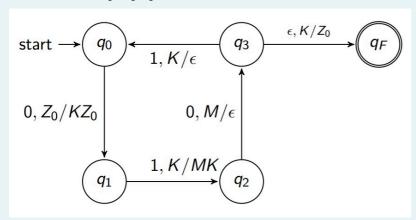
Started on	Wednesday, 15 March 2023, 12:52 PM
State	Finished
Completed on	Wednesday, 15 March 2023, 1:32 PM
Time taken	39 mins 21 secs
Grade	<b>11.00</b> out of 20.00 ( <b>55</b> %)
Question <b>1</b> Correct Mark 1.00 out of 1.00	
What is a regular la	nguage?
<ul> <li>a. A formal language that can be defined by a regular expression</li> </ul>	
O b. A formal language that can be recognized by a finite state machine	
○ c. A formal language that can be generated by a regular grammar	
<ul><li>d. All of the ab</li></ul>	
G. All of the at	
Your answer is correct answer All of the above	
Question 2 Incorrect Mark 0.00 out of 1.00	
What is a context-fr	ee language?
a. A language	that has no grammar rules
<ul><li>b. A language</li></ul>	that can be spoken in any context ×
oc. A language	requiring an unbounded memory to be recognized
○ d. A language	for which you can apply Pumping Lemma
Your answer is inco	rrect.
The correct answer	is:
	g an unbounded memory to be recognized

Which of the following languages suits the drawn PDA?



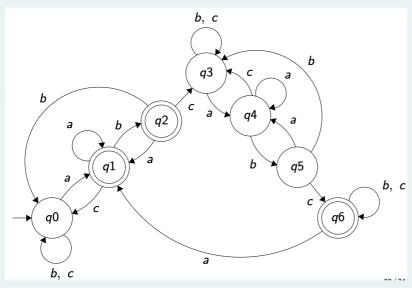
- a.  $L = \{(01)^n0 \mid n>0 \& n \text{ is odd}\} \checkmark$
- b.  $L = {(01)^n1 \mid n>0 \& n \text{ is odd}}$
- C. L = {(01)<sup>n</sup>0 | n>0 & n is even}
- Od.  $L = \{(01)^n 1 \mid n > 0 \& n \text{ is even}\}$

Your answer is correct.

The correct answer is:

 $L = \{(01)^n0 \mid n>0 \& n \text{ is odd}\}$ 

Which of the following languages with  $\Sigma$  = {a, b, c} suits the drawn FSA?



- $\bigcirc$  a. L = {w  $\in \Sigma^*$  | the substring abc in w occurs an odd number of times}
- $\bigcirc$  b. L = {w  $\in \Sigma^*$  | the substring abc in w occurs an even number of times}
- $\circ$  c. L = {w  $\in \Sigma^*$  | the substring abc in w occurs an even number of times  $\Lambda$  |w| > 0}
- $\bigcirc$  d. L = {w  $\in \Sigma^*$  | the substring abc in w occurs an odd number of times  $\land$  |w| > 0}
- e. None of the above

Your answer is correct.

The correct answer is:

None of the above

#### Question 5

Correct

Mark 1.00 out of 1.00

Relate the following statement: All sufficiently long words in a regular language can have a middle section of words repeated a number of times to produce a new word which also lies within the same language.

- a. Halting Problem
- b. Pumping Lemma ✓
- o. Myhill-Nerode theorem
- od. Rice Theorem

Your answer is correct.

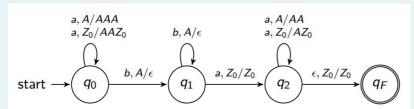
The correct answer is:

**Pumping Lemma** 

Incorrect		
Mark 0.00 out of 1.00		
Which of the following languages is/are NOT regular?		
☑ a. a <sup>n</sup> b <sup>n</sup> ✓		
$ □ $ b. L = {wcw <sup>R</sup>   w ∈ {a, b}* Λ  w  > 0} ✓		
☑ c. a <sup>n</sup> b <sup>m</sup> ×		
$\square$ d. $a^{n!}$		
Your answer is incorrect.		
The correct answers are:		
a <sup>n</sup> b <sup>n</sup> ,		
$L = \{wcw^R \mid w \in \{a, b\} * \land  w  > 0\},$		
$a^{n!}$		
Question 7		
Correct		
Mark 1.00 out of 1.00		
Which of the following one can relate to the following statement: if n items are put in m containers with n>m, the at least one		
container will contain more than one item?		
O a Muhill Navada theorem		
a. Myhill–Nerode theorem		
○ c. None of the mentioned		
○ d. Pumping Lemma		
a. Fullipling Lemma		
a. Tuliping Lemina		
Your answer is correct.		
Your answer is correct.		
Your answer is correct. The correct answers are:		
Your answer is correct. The correct answers are: Pumping Lemma,		
Your answer is correct. The correct answers are: Pumping Lemma,		

Question 6

Which of the following languages suits the drawn PDA?



- a.  $L = \{a^n b^{2n} a \mid n \in N\}$
- b.  $L = \{a^{2n}b^na^n \mid n \in N\}$
- c.  $L = \{a^n b^{2n} a^n \mid n \in \mathbb{N}\} \times$
- $\bigcirc$  d. L = { $a^{2n}b^na \mid n \in \mathbf{N}$ }

Your answer is incorrect.

The correct answer is:

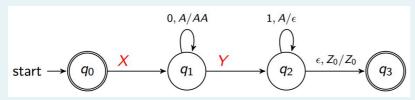
 $L = \{a^n b^{2n} a \mid n \in \mathbf{N}\}$ 

#### Question 9

Incorrect

Mark 0.00 out of 1.00

Find the rule transitions X and Y such that the following PDA recognizes the language L=  $\{0^n1^{n+1} \mid n \in N\}$ 



- a.  $X = "0, Z_0/Z_0", Y = "1, A/\epsilon"$
- O b.  $X = "0, Z_0/AZ_0", Y = "1, A/A"$
- $\circ$  c.  $X = "0, Z_0/Z_0", Y = "1, A/A"$
- ⊚ d.  $X = "0, Z_0/AZ_0", Y = "1, A/\epsilon" ×$

Your answer is incorrect.

The correct answer is:

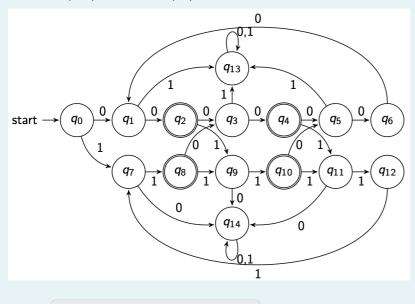
 $X = "0, Z_0/AZ_0", Y = "1, A/A"$ 

Incorrect		
Mark 0.00 out of 1.00		
Which of the following statements about FSA is/are TRUE?		
☑ a. Each complete FSA is also deterministic   X		
☐ b. The set of final states may be empty		
extstyle  ext		
☐ d. The sets of initial and accepting states may not intersect		
<ul> <li>e. In complete FSA the number of transitions from each state matches the number of symbols in the alphabet (assume no ε-moves)</li> </ul>		
Your answer is incorrect.		
The correct answers are:		
An FSA with a total transition function is called complete,		
In complete FSA the number of transitions from each state matches the number of symbols in the alphabet (assume no ε-moves)		
Question 11		
Correct		
Mark 1.00 out of 1.00		
While applying Pumping lemma over a language, we consider a string w that belongs to L and fragment it into parts.		
○ a. 9		
O b. 7		
○ c. 5		
⊚ d. 3 ✓		
Very annual in a second		
Your answer is correct.		
The correct answer is:  3		

Question **10** 

Mark 0.00 out of 1.00

Consider the given language  $L = \{w \in \Sigma^* \mid \text{each 0 in } w \text{ is doubled, each 1 is also doubled } \Lambda \mid w \mid \text{is NOT divisible by 3} \}$  with  $\Sigma = \{0, 1\}$ . Given the incomplete FSA for this language type the missing transitions with its input and output state indices, e.g. for missing transition qy = (qx, 1) you should type x1y. If multiple transitions are missing, do NOT use spaces and use numerical order for both initial states (0-14) and transitions (0-1).



The correct answer is: 6171201

### Question 13

Correct

Mark 1.00 out of 1.00

Answer: 619

Languages are proved to be regular or non regular using Pumping Lemma

- a. Often true
- ob. False
- oc. True
- d. Almost always false

Your answer is correct.

The correct answers are:

True,

False,

Often true,

Almost always false

Question 14 Incorrect Mark 0.00 out of 1.00		
Which of the following statements about FST is/are TRUE?		
a. Also called finite-state deceptor		
extstyle  ext		
☑ c. It is an FSA with input and output tapes ✓		
☑ d. Input and output alphabets have to match ×		
Your answer is incorrect.		
The correct answers are:		
It is an FSA with input and output tapes,		
The translation may happen only to the accepted strings		
Question 15		
Correct		
Mark 1.00 out of 1.00		
Which of the following languages require(s) unbounded memory to be recognized?		
$\square$ a. L = { $a^nb^n$ } $\checkmark$		
$\square$ b. L = { $a^nb^m   n \neq m$ } $\times$		
$\square$ c. L = { $a^{n!}$ } $\checkmark$		
☑ d. $L = \{ wcw^R \mid w \in \{a, b\} * \land  w  > 0 \} \checkmark$		
Your answer is correct.		
The correct answers are:		

 $L = \{ a^n b^n \},$ 

 $L = \{ a^{n!} \}$ 

 $L = \{ \ wcw^R \ | \ w \in \{a, \, b\} * \ \land \ |w| > 0 \ \},$ 

Which of the following languages suits the drawn PDA?

$$a, Z_0/AZ_0$$

$$b, Z_0/BZ_0$$

$$a, A/AA$$

$$a, B/AB$$

$$b, A/BA$$

$$b, B/BB$$

$$a, A/\epsilon$$

$$b, B/BB$$

$$b, B/\epsilon$$

$$c, B/B$$

$$c, A/A$$

$$q_1$$

$$\epsilon, Z_0/Z_0$$

- a.  $L = \{wcw^R \mid w \in \{a, b\}^* \& |w| \ge 0\}$
- b.  $L = \{wcw^R \mid w \in \{a, b\} \& |w| > 0\}$
- ⊚ c.  $L = \{wcw^R \mid w \in \{a, b\} * \& |w| > 0\}$  ✓
- d.  $L = \{wcw^R \mid w \in \{a, b\} * \& |w| = 0\}$

Your answer is correct.

The correct answer is:  $L = \{wcw^R \mid w \in \{a, b\} * \& |w| > 0\}$ 

## Question 17

Incorrect

Mark 0.00 out of 1.00

Pumping lemma for context free grammar is used for:

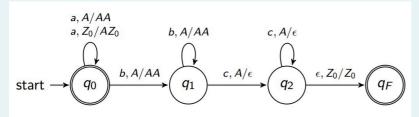
- a. Proving certain languages are not context free
- b. Proving language is infinite
- $\,\,\bigcirc\,$  c. Proving certain languages are not context free and that language is infinite
- d. None of the above X

Your answer is incorrect.

The correct answer is:

Proving certain languages are not context free

The following PDA recognizes L= $\{a^ib^jc^k\mid RULL\}$ . Find "RULL".



- a. RULL = "i + k = j"
- b. RULL = "k+j=i"
- o. RULL = "i+i=k"
- d. RULL = "j + i = k "

Your answer is correct.

The correct answer is:

RULL = "j + i = k"

### Question 19

Incorrect

Mark 0.00 out of 1.00

What is the Kleene star operation on a formal language?

a. The operation that returns the set of all strings that can be obtained by inserting any symbol from the alphabet between any two symbols of any string from the language

×

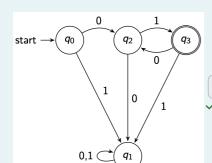
- b. The operation that returns the set of all strings that can be obtained by repeating any string from the language any number of times, including zero times
- c. The operation that returns the complement of a given language
- d. It is the free monoid on that language

Your answer is incorrect.

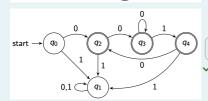
The correct answer is:

It is the free monoid on that language

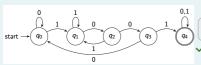
#### Match the FSAs with the languages that may be accepted by them ( $\Sigma = \{0, 1\}$ )



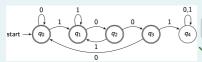
 $L = \{w \in \Sigma * \mid w \text{ starts from 0 } \land w \text{ ends with 1 } \land w \text{ does NOT contain substrings containing the same symbol starts} \}$ 



 $L = \{w \in \Sigma * \mid \text{in } w \text{ each } 1 \text{ is preceded by double } 0 \text{ } \Lambda \text{ } w \text{ starts from } 0\}$ 



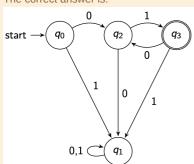
 $L = \{w \in \Sigma * \mid w \text{ contains the substring 1001}\}$ 



 $_{-}$  = {w  $\in \Sigma * | w \text{ does NOT contain the substring 1001}}$ 

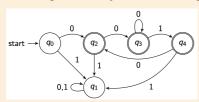
#### Your answer is correct.

# The correct answer is:

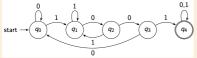


 $_{\rightarrow}$  L = {w  $\in \Sigma *$  | w starts from 0  $_{\Lambda}$  w ends with 1  $_{\Lambda}$  w does NOT contain substrings

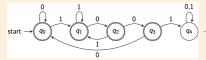
containing the same symbol with the length 2  $\Lambda$  |w| > 1,



 $_{\rightarrow}$  L = {w ∈ Σ\* | in w each 1 is preceded by double 0 ∧ w starts from 0},



 $\rightarrow$  L = {w  $\in \Sigma * | w \text{ contains the substring 1001}},$ 



 $L = \{w \in \Sigma * \mid w \text{ does NOT contain the substring 1001}\}$ 

