

Basic mathematical question: *existence* and *uniqueness*

For $Ax=b$, one solution, no solution or infinitely many?

All start with “Elimination”

Gaussian Elimination

Example:

$$2u + v + w = 5$$

$$4u - 6v = -2$$

$$-2u + 7v + 2w = 9$$

Subtracting multiples of the first equation from the others, so as to eliminate u from the last two equations.

$$2u + v + w = 5$$

$$-8v - 2w = -12$$

$$8v + 3w = 14$$

2: the first pivot

-8: the second pivot

We continue to eliminate v from the third equation.

$$2u + v + w = 5$$

$$-8v - 2w = -12$$

$$1w = 2$$

Pivots: 2, -8, 1

Back-substitution: solve w then v then u .

Breakdown of Gaussian Elimination

- n equations $\rightarrow n$ pivots: nonsingular, only one solution
- a zero appears in a pivot position \rightarrow some problem! May be curable or incurable.
- Curable case (cured by row exchange): nonsingular case

$$u + v + w = \quad \quad u + v + w = \quad \quad u + v + w =$$

$$2u + 2v + 5w = \rightarrow 3w = \rightarrow 2v + 4w =$$

$$4u + 6v + 8w = \quad 2v + 4w = \quad 3w =$$

- Incurable case: singular case

$$u + v + w = \quad \quad u + v + w =$$

$$2u + 2v + 5w = \rightarrow 3w =$$

$$4u + 4v + 8w = \quad 4w =$$

Why is this system singular?

If $3w=6$ and $4w=7$ then inconsistent \rightarrow no solution

If $3w=6$ and $4w=8$ then $u+v = \text{constant}$ \rightarrow many solutions

Cost of Gaussian Elimination

- How many separate arithmetical operations (cost) does elimination require, for n equations in n unknowns?
- **Operations: division and multiply-subtract**
- To produce a zero in the first column: 1 division and $n-1$ multiply-subtract are needed \rightarrow total n operations needed
- There are $n-1$ rows in the first column to be operated on: $n(n-1)=n^2-n$ operations are needed to produce zeros in the first column
- For the second column: $(n-1)^2-(n-1)$ and so on
- There are n columns:

$$(n^2 + \dots + 1^2) - (n + \dots + 1) = \frac{n(n+1)(2n+1)}{6} - \frac{n(n+1)}{2} = \frac{n^3 - n}{3}$$

- When n is large, $n^3/3$ is a pretty good estimate
- For the right side: about $n^2/2$ is required. Why?
- Operations required in back-substitution:

$$1 + 2 + \dots + n = \frac{n(n+1)}{2} \approx \frac{n^2}{2}$$

Why?

- The right side is responsible for a total of n^2 operations
- Minimum operations required? $Cn^{\log_2 7} = Cn^{2.8}$
($\log_2 7 \approx 2.8$). IBM research center has find some the power of n below 2.5!!

Multiplication of a Matrix and a Vector

- Inner product of two vectors:

$$\begin{bmatrix} 2 & 1 & 1 \end{bmatrix} \begin{bmatrix} u \\ v \\ w \end{bmatrix} = [2u + v + w]$$

- **By rows:**

$$Ax = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} \begin{bmatrix} u \\ v \\ w \end{bmatrix} = \begin{bmatrix} 2u + v + w \\ 4u - 6v \\ -2u + 7v + 2w \end{bmatrix}$$

- Let a_{ij} be the entry in the i th row and j th column and x_j be the j th component of x

$$\sum_{j=1}^n a_{ij} x_j \text{ is the } i\text{th component of } Ax$$

Each row is the linear combination of components of x

- **By columns:**

$$Ax = u \begin{bmatrix} 2 \\ 4 \\ -2 \end{bmatrix} + v \begin{bmatrix} 1 \\ -6 \\ 7 \end{bmatrix} + w \begin{bmatrix} 1 \\ 0 \\ 2 \end{bmatrix}$$

Ax is a combination of the columns of A . The coefficients which multiply the columns are the components of x .

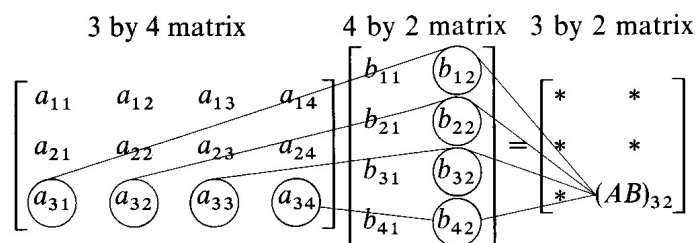
- **Identity matrix I** leaves every vector unchanged.

$$Ib = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix} = \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix}$$

Matrix Multiplication

- Most common and least used: Each entry of AB is the product of a row and a column: $(AB)_{ij}$ = inner product of row i of A and column j of B

Example: Multiplication of matrix $A_{3 \times 4}$ and matrix $B_{4 \times 2}$



- Each column of AB is the product of A and a column of B
Column j of AB = A times column j of B

Example:

$$AB = \begin{bmatrix} 1 & 0 \\ 2 & 3 \end{bmatrix} \begin{bmatrix} a & b \\ c & d \end{bmatrix} = \begin{bmatrix} a & b \\ 2a + 3c & 2b + 3d \end{bmatrix} \Rightarrow$$

Each column of AB is a combination of the columns of A with coefficients coming from each column of B

- Each row of AB is the product of a row of A and B :
Row i of AB = row i of A times $B \Rightarrow$

Each row of AB is a combination of the rows of B with coefficients coming from each row of A

Elimination Steps and Elementary Matrices

- An Elimination step can be a Matrix: Elementary Matrix
HOW? Subtracting 2 times the first row of A from the second row of A

$$EA = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ -2 & 7 & 2 \end{bmatrix}$$

The matrix that subtracts a multiple l_{ij} of row j from row i is the elementary matrix E_{ij} with 1's on the diagonal and the number $-l_{ij}$ in row i , column j .

Example: $E_{31} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -l_{31} & 0 & 1 \end{bmatrix}$

- Gaussian elimination for

$$Ax = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} \begin{bmatrix} u \\ v \\ w \end{bmatrix} = \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix} = b$$

Step 1: Subtract 2 times the 1st equation from the 2nd (E)

Step 2: Subtract -1 times the 1st equation from the 3rd (F)

Step 2: Subtract -1 times the 2nd equation from the 3rd (G)

- Matrix form of Gaussian elimination: $GFEA = U$

$$GFE = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Matrix Multiplication (Cont'd)

- Matrix multiplication is associative $(AB)C=A(BC)$

- Matrix multiplication is distributive

$$A(B+C)=AB+AC \quad \text{and} \quad (B+C)D=BD+CD$$

- Matrix multiplication is not commutative. Usually,

$$AB \neq BA$$

Example: let G be the elementary matrix that adds row 2 to row 3 and E be the matrix that subtracts 2 times row 1 from row 2. Then,

$$EG = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix} \neq GE = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ -2 & 1 & 1 \end{bmatrix}$$

That is, the order of elimination steps matters

Inverse of Elimination Steps

$$Ax = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} \begin{bmatrix} u \\ v \\ w \end{bmatrix} = \begin{bmatrix} 5 \\ -2 \\ 9 \end{bmatrix} = b$$

- Gaussian elimination leads to:

$$GF EA x = Ux = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} u \\ v \\ w \end{bmatrix} = GF Eb = \begin{bmatrix} 5 \\ -12 \\ 2 \end{bmatrix} = c$$

U is called *upper triangular*

After Gaussian elimination $Ax=b \Rightarrow Ux=c$

- Undo the steps of Gaussian elimination:

Example: $E = \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$ then $E^{-1}=?$

$$E^{-1}E = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

If the elementary matrix E has the number $-l_{ij}$ in the (i,j) position, then its inverse has $+l_{ij}$ in that position

- The 3rd step (G) was last in going from A to U . Its matrix G must be the first to be inverted in the reverse direction. That is:

$$E^{-1}F^{-1}G^{-1}U=A$$

Inverse of Elimination Steps in Matrix Form: L

- The matrix that takes U back to A : $L \Rightarrow LU=A$

$$E^{-1}F^{-1}G^{-1} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -1 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -1 & 1 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & -1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1 \end{bmatrix} = L$$

$$\begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \bullet \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ l_{31} & 0 & 1 \end{bmatrix} \bullet \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & l_{32} & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1 \end{bmatrix}$$

\Rightarrow Expressing Inverse of Elimination Steps is EASY!

How about Gaussian Elimination in matrix form? It's NOT EASY!

$$GFE=L^{-1}=$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -l_{32} & 1 \end{bmatrix} \bullet \begin{bmatrix} 1 & 0 & 0 \\ -l_{21} & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \bullet \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -l_{31} & 0 & 1 \end{bmatrix} \neq \begin{bmatrix} 1 & 0 & 0 \\ -l_{21} & 1 & 0 \\ -l_{31} & -l_{32} & 1 \end{bmatrix}$$

Triangular Factorization of A : $A=LU$

- Example:

$$\begin{array}{rrcrcl} 2 & u & + & v & + & w & = \\ & 4 & u & - & 6 & v & = \\ -2 & u & + & 7 & v & + 2 & w & = \end{array}$$

Step 1: Subtract 2 times the 1st equation from the 2nd $\Rightarrow l_{21}=2$

Step 2: Subtract -1 times the 1st equation from the 3rd $\Rightarrow l_{31}=-1$

Step 2: Subtract -1 times the 2nd equation from the 3rd $\Rightarrow l_{32}=-1$

$$GFEA = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} = \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = U$$

$$L=E^{-1}F^{-1}G^{-1} = \begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & -1 & 1 \end{bmatrix}$$

$$A = \begin{bmatrix} 2 & 1 & 1 \\ 4 & -6 & 0 \\ -2 & 7 & 2 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & -1 & 1 \end{bmatrix} \begin{bmatrix} 2 & 1 & 1 \\ 0 & -8 & -2 \\ 0 & 0 & 1 \end{bmatrix} = LU$$

- Triangular Factorization $A=LU$

Without row exchanges required, the original matrix A can be written as a product $A=LU$. The matrix L is lower triangular, with 1's on the diagonal and the multipliers l_{ij} (taken from the elimination) below the diagonal. U is the upper triangular matrix which appears after forward elimination and before back-substitution; its diagonal entries are the pivots.

Examples of Triangular Factorization

- The steps of elimination are actually L^{-1} that takes A to U and reduce L to I : $L^{-1}A = L^{-1}LU = U$.

- Another way of looking at it:

$$\begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1 \end{bmatrix} \begin{bmatrix} \text{row 1 of } U \\ \text{row 2 of } U \\ \text{row 3 of } U \end{bmatrix} = \text{original } A$$

The matrix L , applied to U , brings back A !

Example: When U is the identity then L is the same as A

$$A = \begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ l_{31} & l_{32} & 1 \end{bmatrix} I = \begin{bmatrix} 1 & 0 & 0 \\ l_{21} & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ l_{31} & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & l_{32} & 1 \end{bmatrix} I$$

Example:

$$A = \begin{bmatrix} 1 & -1 & & \\ -1 & 2 & -1 & \\ & -1 & 2 & -1 \\ & & -1 & 2 \end{bmatrix} = \begin{bmatrix} 1 & & & \\ -1 & 1 & & \\ & -1 & 1 & \\ & & -1 & 1 \end{bmatrix} \begin{bmatrix} 1 & -1 & & \\ & 1 & -1 & \\ & & 1 & -1 \\ & & & 1 \end{bmatrix}$$

Solving a Linear System

- c is the right side after Gaussian elimination. That is,

$$c = L^{-1}b \Rightarrow c \text{ can be solved by } Lc = b$$

- When $A = LU$ and L and U are known,

$$Ax = b \Rightarrow \boxed{Lc = b \text{ and } Ux = c}$$

- Given $A = LU$

One Linear System = Two Triangular Systems

- We first solve a lower triangular system for c ($Lc = b$) then the upper triangular system for x ($Ux = c$)
- Multiply $Ux = c$ by L . We have $LUx = Lc \Rightarrow Ax = b$.

Example: Band matrix in the previous example

- U can be further rewritten as:

$$U = \begin{bmatrix} d_1 & & & \\ & d_2 & & \\ & & \ddots & \\ & & & d_n \end{bmatrix} \begin{bmatrix} 1 & u_{12}/d_1 & u_{13}/d_1 & \cdot \\ & 1 & u_{23}/d_2 & \cdot \\ & & \ddots & \cdot \\ & & & 1 \end{bmatrix}$$

The triangular factorization is often written $A = LDU$,

where L and U have 1's on the diagonal

and D is the diagonal matrix of pivots.

- Some freedom in the elimination steps and calculations, but absolutely no freedom in the final L , D and U . That is,

L , D and U are uniquely determined by A

Row Exchanges and Permutation Matrices

- Remember what we do if a zero appears in the pivot location? Problem may be curable or incurable (singular)
- How can the problem be cured?
Looking below the zero to seek out nonzero entry lower down in the same column. Then, a row exchange is carried out; the nonzero entry becomes the needed pivot and elimination can get going again.

- Row exchanges in matrix forms: Permutation matrices**

Example: an exchange of rows 1 and 3 and exchange of rows 2 and 3:

$$P_{13} = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix} \quad P_{23} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

Why? A trick: $P_{13}I = P_{13}$

- Do both of the row exchanges at once:

$$P_{23}P_{13} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} = P$$

- In a nonsingular case, $Ax=b$ has a unique solution
 - It is found by elimination with row exchanges
 - With the rows reordered in advance, PA can be factored into LU .
- In a singular case, no reordering can produce a full set of pivots.

Inverse of A

The matrix A is invertible if there exists a matrix B such that $BA=I$ and $AB=I$. There is at most one such B , called the inverse of A and denoted by A^{-1} :

$$\underline{A^{-1}A=I} \quad \text{and} \quad \underline{AA^{-1}=I}$$

- There could not be two different inverses, because if $BA=I$ and $AC=I$ then

$$B=BI=B(AC)=(BA)C=IC=C$$

- $(A^{-1})^{-1} = A$
- $\begin{bmatrix} a & b \\ c & d \end{bmatrix}^{-1} = \frac{1}{ad-bc} \begin{bmatrix} d & -b \\ -c & a \end{bmatrix}$
- if $A = \begin{bmatrix} d_1 & & \\ & \ddots & \\ & & d_n \end{bmatrix}$ then $A^{-1} = \begin{bmatrix} 1/d_1 & & \\ & \ddots & \\ & & 1/d_n \end{bmatrix}$
- $(A+B)^{-1}=?$ Nothing. No quick solution. But $(AB)^{-1}=?$
$$(AB)^{-1}=B^{-1}A^{-1}$$

Proof:

$$(AB)(B^{-1}A^{-1})=ABB^{-1}A^{-1}=AIA^{-1}=AA^{-1}=I$$

$$(B^{-1}A^{-1})(AB)=BAA^{-1}B^{-1}=BIB^{-1}=BB^{-1}=I$$

- Similarly, $(ABC)^{-1}=C^{-1}B^{-1}A^{-1}$
- Remember $GFEA=U$ and $A=E^{-1}F^{-1}G^{-1}U=LU$!

Gauss-Jordan Method to Find A^{-1}

- We first go from A to U using Gaussian elimination

$$\begin{aligned}
 [A \quad e_1 \quad e_2 \quad e_3] &= \begin{bmatrix} 2 & 1 & 1 & 1 & 0 & 0 \\ 4 & -6 & 0 & 0 & 1 & 0 \\ -2 & 7 & 2 & 0 & 0 & 1 \end{bmatrix} \\
 &\rightarrow \begin{bmatrix} 2 & 1 & 1 & 1 & 0 & 0 \\ 0 & -8 & -2 & -2 & 1 & 0 \\ 0 & 8 & 3 & 1 & 0 & 1 \end{bmatrix} \\
 &\rightarrow \begin{bmatrix} 2 & 1 & 1 & 1 & 0 & 0 \\ 0 & -8 & -2 & -2 & 1 & 0 \\ 0 & 0 & 1 & -1 & 1 & 1 \end{bmatrix} = [U \quad L^{-1}]
 \end{aligned}$$

$$[A \quad I] \rightarrow L^{-1}[A \quad I] = [L^{-1}A \quad L^{-1}I] = [U \quad L^{-1}]$$

- Then from U to I : creating zeros above the pivots

$$\begin{aligned}
 [U \quad L^{-1}] &\rightarrow \begin{bmatrix} 2 & 1 & 0 & 2 & -1 & -1 \\ 0 & -8 & 0 & -4 & 3 & 2 \\ 0 & 0 & 1 & -1 & 1 & 1 \end{bmatrix} \\
 &\rightarrow \begin{bmatrix} 2 & 0 & 0 & \frac{12}{8} & -\frac{5}{8} & -\frac{6}{8} \\ 0 & -8 & 0 & -4 & 3 & 2 \\ 0 & 0 & 1 & -1 & 1 & 1 \end{bmatrix} \\
 &\rightarrow \begin{bmatrix} 1 & 0 & 0 & \frac{12}{16} & -\frac{5}{16} & -\frac{6}{16} \\ 0 & 1 & 0 & \frac{4}{8} & -\frac{3}{8} & -\frac{2}{8} \\ 0 & 0 & 1 & -1 & 1 & 1 \end{bmatrix} = [I \quad U^{-1}L^{-1}]
 \end{aligned}$$

$$[U \quad L^{-1}] \rightarrow U^{-1}[U \quad L^{-1}] = [I \quad U^{-1}L^{-1}] = [I \quad A^{-1}]$$

Operations required for calculating A^{-1} ? $n^3/6 + n^3/3 + n(n^2/2)$ How?

- Are Gauss-Jordan method the only order of elimination?

We could have produced zeros above and below a pivot!

Invertible = Nonsingular

- Nonsingular \rightarrow Invertible

- A matrix has a full set of n pivots \rightarrow nonsingular by definition. (There exists a unique solution for any b)
- $AA^{-1}=I$ can then be viewed as n separate $Ax_i=e_i$ systems where x_i 's are columns of A^{-1} .
- x_i 's can be solved because A is nonsingular and therefore A^{-1} can be determined uniquely.
- $A^{-1}A=I$? A^{-1} is actually the multiplication of many elementary matrices formed by the Gauss-Jordan steps.

- Invertible \rightarrow Nonsingular

- If A has an inverse, Gauss-Jordan method must not break down. That is, the following should not occur.

$$A' = \begin{bmatrix} d_1 & x & x & x & x \\ 0 & d_2 & x & x & x \\ 0 & 0 & 0 & x & x \\ 0 & 0 & 0 & x & x \\ 0 & 0 & 0 & x & x \end{bmatrix}$$

- No matter what x 's are, the Gauss-Jordan method cannot be carried out. That is, A must have a full set of pivots and is nonsingular.

A square matrix is invertible if and only if it is nonsingular

Transposes

- The transpose of $A = A^T$
- The i th row of A becomes the i th column of A^T
- The entry $(A^T)_{ij} = A_{ji}$

$$\left(\begin{bmatrix} 2 & 1 & 4 \\ 0 & 0 & 3 \end{bmatrix} \right)^T = \begin{bmatrix} 2 & 0 \\ 1 & 0 \\ 4 & 3 \end{bmatrix}$$

- $(AB)^T = B^T A^T$

$$AB = \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} 3 & 3 & 3 \\ 2 & 2 & 2 \end{bmatrix} = \begin{bmatrix} 3 & 3 & 3 \\ 5 & 5 & 5 \end{bmatrix}$$

$$B^T A^T = \begin{bmatrix} 3 & 2 \\ 3 & 2 \\ 3 & 2 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} 3 & 5 \\ 3 & 5 \\ 3 & 5 \end{bmatrix}$$

- $(A^{-1})^T = (A^T)^{-1}$ not trivial!

$$(A^{-1})^T A^T = (A A^{-1})^T = I^T = I \quad \text{and} \quad A^T (A^{-1})^T = (A^{-1} A)^T = I^T = I$$

- A symmetric matrix: $A^T = A$ or $a_{ij} = a_{ji}$
- A symmetric matrix's inverse, if there is one, is also symmetric
- If A is symmetric and can be factored into LDU without row exchanges to destroy the symmetry, then the upper triangular U is the transpose of the lower triangular L . That is, the factorization becomes $A = LDL^T$

$$A = LDU = A^T = U^T D^T L^T$$

Roundoff Error

- A typical computer round-off:
 $.345 + .00123 \rightarrow .346$ (keep three digits)
- Roundoff error: the digits lost (0.00023)
- Example:

$$A = \begin{bmatrix} 1. & 1. \\ 1. & 1.0001 \end{bmatrix} \quad \text{and} \quad A' = \begin{bmatrix} .0001. & 1. \\ 1. & 1. \end{bmatrix}$$

- Some matrices are extremely sensitive to small changes and others are not. The matrix A is ill-conditioned (sensitive) and A' is well-conditioned. In A , if a_{ij} (in A) becomes 1 due to roundoff error the system becomes singular

Example

Original: $u +$	$v = 2$	Roundoff: $u +$	$v = 2$
$u + 1.0001v = 2.0001$		$u + 1.0001v = 2$	

solution: $u = v = 1$,	solution: $u = 2 \quad v = 0$
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A change in the fifth digit of b was amplified to a change in the first digit of the solution!!

Partial Pivoting

- Even a well-conditioned matrix can be ruined by a poor algorithm. *Straightforward Gaussian elimination is among the poor algorithms!!*

Example: the well-condition matrix A' (symmetric)

$$.0001u + v = 1$$

$$u + v = 2$$

after elimination: $-9999v = -9998 \rightarrow v = .99990$

original: $.0001u + .9999 = 1 \rightarrow u = 1.0$

round-off: $v = .99990 \rightarrow v \sim 1.0$

$$.0001u + 1 = 1 \rightarrow u \sim 0 !!!$$

$$A' = \begin{bmatrix} 1 & 0 \\ 10,000 & 1 \end{bmatrix} \begin{bmatrix} .0001 & 0 \\ 0 & -9999 \end{bmatrix} \begin{bmatrix} 1 & 10,000 \\ 0 & 1 \end{bmatrix} = LDU$$

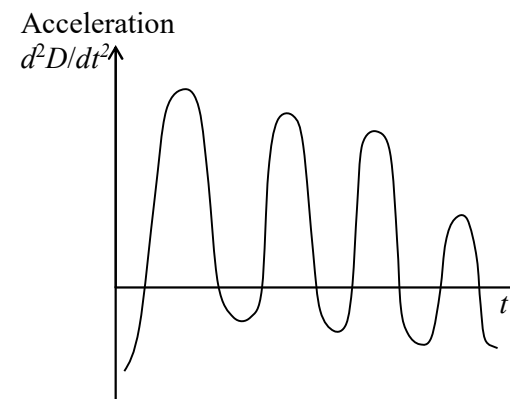
- *The small pivot 0.0001 brought instability!!*
- Remedy: row exchanges to remove small pivots.
- A computer must compare each pivot with all the other possible pivots in the same column. *Choosing the largest of these candidates, and exchanging the corresponding rows so as to make this largest value the pivot is called Partial Pivoting.*

$$A'' = \begin{bmatrix} 1 & 1 \\ .0001 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ .0001 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & .9999 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} = LDU$$

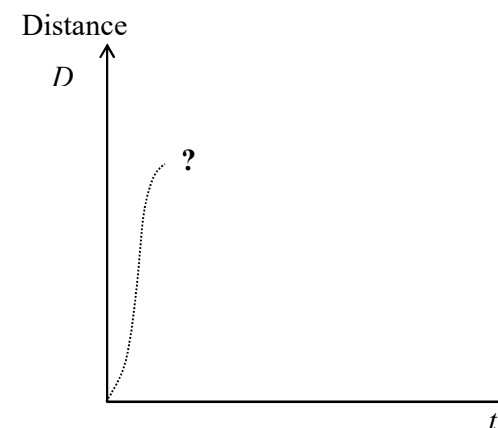
A symmetric A' is not necessarily better for a computer!!

An Application: Differential Equation

- A continuous problem cannot be solved by computers. Computer can only approximate a continuous problem by a discrete problem
- Example: riding motorcycles with only acceleration known as a function of time



Question: distance as a function of time?



Continuous → Discrete

$$\frac{d^2 D}{dt^2} = f(t), \quad 0 \leq t \leq 1 \Rightarrow D(t) = ?$$

- **Uncertainty in the problem** since $a+bt$ for any a and b after second derivative contributes nothing to $f(t)$.
- **Usually, initial conditions**, such as $D(0)=0$ (starting from the origin) and $dD(0)/dt=0$ (initial speed is zero), are added to remove the uncertainty.

- **Discrete approximation of differentiation:**

$$\frac{dD}{dt} \approx \frac{D(t+h) - D(t)}{h} \Rightarrow \frac{d^2 D}{dt^2} = \frac{d(dD/dt)}{dt} \approx \frac{\frac{D(t+h) - D(t)}{h} - \frac{D(t) - D(t-h)}{h}}{h} \approx \frac{D(t+h) - 2D(t) + D(t-h)}{h^2} = f(t)$$

Let time interval = $h \Rightarrow D_{j+1} - 2D_j + D_{j-1} = h^2 f(jh) \quad j=1, \dots, n$

- Let $n=5, D_0=0$ and $dD_0/dt=0 \Rightarrow (D_1 - D_0)/h=0 \Rightarrow D_1 = D_0 = 0$

$$D_0 - 2D_1 + D_2 = D_2 = h^2 f(h)$$

$$D_1 - 2D_2 + D_3 = -2D_2 + D_3 = h^2 f(2h)$$

$$D_{j-1} - 2D_j + D_{j+1} = h^2 f(jh) \quad \text{for } j=3, \dots, 5$$

$$\begin{bmatrix} 1 & 0 & & & \\ -2 & 1 & 0 & & \\ 1 & -2 & 1 & 0 & \\ & 1 & -2 & 1 & 0 \\ & & 1 & -2 & 1 \end{bmatrix} \begin{bmatrix} D_2 \\ D_3 \\ D_4 \\ D_5 \\ D_6 \end{bmatrix} = h^2 \begin{bmatrix} f(h) \\ f(2h) \\ f(3h) \\ f(4h) \\ f(5h) \end{bmatrix} \Rightarrow \text{Lower Triangular}$$

Matrix

Continuous → Discrete

- Let $n=5, D_0=0$ and $dD_5/dt=0 \Rightarrow (D_6 - D_5)/h=0 \Rightarrow$

$$D_0 - 2D_1 + D_2 = -2D_1 + D_2 = h^2 f(h)$$

$$D_{j-1} - 2D_j + D_{j+1} = h^2 f(jh) \quad \text{for } j=2, \dots, 4$$

$$D_4 - 2D_5 + D_6 = D_4 - 2D_5 + D_5 = D_4 - D_5 = h^2 f(5h)$$

$$\begin{bmatrix} -2 & 1 & & & \\ 1 & -2 & 1 & & \\ & 1 & -2 & 1 & \\ & & 1 & -2 & 1 \\ & & & 1 & -1 \end{bmatrix} \begin{bmatrix} D_1 \\ D_2 \\ D_3 \\ D_4 \\ D_5 \end{bmatrix} = h^2 \begin{bmatrix} f(h) \\ f(2h) \\ f(3h) \\ f(4h) \\ f(5h) \end{bmatrix}$$

\Rightarrow **Tri-diagonal Matrix**

- Let $n=5, D_0=0$ and $D_6=0$ (Boundary Conditions)

$$D_0 - 2D_1 + D_2 = -2D_1 + D_2 = h^2 f(h)$$

$$D_{j-1} - 2D_j + D_{j+1} = h^2 f(jh) \quad \text{for } j=2, \dots, 4$$

$$D_4 - 2D_5 + D_6 = D_4 - 2D_5 = h^2 f(5h)$$

$$\begin{bmatrix} -2 & 1 & & & \\ 1 & -2 & 1 & & \\ & 1 & -2 & 1 & \\ & & 1 & -2 & 1 \\ & & & 1 & -2 \end{bmatrix} \begin{bmatrix} D_1 \\ D_2 \\ D_3 \\ D_4 \\ D_5 \end{bmatrix} = h^2 \begin{bmatrix} f(h) \\ f(2h) \\ f(3h) \\ f(4h) \\ f(5h) \end{bmatrix}$$

\Rightarrow **Tri-diagonal Matrix**

Tridiagonal Matrices

$$\begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & -1 & 2 & -1 & \\ & & -1 & 2 & -1 \\ & & & -1 & 2 \end{bmatrix}$$

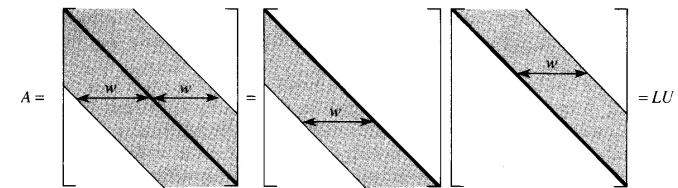
- **Tridiagonal matrix:** $a_{ij}=0$ if $|i-j|>1$; that is, all nonzero entries lie on the main diagonal and the two adjacent diagonals.
- **Symmetric matrix:** $U=L^T$
- **Positive definite matrix:** all pivots are positive. *A symmetric matrix with all positive pivots does not require row exchanges.*
- For a tridiagonal matrix,
 - (a) only one nonzero entry below the pivot
 - (b) operations are carried out on a very short row.

$$A = \begin{bmatrix} 1 & & & & \\ -\frac{1}{2} & 1 & & & \\ & -\frac{2}{3} & 1 & & \\ & & -\frac{3}{4} & 1 & \\ & & & -\frac{4}{5} & 1 \end{bmatrix} = \begin{bmatrix} \frac{2}{1} & & & & \\ & \frac{3}{2} & & & \\ & & \frac{4}{3} & & \\ & & & \frac{5}{4} & \\ & & & & \frac{6}{5} \end{bmatrix} \begin{bmatrix} 1 & -\frac{1}{2} & & & \\ & 1 & -\frac{2}{3} & & \\ & & 1 & -\frac{3}{4} & \\ & & & 1 & -\frac{4}{5} \\ & & & & 1 \end{bmatrix}$$

L and U are *bi-diagonal*. The pivots are converging to +1.

Operations Required by Band Matrices

- For a tridiagonal matrix, every forward elimination stage requires 2 multiply-subtract operations and there are n such stages \rightarrow we need only $2n$ operations instead of the usual $n^3/3$.
- **Back-substitution:** again only $2n$, instead of $n^2/2$, operations are needed.
- No. of operations for a tridiagonal matrix is proportional to n instead of to a higher power of n .



- In general, a band matrix with nonzero entries only on the band $|i-j|<w$. A tridiagonal matrix has $w=2$.
- Each stage of elimination requires $w(w-1)$ operations
- A total of about w^2n operations are needed.
- The exact number of divisions and multiply-subtracts is $w(w-1)(3n-2w+1)/3$ ($=n(n-1)(n+1)/3$ when $w=n$)
- In band matrix problems, calculating A^{-1} to solve $x=A^{-1}b$ is not more efficient than calculating L and U to solve $Lc=b$ and $Ux=c$