

Thompson's Group F

Yarik Vitovsky

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Main Talking Points

- Thompson's Group Introduction
- Binary Tree Mirrored Image and General Expression
- \overline{X} Encryption

Thompson's Group F

Basic Definition

Thompson's group F is defined as the group of piecewise linear maps:

$$F : [0, 1] \rightarrow [0, 1]$$

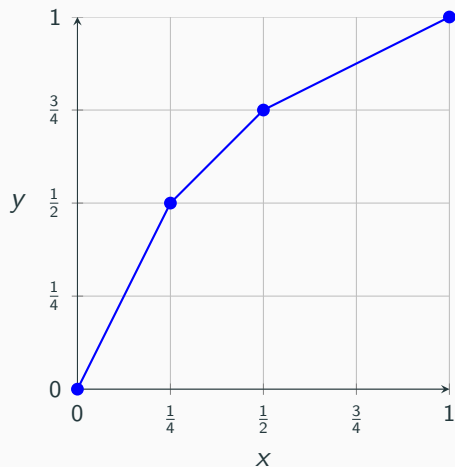
These maps are **homeomorphisms** (bijective and continuous functions with continuous inverses) that **preserve orientation**.

Definition of an Element

Each element in the group is defined by how it partitions the interval. The partition is made by dyadic rational breakpoints of the form:

$$\frac{m}{2^n}, \quad m \in \mathbb{Z}, \quad n \in \mathbb{N}$$

Example X_0



The element x_0 is defined by the following breakpoints:

$$\left[\left(\frac{1}{4}, \frac{1}{2}\right), \left(\frac{1}{2}, \frac{3}{4}\right)\right]$$

with the slopes of $f(t)$ in each interval:

$$f(t) = \begin{cases} 2t & \text{for } 0 \leq t \leq \frac{1}{4} \\ t + \frac{1}{4} & \text{for } \frac{1}{4} \leq t \leq \frac{1}{2} \\ \frac{t+1}{2} & \text{for } \frac{1}{2} \leq t \leq 1 \end{cases}$$

Generators in Thompson's Group F

Generators are the fundamental building blocks of the group. In Thompson's group F , all elements can be written as finite compositions of:

x_0, x_1, x_2, \dots and their inverses

Each generator x_i corresponds with:

- finite amount of unique breakpoints (all dyadic rationals)
- slopes that are powers of 2
- a well-defined inverse x_i^{-1} (also a homeomorphism)

Algebraic Form of F

These generators arise from a broader algebraic definition of Thompson's group F , known as its infinite presentation, which is given as:

$$\langle x_0, x_1, x_2, \dots \mid x_i^{-1} x_j x_i = x_{j+1}, \text{ for } i < j \rangle$$

meaning that, due to its non-abelian structure and the function of composition, Thompson's group contains an infinite number of elements.

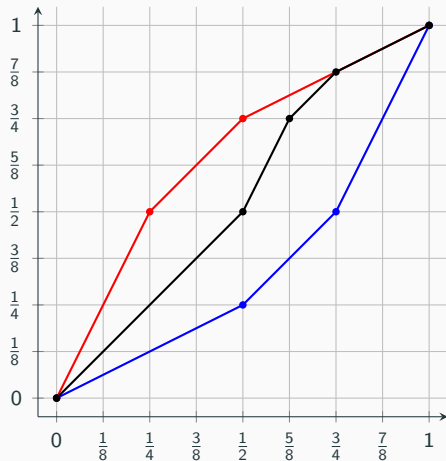
For every X_n , each breakpoint is generated by the formula:

$$\left[\left(1 - \frac{1}{2^n}, 1 - \frac{1}{2^{n+1}} \right), \left(1 - \frac{3}{2^{n+2}}, 1 - \frac{1}{2^{n+1}} \right), \left(1 - \frac{1}{2^{n+1}}, 1 - \frac{1}{2^n} \right) \right]$$

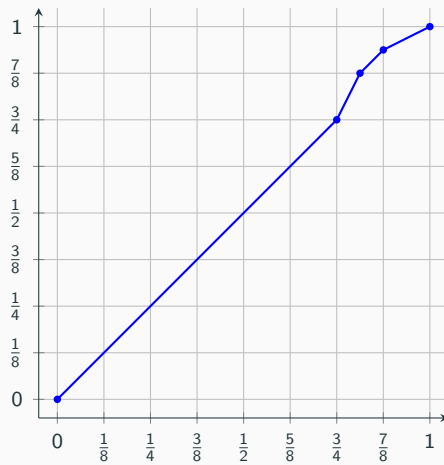
Generating X_2

To generate the element X_2 we apply the conjugation: $X_0^{-1}X_1X_0 = X_2$

X_0^{-1} X_1 X_0

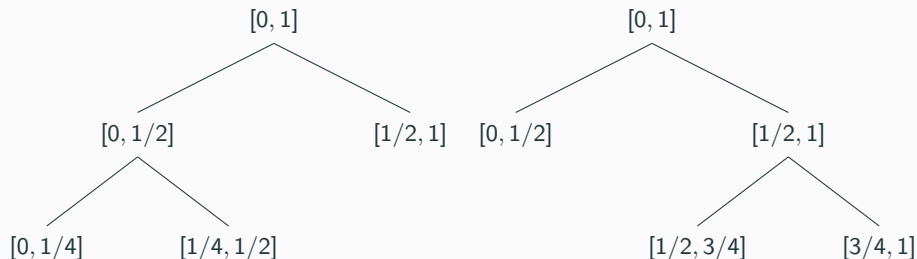


X_2



Binary Trees

Since dividing the interval $[0,1]$ defines a mapping function, it can be equivalently represented as a binary tree



This binary trees correspond to x_0 with the following breakpoints: $[(\frac{1}{4}, \frac{1}{2}), (\frac{1}{2}, \frac{3}{4})]$

Mirroring for Encryption

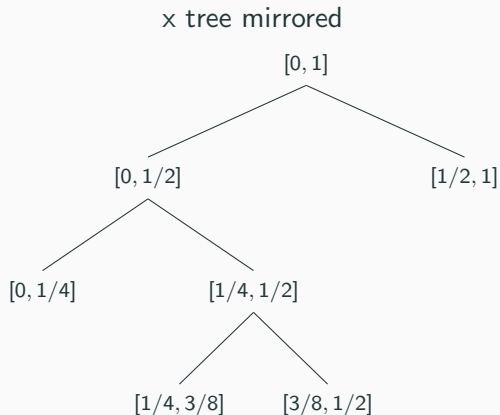
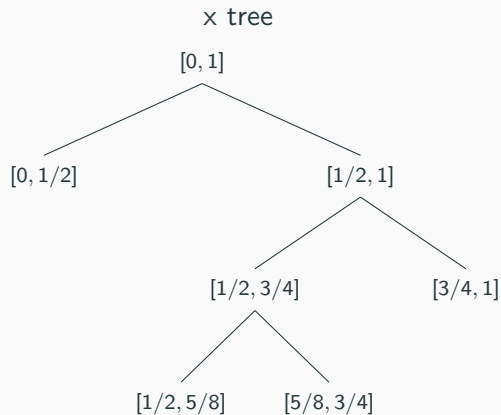
In our project, we studied mirrored binary trees as a part of the RSA public key protocol.

The key idea was to encrypt a message or generate a key using the mirroring operation. We took the tree representation of a known generator (e.g., X_1) and applied a structured mirroring process to produce a transformed version, $\overline{X_1}$.

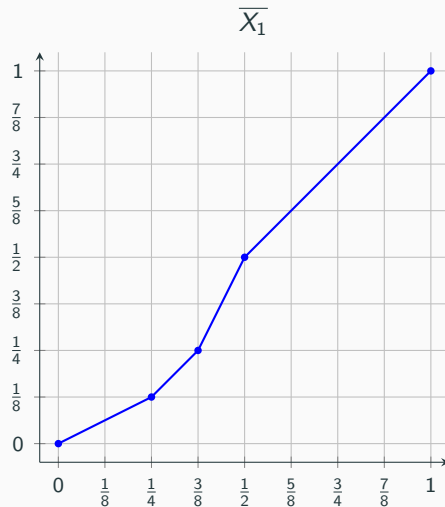
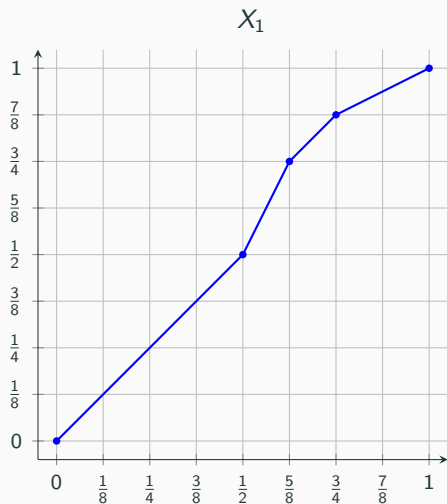
This mirrored version preserves important algebraic properties while being structurally different, making it a strong candidate for use in encoding, decoding, or generating secure elements in group-based cryptographic systems.

Mirroring Operation

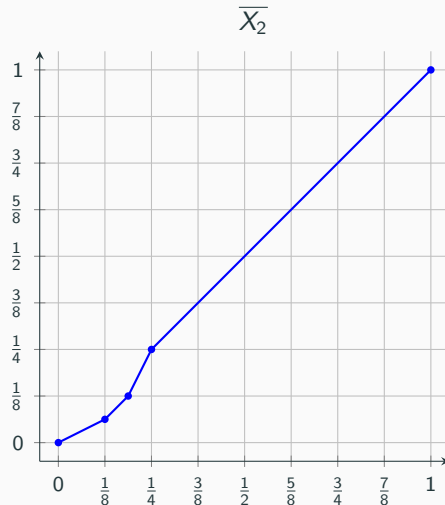
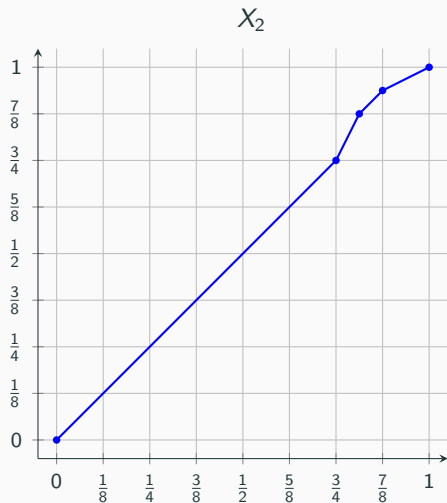
X_1



Mirrored Image Graph



The Pattern Continues



Interval Reflection in $\overline{X_n}$

Let the ordered sequence of subintervals that partition $[0, 1]$ for X_n be:

$$\{(a_1, b_1), (a_2, b_2), \dots, (a_k, b_k)\}$$

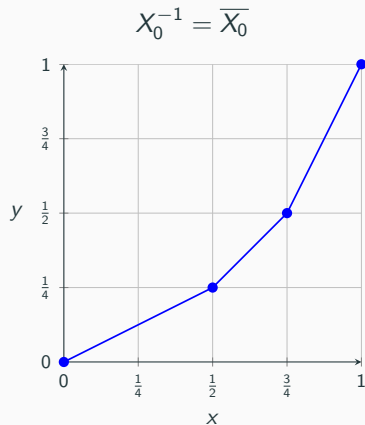
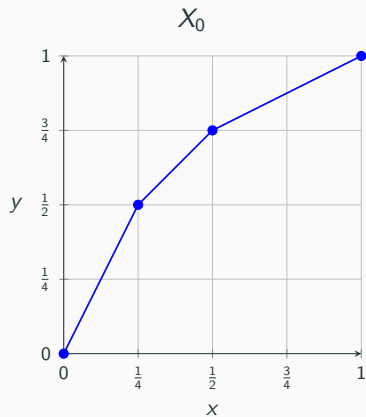
The mirrored function $\overline{X_n}$ reflects each part across the unit interval by the ordered sequence in reverse and by:

$$(a_i, b_i) \mapsto (1 - b_i, 1 - a_i)$$

The resulting sequence becomes:

$$\{(1 - b_k, 1 - a_k), \dots, (1 - b_2, 1 - a_2), (1 - b_1, 1 - a_1)\}$$

Uniqueness of X_0



Among all generators, X_0 is unique in that its inverse coincides with its mirrored version:

New Objective

So far, we've explored:

- The definition and structure of Thompson's group F
- Different representations of group elements with some examples
- The process and patterns that arise when mirroring these elements

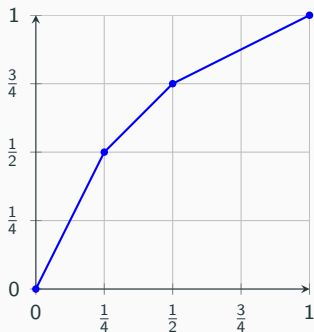
Our next goal:

To find a general algebraic expression for any mirrored element $\overline{X_n}$ using only the known generators and operations of F .

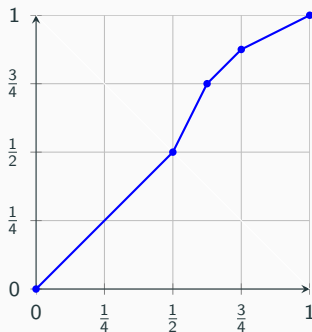
First Impression

To achieve our goal, we first recognize the common pattern of generators in growing iterations.

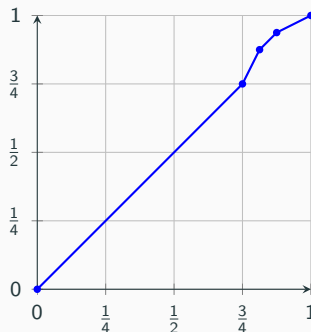
X_0



X_1



$X_0^{-1}X_1X_0 = X_2$



First Expression

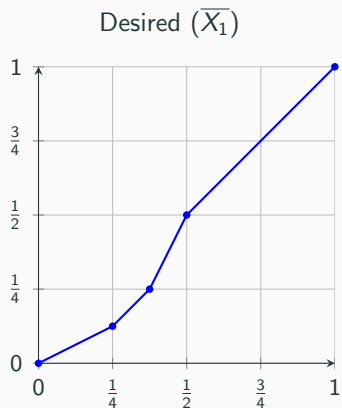
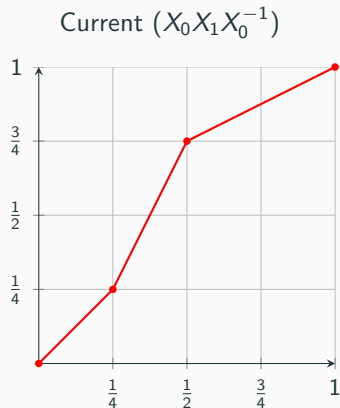
As we move to higher-index generators, we apply more layers of conjugation using X_0 and its inverse (excluding the original generators X_0 and X_1). This repeated conjugation causes the graph to gradually "shrink" towards 1.

$$X_0^{-1}X_1X_0 = X_2 \quad \Rightarrow \quad X_0^{-2}X_1X_0^2 = X_3 \quad \Rightarrow \quad X_0^{-3}X_1X_0^3 = X_4$$

This naturally raises the question: What happens if we reverse the direction of the conjugation? That is:

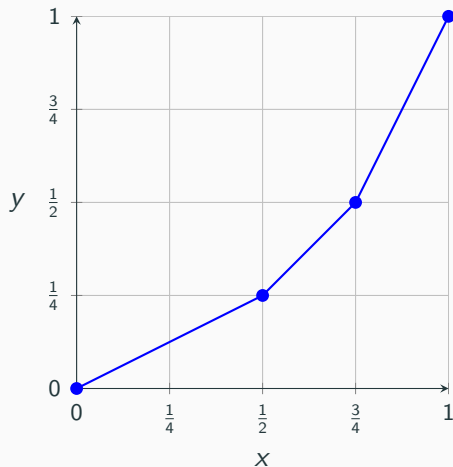
$$X_0X_1X_0^{-1} = ?$$

Current vs. Desired



We're actually quite close - the first and last subintervals, $(0, 1/4)$ and $(1/2, 1)$, already match the mirrored structure. What's left is adjusting the slopes.

Graph of X_0^{-1}



We observe that X_0^{-1} adjusts the slopes exactly as needed.

Slope adjustment:

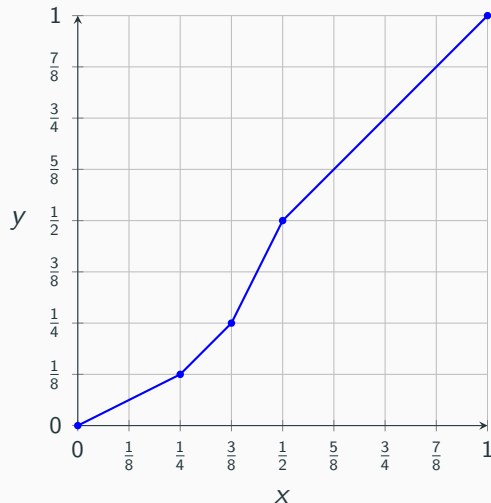
$$\text{Starting slope: } 1 \cdot \frac{1}{2} = \frac{1}{2}$$

$$\text{Ending slope: } \frac{1}{2} \cdot 2 = 1$$

This makes X_0^{-1} the final piece of the puzzle in constructing $\overline{X_1}$.

Result

Graph of $\overline{x_1}$



From our construction, we find:

$$X_0 X_1 X_0^{-2} = \overline{X_1}$$

This suggests the general formula:

$$\overline{X_n} = X_0^n X_1 X_0^{-(n+1)}$$

This expression generates the mirrored version $\overline{X_n}$ for any index n .

Proof By Induction

Goal: Show that

$$\overline{X_n} = X_0^n X_1 X_0^{-(n+1)} \quad \text{for all } n \in \mathbb{N}$$

using induction.

Base Case: $n = 1$

$$\overline{X_1} = X_0 X_1 X_0^{-2}$$

The base case holds.

Inductive Hypothesis: Assume

$$\overline{X_k} = X_0^k X_1 X_0^{-(k+1)} \quad \text{for some } k \in \mathbb{N}.$$

Inductive Step

To prove:

$$\overline{X_{k+1}} = X_0^{k+1} X_1 X_0^{-(k+2)}$$

Recall the recursive rule:

$$\overline{X_{k+1}} = X_0 \cdot \overline{X_k} \cdot X_0^{-1}$$

Substitute the inductive hypothesis:

$$= X_0 \cdot \left(X_0^k X_1 X_0^{-(k+1)} \right) \cdot X_0^{-1}$$

Group powers:

$$= X_0^{k+1} X_1 X_0^{-(k+2)}$$

The formula holds for $k + 1$. Thus, by induction, it holds for all $n \in \mathbb{N}$.