Experiment 26 Introduction to the ARM Microprocessor Submission of Experimental Results

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Experimental Results and Comments (Total 90 marks):

Please provide the required screenshots, photos, code, values highlighted in the script according to the following requirements (screenshots should be clear and readable):

1. Code of the Hello World programme (Section 4) after editing (put as text NOT as a screenshot, screenshots of code will receive zero marks) [5 marks]

```
Code:
```

```
#include "mbed.h"
DigitalOut myled(LED_GREEN);
Serial pc(USBTX, USBRX);
int main()
{
    int i = 0;
    pc.printf("Hello World!\n");

    while (true) {
        wait(2); // wait a small period of time
        pc.printf("%d \ n\ r ", i); // print the value of variable i i++; // increment the variable
        myled = !myled; // toggle a led
    }
}
```

Explanation: Programme after last editing.

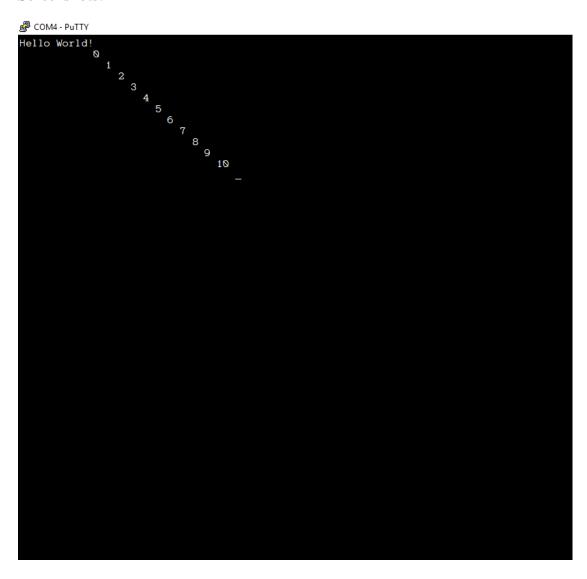
2. Answer to Q1 [5 marks]

Answer: Smallest chosen time: 0.01s (0.02s for a complete cycle), corresponding to 50 Hz.

Explanation: It was determined that the smallest time to allow the lab group to see the LEDs flashing was 0.01s, which means a complete flash and black out cycle takes 0.02s, which corresponds to 50Hz.

3. Screenshot of the result of Section 5 part VIII (the output screen) [5 marks]

Screenshots:



```
## COM4-PUTY
Hello World!
1
2
3
Hello World!
0
1
2
3
4
5
6
6
7
8
9
9
10
11
12
13
14
15
16
17
18
19
20
20
21
22
—
```

Comment/Explanation: In the first screenshot of the programme, the message seems to "walk" across the screen, since the "\n" in C is new line instead of newline and carriage return. The printf was changed to use "\n\r" instead of just "\n" in the later version so as to add a carriage return.

4. The modified code of section 6 (put as text NOT as a screenshot, screenshots of code will receive zero marks) [5 marks]

Code:

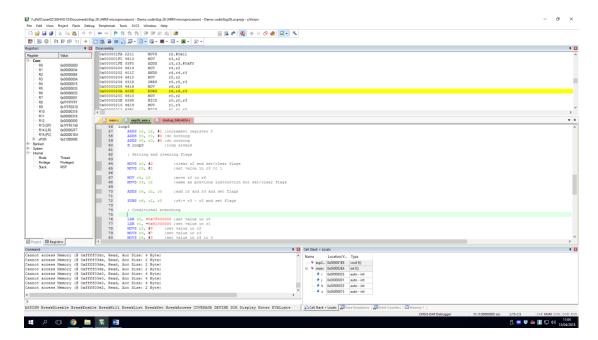
```
DigitalOut gpo(D0);
DigitalOut led(LED_RED);
DigitalOut led2(LED_GREEN);
SLCD slcd;
int main()
{
```

```
while (true) {
    float f=tsi.readPercentage();
    slcd.printf("%1.3f",f);
    if(f>0.5){
        led=true;
        }else{
        led=false;
        }
    if(f<0.5){
        led2=true;
        }else{
        led2=false;
        }
    wait(0.2f);
}</pre>
```

Explanation: As shown above, the code was modified to set the green led illuminating when the read percentage of the LCD display f > 0.5, corresponding to touching the slider anywhere between the midpoint and the right-hand end; and set the red led illuminating when the read percentage of the LCD display f < 0.5, corresponding to touching the slider anywhere between the midpoint and the left-hand end

5. Screenshot of the result of Section 7 part VIII [3 marks]

Screenshots:



Comment/Explanation: Debugging stared. The important windows of the 'Disassembly' and the 'Registers' and all the important values of the variables were displayed in the view correctly as expected.

6. Answer to Q2 [4 marks]

Answer: The bits from the sixth to the eighth bit (start from the left) of the machine code identifies the register used.

Disassembly: 0x000001EA 220E MOVS r2,#0x0E

Mnemonics: MOVS r2, #14 ;move a number into reigister 2

Machine Code 00100 010 00001110

Disassembly: 0x000001EC 2325 MOVS r3,#0x25

Mnemonics: MOVS r3, #37 ;move a 2nd no. into reg. 3

Machine Code 00100 011 00100101

Explanation: $010_2 = 2$ and $011_2 = 3$ justifies the using of corresponding registers, namely, r2 and r3.

7. Answer to Q3 [4 marks]

Answer: The bits from the ninth to the last bit (start from the left) of the machine code gives the number to be moved into the register.

Disassembly: 0x000001EA 220E MOVS r2,#0x0E

Mnemonics: MOVS r2, #14 ;move a number into reigister 2

Machine Code 00100 010 00001110

Disassembly: 0x000001EC 2325 MOVS r3,#0x25

Mnemonics: MOVS r3, #37 ;move a 2nd no. into reg. 3

Machine Code 00100 011 00100101

Explanation: 000011102 = 0x0E and 001001012 = 0x25 justifies that this part of the machine code gives the number, namely, 14 and 37, to be moved into that register.

8. Answer to Q4 [4 marks]

Answer: result -- R4:0x00000015

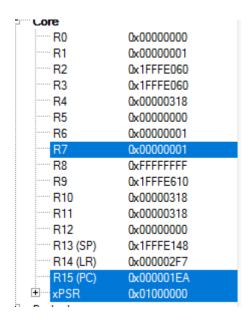
Explanation: The instruction "ADDS r4, r2, #7" put the sum of reg2 and 7 to reg 4, whiche is the addition of 0x0E and 0x07, equalling to 0x15.

9. Answer to Q5 [4 marks]

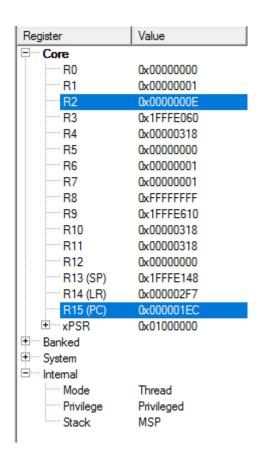
Answer: The order of the registers in the mnemonics is important.

See below the screenshots of the execution results of corresponding instructions:

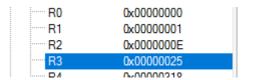
0x000001E8 1847 ADDS r7,r0,r1



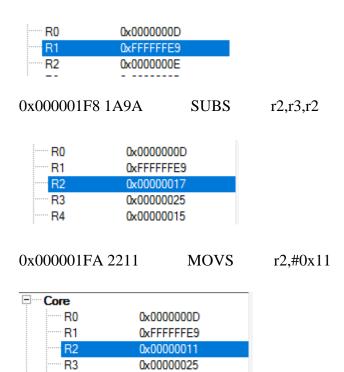
0x000001EA 220E MOVS r2,#0x0E



0x000001EC 2325 MOVS r3,#0x25



0x000001EE 1D	DD4	ADDS	r4,r2,#7
R0	0x00000 0x00000		
R2	0x00000		
R3	0x00000		
R4	0x00000	015	
0x000001F0 18I	D5	ADDS	r5,r2,r3
: K1	0x00000	001	
R2	0x00000	00E	
R3	0x00000		
R4 R5	0x00000		
R6	0x00000		
0x000001F2 189	9E	ADDS	r6,r3,r2
Lore	0.000000		
R0 R1	0x0000000		
R2	0x0000000		
R3	0x0000002		
R4	0x0000001	5	
R5	0x0000003		
R6 R7	0x0000003 0x0000000		
117	uxuuuuuu	,,	
0x000001F4 1E	50	SUBS	r0,r2,#1
			4
Core R0	0x0000	0000	
R1	0x0000		
R2	0x0000		
R3	0x0000	0025	
R4	0x0000		
R5	0x0000 0x0000		
R7	0x0000		
R8	0xFFFF		
R9	0x1FFF	E610	
R10	0x0000		
R11	0x0000		
R12 R13 (SP)	0x0000 0x1FFF		
R14 (LR)	0x0000		
R15 (PC)	0x0000	01F6	
± xPSR	0x2100	0000	
0x000001F6 1A	D1	SUBS	r1,r2,r3



0x00000015

0x00000033

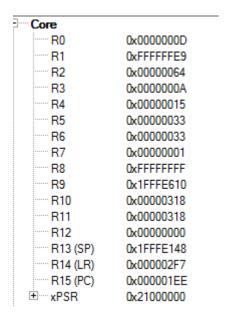
Explanation: As shown above, the order of the registers in mnemonics distinguishes the registers which stores the operands and the registers for storing results, thus important.

10. Answer to Q6 [4 marks]

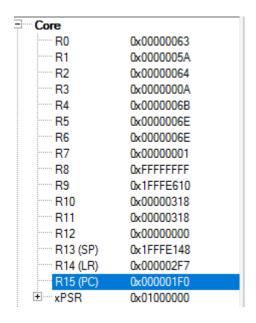
R4

R5

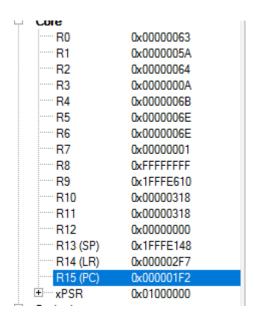
Answer: Values are as expected. Register status before excuation:



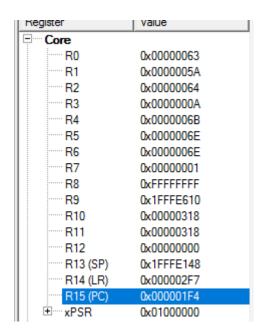
0x000001EE 1DD4 ADDS r4, r2, #7



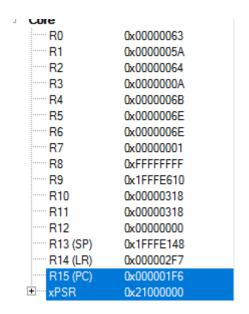
0x000001F0 18D5 ADDS r5, r2, r3



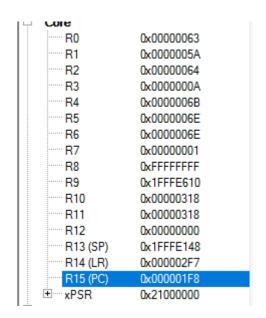
0x000001F2 189E ADDS r6, r3, r2



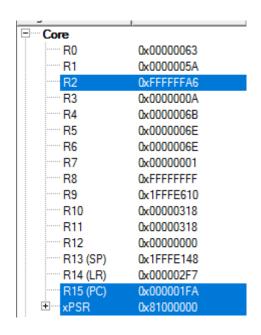
0x000001F4 1E50 SUBS r0, r2, #1



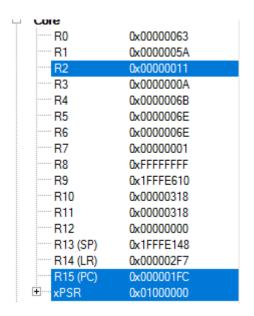
0x000001F6 1AD1 SUBS r1,r2,r3



0x000001F8 1A9A SUBS r2,r3,r2



0x000001FA 2211 MOVS r2,#0x11

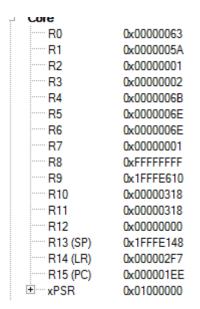


Explanation: The results are shown clearly in the screenshots of the single step execution process.

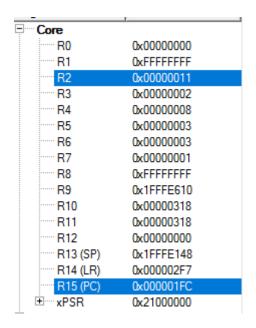
11. Answer to Q7 [4 marks]

Answer: The values are as expected as shown below in the screenshots of the single step results.

Register status before execution:



Result:



Explanation: The results are shown clearly in the screenshots of the single step execution process.

12. Answer to Q8 [4 marks]

Answer: The logic function AND was performed.

0x000001FA 2211	MOVS	r2,#0x11
0x000001FC 4613	MOV	r3,r2
0x000001FE 33F0	ADDS	r3,r3,#0xF0
0x00000200 4614	MOV	r4,r2
0x00000202 401C	ANDS	r4,r4,r3
0x00000204 4615	MOV	r5,r2

Results:

R2	0x00000011
R3	0x00000101
R4	0x00000001

A(R2)	B(R3)	C(R4)
0x00000011	0x00000101	0x00000001

Explanation: C = A AND B, that is R4 = R2 AND R3, that is 0x00000001 = 0x00000011 AND 0x00000101.

13. Answer to Q9 [4 marks]

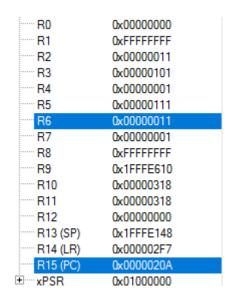
Answer: Yes, besides MOV, AND and SUB, it is also important for functions as BIC, ORR, EOR, etc.

0x00000)204 4615	MOV	r5,r2
R2	0x00000011		
R3	0x00000101		
R4	0x00000001		

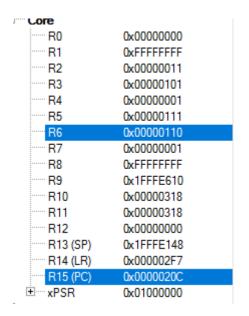
0x00000206 431D	ORRS	r5	.r5	,r3

R0	0x00000000
R1	0xFFFFFFFF
R2	0x00000011
R3	0x00000101
····· R4	0x00000001
R5	0x00000111
R6	0x00000003
R7	0x00000001
R8	0xFFFFFFFF
R9	0x1FFFE610
R10	0x00000318
R11	0x00000318
R12	0x00000000
R13 (SP)	0x1FFFE148
R14 (LR)	0x000002F7
R15 (PC)	0x00000208
±·····xPSR	0x01000000
Papieod	

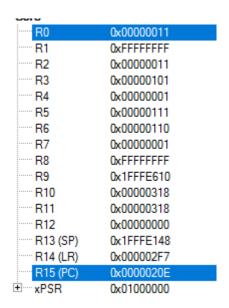
0x00000208 4616 MOV r6,r2



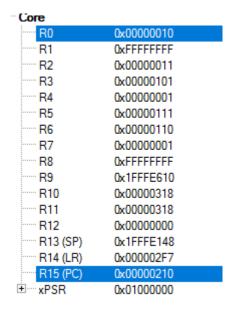
0x0000020A 405E EORS r6,r6,r3



0x0000020C 4610 MOV r0,r2



0x0000020E 4398 BICS r0,r0,r3



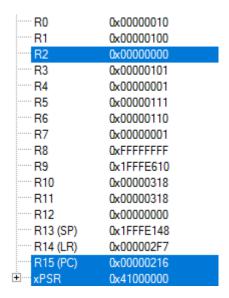
0x00000210 4619 MOV r1,r3

R0	0x00000010
R1	0x00000101
R2	0x00000011
R3	0x00000101
R4	0x00000001
R5	0x00000111
R6	0x00000110
R7	0x00000001
R8	0xFFFFFFFF
R9	0x1FFFE610
R10	0x00000318
R11	0x00000318
R12	0x00000000
R13 (SP)	0x1FFFE148
R14 (LR)	0x000002F7
R15 (PC)	0x00000212
±xPSR	0x01000000
Banked	

0x00000212 4391 BICS r1,r1,r2

R0	0x00000010
R1	0x00000100
R2	0x00000011
R3	0x00000101
····· R4	0x00000001
R5	0x00000111
R6	0x00000110
R7	0x00000001
R8	0xFFFFFFFF
R9	0x1FFFE610
R10	0x00000318
R11	0x00000318
R12	0x00000000
R13 (SP)	0x1FFFE148
R14 (LR)	0x000002F7
R15 (PC)	0x00000214
± ····· xPSR	0x01000000

0x00000214 2200 MOVS r2,#0x00



Explanation:

As shown above, the order of the registers in mnemonics distinguishes which registers to store the operands and which registers for storing results and the order to execute certain order-sensitive functions, thus important.

14. Answer to Q10 [4 marks]

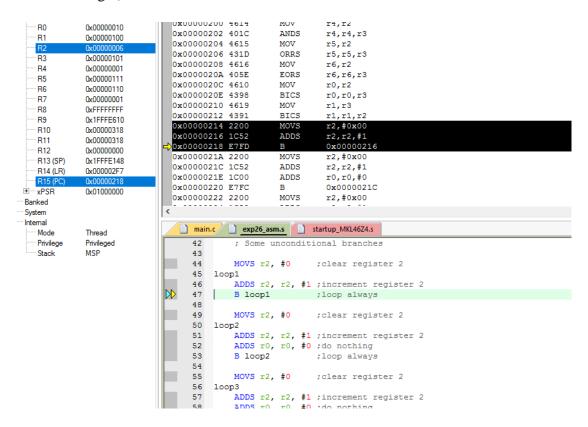
Answer: Branch function was performed, branches to the 0x00000216.

0x00000214 2200	MOVS	r2,#0x00
0x00000216 1C52	ADDS	r2,r2,#1
0x00000218 E7FD	В	0x00000216

Explanation: The operand for B, which stands for unconditional branch, dictates that the address of the next instruction to be executed was at 0x00000216.

15. Answer to Q11 [5 marks]

Answer: The contents of R2 will be 5 after the instruction at address is executed 5 times (in the screenshots the instruction at 0x00000216 was executed one more time thus R2 being 6).



Explanation: Every time after the "ADD 1 to r2" instruction at 0x00000216 has been executed, the instruction at 0x00000218 branches the program counter back to the "ADD 1 to r2" instruction at 0x00000216, thus when the branch instruction has been executed *exactly* 5 times, the "ADD 1 to r2" instruction at 0x00000216 will also be executed *exactly* 5 times, yielding the contents of R2 being 5.

16. Answer to Q12 [4 marks]

Answer: The value held by R2 at 40.19s was approximately twice the value after the initial 20.25s, which means that code block of loop1 has been executed approximately twice times given twice execution time, indicating that the execution of the code block of Loop1is approximately the same.

Value of R2 at time 20.25s: 0x10DE1FD6



Value of R2 at time 40.19s: 0x21809FE6



Disassembly for Loop1

0x00000216 1C52 ADDS r2,r2,#1

0x00000218 E7FD B 0x00000216

Explanation: The execution of the same code block of Loop1 is approximately the same.

17. The graph of Section 11 (using MS Excel or MATLAB) [6 marks]

Screenshots:

Loop2

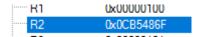
0x0000021A 2200 MOVS r2,#0x00

0x0000021C 1C52 ADDS r2,r2,#1

0x0000021E 1C00 ADDS r0,r0,#0

0x00000220 E7FC B 0x0000021C

Value of R2 at time 20.22s: 0x0CB5486F



Value of R2 at time 39.92s: 0x1919424F



Loop3

0x00000222 2200 MOVS r2,#0x00

0x00000224 1C52 ADDS r2,r2,#1

0x00000226 1C00	ADDS	r0.r0.#0
	ADDO	10.10.#0

0x00000228 1C00 ADDS r0,r0,#0

0x0000022A E7FB B 0x00000224

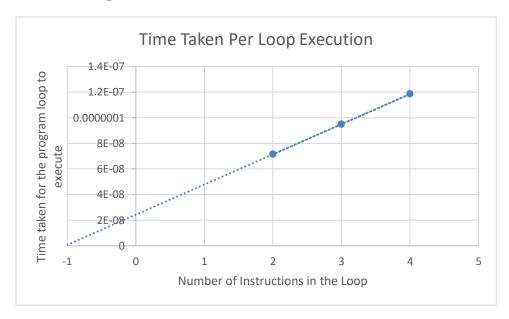
Value of R2 at time 20.22s: 0x0A2AA18C



Value of R2 at time 40.04s: 0x14208A40

R2 0x14208A40

Comment/Explanation:



As can be seen from the backward extrapolation of the plotted graph, every ADD instruction takes 1 clock cycle to execute and the number of instructions in the loop goes to -1 when time goes to 0, indicating the branch instruction takes twice time as the ADD instruction to execute, that is, 2 clock cycles.

18. Answer to **Q13** [4 marks]

Answer: The time taken by one clock cycle is 0.021 μs.

Explanation: It was checked out from the specification that the frequency of the board was 48MHz, which means that one clock cycle takes $0.021 \mu s$.

19. Answer to Q14 [3 marks]

Answer: For the branch instruction, 2 clock cycles are required.

Explanation: As obtained from the specification of the Mbed and verified in the experiment, one branch instruction will take 2 clock cycles.

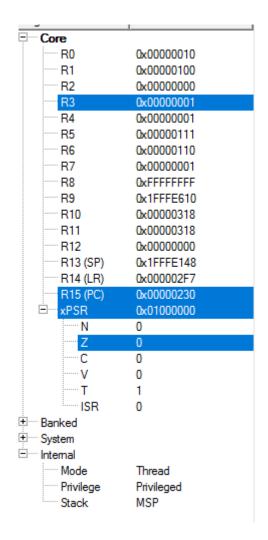
20. Screenshot of the result of Section 12 [3 marks]

Screenshots:

20.1 Results with R2 set as 0x00000000

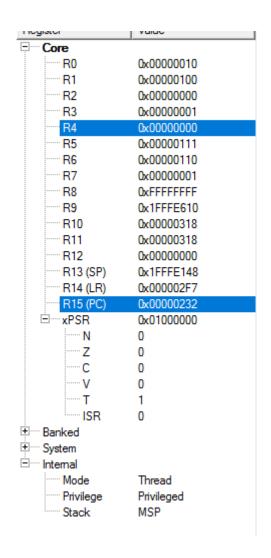
0x0000022C 2200	MOVS	r2,#0x00
0x0000022C 2200	MOVS Core R0 R1 R2 R3 R4 R5 R6 R7 R8 R9 R10 R11 R12 R13 (SI R14 (LF R15 (PC RYSR N Z C V T	0x00000010 0x00000000 0x00000000 0x00000011 0x00000111 0x00000110 0x0000001 0xFFFFFFF 0x1FFFE610 0x00000318 0x00000318 0x000000000 0x1FFFE148 R) 0x000002F7
	ISF	R 0

0x0000022E 2301 MOVS r3,#0x01

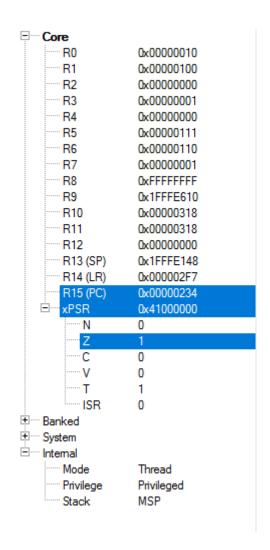


r4,r2

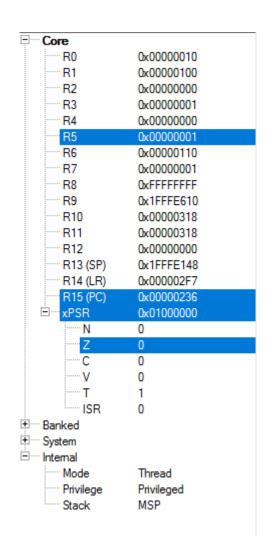
0x00000230 4614 MOV



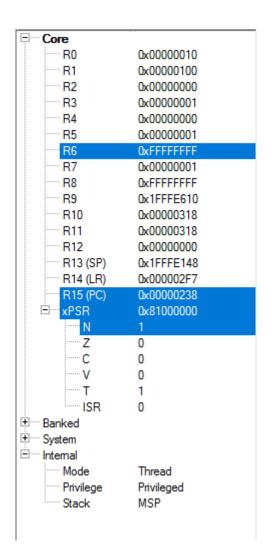
0x00000232 0014 MOVS r4,r2



0x00000234 18D5 ADDS r5,r2,r3

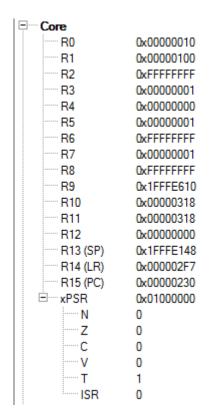


0x00000236 1AD6 SUBS r6,r2,r3



20.2 Results with R2 set to 0xFFFFFFFF

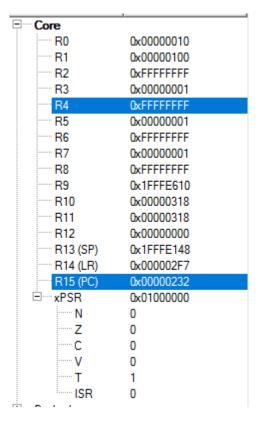
0x0000022E 2301 MOVS r3,#0x01



0x00000230 4614

MOV

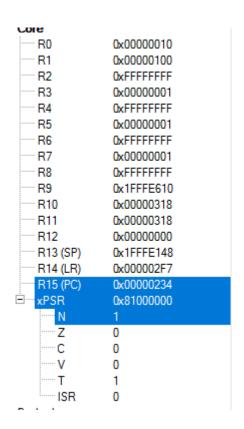
r4,r2



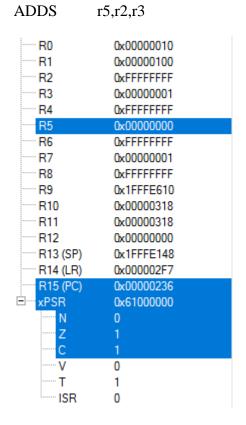
0x00000232 0014

MOVS

r4,r2

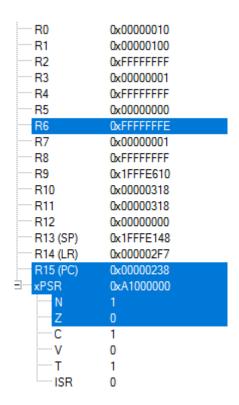


0x00000234 18D5



0x00000236 1AD6

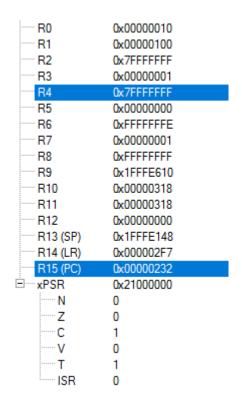
SUBS r6,r2,r3



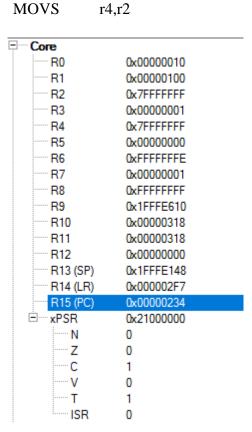
20.3 Results with R3 set to 0x7FFFFFFF

0x0000022E 2301 **MOVS** r3,#0x01 R0 0x00000010 R1 0x00000100 R2 0x7FFFFFFF R3 0x00000001 R4 0xFFFFFFFF R5 0x00000000 R6 0xFFFFFFE R7 0x00000001 R8 0xFFFFFFFF R9 0x1FFFE610 0x00000318 R10 ·R11 0x00000318 R12 0x00000000 R13 (SP) 0x1FFFE148 0x000002F7 R14 (LR) R15 (PC) 0x00000230 xPSR 0x21000000 N Z 0 C 1 ٧ 0 Т 1 ISR 0

0x00000230 4614 MOV r4,r2

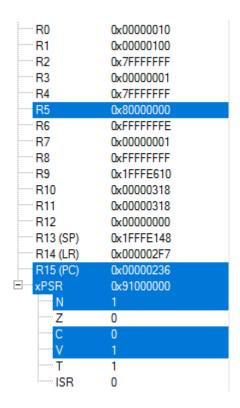


0x00000232 0014 MOVS



0x00000234 18D5 ADDS r5,r2,r3

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r6,r2,r3

SUBS

0x00000236 1AD6

R0	0x00000010
R1	0x00000100
R2	0x7FFFFFFF
R3	0x00000001
····· R4	0x7FFFFFFF
R5	0x80000000
R6	0x7FFFFFFE
R7	0x00000001
R8	0xFFFFFFFF
R9	0x1FFFE610
R10	0x00000318
R11	0x00000318
R12	0x00000000
R13 (SP)	0x1FFFE148
R14 (LR)	0x000002F7
R15 (PC)	0x00000238
Ė·····xPSR	0x21000000
N	0
Z	0
C	1
······ V	0
Т	1
iISR	0

Comment/Explanation: Flags were observed to set only with the execution of instructions with a mnemonic ending with an S.

21. Answer to Q15 [3 marks]

Answer: Branch 1, 3, 6, and the last have been executed. At branch 1 BEQ, zero flag was set, thus branch 1 was executed; at branch 3 BCS, the carry flag was set, thus branch 3 was executed; at branch 6 BPL, the negative flag was clear, thus branch 6 was executed; at the last branch BVC, the overflow flag was clear, thus the last branch was executed;

Code:

LDR r0, =0x7F000000 ;set value in r0

LDR r1, =0x81000000 ;set value in r1

MOVS r2, #3 ;set value in r2

MOVS r3, #7 ;set value in r3

MOVS r4, #0 ;set value in r4 to 0

MOVS r5, #0 ;set value in r5 to 0

MOV r8, r5 ;set value in r8 to 0

MOV r9, r5 ;set value in r9 to 0

ADDS r6, r1, r0 ;add r1 and r0 and set/clear flags

BEQ next1 ;branch if zero flag set

ADD r4, r2 ;add if zero flag clear

next1

BNE next2 ;branch if zero flag clear

ADD r4, r3 ;add if zero flag set

next2

BCS next3 ;branch if carry flag set

ADD r5, r2 ;add if carry flag clear

next3

BCC next4 ;branch if carry flag clear

ADD r5, r3 ;add if carry flag set

next4

BMI next5 ;branch if negative flag set

ADD r8, r2 ;add if negative flag clear

next5

BPL next6 ;branch if negative flag clear

ADD r8, r3 ;add if negative flag set

next6

BVS next7 ;branch if overflow flag set

ADD r9, r2 ;add if overflow flag clear

next7

BVC last ;branch if overflow flag clear

ADD r9, r3 ;add if overflow flag set

Last

Explanation: At branch 1 BEQ, zero flag was set, thus branch 1 was executed; at branch 3 BCC, the carry flag was set, thus branch 3 was executed; at branch 6 BPL,

the negative flag was clear, thus branch 6 was executed; at the last branch, the overflow flag was clear, thus the last branch was executed;

22. Answer to Q16 [3 marks]

Answer: Conditional branches 2, 4, 5, 7 were executed as expected. At branch 2 BNE, zero flag was clear, thus branch 2 was executed; at branch 4 BCC, the carry flag was clear, thus branch 4 was executed; at branch 5 BMI, the negative flag was set, thus branch 5 was executed; at branch 7 BVS, the overflow flag was clear, thus the last branch was executed.

R1 was set as 0x00000238, N and V is set

Explanation:

At branch 2 BNE, zero flag was clear, thus branch 2 was executed; at branch 4 BCC, the carry flag was clear, thus branch 4 was executed; at branch 5 BMI, the negative flag was set, thus branch 5 was executed; at branch 7 BVS, the overflow flag was clear, thus the last branch was executed.

Appendix:

Code:

```
AREA asm_func, CODE, READONLY

THUMB

; Export exp26_asm function location so that C compiler can find it and link

EXPORT exp26_asm

exp26_asm PROC

; Test routine that does something!
```

; Some arithmetic

ADDS r7, r0, r1 ;

MOVS r2, #14 ;move a number into reigister 2

MOVS r3, #37 ;move a 2nd no. into reg. 3

ADDS r4, r2, #7; put sum of reg.2 and 7 in reg.3

ADDS r5, r2, r3; put sum of r2 and r3 in r5

ADDS r6, r3, r2 ; reverse the order

SUBS r0, r2, #1 ;subtract 1 from r2

SUBS r1, r2, r3 ;subtract r3 from r2

SUBS r2, r3, r2 ; reverse the order

; Some logic

MOVS r2, #0x0011 ;set bit pattern 0011 (in hex)

MOV r3, r2

ADDS r3, #0x00F0 ;set bit pattern 0101 (in hex)

MOV r4, r2

ANDS r4, r3 ;AND

MOV r5, r2

ORRS r5, r3 ;OR

MOV r6, r2

EORS r6, r3 ;exclusive OR

MOV r0, r2

BICS r0, r3 ;bit clear

MOV r1, r3

BICS r1, r2 ;bit clear in reverse order

; Some unconditional branches

MOVS r2, #0 ;clear register 2

loop1

ADDS r2, r2, #1; increment register 2

B loop1 ;loop always

MOVS r2, #0 ;clear register 2

loop2

ADDS r2, r2, #1; increment register 2

ADDS r0, r0, #0 ;do nothing

B loop2 ;loop always

MOVS r2, #0 ;clear register 2

loop3

ADDS r2, r2, #1; increment register 2

ADDS r0, r0, #0; do nothing

ADDS r0, r0, #0 ;do nothing

B loop3 ;loop always

; Setting and clearing flags

MOVS r2, #0 ;clear r2 and set/clear flags

MOVS r3, #1 ;set value in r3 to 1

MOV r4, r2 ;move r2 to r4

MOVS r4, r2 ;same as previous instruction but set/clear flags

ADDS r5, r2, r3 ;add r2 and r3 and set flags

SUBS r6, r2, r3 ;r6:= r3 - r2 and set flags

; Conditional branching

LDR r0, =0x7F000000 ;set value in r0

LDR r1, =0x81000000 ;set value in r1

MOVS r2, #3 ;set value in r2

MOVS r3, #7 ;set value in r3

MOVS r4, #0 ;set value in r4 to 0

MOVS r5, #0 ;set value in r5 to 0

MOV r8, r5 ;set value in r8 to 0

MOV r9, r5 ;set value in r9 to 0

ADDS r6, r1, r0 ;add r1 and r0 and set/clear flags

BEQ next1 ;branch if zero flag set

ADD r4, r2 ;add if zero flag clear

next1

BNE next2 ;branch if zero flag clear

ADD r4, r3 ;add if zero flag set

next2

BCS next3 ;branch if carry flag set

ADD r5, r2 ;add if carry flag clear

next3

BCC next4 ;branch if carry flag clear

ADD r5, r3 ;add if carry flag set

next4

BMI next5 ;branch if negative flag set

ADD r8, r2 ;add if negative flag clear

next5

BPL next6 ;branch if negative flag clear

ADD r8, r3 ;add if negative flag set

next6

BVS next7 ;branch if overflow flag set

ADD r9, r2 ;add if overflow flag clear

next7

BVC last ;branch if overflow flag clear

ADD r9, r3 ;add if overflow flag set

last

; return value passed back to C in r0

MOVS r0, r7

; Return to	OC using link register (Branch indirect using LR - a return)
BX	LR
ENIDD	
ENDP	
END	