

United International University (UIU)

Dept. of Computer Science & Engineering (CSE)

Course Code: CSI 228 Section: A
Course Title: Algorithms Lab

1. Quick Sort

The quick sort uses divide and conquer just like merge sort but without using additional storage.

The steps are:

- 1. Select an element q, called a *pivot*, from the array. In this algorithm, we have chosen the last index as the pivot.
- 2. The **PARTITION** function finds the location of the pivot in such a way that all the elements smaller than the pivot is on the left side and all the elements greater than the pivot is on the right-hand side. (Items with equal values can go either way).
- 3. Recursively call the **QUICKSORT** function which performs quicksort on the array on the left side of the pivot and then on the array on the right side, thus, dividing the task into sub-tasks. This is carried out until the arrays can no longer be split.

Implement Quick sort algorithm. The pseudo code is given below:

```
QUICKSORT(A, p, r)
                               PARTITION (A, p, r)
1 if p < r
      q = PARTITION(A, p, r)
2
                               1 \quad x = A[r]
3
       QUICKSORT(A, p, q - 1)
                               2 i = p - 1
       QUICKSORT(A, q + 1, r)
4
                                 for j = p to r - 1
                               3
                                      if A[j] \leq x
                               4
                               5
                                          i = i' + 1
                                           exchange A[i] with A[j]
                               6
                                  exchange A[i + 1] with A[r]
                                  return i + 1
int A [] = \{10, 15, 5, 19, 16, 188, 99, 14, 27, 35\};
```

Hint:

- \circ p = low
- o q = position of pivot after partitioning
- \circ r = high

the way we used in binary search and merge sort

2. Compute how many times a number x occurs in an array which is sorted using divide and conquer.

Hint: [First, find x using binary search. Then, count the number of x in the array.]

```
Boolean BS(A, key, start, end)
    mid = (start+end)/2
    if(A[mid] == key)
        return true
    else
        if(end <= start)
            return false
        else
        if (A[mid] > key)
            return BS(A, key, start, mid-1)
        else
            return BS(A, key, mid+1, end)
```

Figure 1: Binary Search Pseudocode

Sample Input array	Sample Output
x	
{1, 1, 2, 2, 2, 30} 2	4
{1, 1, 2, 2, 2, 30} 1	2
{1, 1, 2, 2, 2, 30} 30	1

3. Compute a^b using recursion.

Sample Input	Sample Output
a b	
3 5	243
2 16	256
5 5	3125