

**Aim:** Implementation of Job sequencing with Deadlines

**Theory:**

In job sequencing problem, the objective is to find a sequence of jobs, which is completed within their deadlines and gives maximum profit.

Solution

Let us consider, a set of  $n$  given jobs which are associated with deadlines and profit is earned, if a job is completed by its deadline. These jobs need to be ordered in such a way that there is maximum profit.

It may happen that all of the given jobs may not be completed within their deadlines.

Assume, deadline of  $i^{\text{th}}$  job  $J_i$  is  $d_i$  and the profit received from this job is  $p_i$ . Hence, the optimal solution of this algorithm is a feasible solution with maximum profit.

Thus,  $D(i) > 0$  for  $1 \leq i \leq n$ .

Initially, these jobs are ordered according to profit, i.e.  $p_1 \geq p_2 \geq p_3 \geq \dots \geq p_n$

**Complexity Analysis :**

In this algorithm, we are using two loops, one is within another. Hence, the complexity of this algorithm is  $O(n^2)$

**Code:**

```
#include <stdio.h>

#include <stdlib.h>

struct Job
{
    int jobid;
    float profit;
    int deadline;
};

int comparator(const void *a, const void *b)
{
    struct Job *j = (struct Job *)a;
    struct Job *j1 = (struct Job *)b;
    return (j1->profit - j->profit);
}
```

```
}  
  
int minValue(int x, int y)  
{  
    if (x < y)  
        return x;  
    return y;  
}  
  
int main()  
{  
    int i, j, n, dmax = 0, counter = 0;  
    float total_profit = 0;  
    printf("Enter number of processes:");  
    scanf("%d", &n);  
    struct Job *arr = (struct Job *)malloc(n * sizeof(struct Job));  
    int alloc[n];  
    for (i = 0; i < n; i++)  
    {  
        printf("Enter profit and deadline for job %d :", i + 1);  
        scanf("%f %d", &arr[i].profit, &arr[i].deadline);  
        arr[i].jobid = i + 1;  
        alloc[i] = -1;  
        if (arr[i].deadline > dmax)  
        {  
            dmax = arr[i].deadline;  
        }  
    }  
    qsort(arr, n, sizeof(struct Job), comparator);  
    for (i = 1; i <= n; i++)  
    {  
        for (j = minValue(dmax, arr[i - 1].deadline); j >= 1; j--)  
        {
```

```
    if (alloc[j] == -1)
    {
        alloc[j] = i - 1;
        counter++;
        total_profit += arr[i - 1].profit;
        break;
    }
}
if (counter == dmax)
{
    break;
}
printf("Sequence:");
for (i = 1; i <= dmax; i++)
{
    printf("Job %d ", arr[alloc[i]].jobid);
}
printf("\nTotal profit= %f", total_profit);
}
```

## Problem

Consider the following 5 jobs and their associated deadline and profit.

index	1	2	3	4	5
JOB	j1	j2	j3	j4	j5
DEADLINE	2	1	3	2	1
PROFIT	60	100	20	40	20

## Sort the jobs according to their profit in descending order

Note! If two or more jobs are having the same profit then sort them as per their entry in the job list.

index	1	2	3	4	5
JOB	j2	j1	j4	j3	j5
DEADLINE	1	2	2	3	1
PROFIT	100	60	40	20	20

## Find the maximum deadline value

Looking at the jobs we can say the max deadline value is 3.  
So,  $d_{\max} = 3$

From this set of jobs, first we select  $J_2$ , as it can be completed within its deadline and contributes maximum profit.

- Next,  $J_1$  is selected as it gives more profit compared to  $J_4$ .
- In the next clock,  $J_4$  cannot be selected as its deadline is over, hence  $J_3$  is selected as it executes within its deadline.
- The job  $J_5$  is discarded as it cannot be executed within its deadline.

Thus, the solution is the sequence of jobs ( $J_2, J_1, J_3$ ), which are being executed within their deadline and gives maximum profit.

Total profit of this sequence is  $100 + 60 + 20 = 180$ .

## Output:

```
Enter number of processes:5
Enter profit and deadline for job 1 :60 2
Enter profit and deadline for job 2 :100 1
Enter profit and deadline for job 3 :20 3
Enter profit and deadline for job 4 :40 2
Enter profit and deadline for job 5 :20 1
Sequence:Job 2 Job 1 Job 3
Total profit= 180.000000
```