$
 (7)

• The quantum capacity of  $\Phi$  is given by the regularized formula:

$$Q(\Phi) = \lim_{n \to \infty} \frac{Q^{(1)}(\Phi^{\otimes n})}{n} = \sup_{n \in \mathbb{N}} \frac{1}{n} Q^{(1)}(\Phi^{\otimes n})$$
(8)

where the second equality is due to the Fekete's lemma.

• The coherent information is also **super-additive**, i.e., for n independent quantum channels  $\{\Phi_i\}_{i=1}^n$ , and  $\mathcal{N} := \Phi_1 \otimes \cdots \otimes \Phi_n$ , we have the inequality

$$Q^{(1)}(\mathcal{N}) \ge \sum_{i=1}^{n} Q^{(1)}(\Phi_i) \tag{9}$$

which can be strict.

#### Weak Additivity

- The super-additivity can be divided into two categories based on the violation of **weak and strong additivity**. Take coherent information as example.
- If

$$Q^{(1)}(\Phi^{\otimes n}) = n Q^{(1)}(\Phi) \quad \forall n \in \mathbb{N}, \tag{10}$$

i.e., in the setting of multiple independent and parallel uses of  $\Phi$ , then the coherent information of  $\Phi$  satisfies weak additivity.

• Once weak additivity holds, the quantum capacity simplifies to

$$Q(\Phi) = Q^{(1)}(\Phi) \tag{11}$$

 $\bullet$  Violation of weak additivity  $\Rightarrow$  weak superadditivity: There exists channels such that

$$\mathcal{Q}^{(1)}(\Phi^{\otimes n}) > n \, \mathcal{Q}^{(1)}(\Phi). \tag{12}$$

#### Strong Additivity

• For a pair of channel  $\Phi_1, \Phi_2$ , if we fix  $\Phi_1$  and the equality

$$Q^{(1)}(\Phi_1 \otimes \Phi_2) = Q^{(1)}(\Phi_1) + Q^{(1)}(\Phi_2). \tag{13}$$

holds for all  $\Phi_2$ , then the coherent information of  $\Phi_1$  satisfies strong additivity.

• Hence, the quantum capacity satisfies:

$$\mathcal{Q}(\Phi_1 \otimes \Phi_2) = \mathcal{Q}(\Phi_1) + \mathcal{Q}(\Phi_2). \tag{14}$$

• Violation of strong additivity ⇒ strong superadditivity: There exists a pair of channels such that

$$Q^{(1)}(\Phi_1 \otimes \Phi_2) > Q^{(1)}(\Phi_1) + Q^{(1)}(\Phi_2).$$
 (15)

This implies that using different channels in parallel can transmit more information than using them individually.

## Private Information and Private Capacity

- The private capacity  $\mathcal{P}(\Phi)$  of a quantum channel  $\Phi$  is defined as the largest rate at which classical information can be faithfully sent through the channel, such that the environment gains no meaningful knowledge about the information.
- The channel's private information is defined as:

$$\mathcal{P}^{(1)}(\Phi) = \max_{\{p_x, \rho_x\}} \left[ I_c(\sum_x p_x \rho_x, \Phi) - \sum_x p_x I_c(\rho_x, \Phi) \right]$$
(16)

which satisfies **supper-additivity**.

• The channel's private capacity is given by the regularized expression:

$$\mathcal{P}(\Phi) = \lim_{n \to \infty} \frac{\mathcal{P}^{(1)}(\Phi^{\otimes n})}{n} = \sup_{n \in \mathbb{N}} \frac{1}{n} \mathcal{P}^{(1)}(\Phi^{\otimes n})$$
 (17)

#### Equivalent Expression of Private Information

• Denote the quantum channel and its complementary respectively as  $\mathcal{B}, \mathcal{C}.\bar{\rho}_a := \sum_x p_x \rho_x$ . We have:

$$\sigma_B = \mathcal{B}(\bar{\rho}_a), \quad \sigma_C = \mathcal{C}(\bar{\rho}_a), \quad \sigma_B^x = \mathcal{B}(\rho_x), \quad \sigma_C^x = \mathcal{C}(\rho_x)$$
 (18)

• We can derive:

$$H(\sigma_{XB}) = H(p_x) + \sum_{x} p_x H(\sigma_B^x), \quad H(\sigma_X) = H(p_x)$$
(19)

Thus,

$$I(X;B)_{\sigma} = H(\sigma_X) + H(\sigma_B) - H(\sigma_{XB}) = H(\sigma_B) - \sum_{x} p_x H(\sigma_B^x)$$
 (20)

$$=H(\mathcal{B}(\bar{\rho}_a))-\sum p_x H(\mathcal{B}(\rho_x)) \tag{21}$$

Similarly:

$$I(X;C)_{\sigma} = H(\mathcal{C}(\bar{\rho}_a)) - \sum_{x} p_x H(\mathcal{C}(\rho_x))$$
 (22)

#### Equivalent Expression of Private Information (cont'd)

So we have:

$$I(X;B)_{\sigma} - I(X;C)_{\sigma} = \underbrace{\left[H(\mathcal{B}(\bar{\rho}_{a})) - H(\mathcal{C}(\bar{\rho}_{a}))\right]}_{=I_{c}(\bar{\rho}_{a},\mathcal{B})} - \sum_{x} p_{x} \underbrace{\left[H(\mathcal{B}(\rho_{x})) - H(\mathcal{C}(\rho_{x}))\right]}_{=I_{c}(\rho_{x},\mathcal{B})}$$
(23)  
$$= I_{c}(\bar{\rho}_{a},\mathcal{B}) - \sum_{x} p_{x} I_{c}(\rho_{x},\mathcal{B})$$
(24)

At last we obtain an equivalent expression for (16):

$$\mathcal{P}^{(1)}(\Phi) = \max_{\{p_x, \rho_x\}} \left[ I(X; \Phi)_{\sigma} - I(X; \Phi^c)_{\sigma} \right]$$
 (25)

Full derivation details can be found in the "Quantum Information Theory
Preliminaries" section of my notes Quantum SWITCH for Communication
Enhancement.

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