```
using the multiple instance of each revource type Q1. Available at time t: [0,0,1] deadlack detection algorithm:
                                               deadlock detection
          1 P3 is not deadlocked at the moment. Pretend to run
         completion Available becomes 1,2,1
         Then P4 can be sortisfied. Pretend to run completion
          Available becomes. 2,2,2
          Then P5. Available becomes: [3,3,5]
         Then P1 Available becomes: [4,3,5]
         Then P2 Available becomes 5,3,5
       Sequence < Ps, P4, P5, P1, P2 > will result Finish[i] = +rue
       For all 7 > NO DEADLOCK
  Q2. Using the Banker's Algorithm
    Need = Max Demand - Alloc Available initally: 1, 0,2
    ABC
   4 12 PH. pretend to run completion Available = [2,2,3]
   0 1 0 Then P2 Available = [3, 3, 4]
P3 1 4 2
P4 1 0 2 Then P5. Available = [5,3,5]
P5 1 3 0 Then P1 Available = [6,3,8]
              P3 can never run. So, state is not safe.
      Page 8120 = 212 = 4096 bytes = 1/2B
Q3.
     > (10,10,12) > 4 bytes (each page table entry)

page affect

2nd level page table index

184 level page table index
          210 = 1024 entries per each page table
          SK 1024KB
SMB Virtual men. = 2048 pages > 2048
          AKB
       Memory required for 1st level PT 1024 x 4byte = 4KB
          " 2nd " 2x 1024 x 4 = 8192 = 8kB
```

122B

0x A 30D = 1010 0011 0000 1101 value = 3 Physical address = 0x03300 0 × 5 0 23 = 0101 0000 00 5 -> value = 17 0×801f = 1000 1010 0001 1110 8 - value = 95 PA= 0×5FA1F L, youlde = DNE! INVALID ADDRESS