

(Please write your Exam Roll No.)

# END TERM EXAMINATION

FIFTH SEMESTER [B.TECH] DECEMBER 2019

Paper Code: IT 301

Time : 3 Hours

Subject: Theory of computation

Maximum Marks :75

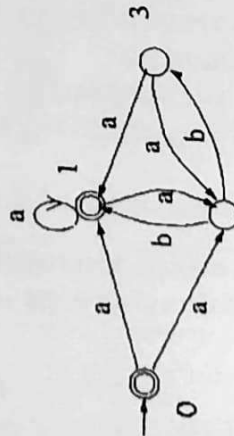
Note: Attempt five questions including Q. NO. 1 which is compulsory. Select one question from each unit.

Q1. Attempt any five of the following:

- Prove that  $L = \{a^n b^n \mid n \geq 1\}$  is non-regular. (5\*5=25)
- Show that the context free languages are not closed under intersection
- What is probabilistic Turing machine?
- Prove that a problem solvable in the space of  $O(f(n))$  requires worst case time of the order of  $O(2^{f(n)})$  [Make necessary assumptions].
- Prove that vertex cover problem is poly-time reducible to clique problem.
- What is parsing? Define LL(1) parsing technique.
- Define any two variants of standard Turing machine.

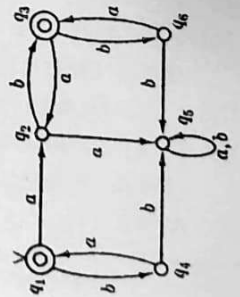
## Unit-I

- Q2. Differentiate in between deterministic and non-deterministic finite automata. Convert following NDFA into DFA. (12.5)



2

- Q3. Explain the process of minimizing number of states of a DFA. Minimize the number of states of following DFA. (12.5)



[2]

## Unit II

- Q4. What is context free grammar? Explain pumping lemma for context free language through an example. (12.5)
- Q5. Define Pushdown automata (PDA). Create a pushdown automaton that accepts the language  $\{0^n 1^n \mid n > 0\}$ . Show that your PDA accepts 000011 and that it rejects 00001. (12.5)

## Unit III

- Q6. Can you write a program which outputs itself? if 'yes' then give an example. Define Recursion theorem and show that construction of 'SELF' Turing machine is possible. (12.5)
- Q7. Differentiate in between computationally intractable and Undecidable problems. Prove that Halting problem is undecidable. (12.5)

## Unit IV

- Q8. Define IP and BPP complexity classes? Prove that  $NSPACE(f(n)) = SPACE(f(n)^2)$ . (12.5)
- Q9. Discuss the proof outline of Cook-Levin theorem. State whether following statements are TRUE or FALSE with justifications. (12.5)
- Some problems in NP complete can not be transformed into satisfiability problem in Polynomial time.
  - Non deterministic RAM may give different results for the same decision problem.
  - A problem with exponentially possible solutions can only be in P if  $P=NP$ .
  - Every problem whose solution requires exponential time on the deterministic RAM can be made to run in polynomial time on non-deterministic RAM.

\*\*\*\*\*