


# RECURSION<sup>1</sup>

## Agenda :

- Recursion (Basics)
- steps
- problems (fact, fibo, power, sum)

## Why recursion ?

- Decision Trees 
- Dynamic programming
- Merge sort / Quick sort
- Graphs
- Backtracking

## Disclaimer :

- It might take some time to digest this topic.
- A lot of dry runs are required.

★ A function is calling inside itself.

### Observations :

- i) Smaller dolls (sub problem)
- ii) Similar dolls (same problem)
- iii) Tinniest doll (end case stop / base condition / stop cond<sup>n</sup>)

### Recursion :

- i) Solving problem using smaller instance of same problem.
- ii)  $f(n) \rightarrow f(n-1) \rightarrow f(n-2) \dots$  base cond<sup>n</sup>

Problem : Find sum of first  $N$ -natural numbers

→ Iteration  
→ Recursion

$$\Rightarrow \text{sum}(5) \Rightarrow 1 + 2 + 3 + 4 + 5$$

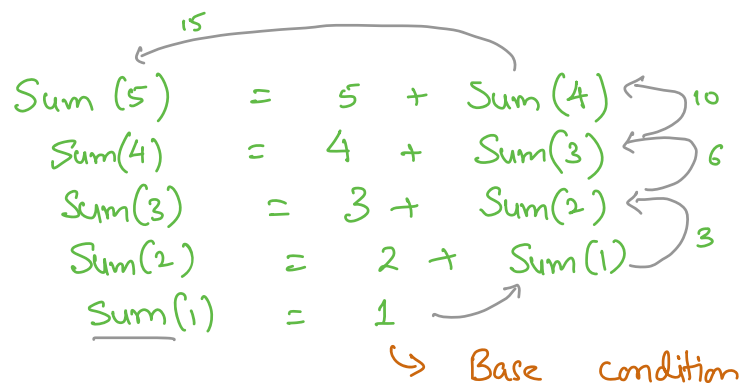
$\underbrace{\hspace{10em}}_{\text{sum}(4) + 5}$

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★

## Steps of recursion:

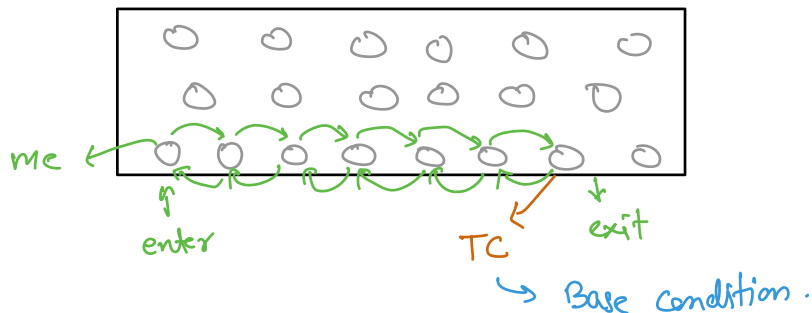
- i) Assumption : Decide what you want from function.
- ii) Main logic : Solving assumption using smaller instance.



- iii) Base condition : Once this base condition is hit we need not go further into recursive calls.

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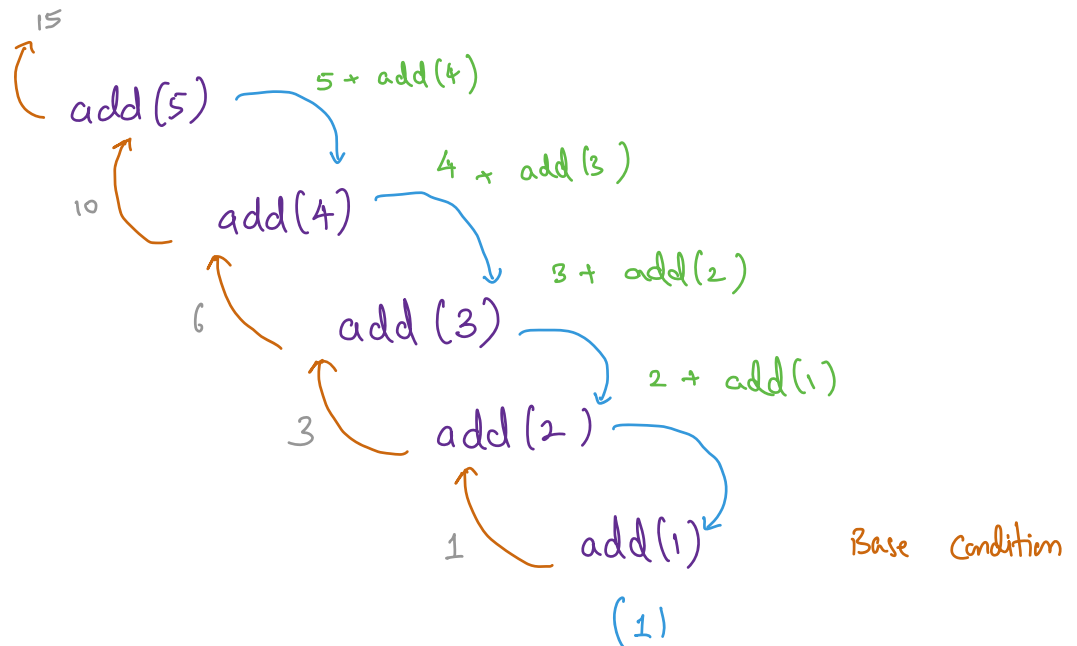
## DTC:



★ Sum of first  $N$ - natural numbers :

Ans  $\text{add}(n) \Rightarrow n + \text{add}(n-1)$

Recursive tree :



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$$\text{add}(5) = n + \text{add}(n-1)$$

↳ Recurrence relation

#

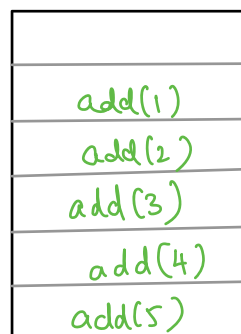
def  $\text{add}(5)$  :

# Base condition

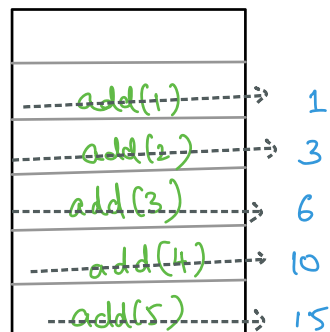
```
if n == 1:
    return 1
return n + add(n-1)
```

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★ Memory stack : →



stack building



stack falling

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## Factorial

$$\# \quad \text{fact}(0) = 1$$

$$\# \quad \text{fact}(1) = 1$$

$$\# \quad \text{fact}(3) = 3 \times 2 \times 1 = 6$$

$$\# \quad \text{fact}(5) = 5 \times 4 \times 3 \times 2 \times 1 = 120$$

$$\# \quad \text{fact}(5) = 5 \times \text{fact}(4)$$

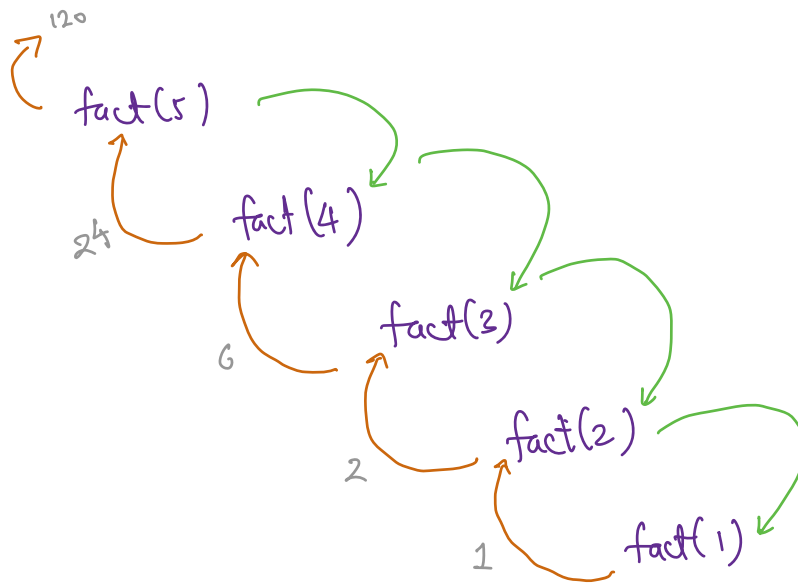
$$\# \quad \text{fact}(n) = n \times \text{fact}(n-1)$$

↳ Recurrence relation

$$\begin{aligned} \# \quad \text{fact}(5) &= 5 \times \text{fact}(4) \\ \text{fact}(4) &= 4 \times \text{fact}(3) \\ \text{fact}(3) &= 3 \times \text{fact}(2) \\ \text{fact}(2) &= 2 \times \text{fact}(1) \\ \text{fact}(1) &= 1 \end{aligned}$$

→ Base condition

## ★ Recursive Tree



# HW : Draw memory stack

```
def factorial(n):  
    # Base condition  
    if n == 1 or n == 0 :  
        return 1  
    return n x factorial(n-1)
```

## ★ fibonacci Series :

N	=	0	1	2	3	4	5	6	7
fib(N)	=	0	1	1	2	3	5	8	13

Detailed description: The table lists the first 8 Fibonacci numbers. Below the table, dashed arrows show the recursive calculation: an arrow from fib(1) and fib(0) points to fib(2), and an arrow from fib(5) and fib(4) points to fib(6).

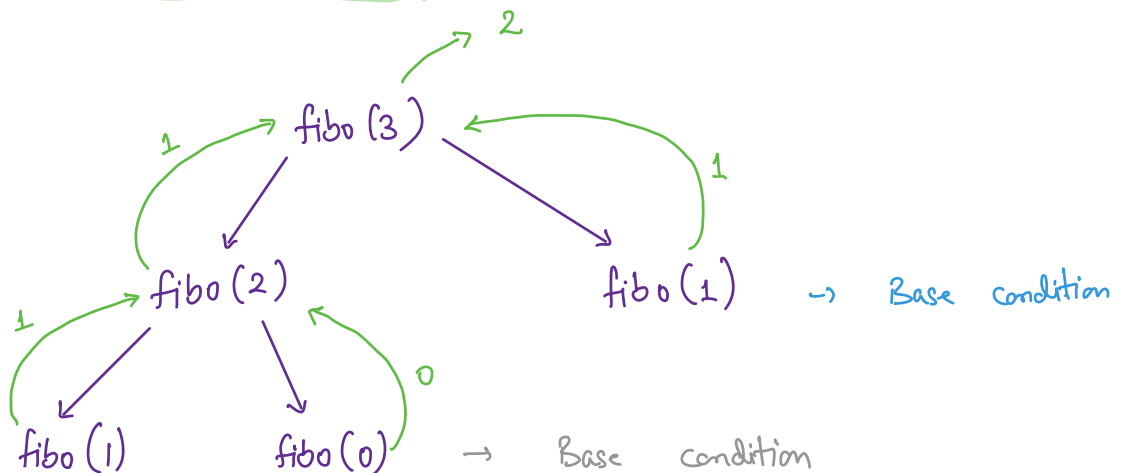
## ★ formula :

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$$

## ★ Base Condition :

if  $n == 0$  or  $n == 1$  :  
return  $n$

## ★ Recursive Tree :





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## Power

# power (a, n)

# power (2, 3)

$$\text{power}(2, 3) = 2 * \text{power}(2, 2)$$

$$\text{power}(2, 2) = 2 * \text{power}(2, 1)$$

$$\text{power}(2, 1) = 2 * \text{power}(2, 0)$$

$$\text{power}(2, 0) = 1$$

↳ Base condition

★ Recurrence relation :

$$\text{power}(a, n) = a * \text{power}(a, n-1)$$

★ power (2, 3)

$$\text{power}(2, 3) = 2 \times 2^2 \rightarrow \text{power}(2, 2)$$

$$\text{power}(2, 2) = 2 \times 2^1 \rightarrow \text{power}(2, 1)$$

$$\text{power}(2, 1) = 2 \times 2^0 \rightarrow \text{power}(2, 0)$$

↳ 1

Ans

$$\text{power}(2, 1) = 2$$

$$\text{power}(2, 2) = 2 \times 2 = 4$$

$$\text{power}(2, 3) = 2 \times 4 = 8$$