

CSCI 335 Theory of Computing



Lecture 14 – Turing Machines IV: Church-Turing Thesis

Learning Outcomes

By the end of this section you will be able to:

- Understand Church-Turing thesis
- Get acquainted with an undecidable problem among Hilbert's problems
- Get familiar with string encoding of different objects for TM simulations

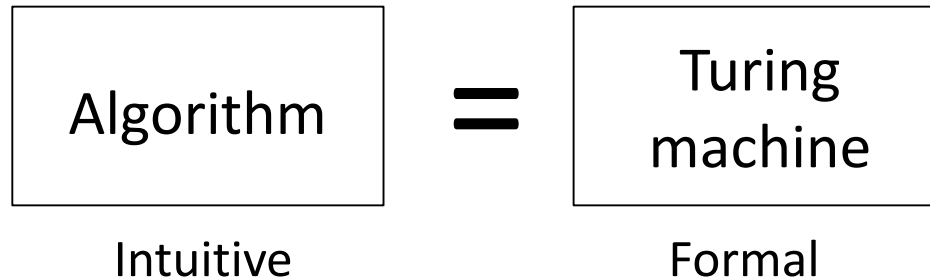
- Church-Turing thesis
- Hilbert's problems
- String representation of objects

Church–Turing thesis

Church–Turing thesis



Alonzo Church
1903–1995



Instead of Turing machines,
can use any other “reasonable” model
of unrestricted computation:
 λ -calculus, random access machine,
your favorite programming language, ...



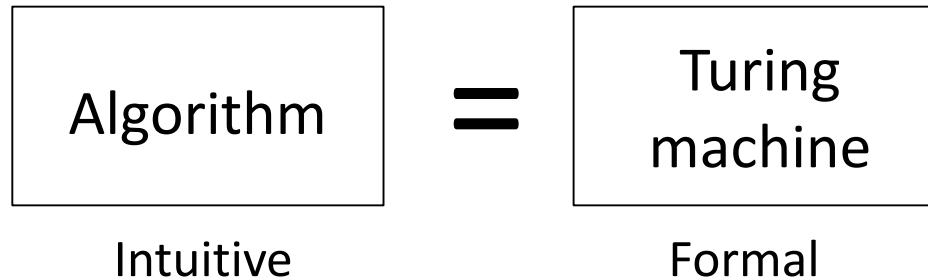
Alan Turing
1912–1954

Big impact on mathematics, computer science, and artificial intelligence.

Church–Turing thesis



Alonzo Church
1903–1995



- This thesis states that **all possible models of computation, if they are sufficiently broad, must be equivalent.**
- It also implies that **there is an inherent limitation in this and that there are functions that cannot be expressed in any way that gives an explicit method for their computation.**
- A general principle for algorithmic computation and, while not provable, gives strong evidence that no more powerful models can be found.



Alan Turing
1912–1954

Hilbert's problems

Hilbert's 10th problem

NEW GIZA UNIVERSITY

In 1900 David Hilbert posed 23 problems

#1) Problem of the continuum (Does set A exist where $|\mathbb{N}| < |A| < |\mathbb{R}|$?).

#2) Prove that the axioms of mathematics are consistent.

#10) Give an algorithm for solving *Diophantine equations*.

Diophantine equations:

Equations of polynomials where solutions must be integers.

Example: $3x^2 - 2xy - y^2z = 7$ solution: $x = 1, y = 2, z = -2$

Let $D = \{p \mid \text{polynomial } p(x_1, x_2, \dots, x_k) = 0 \text{ has a solution in integers}\}$

Hilbert's 10th problem: Give an algorithm to decide D .

Matiyasevich proved in 1970: D is not decidable.

Note: D is T-recognizable.



David Hilbert
1862—1943

String representation of objects

String representation of objects

Notation for encoding objects into strings

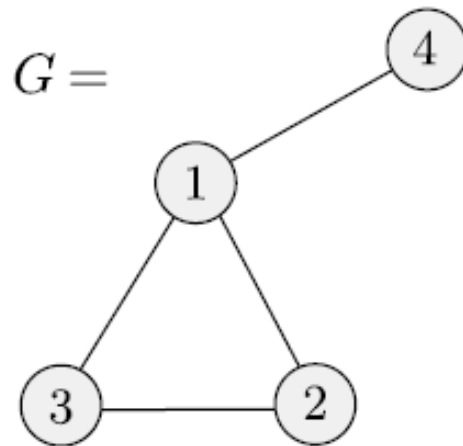
- If O is some object (e.g., polynomial, automaton, graph, etc.), we write $\langle O \rangle$ to be an encoding of that object into a string.
- If O_1, O_2, \dots, O_k is a list of objects then we write $\langle O_1, O_2, \dots, O_k \rangle$ to be an encoding of them together into a single string.

Notation for writing Turing machines

We will use high-level English descriptions of algorithms when we describe TMs, knowing that we could (in principle) convert those descriptions into states, transition function, etc. Our notation for writing a TM M is

$M =$ “On input w
[English description of the algorithm]”

Example: String representation of graphs



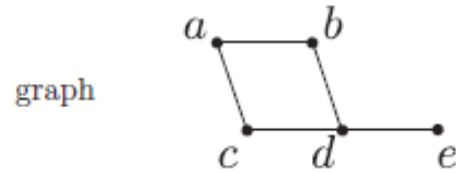
$\langle G \rangle =$

$$(1,2,3,4)((1,2), (2,3), (3,1), (1,4))$$

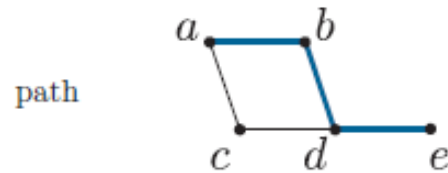
FIGURE 3.24

A graph G and its encoding $\langle G \rangle$

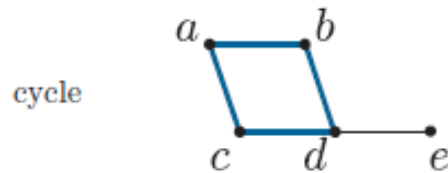
Example: String representation of graphs



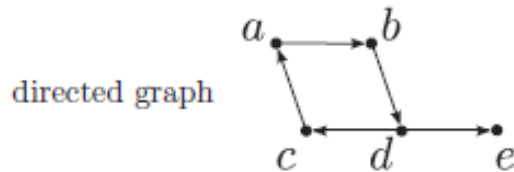
"a,b a,c b,d c,d d,e"



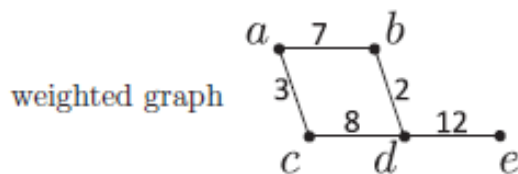
"a,b,d,e"



"a,b,d,c"



"a,b c,a b,d d,c d,e"



"a,b,7 a,c,3 b,d,2
c,d,8 d,e,12"

The right-hand column shows possible ASCII string representations of the corresponding examples of graph concepts.

Page 48, John MacCormick: What Can Be Computed?: A Practical Guide to the Theory of Computation. Princeton University Press, 2018

Summary

- Church-Turing thesis
- Hilbert's problems
- String representation of objects

- Church-Turing thesis: Section 3.3 (Sipser 2013)