7.4-1.

Show that in the recurrence

$$T\left(n\right) = \max_{0 \leq q \leq n-1} \left(T\left(q\right) + T\left(n-q-1\right)\right) + \Theta\left(n\right)$$

$$T(n) = \Omega(n^2)$$
.

Proof.

We begin the argument with a guess that $T(n) \ge c n^2$ for some constant c. Substituting this guess into recurrence (7.1) above, we have

$$\begin{array}{ll} T\left(n \right) & \geq & \mathop {\max }\limits_{0 \le q \le n - 1} \left({c\,{q^2} + c\left({n - q - 1} \right)^2} \right) + \Theta \left(n \right) \\ & = & c \cdot \mathop {\max }\limits_{0 \le q \le n - 1} \left({{q^2} + \left({n - q - 1} \right)^2} \right) + \Theta \left(n \right) \end{array}$$

The expression $q^2 + (n-q-1)^2$ achieves a maximum over that parameter's range $0 \le q \le n-1$ at either endpoint. This indicates that $\max_{0 \le q \le n-1} (q^2 + (n-q-1)^2) \le (n-1)^2 = n-2$ n+1. Going back to the bounding of T(n), we obtain

$$\begin{array}{ll} T\left(n\right) & \geq & c\,n^2 - c\left(2\,n - 1\right) + \Theta\left(n\right) \\ & \geq & c\,n^2 \end{array}$$

since we can pick the constant large enough so that the c(2n-1) term dominates the $\Theta(n)$ term. Therefore, $T(n) = \Omega(n^2)$.

^{*.} Creative Commons © 1000 2014, Lawrence X. Amlord (颜世敏, aka 颜序). Email address: informlarry@gmail.com