

Functional Pearls

Christophe Scholliers

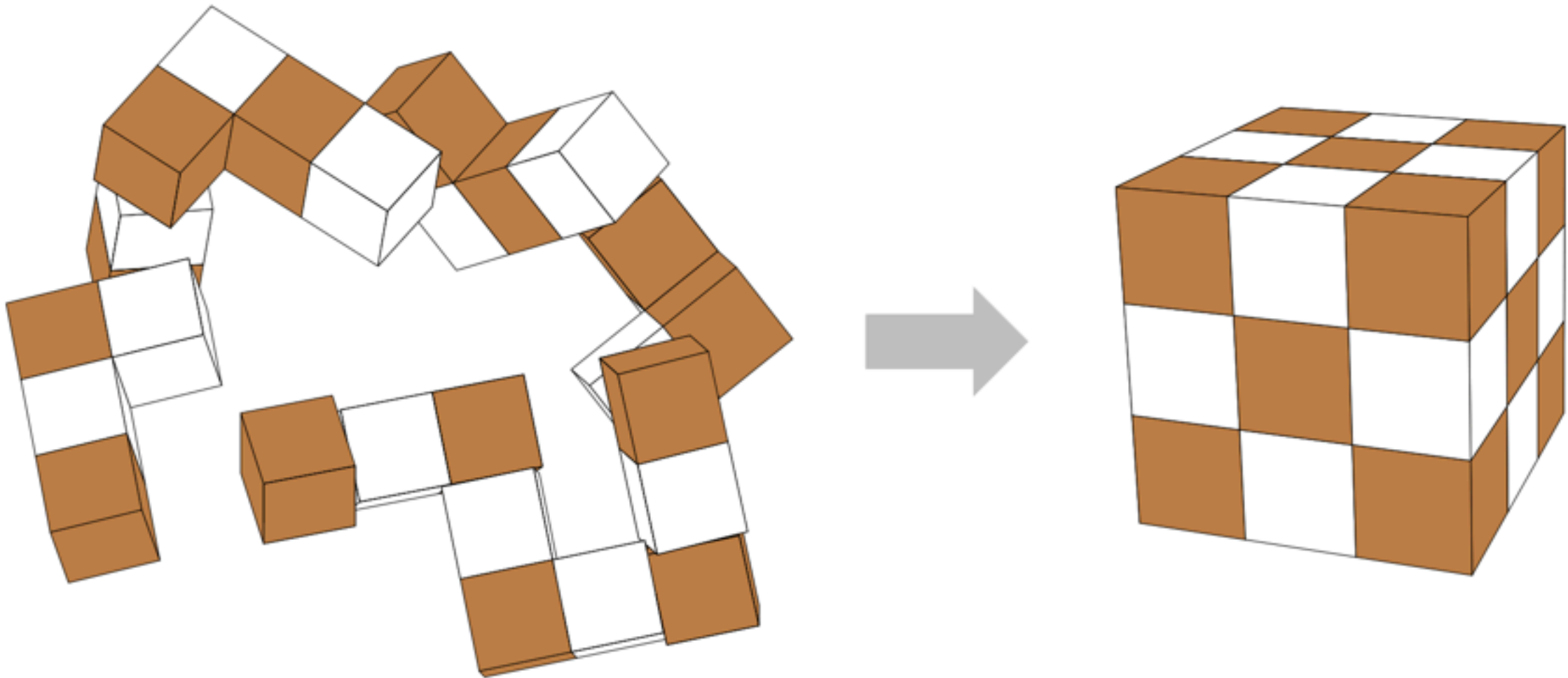
What are Functional Pearls ?

Functional pearls are elegant, instructive examples of functional programming. They are supposed to be fun, and they teach important programming techniques and fundamental design principles. They traditionally appear in [The Journal of Functional Programming](#), and at [ICFP](#) and affiliated workshops.

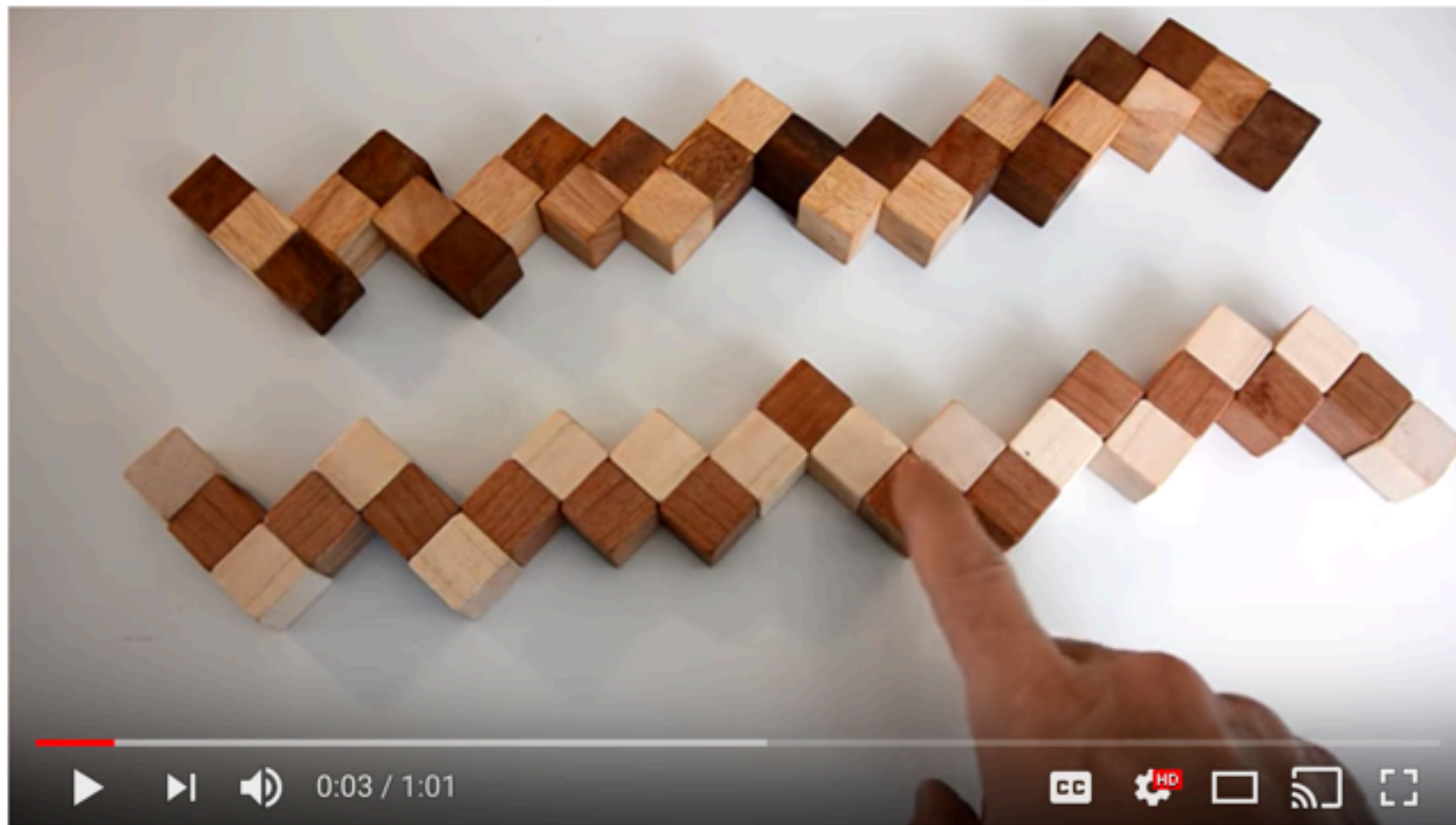
Solving the Snake Cube Puzzle in Haskell

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Snake Cube

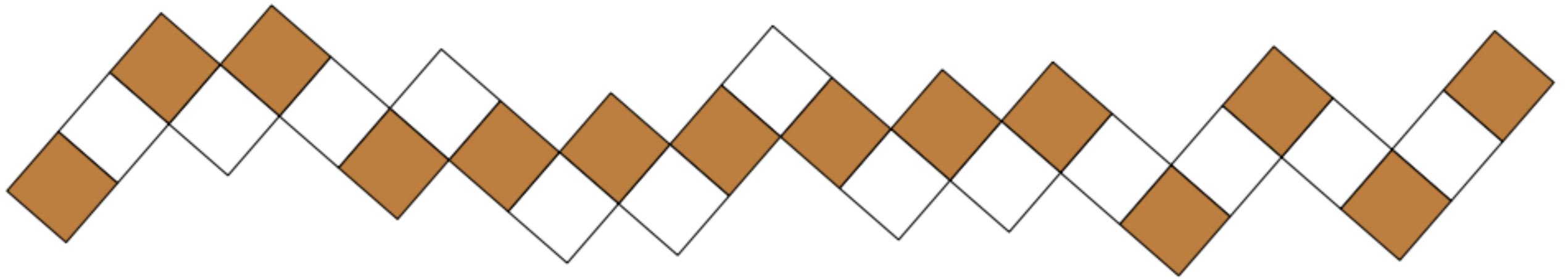


How to solve one



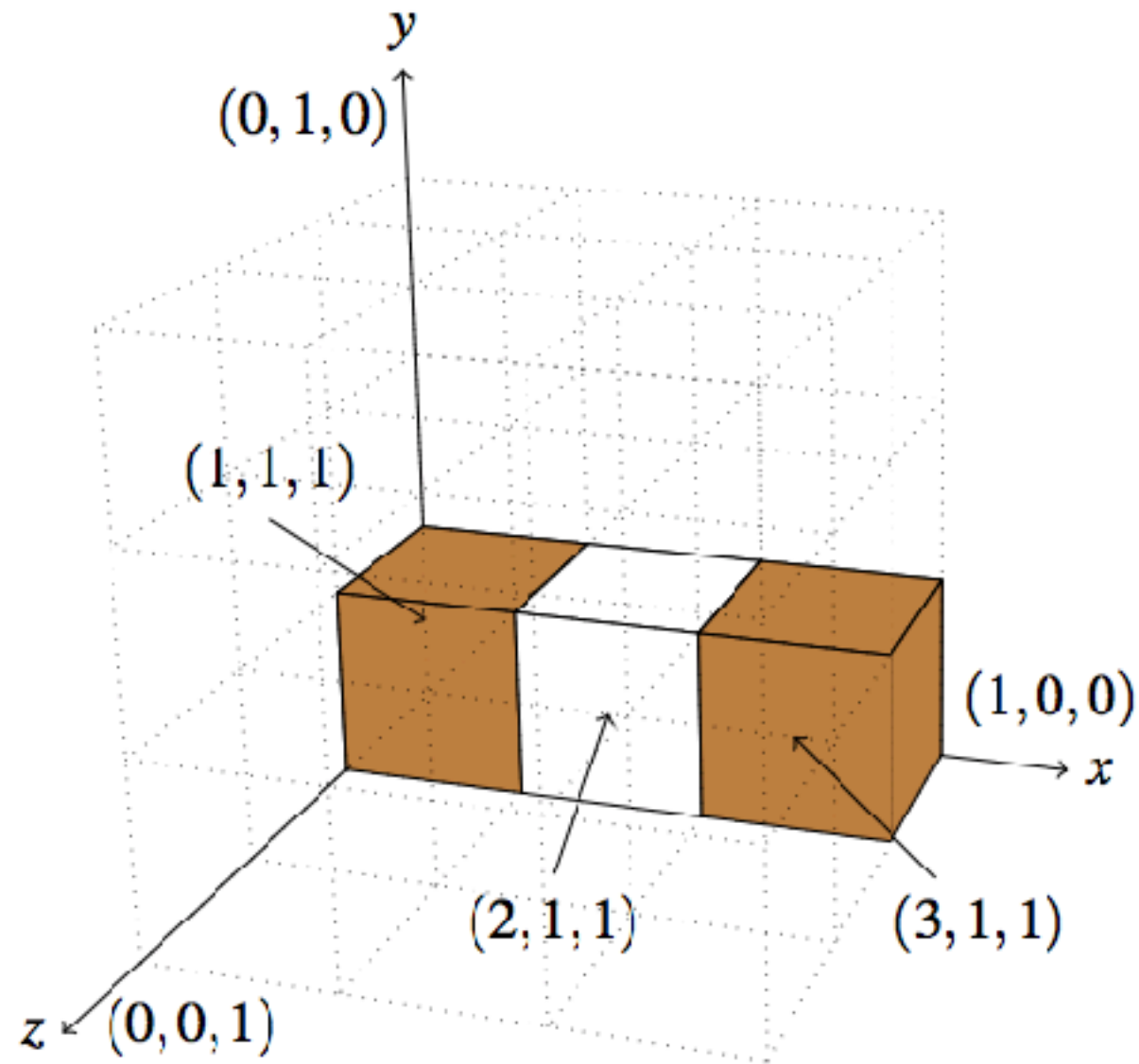
<https://www.youtube.com/watch?v=iTzVPgFjE9c>

Construction



```
snake :: [Int]
snake = [ 3, 2, 2, 3, 2, 3, 2, 2, 3, 3, 2, 2, 2, 3, 3, 3, 3 ]
```

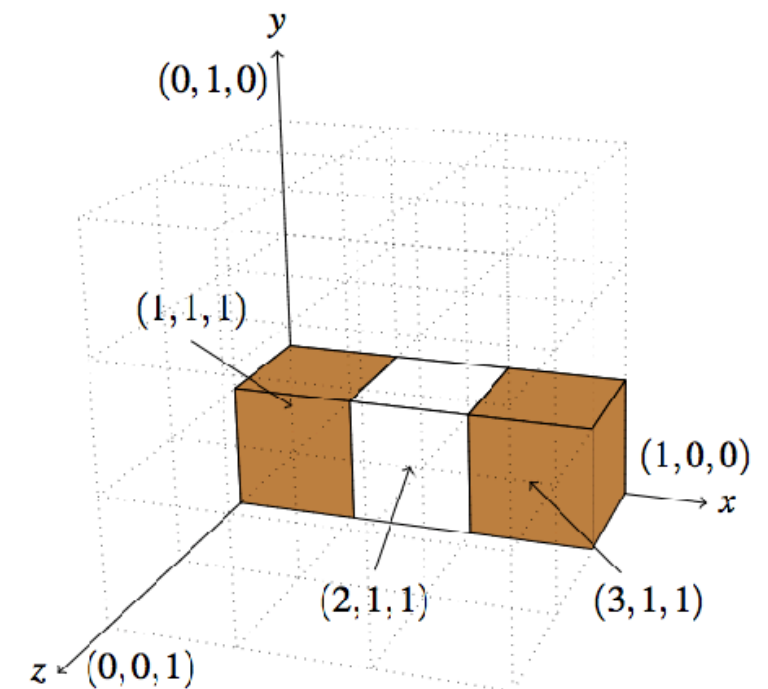
Moving in to Three Dimensions



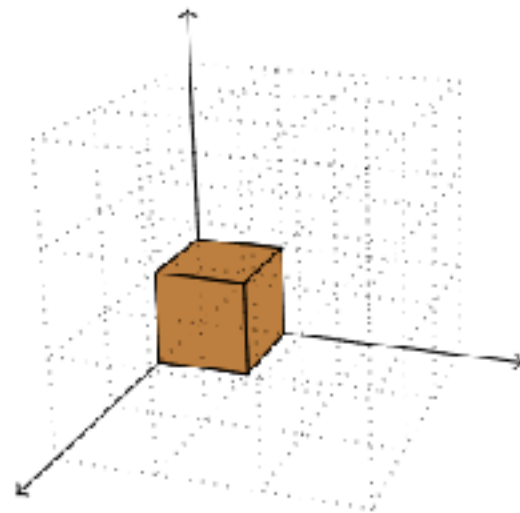
Moving in to Three Dimensions

```
type Position = (Int, Int, Int)
type Direction = (Int, Int, Int)
```

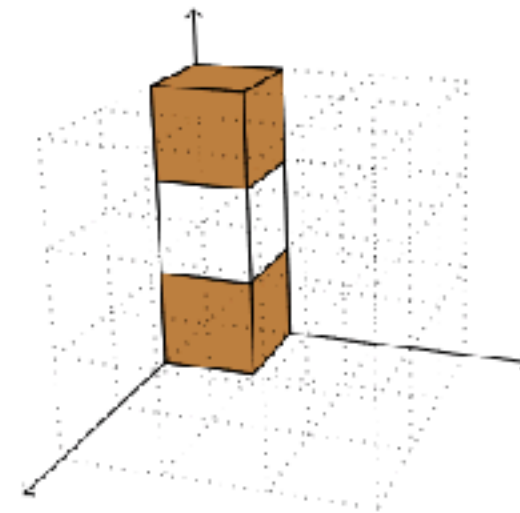
```
inCube :: Int -> Position -> Bool
inCube n (x, y, z) = valid x && valid y && valid z
                    where valid i = 1 <= i && i <= n
```



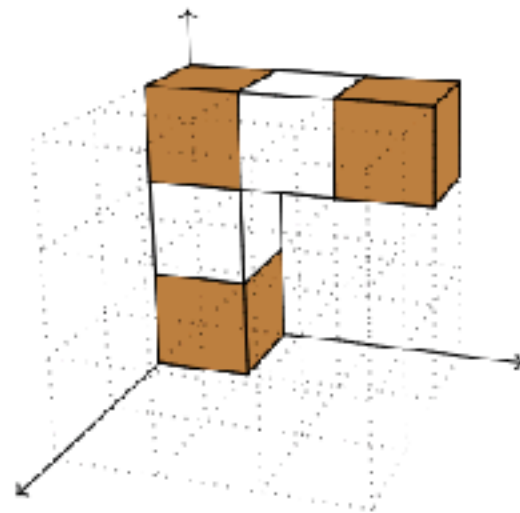
Describing Solutions



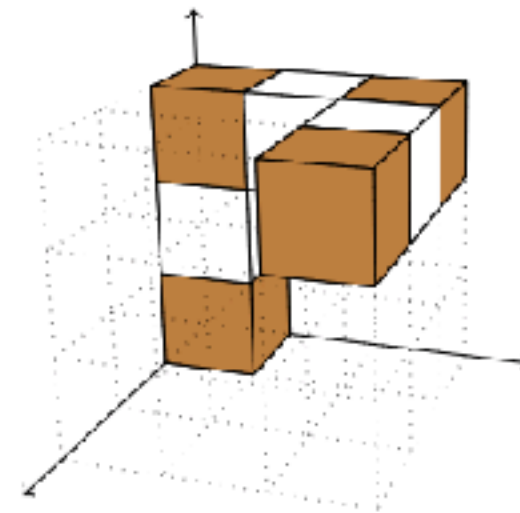
$$soln_0 = [(1, 1, 1)]$$



$$soln_1 = [(1, 3, 1), (1, 2, 1)] : soln_0$$



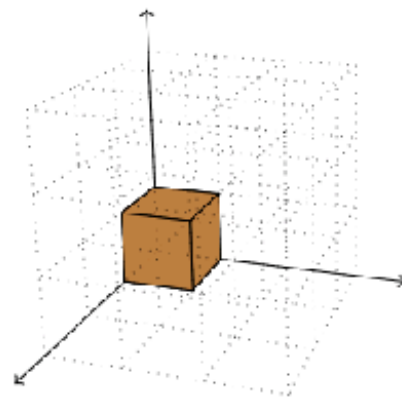
$$soln_2 = [(3, 3, 1), (2, 3, 1)] : soln_1$$



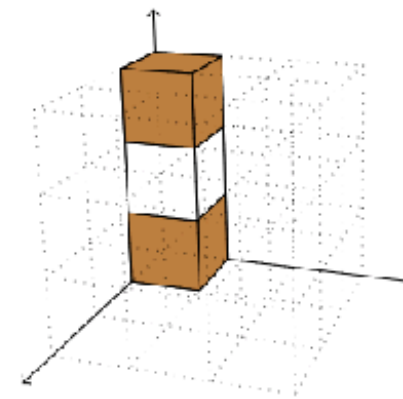
$$soln_3 = [(3, 3, 3), (3, 3, 2)] : soln_2$$

Describing Solutions

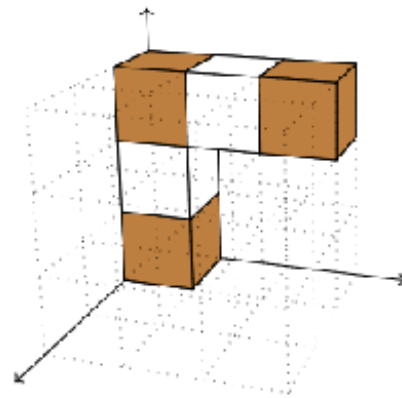
```
type Solution = [Section]  
type Section  = [Position]
```



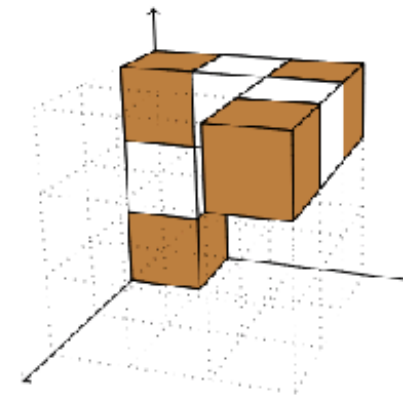
$soln_0 = [(1, 1, 1)]$



$soln_1 = [(1, 3, 1), (1, 2, 1)] : soln_0$



$soln_2 = [(3, 3, 1), (2, 3, 1)] : soln_1$



$soln_3 = [(3, 3, 3), (3, 3, 2)] : soln_2$

Building Sections

```
iterate :: (a -> a) -> a -> [a]
```

[Source](#)

#

`iterate` `f` `x` returns an infinite list of repeated applications of `f` to `x`:

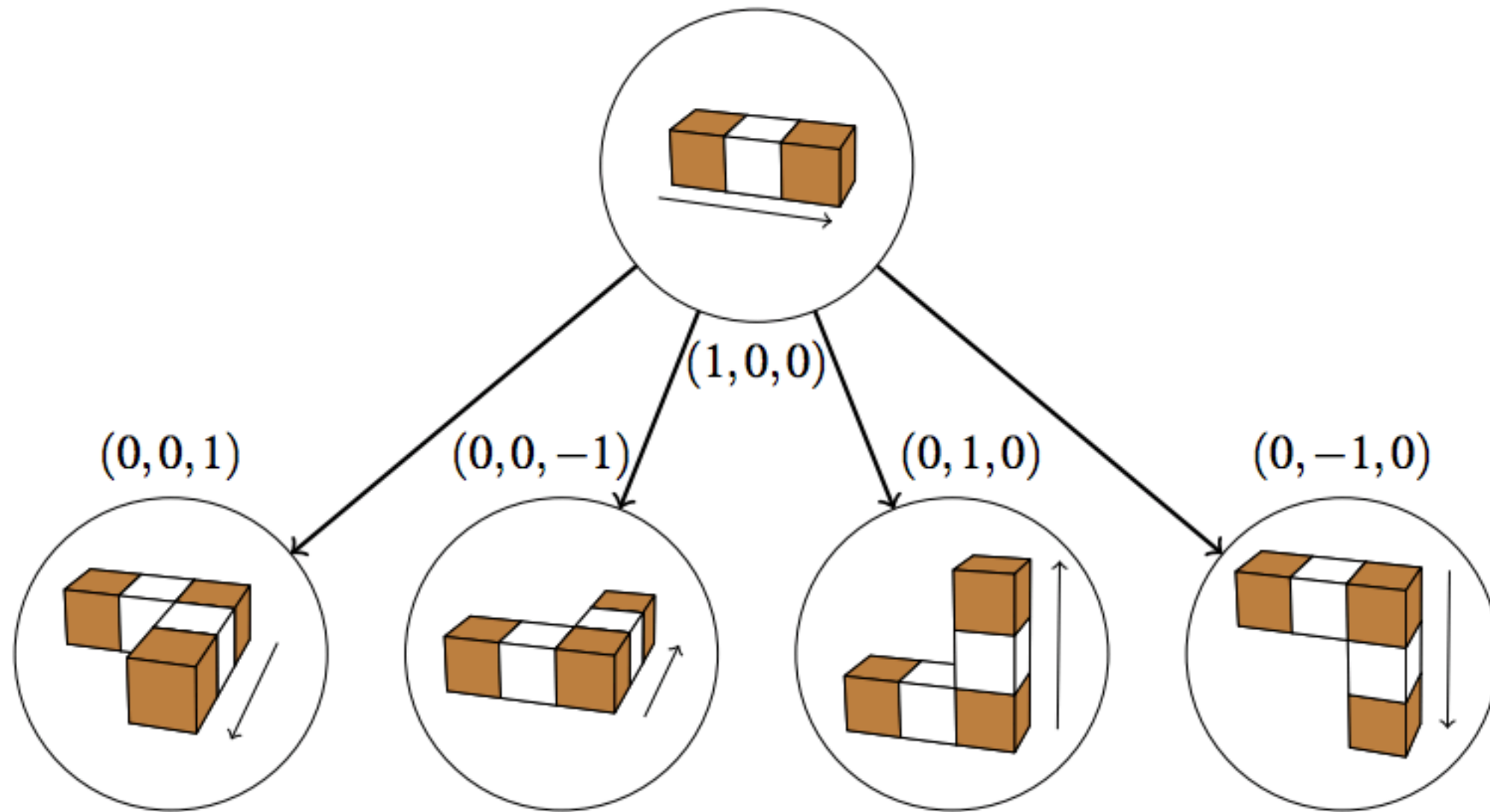
```
iterate f x == [x, f x, f (f x), ...]
```

```
section :: Position -> Direction -> Int -> Section
section start (u,v,w) len = reverse (tail (take len pieces))
  where pieces = iterate (\ (x, y, z) -> (x+u, y+v, z+w)) start
```

Example:

```
section (1, 1, 1) (0, 1, 0) 3 = [(1, 3, 1), (1, 2, 1)]
section (1, 3, 1) (1, 0, 0) 3 = [(3, 3, 1), (2, 3, 1)]
```

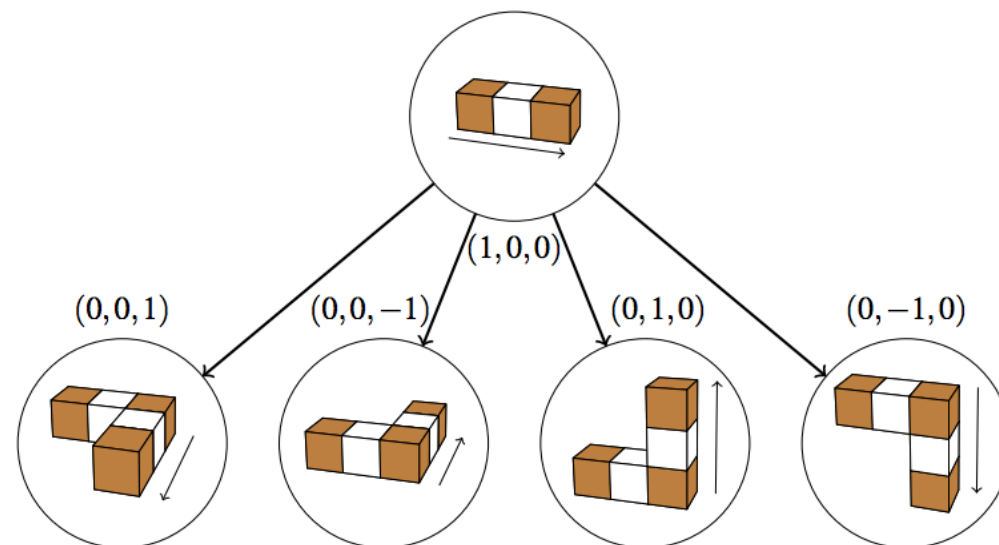
Changing Direction



Changing Direction

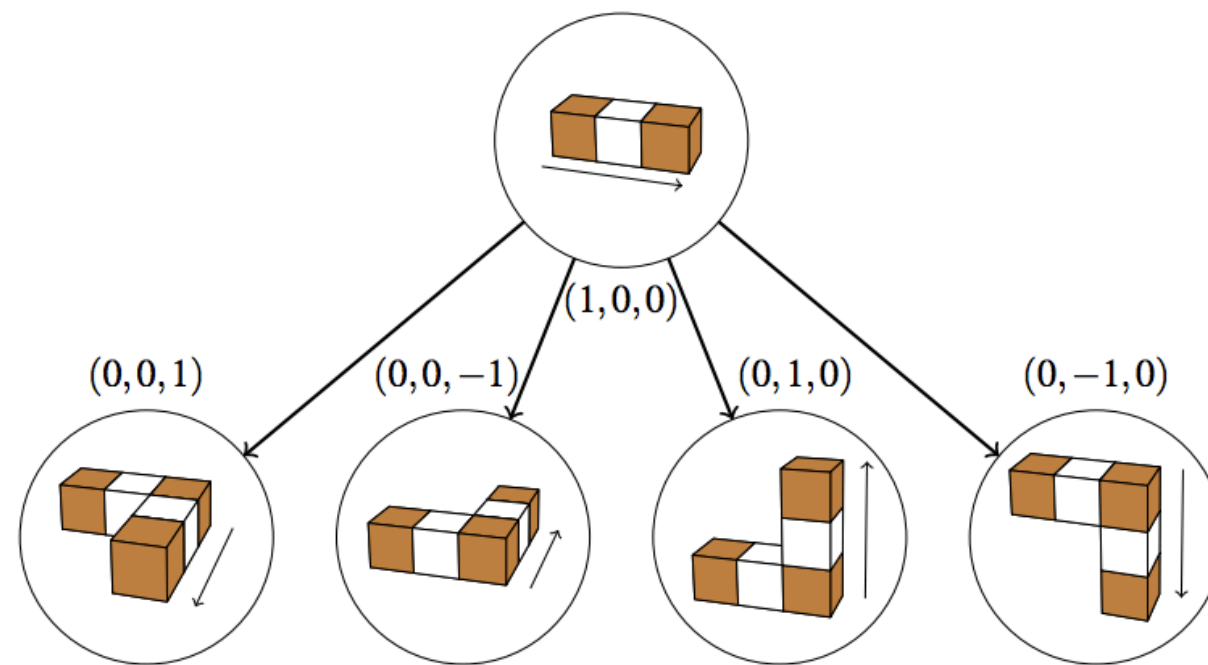
```
newDirs    :: Direction -> [Direction]
newDirs dir = [ rotate (sign dir) | rotate <- [left, right],
               sign    <- [id, inv] ]

where left, right, inv :: Direction -> Direction
left   (x, y, z)      = ( y,  z,  x)
right  (x, y, z)      = ( z,  x,  y)
inv    (x, y, z)      = (-x, -y, -z)
```



Changing Direction

```
newDirs (x, y, z) = [ (y, z, x),  
                      (-y, -z, -x),  
                      (z, x, y),  
                      (-z, -x, -y) ]
```



Describing Complete Puzzles

```
data Puzzle = Puzzle { sections :: [Int],  
                      valid      :: Position -> Bool,  
                      initSoln  :: Solution,  
                      initDir   :: Direction }
```

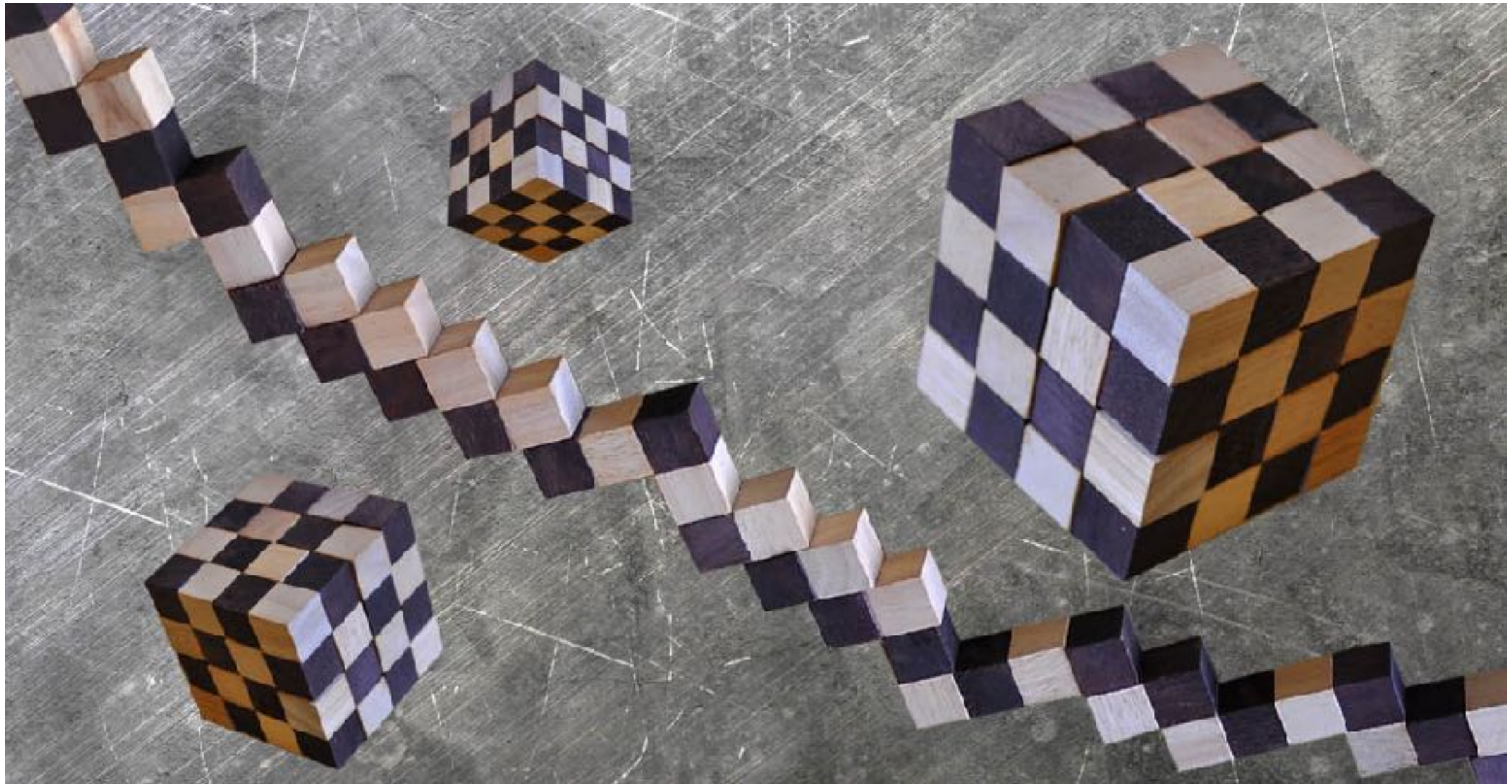
Describing Complete Puzzles

```
standard    :: Puzzle
standard    = Puzzle { valid      = inCube 3,
                       initSoln   = [[(1,1,1)]],
                       initDir    = (0,0,1),
                       sections   = snake }
```


Describing Complete Puzzles

```
meanGreen :: Puzzle
meanGreen = standard { sections = [ 3, 3, 2, 3, 2, 3, 2, 2,
                                     2, 3, 3, 3, 2, 3, 3, 3 ] }
```

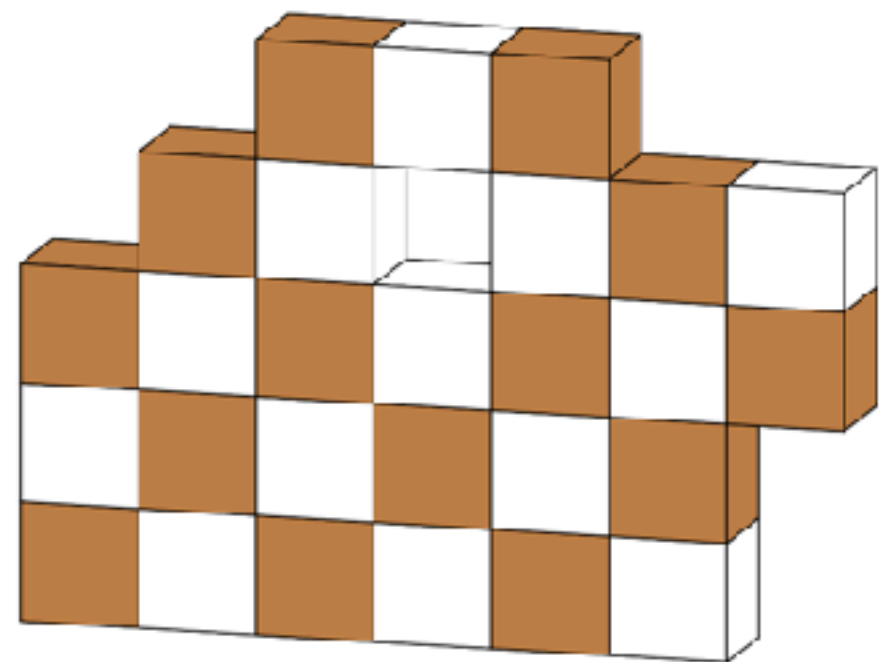
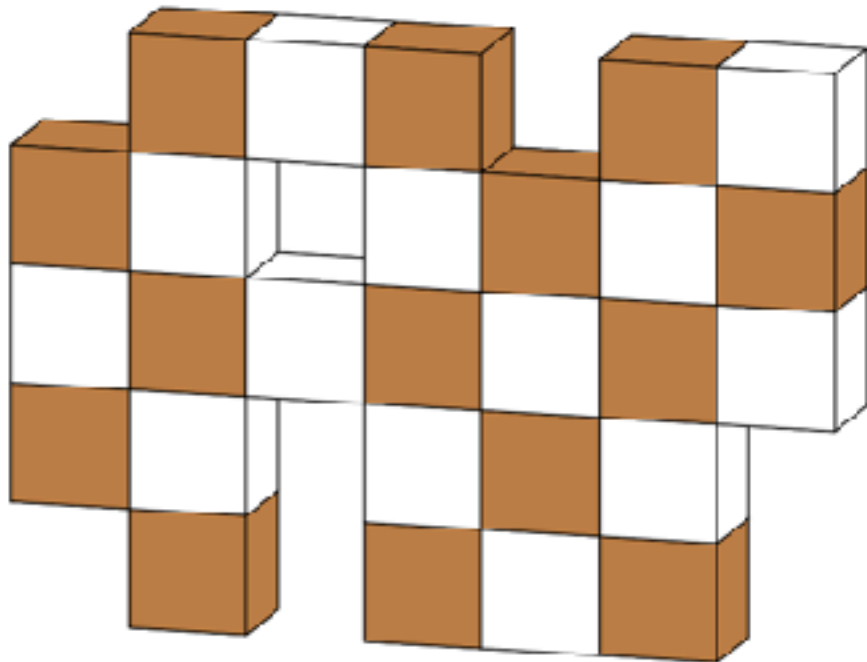
Describing Complete Puzzles



Describing Complete Puzzles

```
king :: Puzzle
king = standard {
  valid      = inCube 4,
  sections = [ 3, 2, 3, 2, 2, 4, 2, 3, 2, 3, 2, 3, 2, 2, 2,
               2, 2, 2, 2, 2, 3, 3, 2, 2, 2, 2, 3, 4, 2,
               2, 2, 4, 2, 3, 2, 2, 2, 2, 2, 2, 2, 2, 2, 4, 2 ] }
```

Flat Configurations



```
flat    :: Puzzle -> Puzzle
flat p = p { valid = ( \ (x,y,z) -> z==1) }
```

Solving Puzzles

```
extend :: Puzzle -> Solution -> Direction -> Int -> [Solution]
extend p soln dir len = [ sect:soln | let start = head (head soln)
                                     sect  = section start dir len,
                                     all (valid p) sect,
                                     all (`notElem` concat soln) sect ]
```

Replace failure by a list of success!

```
all :: Foldable t => (a -> Bool) -> t a -> Bool
```

Source

#

Determines whether all elements of the structure satisfy the predicate.

Solving Puzzles

```
solutions :: Puzzle -> [Solution]
solutions p = solve (initSoln p) (initDir p) (sections p)
  where solve :: Solution -> Direction -> [Int] -> [Solution]
        solve soln dir [] = [soln]
        solve soln dir (len:lens)
          = concat [ solve soln' newdir lens
                    | newdir <- newDirs dir,
                      soln'   <- extend p soln newdir len ]
```

Solving Puzzles

Alternative

```
solve soln dir (len:lens)
  = do newdir <- newDirs dir
      soln'   <- extend p soln newdir len
      soln''  <- solve soln' newdir lens
      return soln''
```

Solving Puzzles

```
head (solutions standard)
= [[(3, 3, 3), (2, 3, 3)], [(1, 3, 3)],
   [(1, 2, 3)], [(2, 2, 3), (2, 2, 2)],
   [(2, 2, 1)], [(2, 1, 1), (2, 1, 2)],
   [(2, 1, 3)], [(1, 1, 3)], [(1, 1, 2), (1, 2, 2)],
   [(1, 3, 2), (2, 3, 2)], [(3, 3, 2)], [(3, 2, 2)],
   [(3, 2, 3)], [(3, 1, 3), (3, 1, 2)],
   [(3, 1, 1), (3, 2, 1)], [(3, 3, 1), (2, 3, 1)],
   [(1, 3, 1), (1, 2, 1)], [(1, 1, 1)]]
```




The Countdown Problem

Graham Hutton. Journal of Functional Programming, 12(6):609-616, Cambridge University Press, November 2002.

Slides adopted from <http://www.cs.nott.ac.uk/~pszgmh/bib.html#countdown>

What Is Countdown?

- A popular quiz programme on British television that has been running for almost 20 years.
- Based upon an original French version called "Des Chiffres et Des Lettres".
- Includes a numbers game that we shall refer to as the countdown problem.

<https://www.youtube.com/watch?v=pfa3MHLLSWI>

Rules

Using the numbers

1 3 7 10 25 50

and the arithmetic operators



construct an expression whose value is

765

Rules

- All the numbers, including intermediate results, must be integers greater than zero.
- Each of the source numbers can be used at most once when constructing the expression.
- We abstract from other rules that are adopted on television for pragmatic reasons.

For our example, one possible solution is

$$(25-10) * (50+1) = 765$$

Notes:

- There are 780 solutions for this example.
- Changing the target number to **831** gives an example that has no solutions.

Evaluating Expressions

Operators:

```
data Op = Add | Sub | Mul | Div
```

Apply an operator:

```
apply :: Op → Int → Int → Int
apply Add x y = x + y
apply Sub x y = x - y
apply Mul x y = x * y
apply Div x y = x `div` y
```

Decide if the result of applying an operator to two integers greater than zero is another such:

```
valid      :: Op → Int → Int → Bool
valid Add  _ _    = True
valid Sub  x y    = x > y
valid Mul  _ _    = True
valid Div  x y    = x `mod` y == 0
```

Expressions:

```
data Expr = Val Int | App Op Expr Expr
```

Return the overall value of an expression, provided that it is an integer greater than zero:

```
eval :: Expr → [Int]
eval (Val n)    = [n | n > 0]
eval (App o l r) = [apply o x y | x ← eval l
                                   , y ← eval r
                                   , valid o x y]
```

Either succeeds and returns a singleton list, or fails and returns the empty list.

Specifying The Problem

Return a list of all possible ways of selecting zero or more elements from a list:

```
subbags :: [a] → [[a]]
```

For example:

```
> subbags [1,2]
```

```
[[],[1],[2],[1,2],[2,1]]
```

Return a list of all the values in an expression:

```
values          :: Expr → [Int]
values (Val n)   = [n]
values (App _ l r) = values l ++ values r
```

Decide if an expression is a solution for a given list of source numbers and a target number:

```
solution        :: Expr → [Int] → Int → Bool
solution e ns n = elem (values e) (subbags ns)
                  && eval e == [n]
```

Brute Force Implementation

Return a list of all possible ways of splitting a list into two non-empty parts:

```
nesplit :: [a] → [([a],[a])]

```

For example:

```
> nesplit [1,2,3,4]

```

```
[[([1],[2,3,4]),([1,2],[3,4]),([1,2,3],[4])]
```

Return a list of all possible expressions whose values are precisely a given list of numbers:

```
exprs      :: [Int] → [Expr]
exprs []    = []
exprs [n]   = [Val n]
exprs ns    = [e | (ls,rs) ← nesplit ns
                  , l      ← exprs ls
                  , r      ← exprs rs
                  , e      ← combine l r]
```

The key function in this lecture.

Combine two expressions using each operator:

```
combine      :: Expr → Expr → [Expr]
combine l r =
    [App o l r | o ← [Add, Sub, Mul, Div]]
```

Return a list of all possible expressions that solve an instance of the countdown problem:

```
solutions    :: [Int] → Int → [Expr]
solutions ns n = [e | ns' ← subbags ns
                      , e   ← exprs ns'
                      , eval e == [n]]
```

Correctness Theorem

Our implementation returns all the expressions that satisfy our specification of the problem:

`elem e (solutions ns n)`



`solution e ns n`

How Fast Is It?

1GHz Pentium-III laptop

GHC version 5.00.2

solutions [1,3,7,10,25,50] 765

0.89 seconds

113.74 seconds

Can We Do Better?

- Many of the expressions that are considered will typically be invalid - fail to evaluate.
- For our example, only around 5 million of the 33 million possible expressions are valid.
- Combining generation with evaluation would allow earlier rejection of invalid expressions.

Applying Program Fusion

Valid expressions and their values:

```
type Result = (Expr, Int)
```

Specification of a function that fuses together the generation and evaluation of expressions:

```
results    :: [Int] → [Result]
results ns = [(e,n) | e ← exprs ns
                    , n ← eval e]
```

We can now calculate an implementation:

```
results [] = []
results [n] = [(Val n,n) | n > 0]
results ns =
    [res | (ls,rs) ← nesplit ns
    , lx ← results ls
    , ry ← results rs
    , res ← combine' lx ry]
```

where

```
combine' :: Result → Result → [Result]
```

Combine'

```
ops :: [Op]
ops = [Add, Sub, Mul, Div]

combine' :: Result -> Result -> [Result]
combine' (l,x) (r,y) = [(App o l r, apply o x y)
                        | o <- ops, valid o x y]
```

Return a list of all possible expressions that solve an instance of the countdown problem:

```
solutions'      :: [Int] → Int → [Expr]
solutions' ns n =
    [ e | ns' ← subbags ns
        , (e,m) ← results ns'
        , m == n ]
```

Correctness Theorem:

solutions'

=

solutions

How Fast Is It Now?

Example:

```
solutions' [1,3,7,10,25,50] 765
```

One solution: 0.08 seconds

All solutions: 5.76 seconds

Can We Do Better?

- Many expressions will be essentially the same using simple arithmetic properties, such as:

$$x * y = y * x$$

$$x * 1 = x$$

- Exploiting such properties would considerably reduce the search and solution spaces.

Exploiting Properties

```
valid      :: Op → Int → Int → Bool
valid Add  x y  = True       $x \leq y$ 
valid Sub  x y  =  $x > y$      $x \leq y$ 
valid Mul  x y  = True       $x \leq y \ \&\& \ x \neq 1 \ \&\& \ y \neq 1$ 
valid Div  x y  =  $x \text{ `mod` } y == 0$ 
               $x \leq y \ \&\& \ x \neq 1 \ \&\& \ y \neq 1$ 
```

Using this new predicate gives a new version of our specification of the countdown problem.

Notes:

- The new specification is sound with respect to our original specification.
- It is also complete up to equivalence of expressions under the exploited properties.

Using the new predicate also gives a new version of our implementation, written as solutions".

Notes:

- The new implementation requires no new proof of correctness, because none of our previous proofs depend upon the definition of valid.
- Hence our previous correctness results still hold under changes to this predicate.

How Fast Is It Now?

Example:

solutions" [1,3,7,10,25,50] 765

Valid:

250,000 expressions

Around 20
times less.

Solutions: 49 expressions

Around 16
times less.

One solution: 0.04 seconds

Around 2
times faster.

All solutions: 0.86 seconds

Around 7
times faster.

More generally, our program usually produces a solution to problems from the television show in an instant, and all solutions in under a second.

Most important points

Program Fusion

Efficiency

Replace failure by a list of success !

Elegancy

Is Carol Now Redundant?

Vorderman said of the computer program:

"It's a nice idea, but thankfully the bosses of Countdown don't want to do away with me and would prefer to keep it human against human!"

How the computer program would handle a makeover show, or dressing up for an awards ceremony remains to be seen.

(London Evening Standard, 1/11/2001)