

Assignment9 Report

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实验课时段: 周五5-6节

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1. Disk scheduling

(1) READ/WRITE data time = Seek Time + Rotational Latency + Transfer Time

(2)

a. Track Access Sequence:

- FIFO: 70 => 30 => 90 => 120 => 60 => 20
- SSTF: 90 => 70 => 60 => 30 => 20 => 120
- SCAN: 120 => 90 => 70 => 60 => 30 => 20
- CSCAN: 120 => 20 => 30 => 60 => 70 => 90

b. Time Calculation

For the seek time, first calculate their tracks:

- FIFO: $(100-70) + (70-30) + (90-30) + (120-90) + (120-60) + (60-20) = 260$
- SSTF: $(100-90) + (90-70) + (70-60) + (60-30) + (30-20) + (120-20) = 180$
- SCAN: $(120-100) + (199-120) + (199-90) + (90-70) + (70-60) + (60-30) + (30-20) = 278$
- CSCAN: $(120-100) + (199-120) + (199-0) + (20-0) + (30-20) + (60-30) + (70-60) + (90-70) = 388$

Thus time is:

- FIFO: 260ms

- SSTF: 180ms
- SCAN: 278ms
- CSCAN: 388ms

For FIFO\SSTF\SCAN\CSCAN algorithm, their rotational latency is the same.

12000r/min => 200r/s => 0.2r/ms => 5ms/r

Since it is randomly distributed access, we treat it as half round: 2.5ms

Since 6 accesses, total time = 2.5ms*6 = 15ms

The question does not tell us the transfer time, thus we omitted.

Total Time:

- FIFO: 260+15=275ms
- SSTF: 180+15=195ms
- SCAN: 278+15=293ms
- CSCAN: 388+15=403ms

2. Simple File System

Consider adding one statement in `tools/mksfs.c` :

```

428 static void
429 add_entry(struct sfs_fs *sfs, struct cache_inode *current, struct cache_inode *file, const char *name) {
430     static struct sfs_entry __entry, *entry = &__entry;
431     assert(current->inode.type == SFS_TYPE_DIR && strlen(name) <= SFS_MAX_FNAME_LEN);
432     entry->ino = file->ino, strcpy(entry->name, name);
433     uint32_t entry_ino = sfs_alloc_ino(sfs);
434     write_block(sfs, entry, sizeof(struct sfs_entry), entry_ino);
435     append_block(sfs, current, sizeof(struct sfs_entry), entry_ino, name);
436     printf("entry is %d, name is %s, file->ino is %d\n", entry_ino, name, file->ino);
437     file->inode.nlinks ++;
438 }
439 }
440

```

And the partial output of `make qemu` is:

```
48016 records out
480000 bytes (480 kB, 469 KiB) copied, 0.00130072 s, 369 MB/s
entry is 3, name is ., file->ino is 1
entry is 4, name is .., file->ino is 1
entry is 7, name is test, file->ino is 5
entry is 12, name is hello, file->ino is 8
entry is 17, name is sh, file->ino is 13
entry is 19, name is ., file->ino is 18
entry is 20, name is .., file->ino is 1
entry is 23, name is file1, file->ino is 21
entry is 24, name is dir1, file->ino is 18
create bin/sfs.img (disk0) successfully.
+ ld bin/kernel
riscv64-unknown-elf-objcopy bin/kernel --strip-all -O binary bin/ucore.img
```

Thus we draw the diagram according to the running result:

0	superblock		0	3		
1	inode	rootdir	1	4		
2		freemap	2	7		
3	entry	name='.'	3	12		
4	entry	name='..'	4	17		
5	inode	test文件的inode	5	24		
6		test data				
7	entry	name='test'	0	6		
8	inode	hello文件的inode				
9		hello data 1	0	9		
10		hello data 2	1	10		
11		hello data 3	2	11		
12	entry	name='hello'				
13	inode	sh文件的inode	0	14		
14		sh data 1	1	15		
15		sh data 2	2	16		
16		sh data 3				
17	entry	name='sh'	0	19		
18	inode	dir1的inode	1	20		
19	entry	name='.'(一个点)	2	23		
20	entry	name='..'(两个点)				
21	inode	file1的inode				
22		file1 data	0	22		
23	entry	name='file1'				
24	entry	name='dir1'				
25-116		空闲状态				

Here is a brief explanation:

第一个 Table (root dir inode)

- 的 entry

- 的 entry

- dir1 的 entry

- test 的 entry

- sh 的 entry

- hello 的 entry

第二个 Table (dir inode)

- 的 entry

- 的 entry

- file1 的 entry

第三个 Table (file1 的 inode)

- file1 的 data x1

第四个 Table' (test 的 inode)

- test 的 data x1

第五个 Table (sh 的 inode)

- sh 的 data x3

第六个 Table (hello 的 ~~data~~ inode)

hello 的 data x3

By 11910104