

Inductive logic programming

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- From previously:
 - 1 Equations
 - 2 Substitution
 - 3 \perp_i
 - 4 Modes
 - 5 Algorithms
- Algorithms:
 - 1 Construct C from \perp_i
 - 2 Cover set algorithm
- Difficulties:
 - 1 \perp_i
 - 2 Progol

Previously: equations

$$B \wedge H \models E \quad (1)$$

$$B \wedge \overline{E} \models \overline{H} \quad (2)$$

$$B \wedge \overline{E} \models \overline{\perp} \quad (3)$$

$$\overline{\perp} \models \overline{H} \quad (4)$$

$$H \models \perp \quad (5)$$

$$B \wedge \overline{e} \wedge \perp \vdash_h \square \quad (6)$$

- B : background knowledge
- H : hypothesis
- E : examples
- $\overline{\perp}$: set of all true literals (wrt. $B \wedge \overline{E}$)
- \perp : most specific clause

Previously: substitution

Substitution

Let θ be a substitution of the form $\{v_1/t_1, \dots, v_n/t_n\}$.

Let F be an arbitrary atom.

$F\theta$ is the atom F where each of its variable v_i have been replaced by t_i .

Ex:

- $F : \text{parent}(X, Y)$
- $\theta : \{X/\text{jean}, Y/\text{bob}\}$
- $F\theta : \text{parent}(\text{jean}, \text{bob})$

Clause substitution

The same can be done for a clause C :

- $F : \text{parent}(X, Y) : \neg \text{father}(X, Y).$
- $\theta : \{X/\text{jean}, Y/\text{bob}\}$
- $F\theta : \text{parent}(\text{jean}, \text{bob}) : \neg \text{father}(\text{jean}, \text{bob}).$

Previously: \perp_i

Depth $d(v)$

$$d(v) = \begin{cases} 0, & \text{if } v \text{ is in the head of } C \\ (\min_{u \in U_v} d(u)) + 1, & \text{otherwise} \end{cases}$$

where U_v are the variables in atoms in the body of C containing v .

Ex:

- $C: p(A) :- p(A,B), g(B,C), f(C,D)$
- $d(A) = 1, d(B) = 2, d(C) = 3, d(D) = 3$
- \perp can have an infinite cardinality
- \perp_i is more restrained
 - the distance of its variables is $\leq i$

Previously: modes

Horn clause

$C : A \leftarrow B_1, \dots, B_n$

- A is the head of clause C
- B_1, \dots, B_n is the body of clause C

Mode declaration

- $\text{modeh}(n, \text{atom}) \mid \text{modeb}(n, \text{atom})$
- $\text{modeh}(*, f(+\text{int}, -\text{int}))$, $\text{modeb}(*, d(+\text{int}, -\text{int}))$
- Use *modeh* for the head of a clause, *modeb* for its body

Previously: modes (ii)

Instantiation

- Let M be a set of modes
- $m \in M$ is a mode declaration
- $a(m)$ is the atom of m with place markers replaced

Ex

$$\begin{aligned} m &= \text{modeh}(*, f(+\text{int}, -\text{int})) \\ a(m) &= f(A, B) \end{aligned}$$

Previously: algorithms

- Three algorithms:
 - Cover set algorithm
 - Construct \perp_i
 - Construct C form \perp_i

Algorithm 1: Cover set algorithm

input: h, i, B, M, E

```
1 forall  $e \in E$  do
2   Construct  $\perp_i$  for  $e$  using Algorithms 3 and 4
3   Construct state  $s$  from  $\perp_i$  using Algorithm 2
4   Let  $C'$  be the unflattening of  $C(s)$ 
5    $B \leftarrow B \cup C'$ 
6    $E \leftarrow E - \{e : e \in E, B \wedge \bar{e} \vdash_h \emptyset\}$ 
```

Algorithm 2: Construct C

input: h, B, e, \perp_i

```
1 Open  $\leftarrow \{\langle \square, \emptyset, 1 \rangle\}$ , Closed  $\leftarrow \emptyset$ 
2 while True do
3    $s \leftarrow \text{best}(\text{Open})$ 
4   Open  $\leftarrow \text{Open} - \{s\}$ , Closed  $\leftarrow \text{Closed} \cup \{s\}$ 
5   if  $\neg \text{prune}(s)$  then
6     Open  $\leftarrow (\text{Open} \cup \rho(s)) - \text{Closed}$ 
7   if terminated(Closed, Open) then
8     return  $\text{best}(\text{Closed})$ 
9   if Open =  $\emptyset$  then
10    print 'no compression'
11    return  $\langle e, \emptyset, 1 \rangle$ 
```

- ρ is the refinement operator

Algorithms: Constructing \perp_i (i)

Algorithm 3: Construct \perp_i - Part 1

```
1 Get  $m \in M$ , modelh such that  $a(m) \preceq a$  with substitution  $\theta_h$ 
2 if  $\#m$  then
3    $\perp$  return  $\square$ 
4  $a_h \leftarrow a(m)$ 
5 for  $v/t \in \theta_h$  do
6   if  $v$  corresponds to  $\#type$  in  $m$  then
7      $\perp$  Replace  $v$  by  $t$  in  $a_h$ 
8   else
9      $\perp$  Replace  $v$  by  $v_k$  in  $a_h$ , with  $k = \text{hash}(t)$ 
10  if  $v$  corresponds to  $+type$  in  $m$  then
11     $\perp$  Add  $v$  to InTerms
12 Add  $a_h$  to  $\perp_i$ 
```

Algorithms: Constructing \perp_i (ii)

Algorithm 4: Construct \perp_i - Part 2

```
1 for  $k \leftarrow 1, \dots, i$  do
2   forall modeb  $m \in M$  do
3     Let  $\{v_1, \dots, v_n\}$  be the variables corresponding to +type in  $a(m)$ 
4     Let  $T_i$  be the set of all terms of the type associated with  $v_i$  in  $m$ 
5      $T(m) \leftarrow T_1 \times \dots \times T_n$ 
6     forall  $\langle t_1, \dots, t_n \rangle \in T(m)$  do
7        $a_b \leftarrow a(m)$ 
8        $\theta \leftarrow \{v_1/t_1, \dots, v_n/t_n\}$ 
9       if Prolog succeeds on goal  $a_b\theta$  then
10        Let  $\Theta_b$  be the set of answer substitutions
11        forall  $\theta_b \in \Theta_b$  do
12          forall  $v/t \in \theta_b$  do
13            if  $v$  corresponds to #type in  $m$  then
14              Replace  $v$  by  $t$  in  $a_b$ 
15            else
16              Replace  $v$  by  $v_k$  in  $a_b$ , with  $k = \text{hash}(t)$ 
17            if  $v$  corresponds to -type then
18              Add  $v$  to InTerms
19          Add  $\overline{a_b}$  to  $\perp_i$ 
20 return  $\perp_i$ 
```

Important numbers

f = Number of positive examples covered -
Number of negative examples covered -
Number of literals in body of clause -
Optimistic estimate of literals needed
p = Number of positive examples covered
n = Number of negative examples covered
h = Optimistic estimate of literals needed

Example of execution

```
[:- modeb(1,dec(+int,-int))? - Time taken 0.00s]
[:- modeb(1,plus(+int,+int,-int))? - Time taken 0.00s]
[:- modeb(1,inc(+int,-int))? - Time taken 0.00s]
[:- modeb(1,mult(+int,+int,-int))? - Time taken 0.00s]
[:- modeb(1,plus(+int,+int,-int))? - Time taken 0.02s]
[:- modeh(1,plus(+int,+int,-int))? - Time taken 0.00s]
[:- modeh(1,mult(+int,+int,-int))? - Time taken 0.00s]
[:- determination(plus/3,dec/2), determination(plus/3,inc/2), determination(plus/3,plus/3)? - Time taken 0.02s]
[:- determination(mult/3,dec/2), determination(mult/3,plus/3), determination(mult/3,mult/3)? - Time taken 0.00s]
[:- commutative(mult/3), commutative(plus/3)? - Time taken 0.00s]
[:- set(c,3)? - Time taken 0.00s]
[Testing for contradictions]
[No contradictions found]
[Generalising plus(4,X,Y) :- inc(X,U), inc(U,V), inc(V,W), inc(W,Y).]
[Most specific clause is]

plus(A,B,C) :- dec(A,D), inc(B,E), dec(D,F), plus(D,B,G), plus(D,
    E,C), inc(D,A), inc(E,H), dec(F,I), plus(F,B,H), plus(F,
    E,G), plus(F,F,A), plus(F,H,C), inc(F,D), inc(G,C), inc(H,
    G).

[C:-2,6,5,3 plus(A,B,C).]
[C:-3,4,4,2 plus(A,B,C) :- dec(A,D).]
[C:-3,4,4,1 plus(A,B,C) :- dec(A,D), inc(B,E).]
[C:1,4,0,0 plus(A,B,C) :- dec(A,D), inc(B,E), plus(D,E,C).]
[C:-3,5,5,2 plus(A,B,C) :- inc(B,D).]
[5 explored search nodes]
[f=1,p=4,n=0,h=0]
[Result of search is]

plus(A,B,C) :- dec(A,D), inc(B,E), plus(D,E,C).
```

Difficulties