Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.

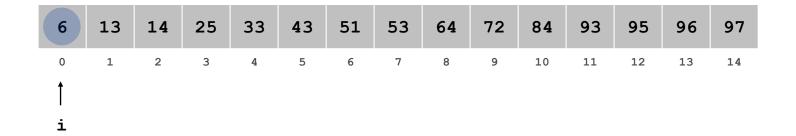
```
public int linear search(int[] a, int key) {
    for(int i=0; i < a.length; i++) {
        if(a[i]== key) return i;
    }
    return -1;
}</pre>
```

Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.

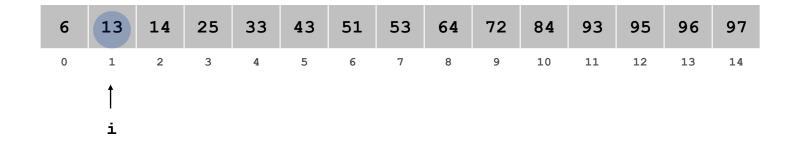
										84					
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	

Linear Search (Average Case)

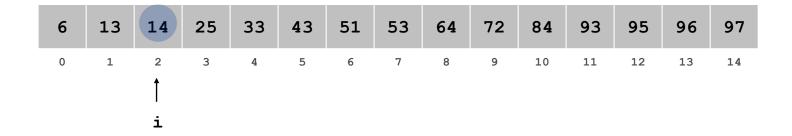
Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.



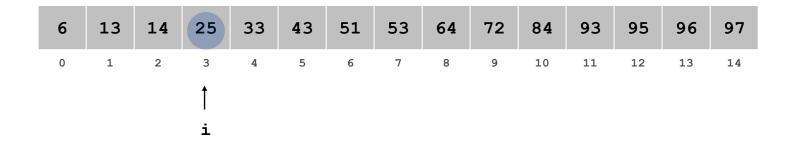
Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.



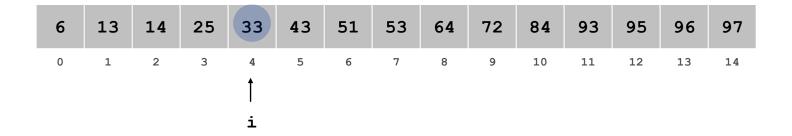
Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.



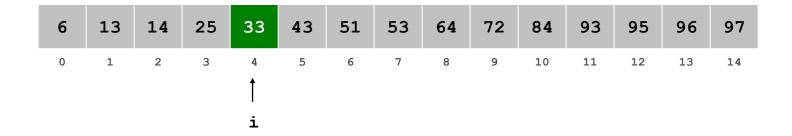
Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.



Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.



Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.



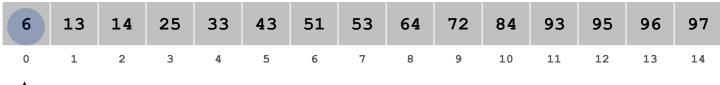
Linear Search: Best Case

Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.

6							53								
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	

Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.

Ex. Linear search for 6.

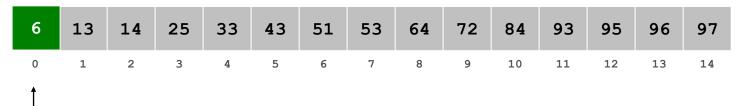


Î

i

Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.

Ex. Linear search for 6.



•

i

Linear Search: Worst Case

Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.

6														97
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

Linear Search: Worst Case

Linear search. Given value and an array a[], find index i such that a[i] = value, or report that no such index exists.

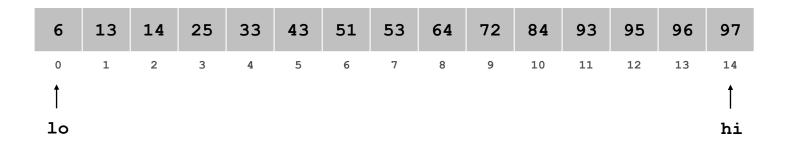
6							53								
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	

Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

```
public int binary search(int[] a, int key) {
       int lo = 0;
       int hi = a.length - 1;
       while (lo <= hi) {
           // Key is in a[lo..hi] or not present.
           int mid = lo + (hi - lo) / 2;
           if
                   (key \langle a[mid] \rangle hi = mid - 1;
           else if (key > a[mid]) lo = mid + 1;
           else return mid;
       return -1;
```

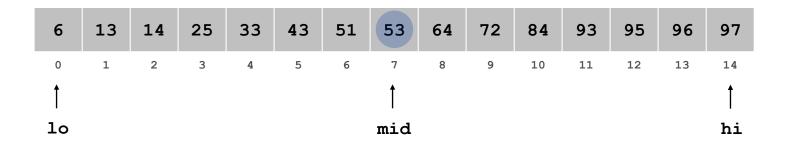
Binary Search: Average case

Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.



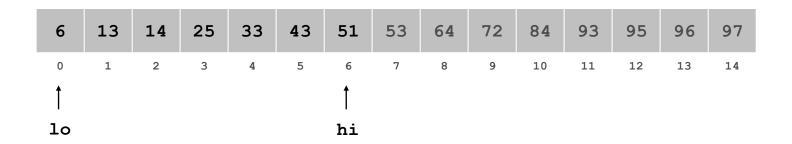
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



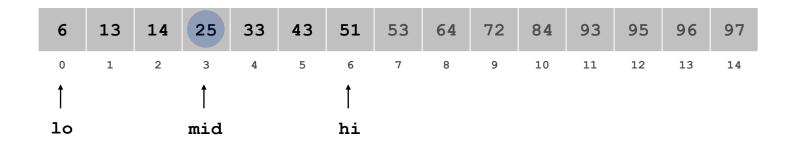
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



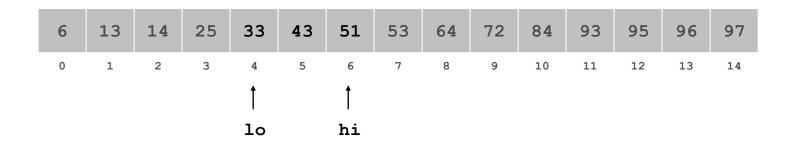
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



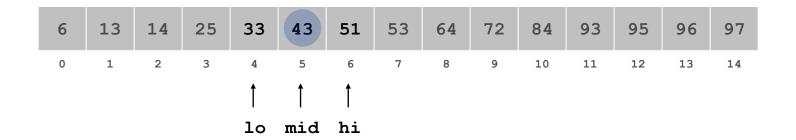
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



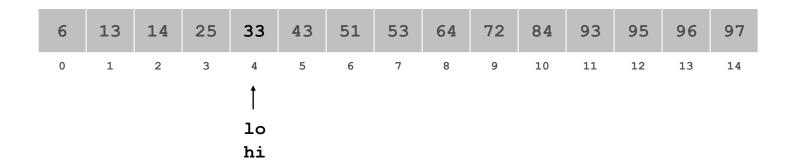
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



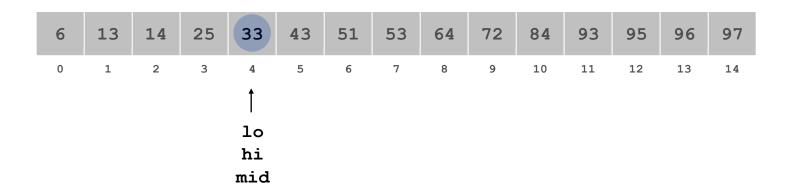
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



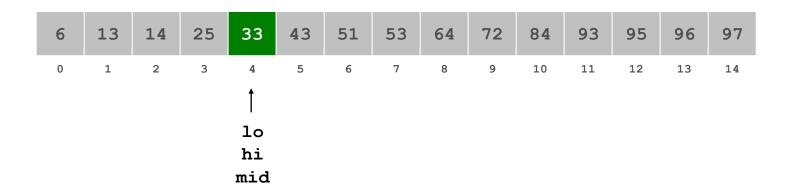
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



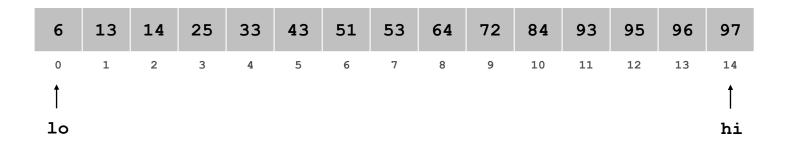
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



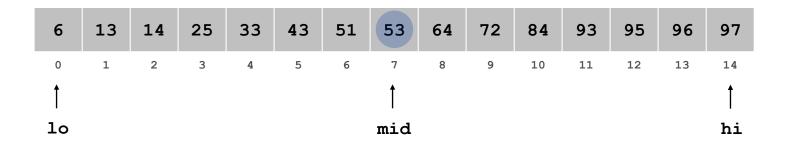
Binary Search: Best case

Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.



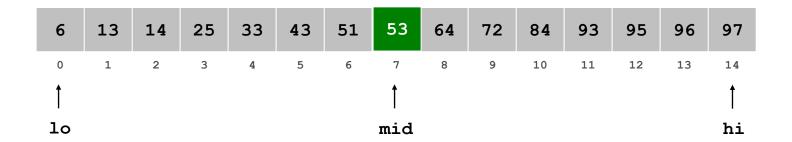
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



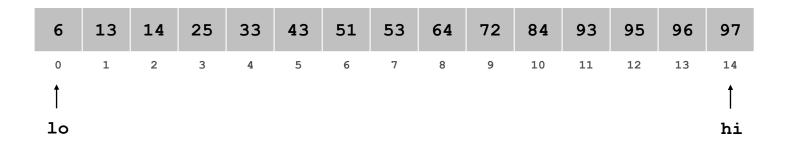
Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.

Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



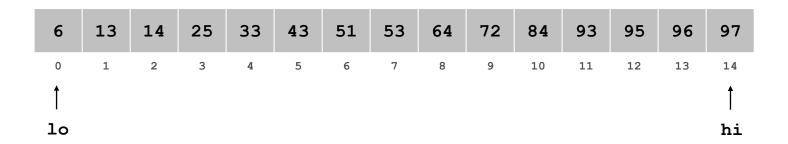
Binary Search: Worst case

Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.



Binary Search: Worst case

Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.



Finding time complexity of Linear Search

```
public int linear search(int[] a, int key) {
    for(int i=0; i < a.length; i++) {
        if(a[i]== key) return i;
        }
    return -1;
}</pre>
Approximately 3 operations per iteration
    f(n) = 3n + 2
```

Finding time complexity of Linear Search: Average case

$$f(n) = 3n + 2$$

n = no of iterations/ location of key

$$f(1) = 3.1 + 2$$

$$f(2) = 3.2 + 2$$

$$f(3) = 3.3 + 2$$

$$\vdots$$

$$\vdots$$

$$f(n) = 3.n + 2$$

$$f(1) + f(2) + ... + f(n) = 3.(1 + 2 + ... + n) + (2 + 2 + ... + 2)$$

$$= 3 n(n+1) / 2 + 2.n$$

$$= \frac{1}{2} (3n^2 + 7n)$$

No of cases = n Average case time complexity = $\frac{1}{2}$ (3n² + 7n) /n = $\frac{1}{2}$ (3n + 7)

Big O notation: O(n)

Finding time complexity of Binary Search: Average case

```
public int binary search(int[] a, int key) {
        int lo = 0;
                                                   Approximately 3 comparisons per
                                                   iteration
        int hi = a.length - 1;
        while (lo <= hi) {
            // Key is in a[lo..hi] or not present.
            int mid = lo + (hi - lo) / 2;
            if
                    (key < a[mid]) hi = mid - 1;
            else if (key > a[mid]) lo = mid + 1;
            else return mid;
        return -1;
    }
```

Finding time complexity of Binary Search: Average case

```
List size n = 2^k
 k = \log_2 n
```

```
List size No of comparisons
2^{k} = 3
2^{k-1} = 3
2^{k-2} = 3
\vdots
\vdots
1 = 3
= (3+3+...+3)
= 3(k+1)
= 3(\log_{2}n + 1)
```

Big O notation: $O(log_2n)$