



# PROJECT 1

## ● THE FLOW SHOP SCHEDULING PROBLEM

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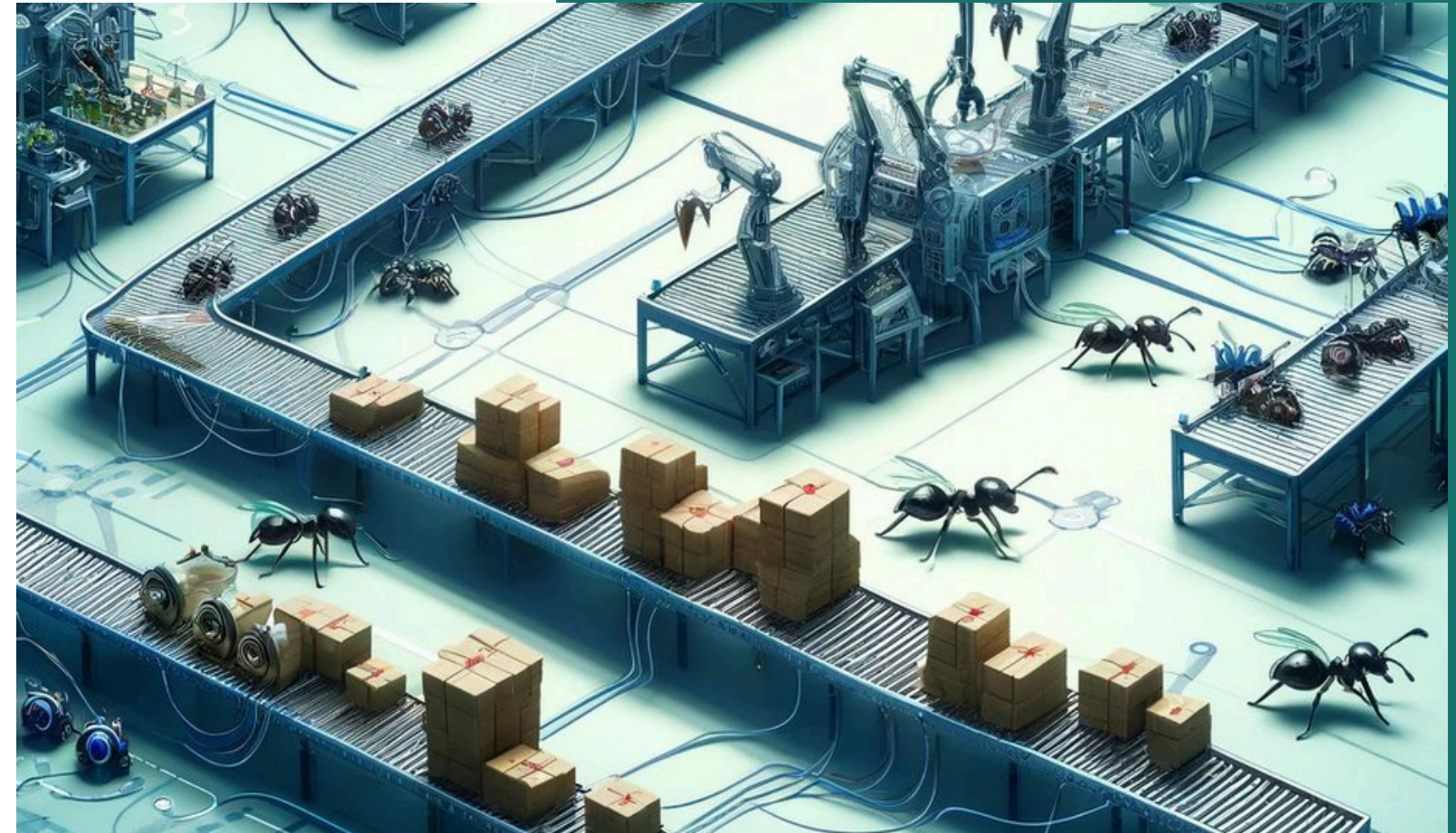
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# TABLE OF CONTENT

SOLVING THE SHOP FLOW SCHEDULING PROBLEM USING ANT COLONY OPTIMIZATION ALGORITHM



**01** Problem introduction

**02** Mathematical Formulation

**03** Complexity

**04** Gantt Diagram

**05** Solution using heuristic

**06** Result analyses



# MATHEMATICAL FORMULATION

## Formulation of The Problem :

### Terminology :

- J = (1,...,N) set representing All jobs, j denotes job number.
- M = (1,...,M) set representing all machines, I denotes machine number.
- P is a matrix M\*N representing Set of Processing time such that  $p_{ij}$  denotes processing time for j in machine I .
- $r_i$  denotes the minimum idle time for each machine
- $d_i$  denotes the maximum idle time for each machine

### Decision Variable :

- S = matrix (M,J), such that  $s_{ij}$  denotes starting time of job j on machine i.
- C = matrix(M,J), such that  $c_{ij}$  denotes finishing time of job j on machine I
- $x_{kj}$  boolean variable presents if j is the successor of job k.

## Objective Function (Minimizing Makespan) :

Min  $z = \max(C_i)$  sd ( We minimize the maximum of all jobs time completions.)

Such That :

for a job k, he has one successor and one predecessor.

$\sum x_{ik} = 1$  ( job k has one predecessor only) for all I in {1,...,J}

$\sum x_{ki} = 1$  (job k has one successor only ) for all I in {1,...,J}

$$y \geq \sum (c_{ij} - s_{ij}) = \sum p_{ij}$$

$$C_{ij} = S_{ij} + p_{ij}$$

Constraint 1 :

$S_{i,k} \geq C_{i,j}$  ( processing of operations can not be interrupted) (job k is successor of job j)

$$x_{kj} = 1$$

Constraint 2 : (idle times are respected)

$d_i \geq S_{i,k} - C_{i,j} \geq r_i$  , I in {1,2,...,N} and j in {2,...,M}.( job k is successor of job j)

$$x_{kj} = 1$$

Constraint 3 : ( technological order of processes is respected)

$S_{i,j} \geq C_{i-1,j}$  ; I in {2,...,N} and j in {2,...,M}.( machines follow same order)

# Complexity of the problem

## NP-HARD!

The Flow Shop Scheduling Problem (FSSP) is known to be NP-hard. This can be explained by the following points:

### 1. **Permutation Flow Shop Scheduling Problem (PFSSP):**

- When the sequence of jobs is fixed across all machines (permutation assumption), the problem remains NP-hard. This is because even with a fixed order, the problem of determining the optimal sequence is combinatorial in nature and grows factorially with the number of jobs.
- For  $n$  jobs, the number of possible permutations is  $n!$ .

### 2. **Non-Permutation Flow Shop Scheduling Problem (NPFSSP):**

- In the non-permutation case, where each machine can process jobs in a different order, the problem is even more complex. The number of possible sequences grows exponentially.
- For each machine, there are  $n!$  possible sequences, leading to a total of  $(n!)^m$  possible schedules for  $m$  machines.

it is considered as an NP-HARD also due to the factorial growth of the problem size.

# GANTT DIAGRAM





# ANT COLONY OPTIMIZATION FOR FLOW SHOP SCHEDULING

## MAIN FUNCTIONS USED IN THE CODE

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### 01 Objective Function:

**Purpose:** Minimize makespan (completion time of the last job on the last machine).

### 02 Ant Colony Optimization (ACO):

**Process:**

- Iteration: Generate and evaluate job schedules.
- Pheromone Update: Adjust pheromone trails based on solution quality.
- Best Solution: Track and improve the best schedule.

### 03 Update Pheromones:

- Evaporation: Reduces pheromones to prevent stagnation.
- Deposit: Increases pheromones for better solutions to guide future ants.

### 04 Gantt Chart Visualization:

**Purpose:** Visualize job schedules and machine usage.

# RESULTS ANALYSIS

## Best solution found:

[[2 3 4 0 1]

[2 3 4 0 1]

[2 3 0 4 1]]

Makespan: 30.0

The best solution found by the Ant Colony Optimization (ACO) algorithm has yielded the following job sequences for each machine:

**Machine 1 processes jobs in the order [2, 3, 4, 0, 1].**

**Machine 2 processes jobs in the same order [2, 3, 4, 0, 1].**

**Machine 3 processes jobs in the order [2, 3, 0, 4, 1].**

The makespan, or the total completion time for all jobs across all machines, is 30.0.

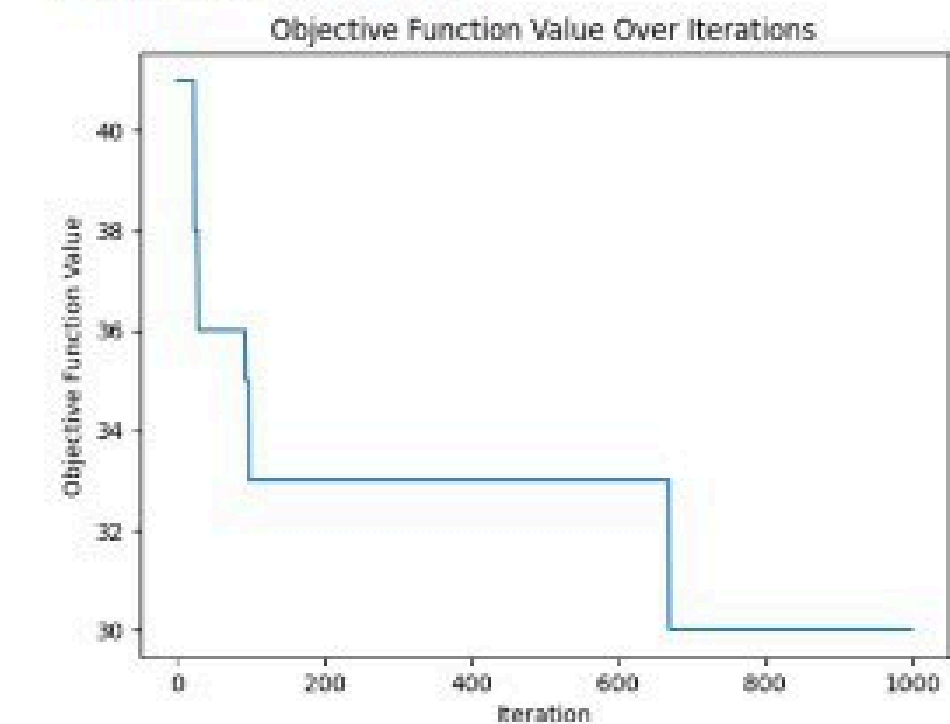
## Variability of Results:

Each run of the ACO algorithm produces different results, including variations in the makespan. This variability is expected due to the stochastic nature of the ACO algorithm, which relies on random initial solutions and probabilistic decisions during the search process.

## Recommendations for Consistency and Improvement:

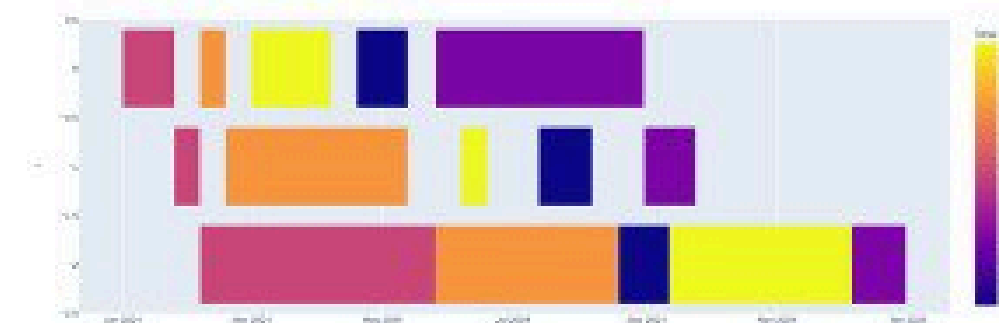
- multiple runs
- parameter tuning
- visualization

Best solution found:  
[[2 3 4 0 1]  
[2 3 4 0 1]  
[2 3 0 4 1]]  
Objective function value: 30.0



[[2 3 4 0 1]  
[2 3 4 0 1]  
[2 3 0 4 1]]  
Makespan: 30.0

Best solution found:  
[[2 3 4 0 1]  
[2 3 4 0 1]  
[2 3 0 4 1]]  
Makespan: 30.0





# THANK YOU

● FOR YOUR NICE ATTENTION

