

Homework 3 - Q8

Team 2

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Question

“ A rod-cutting problem is said that, given a rod and a table of prices. How to cut a steel rod of length n into pieces in order to maximize the revenue r ? Suppose the rod lengths are always an integer. Given a rod of length 10, how to cut it to get the maximum revenue? You need to show your steps and find the maximum revenue.

(Hint: You can use the recursive function in PPT Unit 5.)

”

Concept

For rod cut method $T(n)$ and revenue function R

- Using **divide-and-conquer**:

$$T(n) = \{i_n\} \cup T(n - i_n),$$

where

$$i_n = \min_{i \in \mathbb{Z}_n} R(\{n - i\}) + R(T(i)),$$

and treated $T(n)$ as a sequence.

Algorithm (revenue)

```
def getRevenue(ns):  
    revenues = [2, 7, 9, 11, 13, 15, 17, 17, 19, 20]  
    r = 0  
    for i in ns:  
        if i == 0:  
            r += 0  
        if i > len(revenues) or i < 0:  
            raise ValueError("i out of range.")  
        else:  
            r += revenues[i-1]  
    return r
```

Algorithm (divide-and-conquer)

```
from numpy import argmax

def rodcut(n):
    global cMem # with memoization
    if n == 0: return []
    if cMem[n-1]: return cMem[n-1]
    cuts, revenues = [], []
    for i in range(n):
        c = [n-i] + rodcut(i)
        cuts.append(c)
        revenues.append(getRevenue(c))

    argmaxCuts = cuts[argmax(revenues)]
    cMem[n-1] = argmaxCuts
    return argmaxCuts
```

Algorithm (main function)

```
n = 10
cMem = [None for i in range(n)]

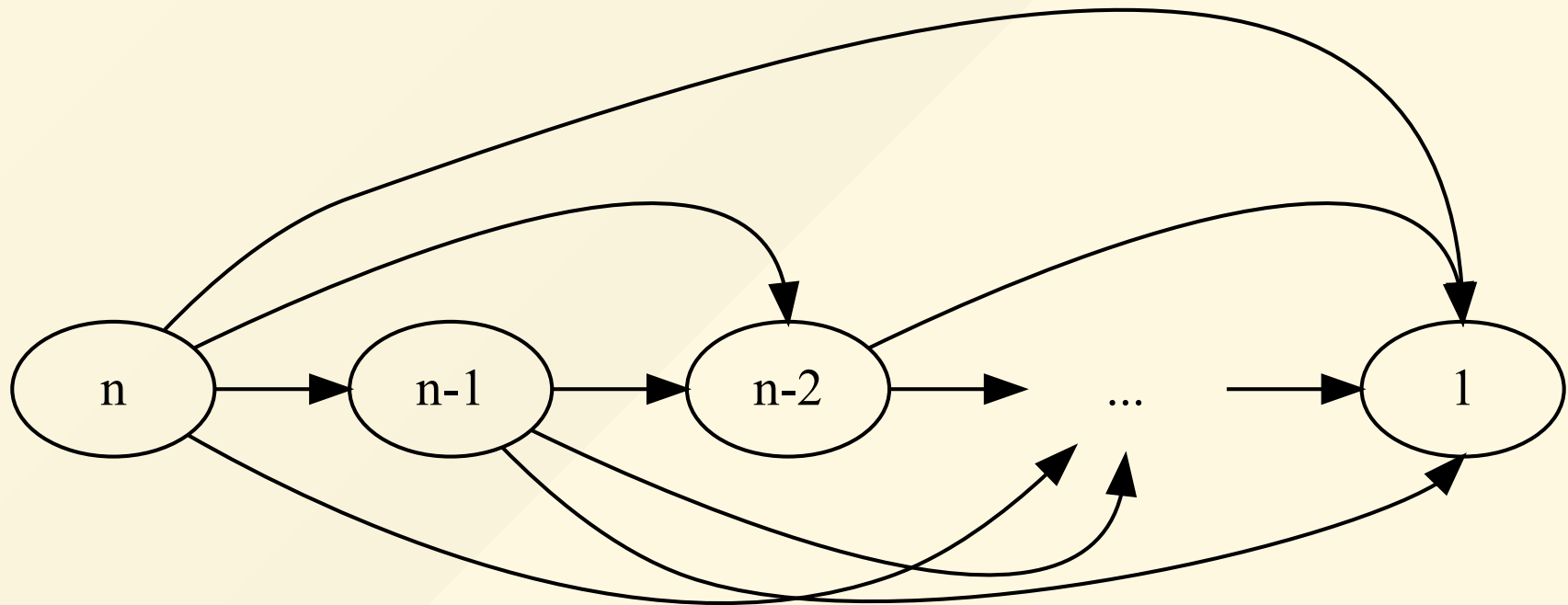
argmax = rodcut(n) # best partition
max = getRevenues(argmax) # best revenue
```

This algorithm is of class $O(n^2)$;

$$(n - 1) + (n - 2) + \dots + 1 = \frac{(n - 1)n}{2} = O(n^2).$$

(see the next page for more information)

$$n = 10$$



Thank you for your kind attention

Questions & comments are welcome.