### **CD Lab Assignment 4**

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#### Question 1:

NOTE 1: The implementation is for non cyclic and no epsilon production grammar

NOTE 2: In the input each non terminal should occur only once in the LHS, write pipe separated productions in such cases

Note 3: Non terminals are from A to Z only, remaining all considered as terminals

Note 4: The terminals of the grammar should not start with capital letters

#### The output of the program for an input is as shown below:

```
yasaswin@LAPTOP-N4V07KP4:/mnt/d/Compiler_Design_Assignment_1/20CS01040_Nalanaga_Yasaswin_A4$ ./q1
NOTE 1: The implementation is for non cyclic and no epsilon production grammar
NOTE 2: In the input each non terminal should occur only once in the LHS, write pipe separated productions in such cases
NOTE 3: Non terminals are from A to Z only, remaining all considered as terminals
NOTE 4: The terminals of the grammar should not start with capital letters
Enter the no.of lines in the grammar
enter a non cyclic, no epsilon production grammar
F -> T*F | F
F -> (E) | id
The processed grammar
E E+T T
F T*F F
F (E) id
**********The Output************
The grammar after removing left recursion is as shown below :
F -> T*FF'
F -> (E) | id
E' -> +ΤΕ' | ε
```

## Question 2:

NOTE 1: General instructions, non terminals are capital letters with single character

NOTE 2: General instructions, terminals are small letters with single character

The output of the program for an input is as shown below:

```
yasaswin@LAPTOP-N4V07KP4:/mnt/d/Compiler_Design_Assignment_1/20CS01040_Nalanaga_Yasaswin_A4$ ./q2

NOTE 1: General instructions, non terminals are capital letters with single character

NOTE 2: General instructions, terminals are small letters with single character

Enter the no.of lines in the grammar

1

Enter grammar productions

S -> aSSbS | aSaSb | abb | b

*************The output********************

The required left factored grammar is:

S -> aB | b

A -> SbS | aSb

B -> SA | bb
```

# Question 2:

I have built a recursive descendant parser for the given grammar, depending on the non terminal I encounter, I will check few conditions and then I call that function corresponding to that non terminal and proceed till the give string is exhausted and we can't move further in sentential form also and then I declared two vector of vector of strings that keep track of these steps

```
yasaswin@LAPTOP-N4V07KP4:/mnt/d/Compiler_Design_Assignment_1/20CS01040_Nalanaga_Yasaswin_A4$ ./q3
Enter the string
*********The output is********
The string is accepted
STEPS OF ACCEPTANCE ARE:
sentential form - Production Applied - Remaining Input
       ------ no production ------ nn+$
------ S -> nB ------ nn+$
                                                                          nn+$

      nnBAB$
      ------

      nnAB$
      ------

      nn+B$
      ------

      nn+$
      ------

      B -> ε
      ------

      +$
      ------

      B -> ε
      ------

                                                                          n+$
OVERALL STEPS ARE AS FOLLOWS:
ind - sentential form - Production Applied - Remaining Input
1 : S$ ----- no production ----- 2 : nB$ ----- S \rightarrow nB ----- nn+$
3 : nnBAB$ ------ B -> nBAB
4 : nnnBABAB$ ----- B -> nBAB
                                                                                     n+$
                                                                                      +$
5 : Backtracked, Back to line : 3
6 : nnAB$ ------ B -> € ------
7 : nn+B$ ------ A -> + ------
                                                                         +$
7 : nn+B$ ------ A -> + ------ 8 : nn+nBAB$ ------ B -> nBAB -----
                                                                          +$
9 : Backtracked, Back to line : 7
10 : nn+$
```