

Subject: \_\_\_\_\_

Date: 1 / 1

## Heap Sorting

① heap sorting algorithm is a comparison based sorting algorithm that uses binary heap (binary tree) to sort a sequence of numbers. It follows 3 steps: ① building a binary heap of  $n$  elements in a where the min (max) element is located at the top of heap.  $O(n \log n)$ .

② Call 'Extract and remove' for extracting the biggest element remaining in the heap and repeat it until heap is empty. Saving these elements in array, it will take  $O(n)$ .

Algorithm step by step:

Heapify: this is a helper function used to maintain the heap property for a subtree rooted at index  $i$ .

② heap sort: the main heap sort function that uses the build-max heap and repeatedly extracts the max/min.

It swaps the last element with the last element and then calls heapify to rebuild the heap. So the time complexity for this step is  $O(n \log n)$ .

③ build max heap: it is a function iterates through the elements from the elements from  $n/2$  to 0 and for each element it calls heapify. Since heapify takes  $O(\log n)$  the overall time complexity for building the heap is  $O(n)$ .

• the overall complexity for the heap sorting process is  $O(n \log n)$ .

Subject: \_\_\_\_\_

Date: 1

19/02/20

### Kruskal's algorithm

Kruskal's algorithm is a greedy algorithm used to find the minimum Spanning Tree (MST) of a connected, weighted graph. It works by sorting the edges of the graph and adding them to the MST one by one, ensuring no cycles are formed.

Sort the edges of the graph in non-decreasing order of their weight and initialize the MST as an empty set.

For each edge in the sorted edge list, check if adding the edge to the MST forms a cycle using a union-find data structure. If no cycle is formed, add the edge to the MST. Repeat until the MST contains  $V-1$  edges, where  $V$  is the number of vertices in the graph.

Sorting the edges: Sorting all edges takes  $O(E \log E)$  where  $E$  is the number of edges.

Union-Find operations: each find, set or union operation runs in almost constant time,  $O(\alpha(V))$  where  $\alpha$  is inverse Ackermann function which grows very slowly.

Overall complexity: We process each edge the overall

complexity is dominated where  $O(E \log E + E \alpha(V))$   
 $= O(E \log E)$ .

• It works for connected, undirected, weighted graphs.