

Linear Search

Linear search. Given `value` and an array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

```
public int linear search(int[] a, int key) {  
    for(int i=0; i < a.length; i++){  
        if(a[i]== key) return i;  
    }  
    return -1;  
}
```

Linear Search

Linear search. Given `value` and an array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Linear search for 33.

6	13	14	25	33	43	51	53	64	72	84	93	95	96	97
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

Linear Search (Average Case)

Linear search. Given `value` and an array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Linear search for 33.

[illegible]

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`i`

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↑
`i`

Linear Search: Best Case

Linear search. Given `value` and an array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Linear search for 6.

6	13	14	25	33	43	51	53	64	72	84	93	95	96	97
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`i`

Linear Search

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Ex. Linear search for 6.

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
i↑	6	13	14	25	33	43	51	53	64	72	84	93	95	96	97

Linear Search: Worst Case

Linear search. Given `value` and an array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Linear search for 97 .

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0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

Linear Search: Worst Case

Linear search. Given `value` and an array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Linear search for 101 .

6	13	14	25	33	43	51	53	64	72	84	93	95	96	97
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Binary Search

Binary search. Given `value` and **sorted** array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

```
public int binary search(int[] a, int key) {  
    int lo = 0;  
    int hi = a.length - 1;  
    while (lo <= hi) {  
        // Key is in a[lo..hi] or not present.  
        int mid = lo + (hi - lo) / 2;  
        if      (key < a[mid]) hi = mid - 1;  
        else if (key > a[mid]) lo = mid + 1;  
        else return mid;  
    }  
    return -1;  
}
```

Binary Search: Average case

Binary search. Given `value` and sorted array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Binary search for 33.

[illegible]

Binary Search

Binary search. Given `value` and sorted array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Invariant. Algorithm maintains $a[\text{lo}] \leq \text{value} \leq a[\text{hi}]$.

Ex. Binary search for 33.

6	13	14	25	33	43	51	53	64	72	84	93	95	96	97
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
↑							↑							↑
lo							mid							hi

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↑			↑			↑								
lo			mid			hi								

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				↑		↑								
				lo		hi								

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				↑	↑	↑								
				lo	mid	hi								

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↑
lo
hi

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↑
lo
hi
mid

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↑
lo
hi
mid

Binary Search: Best case

Binary search. Given `value` and sorted array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Binary search for 53.

[illegible]

Binary Search

Binary search. Given `value` and sorted array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Invariant. Algorithm maintains $a[\text{lo}] \leq \text{value} \leq a[\text{hi}]$.

Ex. Binary search for 53.

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↑							↑							↑
lo							mid							hi

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Binary search. Given `value` and sorted array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Invariant. Algorithm maintains $a[lo] \leq value \leq a[hi]$.

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↑							↑							↑
lo							mid							hi

Binary Search: Worst case

Binary search. Given `value` and sorted array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Binary search for 4.

[illegible]

Binary Search: Worst case

Binary search. Given `value` and sorted array `a[]`, find index `i` such that `a[i] = value`, or report that no such index exists.

Ex. Binary search for 101.

[illegible]

Finding time complexity of Linear Search

```
public int linear search(int[] a, int key) {  
    for(int i=0; i < a.length; i++){  
        if(a[i]== key) return i;  
    }  
    return -1;  
}
```

Approximately 3 operations per
iteration
 $f(n) = 3n + 2$

Finding time complexity of Linear Search: Average case

$$f(n) = 3n + 2$$

n = no of iterations/ location of key

$$f(1) = 3.1 + 2$$

$$f(2) = 3.2 + 2$$

$$f(3) = 3.3 + 2$$

⋮

⋮

⋮

$$f(n) = 3.n + 2$$

$$f(1) + f(2) + \dots + f(n) = 3.(1+2+\dots + n) + (2+ 2+ \dots + 2)$$

$$= 3 n(n+1) / 2 + 2.n$$

$$= \frac{1}{2} (3n^2 + 7n)$$

No of cases = n

$$\text{Average case time complexity} = \frac{1}{2} (3n^2 + 7n) / n$$

$$= \frac{1}{2} (3n + 7)$$

Big O notation: $O(n)$

Finding time complexity of Binary Search: Average case

```
public int binary search(int[] a, int key) {
```

```
    int lo = 0;
```

```
    int hi = a.length - 1;
```

```
    while (lo <= hi) {
```

```
        // Key is in a[lo..hi] or not present.
```

```
        int mid = lo + (hi - lo) / 2;
```

```
        if      (key < a[mid]) hi = mid - 1;
```

```
        else if (key > a[mid]) lo = mid + 1;
```

```
        else return mid;
```

```
    }
```

```
    return -1;
```

```
}
```

Approximately 3 comparisons per iteration

Finding time complexity of Binary Search: Average case

List size $n = 2^k$

$k = \log_2 n$

List size	No of comparisons
-----------	-------------------

$$2^k = 3$$

$$2^{k-1} = 3$$

$$2^{k-2} = 3$$

$$\vdots$$
$$\vdots$$
$$\vdots$$

$$1 = 3$$

$$= (3 + 3 + \dots + 3)$$

$$= 3(k + 1)$$

$$= 3(\log_2 n + 1)$$

Big O notation: $O(\log_2 n)$