

$$F(m) = \min_{0 \leq i < m} \{ F(i) + w(i, m) \}$$

	1	2	3	4	5	6	7	8
a	the	i	and	his	her	arm	but	

$$F(1) = \min\{0 + w(0,1) = 0 + 25 = 25\}$$

$$F(2) = \min \left\{ \begin{array}{l} 0 + w(0,2) = 0 + 4 = 4 \\ F(1) + w(1,2) = 25 + 9 = 34 \end{array} \right\}$$

$$F(3) = \min \left\{ \begin{array}{l} 0 + w(0,3) = 0 + 1 = 1 \\ F(1) + w(1,3) = 25 + 4 = 29 \\ F(2) + w(2,3) = 4 + 25 = 29 \end{array} \right.$$

$$F(i)$$

0	
25	
4	
1	
4	
1	
4	
1	

i

w(i,m) Table

[illegible]

$$F(m) = \min_{0 \leq i < m} \{ F(i) + w(i, m) \}$$

$F(i)$

$w(i, m)$

$OUT(i, m) = F(i) + w(i, m)$

	0	1	2	3	4	5	6	7	8
0		25	4	1	4	25	64	121	196
25	-	-	9	4	1	16	49	100	169
4	-	-	-	25	4	1	16	49	100
1	-	-	-	-	9	0	9	36	81
4	-	-	-	-	-	9	0	9	36
1	-	-	-	-	-	-	9	0	9
4	-	-	-	-	-	-	-	9	0
1	-	-	-	-	-	-	-	-	9
8	-	-	-	-	-	-	-	-	-

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10

1. Add $F(i)$ to each cell (i, m) in row i of $w \rightarrow OUT(i, m)$
2. Compute the column minima of the OUT matrix, such that $F(i) \leftarrow$ the minimum of column i

Convexity/Concavity Properties of the tables W and OUT

(1) Monge (concave)

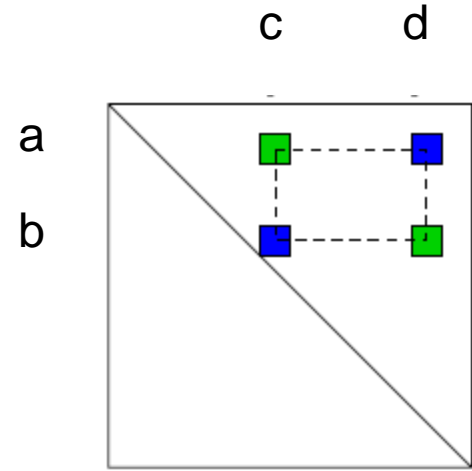
$$W[b, d] - W[a, d] \leq W[b, c] - W[a, c]$$

for all $a < b$ and $c < d$

Monge (convex)

$$W[b, d] - W[a, d] \geq W[b, c] - W[a, c]$$

for all $a < b$ and $c < d$



(2) Total Monotonicity (concave)

$$W[b, c] \leq W[a, c] \Rightarrow W[b, d] \leq W[a, d]$$

for all $a < b$ and $c < d$

Total Monotonicity (convex)

$$W[b, c] \geq W[a, c] \Rightarrow W[b, d] \geq W[a, d]$$

for all $a < b$ and $c < d$

[illegible]

W is Monge (concave) and therefore W is Totally Monotone (concave)

$$W[b, d] - W[a, d] \leq W[b, c] - W[a, c] \text{ for all } a < b \text{ and } c < d$$

$$W[b, c] \leq W[a, c] \Rightarrow W[b, d] \leq W[a, d] \text{ for all } a < b \text{ and } c < d$$

0 1 2 3 4 5 6 7 8

a	the	i	and	his	her	arm	but
---	-----	---	-----	-----	-----	-----	-----

i

	0	1	2	3	4	5	6	7	8
0		25	4	1	4	25	64	121	196
1	-	-	9	4	1	16	49	100	169
2	-	-	-	25	4	1	16	49	100
3	-	-	-	-	9	0	9	36	81
4	-	-	-	-	-	9	0	9	36
5	-	-	-	-	-	-	9	0	9
6	-	-	-	-	-	-	-	9	0
7	-	-	-	-	-	-	-	-	9
8	-	-	-	-	-	-	-	-	-

$w(i, m) = (\# \text{ of characters from } i \text{ to } m - \text{lineopt})^2$

ex: $W(0, 4) = (8 - 6)^2 = 4$, $W(0, 5) = (11 - 6)^2 = 25$

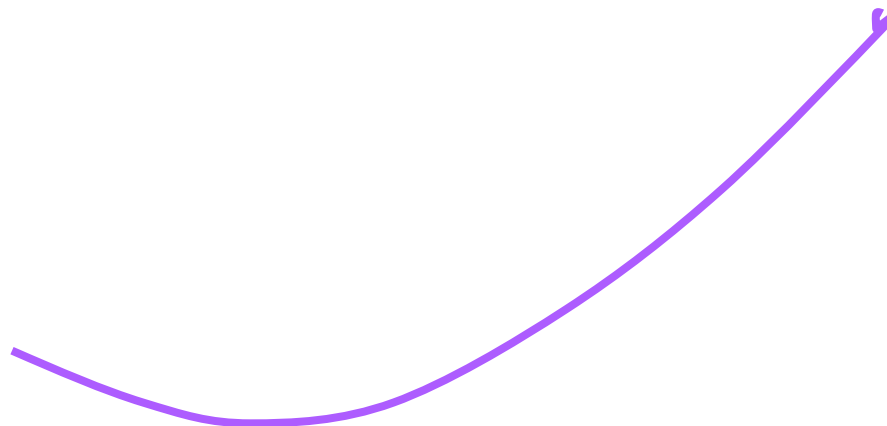
$$1 - 4 = -3$$

$$16 - 25 = -9$$

$W(1, 4) = (7 - 6)^2 = 1$, $W(1, 5) = (10 - 6)^2 = 16$

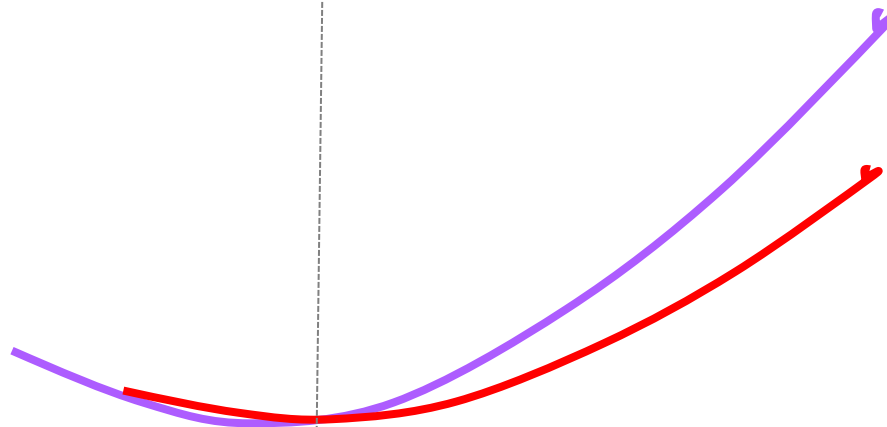
w(i,m) Table

		m								
		0	1	2	3	4	5	6	7	8
i	0		25	4	1	4	25	64	121	196
	1	-	-	9	4	1	16	49	100	169
	2	-	-	-	25	4	1	16	49	100
	3	-	-	-	-	9	0	9	36	81
	4	-	-	-	-	-	9	0	9	36
	5	-	-	-	-	-	-	9	0	9
	6	-	-	-	-	-	-	-	9	0
	7	-	-	-	-	-	-	-	-	9
	8	-	-	-	-	-	-	-	-	-



w(i,m) Table

		m								
		0	1	2	3	4	5	6	7	8
i	0		25	4	1	4	25	64	121	196
	1	-	-	9	4	1	16	49	100	169
	2	-	-	-	25	4	1	16	49	100
	3	-	-	-	-	9	0	9	36	81
	4	-	-	-	-	-	9	0	9	36
	5	-	-	-	-	-	-	9	0	9
	6	-	-	-	-	-	-	-	9	0
	7	-	-	-	-	-	-	-	-	9
	8	-	-	-	-	-	-	-	-	-



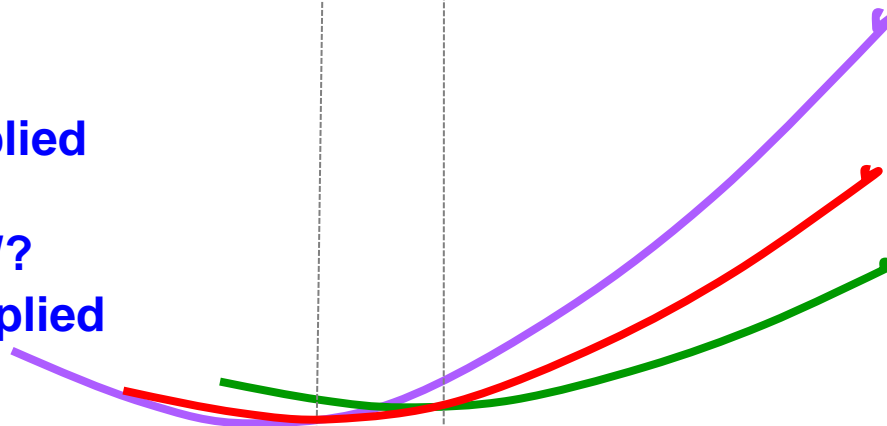
w(i,m) Table

		m								
		0	1	2	3	4	5	6	7	8
i	0		25	4	1	4	25	64	121	196
	1	-	-	9	4	1	16	49	100	169
	2	-	-	-	25	4	1	16	49	100
	3	-	-	-	-	9	0	9	36	81
	4	-	-	-	-	-	9	0	9	36
	5	-	-	-	-	-	-	9	0	9
	6	-	-	-	-	-	-	-	9	0
	7	-	-	-	-	-	-	-	-	9
	8	-	-	-	-	-	-	-	-	-

Remember SMAWK?

Can SMAWK be applied
to find the
column minima of W?

But can SMAWK be applied
to computing $F(n)$?



$$F(m) = \min_{0 \leq i < m} \{ F(i) + w(i, m) \}$$

$F(i)$

$w(i, m)$

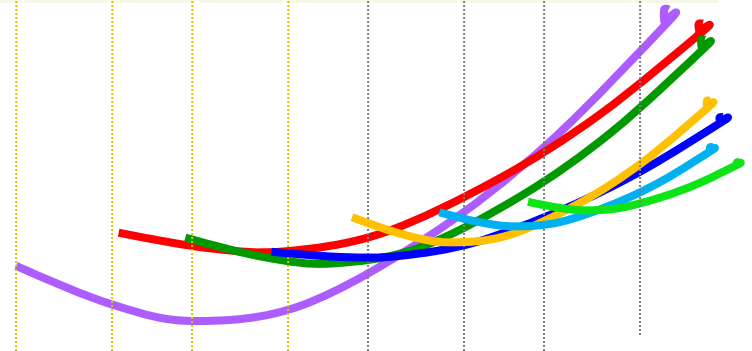
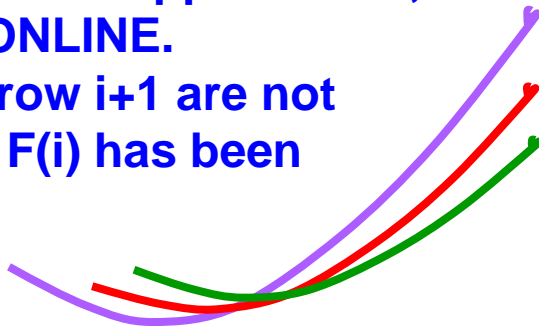
$OUT(i, m) = F(i) + w(i, m)$

		0	1	2	3	4	5	6	7	8
0	0	-	25	4	1	4	25	64	121	196
25	1	-	-	9	4	1	16	49	100	169
4	2	-	-	-	25	4	1	16	49	100
1	3	-	-	-	-	9	0	9	36	81
4	4	-	-	-	-	-	9	0	9	36
1	5	-	-	-	-	-	-	9	0	9
4	6	-	-	-	-	-	-	-	9	0
1	7	-	-	-	-	-	-	-	-	9
	8	-	-	-	-	-	-	-	-	-

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10

SMAWK can not be applied here,
since OUT is ONLINE.

The values of row $i+1$ are not
available until $F(i)$ has been
computed.



W is Monge (concave) and therefore OUT is Totally Monotone (concave)

$OUT[a, c] \geq OUT[b, c] \Rightarrow OUT[a, d] \geq OUT[b, d]$ for all $a < b$ and $c < d$

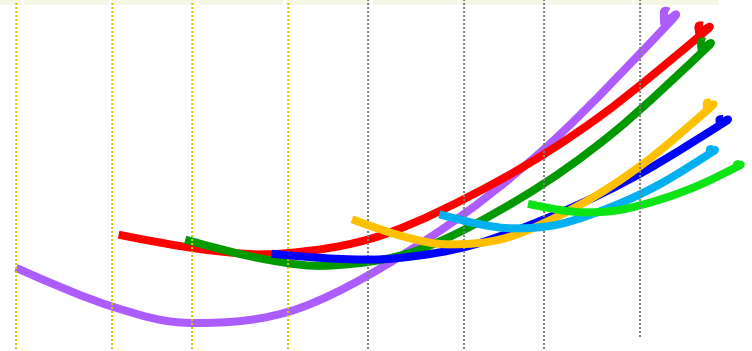
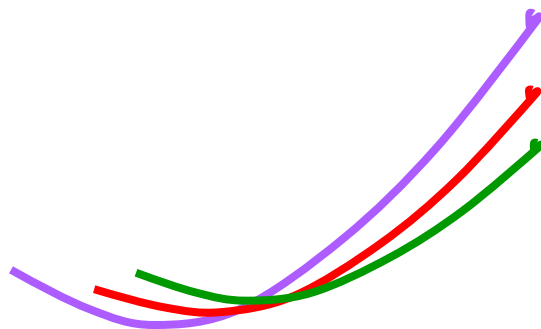
$F(i)$

$w(i, m)$

$OUT(i, m) = F(i) + w(i, m)$

		0	1	2	3	4	5	6	7	8
0	0	-	25	4	1	4	25	64	121	196
25	1	-	-	9	4	1	16	49	100	169
4	2	-	-	-	25	4	1	16	49	100
1	3	-	-	-	-	9	0	9	36	81
4	4	-	-	-	-	-	9	0	9	36
1	5	-	-	-	-	-	-	9	0	9
4	6	-	-	-	-	-	-	-	9	0
1	7	-	-	-	-	-	-	-	-	9
	8	-	-	-	-	-	-	-	-	-

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10

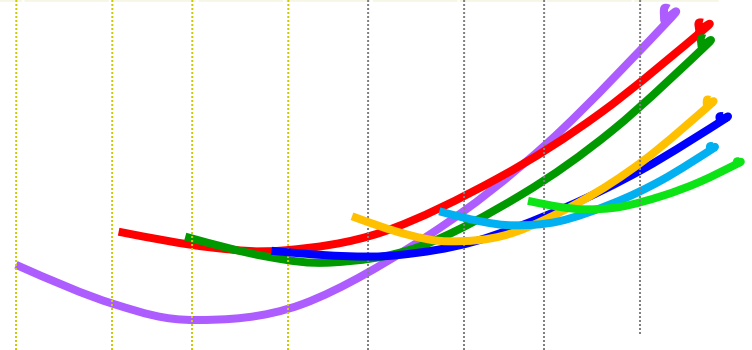


The 1D/1D DP problem:

Definition: *The 1D/1D dynamic programming problem is to find the column minima in an **on-line** upper triangular matrix OUT.*

$$OUT(i, m) = F(i) + w(i, m)$$

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10

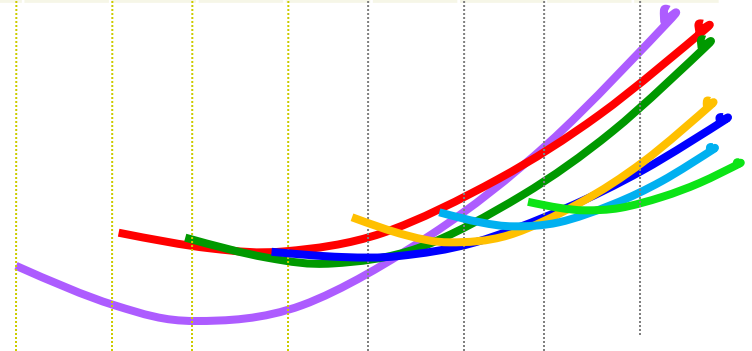


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0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10



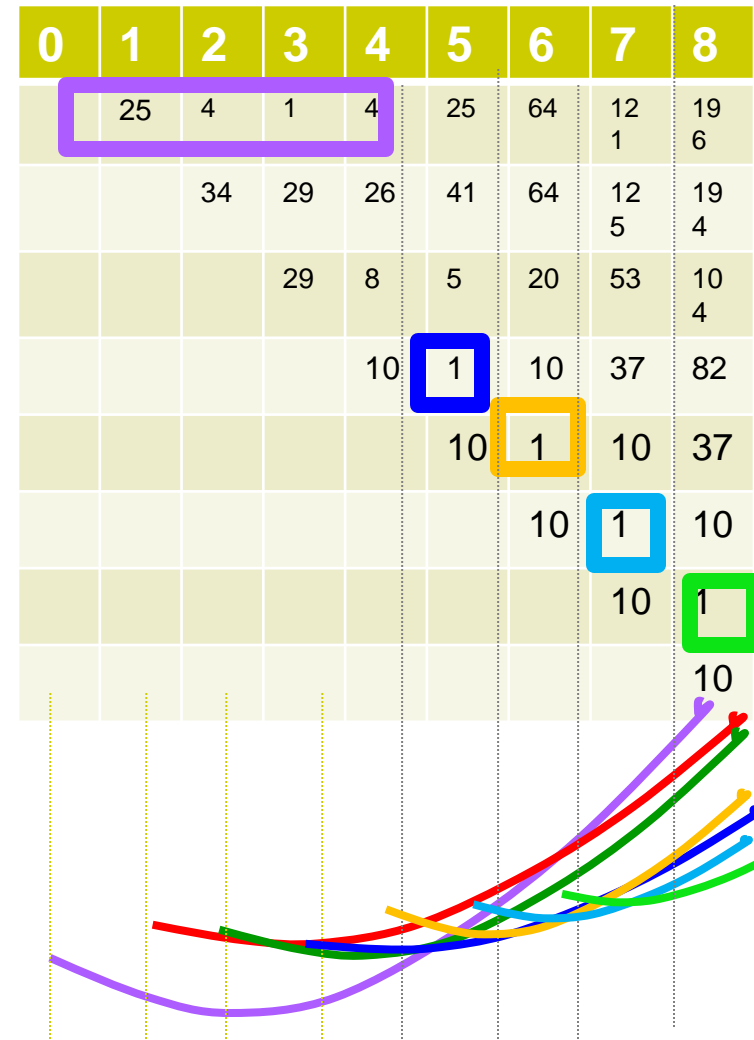
The 1D/1D DP problem:

Procedure 1D/1D

Initialize Queue with row 0

for $m = 2$ to n

{
 find minimum of column m
 update Queue with row $j-1$
}



Example:

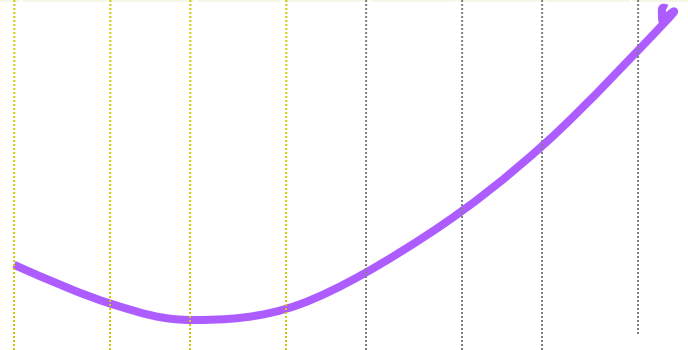
Queue:

$i=1$: $OUT[1,8]$

$m=1$, min in row 1

$$OUT(i, m) = F(i) + w(i, m)$$

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10



Example:

$m=2$, min in row 1

Queue:

$i=1 : OUT[1,8]$

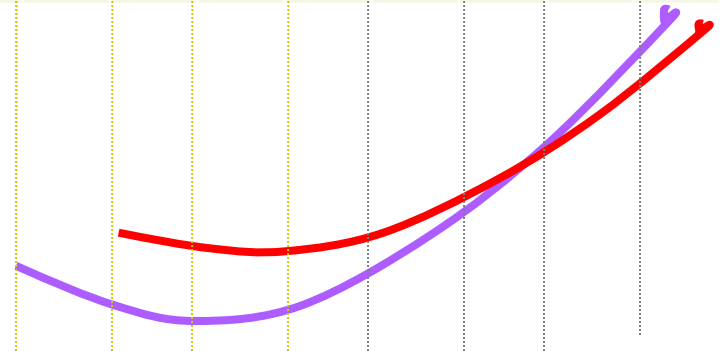
$i=2 : ??$

Where does **row 2**
intersect with **row 1**?

Use binary search:
intersection found
in column 6...

$$OUT(i, m) = F(i) + w(i, m)$$

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10



Example:

Queue:

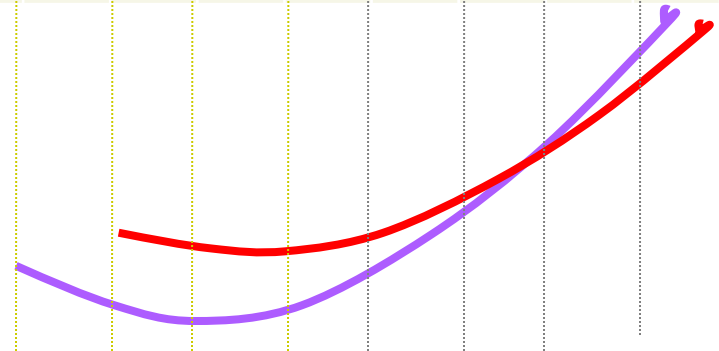
$i=1$: $OUT[1,5]$

$i=2$: $OUT[6,8]$

$m=2$, min in row 1

$$OUT(i, m) = F(i) + w(i, m)$$

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10



Example:

Queue:

$i=1$: $OUT[1,5]$

$i=2$: $OUT[6,8]$

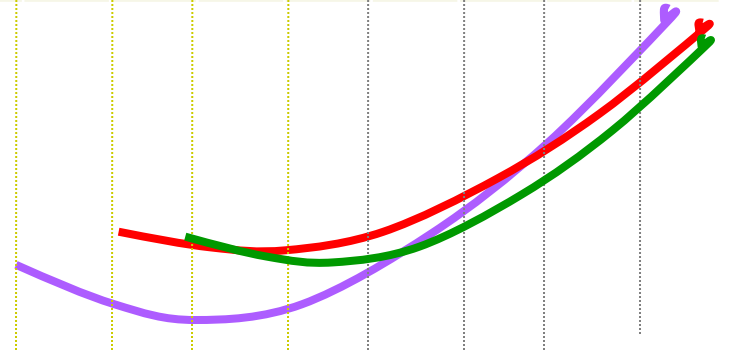
$i=3$:??

Where does **row 3**
intersect with **row 2**?

Use binary search:
intersection at
column 3, even before
Row 2 becomes leader!

$$OUT(i, m) = F(i) + w(i, m)$$

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10



Example:

Queue:

$i=1$: $OUT[1,8]$

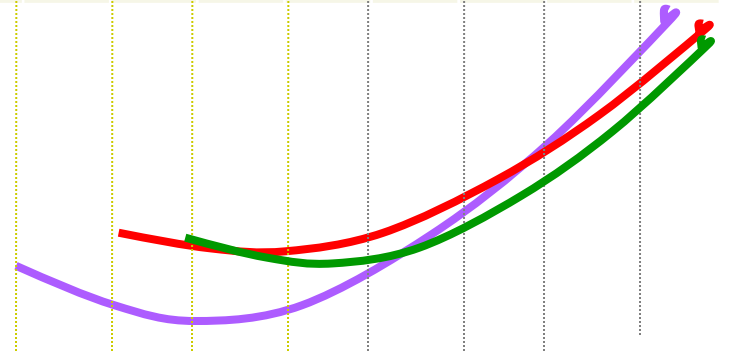
$i=3$: ??

Where does **row 3**
intersect with **row 1**?

Use binary search:
intersection at
Column 5...

$$OUT(i, m) = F(i) + w(i, m)$$

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10



Example:

Queue:

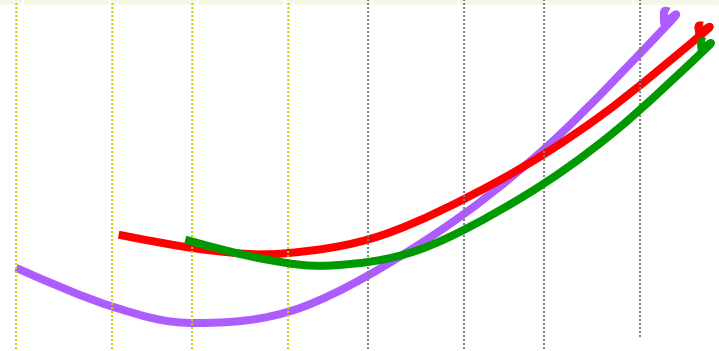
i=1 : **OUT**[1,4]

i=3: **OUT**[5,8]

m=3, min in row 1

$$OUT(i, m) = F(i) + w(i, m)$$

0	1	2	3	4	5	6	7	8
	25	4	1	4	25	64	121	196
		34	29	26	41	64	125	194
			29	8	5	20	53	104
				10	1	10	37	82
					10	1	10	37
						10	1	10
							10	1
								10



Time Complexity Analysis

Each row, for $m = 1..n$, enters the queue at most once, when its column is reached

Each row can get removed from the queue at most once.

Altogether, $O(n)$ rows inserted/removed from the queue.

Computing the “intersection” between the “leadership intervals” of the new row m versus the last row in the queue takes $O(\log n)$ time via binary search.

For each “intersection” computation, either a row is removed from the queue, or the row comparisons for that iteration stop. Therefore, no more than $O(n)$ “intersection” computations.

Total Time Complexity: $O(n \log n)$