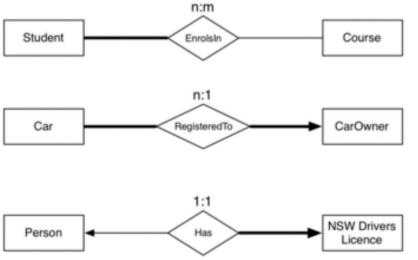
# COMP3311 WEEK 02 LECTURE

#### ER MODEL (cont)



Entities + Relationships + Attributes

Relationship constrains: total / partial, n:m/1:n/1:1

(inheritance hierarchies, weak entities)

# **RELATIONAL DATA MODEL**

- A collection of inter-connected relations (or tables)
  - o Simple, general data modelling formalism
  - Maps easily to file structures (like implementable)
- Mathematical (relation, tuple, attribute) + data-oriented (table, record, field/column)
- In relational model, ROWs represent records + COLUMNs represent attributes.
- RELATION = Name (unique within a given relation) + A set of attributes (column headings)
  - a. Each relation has a KEY (subset of attributes unique for each tuple)

A relation is nothing but a table of values.

Every row in the table represents a collection of related date values (these rows in the table denote a real-world entity or relationship)

The data are represented as a set of relations.

- 2. ATTRIBUTES = Name + An associated domain
  - a. Attribute values are ATOMIC (no composite or multi-valued attributes)
  - b. a distinguished value NULL belongs to all domains

Attribute is each column in a table. They are the properties which define a relation.

1. Many students can enroll in many courses, all the student must enroll in at least one courses.

- 2. Many cars registered to car owner, one car owner can have many cars (but at least one car). Each car has only one car owner.
- 3. One person has one drivers license, each license must be owned by one person.

SUMMARY OF ER

# Table also called Relation Domain **Primary Key** © guru99.com CustomerID CustomerName Status 1 Google Active Tuple OR Row 2 Amazon Active 3 Apple Inactive Total # of rows is Cardinality Column OR Attributes Total # of column is Degree

Database schema = a collection of RELATION SCHEMAS

Database (instance) = a collection of relation instances

#### DATABASE MANAGEMENT SYSTEMS

(Database management system is system software for creating and managing databases. The DBMS provides users and programmers with a systematic way to create, retrieve, update and manage data)

- 1. Relational model is a mathematical construct
  - i. giving a representation for data structures
  - ii. with constrains as logic formulae on relations/tuples
  - iii. an algebra for manipulating (handle or control) relations/tuples
- 2. Relational Database Management systems (RDBMSs)
  - i. RDBMSs is a collection of programs and capabilities that enable IT teams and others to create, update, administer and otherwise interact with a "relational model". Mostly use Structured query language (SQL) to access the database.
  - ii. Provide an implementation of the relational model
  - iii. Using SQL as language for: data definition, query, update
- 3. Approximations made by SQL
  - i. Relations are not required to have a key
  - ii. Relations are bags rather than sets

# **DBMS TERMINOLOGY**

Many DBMSs have multiple namespaces:

- o DBMS-level database names must be unique
- o Database-level schema names must be unique
- Schema-level table names must be unique
- o Table-level attribute name must be unique

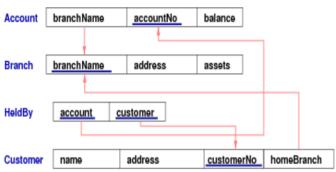
Sometimes it is convenient to use same name in several tables.

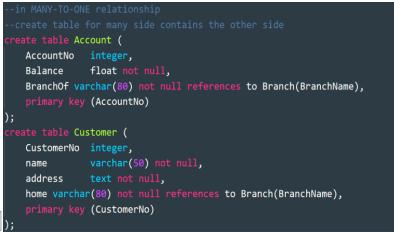
->so we distinguish which attribute we mean using qualified names

(eg. Account.branchname && Branch.branchname)

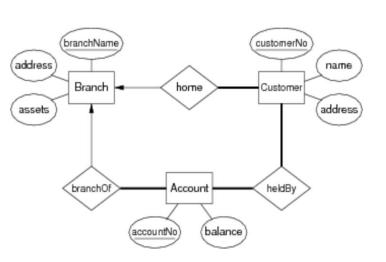
# **Example Database Schema**

Schema with 4 relations:





#### **Example Database (Instance)**



		T	Branch	T	
branchName	accountNo	balance	branchName	address	assets
Downtown	A-101	500	Downtown	Brooklyn	9000000
Mianus	A-215	700	Redwood	Palo Alto	2100000
Perryridge	A-102	400	Perryridge	Horseneck	1700000
Round Hill	A-305	350	Mianus	Horseneck	400000
Brighton	A-201	900	Round Hill	Horseneck	8000000
Redwood	A-222	700	North Town	Rys	3700000
			Brighton	Brooklyn	7100000
Customer				Depositor	
Customer	address	customerNo	homeBranch	Depositor account	customer
	address Rye	customerNo	homeBranch Mianus		customer
name	000,000		111111111111	account	
name Smith	Rye	1234567	Mianus	A-101	1313131
name Smith Jones Smith	Rye Palo Alto	1234567 9876543	Mianus Redwood	A-101 A-215	1313131
name Smith Jones	Rye Palo Alto Brooklyn	1234567 9876543 1313131	Mianus Redwood Downtown	A-101 A-215 A-102	1313131 1111111 1313131

A-102

1234567

# INTEGRITY CONSTRAINSTS

- Relations are used to represents ENTITIES and RELATIONSHIPS
- Domain limits the set of values that attributes can take
- To represents real-world problems, need to describe:
  - What values are/ are not allowed
  - What combinations of values are/ are not allowed
- Constraints are logical statements that do this:
  - o Domain, key, entity integrity, referential integrity...
- 1. NOT NULL ensures that a column cannot have a NULL value
- 2. UNIQUE ensures that all values in a column are different
- 3. PRIMARY KEY not null + unique. Uniquely identifies each row in a table
- 4. FOREIGN KEY uniquely identifies a row/record in another table
- 5. CHECK ensures that all values in a column satisfies a specific condition
- 6. DEFAULT sets a default value for a column when no value is specified
- 7. INDEX used to create and retrieve date from the database very quickly

- Domain constraints example:
  - a. "Employee.age" attribute is typically defined as integer
  - b. Better modeled by adding extra constraint (15<age<66)

```
create table Employee (
   age integer,
   age integer check (between 15 to 66),
);
```

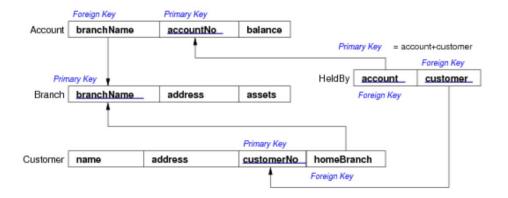
- c. NULL satisfies all domain constraints (except (not null))
- Key constraints example:
  - a. Student (id, ...) is guaranteed unique

```
create table students (
   ID integer not null unique,
   lastName varchar(255) not null,
   firstName varchar(255),
   primary key(ID) --primary key is unique and not null
);
```

- b. Class (..., day, location, time, ...) is unique
- Entity integrity example:
  - a. Class (..., mon, 2pm, lyre, ...) is well defined
  - b. Class (..., NULL, 2pm, lyre, ...) is not well defined

#### REFERENTIAL INTEGRITY

- Describe references between relations (tables)
- Are related to notion of a foreign key.
  - o The foreign key may be null and may have the same value but:
    - The foreign key value must match a record in the table it is referring to
  - Tables that do not follow this are violating the referential integrity rule



A set of attributes F in relation R1 is a foreign key for R2 if:

- the attributes in F correspond to the primary key of R2
- the value of F in each tuple of R1
   Either occurs as a primary key in in R2
   Or is entirely NULL

# **RELATIONAL DATABASES**

A relational database schema is

- A set of relation schemas {R1, R2, ..., Rn}, and
- A set of integrity constraints

#### A relational database instance is

- A set of relation instances {r1(R1), r2(R2), ...}
- Where all of the integrity constraints are satisfied

One of the MOST important functions of relational DBMS: ensure that all data in the database satisfies constraints

# **DESCRIBING RELATIONAL SCHEMAS**

We need a formalism to express relational schemas

SQL provides a Data Definition Language (DDL) for this.

```
CREATE TABLE TableName (

attrName1 domain1 constraints1,

attrName2 domain2 constraints2,

...

PRIMARY KEY (attr1, attr2, ...)

FOREIGN KET (attr5, attr6, ...) REFERENCES OtherTable(attr8, attr9)
);
```

#### **DOMAIN CONSTRAINTS**

CREATE DOMAIN creates new domain

A domain is essentially a DATA TYPE with optional constraints

(eg. Several table might contain email address columns, all requiring the same CHECK constraints to verify the address syntax, define a domain rather than setting up each table's constraint individually.)

```
CREATE DOMAIN CPI_DATA AS REAL CHECK (value >= 0 AND value <=10)

CREATE TABLE student (
    Sid char(9) primary key,
    Name varchar(30),
    Cpi CPI_DATA

);

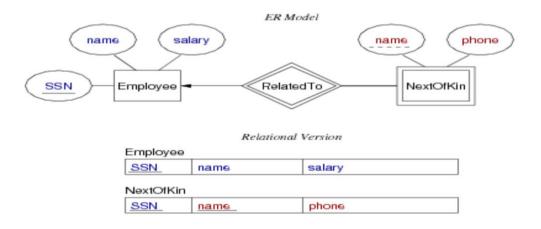
** every time cpi_data will check the constraint when you add data In student table
```

#### MAPPING ER DESIGNS TO RELATIONAL SCHEMAS

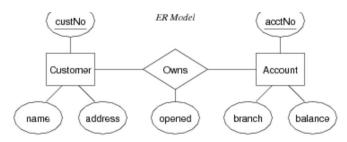
#### **ER TO RELATIONAL MAPPING**

- Perform initial data modelling using ER or OO
- Transform conceptual design into relational model

# WEAK ENTITIES MAPPING



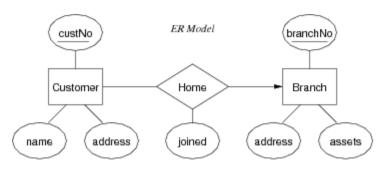
# MANY-TO-MANY MAPPING



#### Relational Version

# Customer custNo name address Account acctNo branch balance Owns custNo acctNo opened

# **ONE-TO-MANY MAPPING**



#### Relational Version

	_					
- 1	٦,	104	toi	רד	0.0	

custNo name	address	branchNo	joined
-------------	---------	----------	--------

#### Branch

branchNo	address	assets	
----------	---------	--------	--