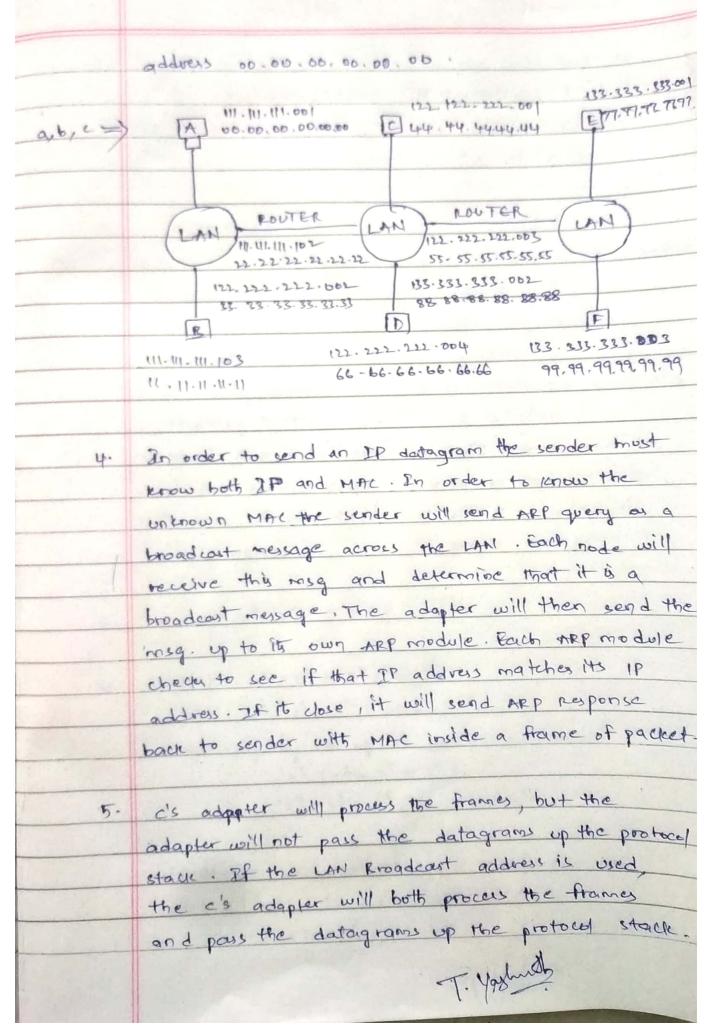
1.	Ethernet à connectionless. Upper layer protocols like
	TCP establishes connection - Connection is implemented
	through software and connection oriented ethernet
	is also there but by default ethernet is connectionless
2.	248 MAC address
	232 IPV4 address.
	2128 IPUG address
3,	
(4)	is forwarding table in A determines that the datagram
	should be routed to interface 11.111.111.002
	(ii) The adapter in A creates and ethernet packed with
	etternet destination address 22-22-22-22
	in The first router receives the packet and extracts
	the datagram. The datagram is routed top 22.22.20
	(by The first nover sends the ethernet payeet with
	the dest add 55-55.55.55-55-55.55 and sie add of
	33.33.33-33.33 via hterface with IP 122.222.2202
	(1) Process continue until packet reaches host F.
	The first of the second
(e)	ARP in A must now determine the LAN address of
	111.111.111.002. Host A sends out an ARP guery packed
	within a broadcast ethernet frame. The first
	router receives the Query packet and sends host
	A and ARP response packed. This ARP Packet is
	corried by externet frame and externet destination
	T. Yaghtot
	170



6.	-) An IP datagram and from the caree hast to the destination host will travel over 8 interferces.
	-) 3 forwarding tables will be indexed to move
	datagram from source to the dufination.
٦,	CSMA/CD CSMA/CA
	1. CSMA/CD is effective after a 1. CSMA/CD is effective before
	collision. collision.
	2. It is used in wired 2. Commonly used in
	networks wireless.
	3. It is only reduced to 3. Minimizes the possibility
	recovery time. of collision.
	4. Resend the data frame 4. Will first transmit the
	whenever conflict occur. intent to send for data.
	5. Is more officient than 5. Will it's similar to
	simple CSMA. Simple CSMA.
0.	(0 to 2 1) X RTT
8,	
10-	-> P(x) = 93 + x2 + x6 (1101)3
10	G(x) = x6+ x3+ x2 (1001/00)
	Multiply by no. of bits in CRC polynomial
	23 (46+ 43 + x2)
	x9+x6+ x5 (1001100000
	We divide and determine the remainder, the
	result is bol. so the transmitted in
	[00 [100]00]
	T. Joshus

11. — suppose two nodes start to transmit at the same time a passet of tenyth L over a broaded channel of talk R.  The propagation delay blue to two nodes as deprop.  — 21 transfer noded at sametime, then receive the bits of passet from another node.  Collision is occurred.		
came time a packet of length L over a breaded channel of tate R.  The propagation delay blue the two modes as deprop.  2) transfer noded at sametime, then receive the bits of packet from another node.  Collision is owned.		
came time a pawet of tength L over a product channel of take R.  The propagation delay blue the two modes as deprop.  2) transfer noded at sametime, then receive the bits of packet from another node.  Collision is owned.		
channel of talk R.  The propagation delay blue the two nodes as deprop.  If transfer noded at sametime, then receive the bits of packet from another node.  Collidor is owned.	n-	
The propagation delay blue the two nodes as deprop.  — 2) transfer noded at sametime, then receive the bits of packet from another node.  .: Collicion is owned.		
Aprop  — 21 transfer noded at sametime, then receive the bits of packet from another node.  .: Collicion is owned.		
The bits of packet from another node.  Collicion is owned.		
the bits of packet from another node.  Collicion is owned.		
Collition is orwined.		
- und A		
- und A		
- Wald		
- und D		
- Waled		
- Waled		
1,70		T. Aost al