Palindrome Partition - QOTD 22 Jan 23

Leetcode Link : Click

Approach-1 (backtrack & recursion)

Time : $O(n^n)$

where n is length of input string

Space : $O(n^n)$

approach :-

```
/* ✓ ★ Appraoch - 1 (recursive & backtracking)
        explanation :-
               -> // Fun.3 - to check if the string we are inserting into 'ans' vector
is palindrome or not
                -> // Fun.2 : this function will fill out result 2d vector with all
the possible /partitioning combinations that can be palindrome
                    step1 : base case- if the string is completely consumed, then push
the ans into result
                    step 2 : run a loop for complete length of string (including i =
length )
                        step 2.1 : if the sub part doesnt exsist then break the loop
                        step 2.2 : fetch the part of string and then check if it is
palindrome or not
                        step 2.3 : if part is palindrome then push it into the ans ,
and recursively call for index = index + i, while returning pop the last element of
the ans
                -> // MAIN Function
                        step 1 : create a 'ans' vector<string> and a 'result' 2d
string vector, index = 0
                        step 2 : call function.2 palindromePart(s, result, ans, index)
        \nabla T : O(n^n)
         S - O(n^n) [worst case]
        ✓ solved at leetcode :
*/
```

code:-

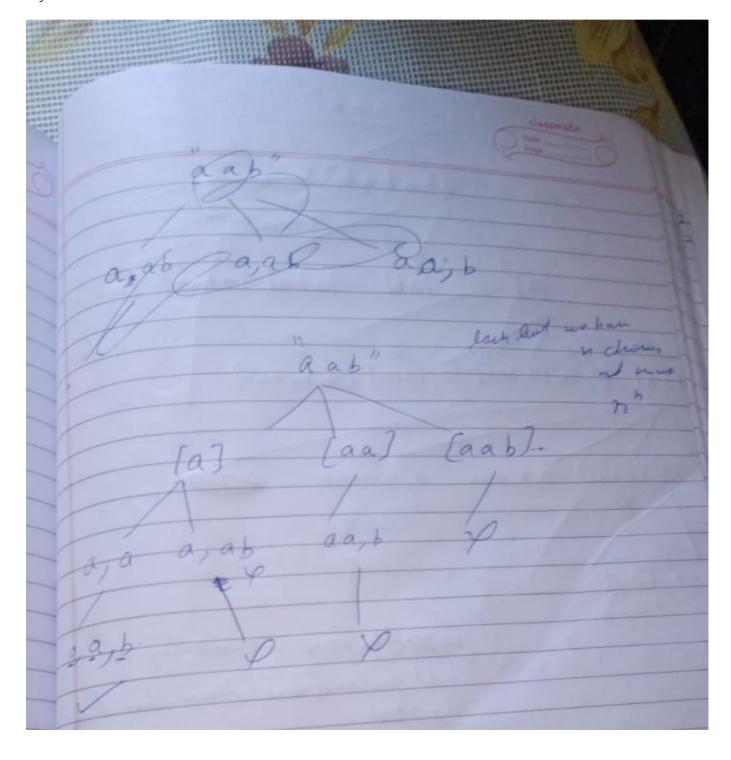
```
private:
    // Fun.3 - to check if the string we are inserting into 'ans' vector is palindrome
    bool isPalindrome(string s){
        int i = 0, j = s.length() -1;
        while(i < j){</pre>
            if(s[i] != s[j]) return false;
            j--;
        return true;
    // Fun.2 : this function will fill out result 2d vector with all the possible
partitioning combinations that can be palindrome
   void palindromePart(string &s, vector<vector<string>> &result, vector<string>
&ans, int index){
        // step1 : base case- if the string is completely consumed, then push the ans
into result
        if(index == s.length()){
            result.push back(ans);
            return;
        }
        // step 2 : run a loop for complete length of string (including i = length )
        for(int i = 1; i <= s.length(); i++){</pre>
            // step 2.1 : if the sub part doesnt exsist then break the loop
            if(i + index > s.length()) break;
            // step 2.2 : fetch the part of string and then check if it is palindrome
or not
            string part = s.substr(index,i);
            // step 2.3 : if part is palindrome then push it into the ans , and
recursively call for index = index + i, while returning pop the last element of the
ans
            if(isPalindrome(part)){
                ans.push back(part);
                palindromePart(s, result, ans, index + i);
                ans.pop_back();
        }
public: // MAIN Function
   vector<vector<string>> partition(string s) {
        // step 1 : create a 'ans' vector<string> and a 'result' 2d string vector,
```

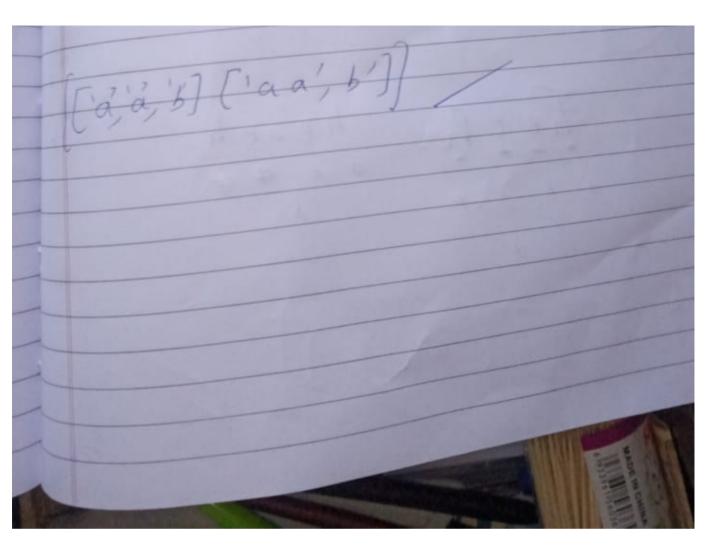
```
index = 0
    int index = 0;
    vector<string> ans;
    vector<vector<string>> result;

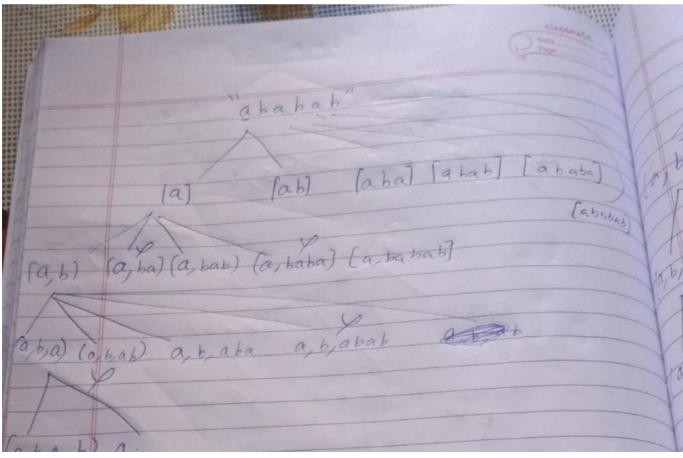
// step 2 : call function.2 palindromePart(s, result, ans, index)
    palindromePart(s, result, ans, index);

    return result;
}
```

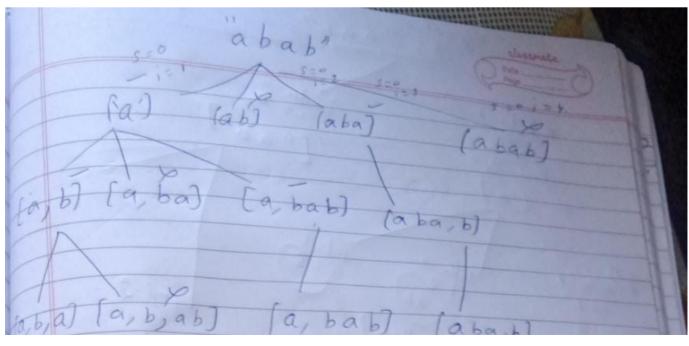
dry run :







not palindrome then don't



ta, b, a, b) (a, bab) what the approach 11 base case (index == length of str) invert ans into result I loop from i = of to lugh of sto. part = 28 Arstr (Andrest, 2) part is palendren then invest it else consume,

