

# [Week 4] Path Planning

授課教師:郭重顯教授

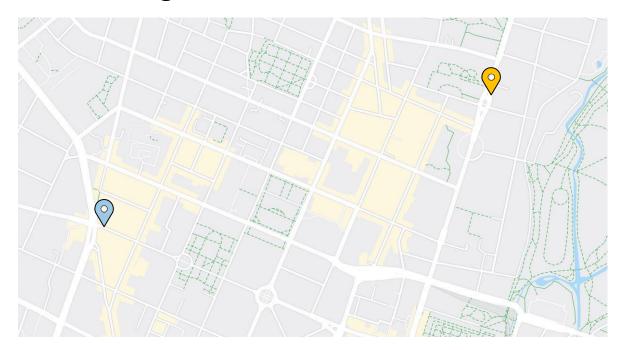
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### **Motivation**



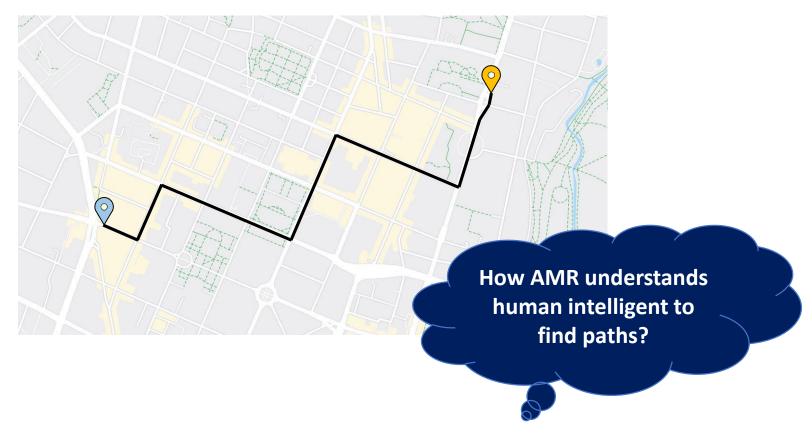
- The human approach to finding routes on a street map is a very visual way
- ❖ Feasible routes/paths are developed by visually tracing different street lines in the general direction of the target



### **Motivation**



- The human approach to finding routes on a street map is a very visual way
- ❖ Feasible routes/paths are developed by visually tracing different street lines in the general direction of the target



### **Definition of terms**



Global Planning	Uses map for planning without information of local environment. Focus: hours to minutes.	
	High-level description of the vehicle motion. Focus: minutes to seconds.	
Local Planning	Consideration of local objects.  Deals with reactive decisions.  Focus: seconds to milliseconds.	Previous path Valid path candidates Invalid path candidates Final path

### **Global vs Local Planning**



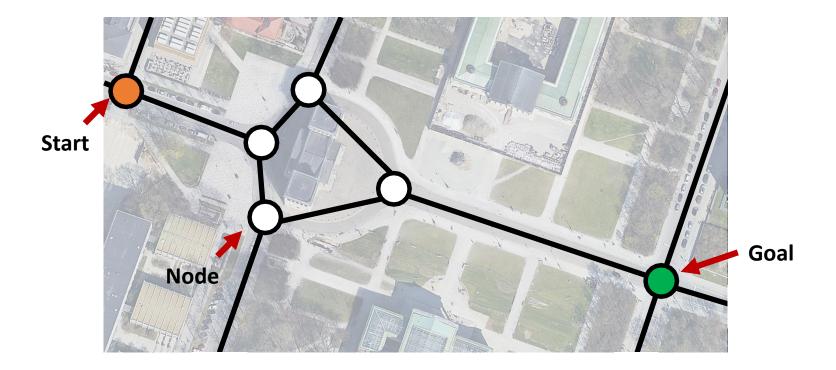
<b>Global Planning</b>	Local Planning
Map based	Sensor based
Relatively slower response	Fast response
Known workspace	Incomplete workspace
Generate path/route before moving	Planning and moving at the same time

# Route/Path/Trajectory Planning



#### **Route Planning**

- ☐ Decision to take a route from source (start) to destination (goal)
- ☐ Sequence of discrete geometrical nodes in map network

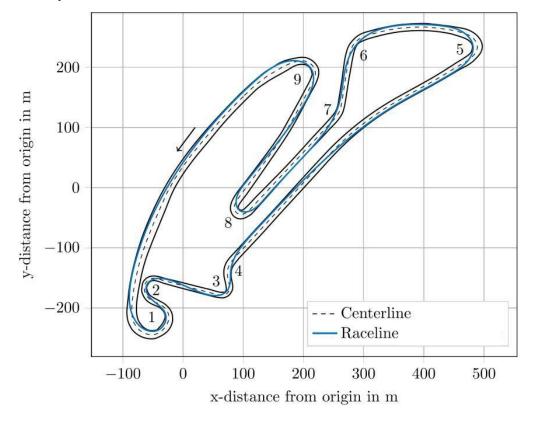


# **Route/Path/Trajectory Planning**



#### Path Planning

- ☐ Continuous curve in spatial domain
- ☐ Consideration of spatial / geometrical boundary conditions possible (e.g., track width, maximum curvature)

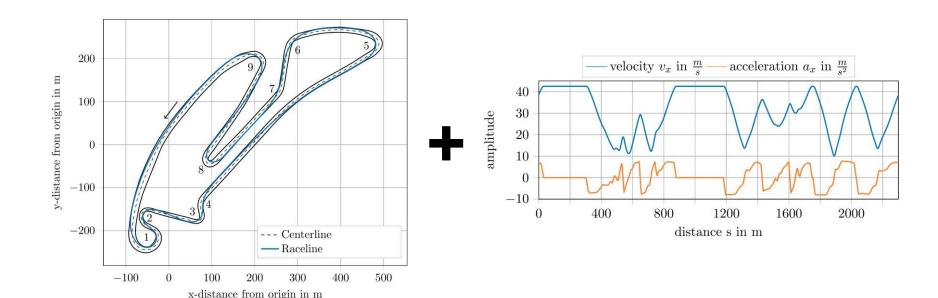


# Route/Path/Trajectory Planning



#### Trajectory Planning

- ☐ Continuous curve in **spatial-temporal** domain
- ☐ Further conditions can be checked by <u>temporal information</u> (e.g., freedom from collisions, compliance with acceleration limits)



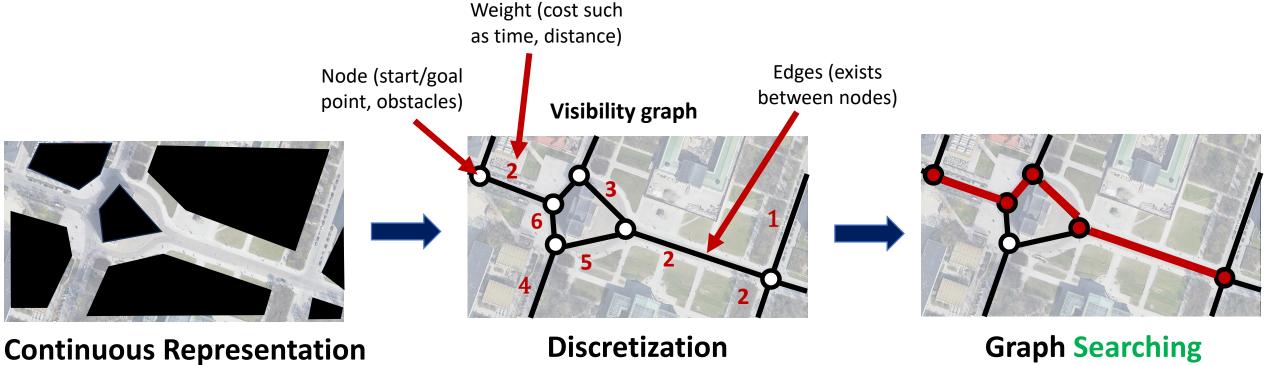
# Global Route Planning vs. Global Path Planning



- Global Route Planning is generally used to generate a route from a source node to a destination node
  - ☐ Certain <u>target</u> variables, such as distance or time, can be <u>minimized</u> in this context
  - ☐ **Exact** geometric curve lines of the vehicle **are not** taken into account
- ❖ Global path planning considers the continuous planning of the vehicle's curve
  - ☐ Spatial and geometrical boundary conditions must be observed (e.g., track width, maximum curvature)

# **Global Planning Framework**





(configuration space formulation) (random sampling, processing critical geometric events)

**Graph Searching** 

(Dijkstra, A\*)



- Dijkstra's Algorithm invented by Edsger W. Dijkstra in 1956
  - ☐ Finding optimal path from source node to destination node
  - ☐ Graph-based algorithm
  - ☐ Informed, complete and optimal path search algorithm

#### **Assumptions**

- ☐ All edges are weighted by costs
- □ No <u>negative</u> weighted costs
- ☐ Path to start node has no costs
- ☐ Starting node has no predecessor
- ☐ For initialization, costs to all other nodes are infinite

Dijkstra, EW 1959, 'A note on two problems in connexion with graphs', Numerische Mathematik, vol. 1, pp. 269-271. https://doi.org/10.1007/BF01386390

You can try: https://www.cs.usfca.edu/~galles/visualization/Dijkstra.html

#### Steps

- 1. Initialization: Starting node=0, other nodes=∞
- 2. Distance of starting node is set to permanent, all other distances are temporarily
- 3. Setting starting node as active
- 4. Calculation of the temporary distances to the current active node:
  - i. Consideration of all reachable neighbor nodes
  - ii. Summing up its distance with the weights of the edges
- 5. If calculated distance is smaller as the current one:
  - i. Update the distance
  - ii. Set the current node as antecessor
- 6. Setting node with minimal temporary distance as **active**. Mark its distance as permanent. .....

Explore all shortest paths until destination is reached

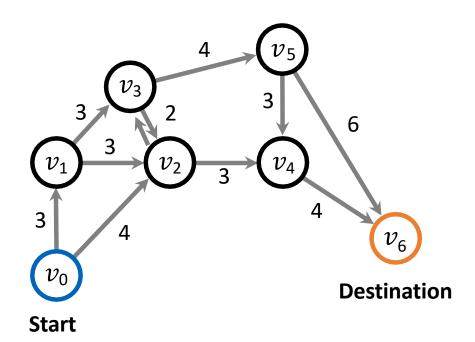
#### **Pseudocode:**

```
function Dijkstra(graph G, source s):
          for each v in G:
             dist[v] \leftarrow \infty
             previous[y] \leftarrow Null
                           u equals s at the 1st time
          dist[s] := 0
   6:
          while G is not empty:
             u := node in G with smallest dist[]
             remove u from G
                                        u is active now
             for each reachable neighbor v of \underline{\mathbf{u}}:
                tmp := dist[u] + dist_between(u, v)
   10:
   11:
                if tmp < dist[v]:
\cdots 12 : \cdots \rightarrow dist[v] := tmp \leftarrow
                                                    5(u) \rightarrow 1(v) feasible (5)
   13:
                   previous[v] := u
                                                    5(u) \rightarrow 7(v) feasible (5)
          return dist[], previous[]
   14:
```

**Hint:** If there are some nodes never reach the goal, and they will go through the final selection of G in step 7 of the minimum distance, and that node will be removed in step 8; however, the steps 9-13 will never be executed because the distance is large/ infinite (beginning setting of **dist[]**).



for each v in G: dist[v]  $\leftarrow \infty$ previous[v]  $\leftarrow Null$ dist[s] := 0



s = 0 in this case

d = distance

P = predecessor c = considered

#### Step 0:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
p[v]	-	-	-	-	-	-	-
c[ <i>v</i> ]	-	-	-	-	-	-	-

#### **Priority Queue:**

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$



```
while G is not empty:
    u := node in G with smallest dist[]
    remove u from G
    for each neighbor v of u:
        alt := dist[u] + dist_between(u, v)
        if alt < dist[v]:
        dist[v] := alt
        previous[v] := u
return dist[], previous[]</pre>
```

Active node stays in *G*; otherwise remove inactive node from *G* 

d = distance

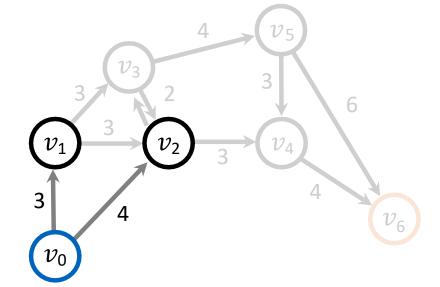
P = predecessor

c = considered

#### Step 1:

	v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
_	d[v]	0	3	4	$\infty$	$\infty$	$\infty$	$\infty$
_	p[ <i>v</i> ]	-	$v_0$	$v_0$	-	-	-	-
	c[ <i>v</i> ]	<b>/</b>	-	-	-	-	-	-
		·						

$\underline{\hspace{1cm}} v$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$	
d[v]	3	4	$\infty$	$\infty$		$\infty$	

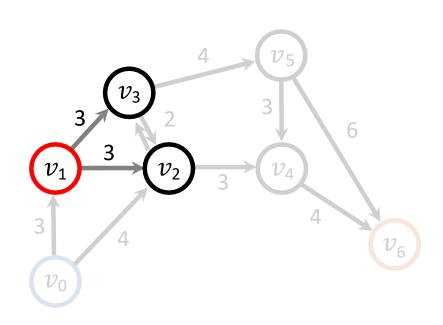




d = distance

P = predecessor

c = considered



#### Step 2:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	3	4	6	$\infty$	$\infty$	$\infty$
p[ <i>v</i> ]	-	$v_0$	$v_0$	$v_1$	-	-	-
c[v]	<b>/</b>	<b>✓</b>	-	-	-	-	-

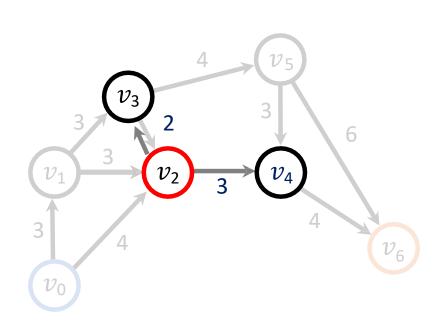
$$3+3=6$$



d = distance

P = predecessor

c = considered



#### Step 3:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	3	4	6	7	$\infty$	$\infty$
p[v]	-	$\overline{v_0}$	$v_0$	$v_1$	$v_2$	-	-
c[ <i>v</i> ]	<b>/</b>	<b>/</b>	<b>/</b>	-	-	-	-

$$\begin{array}{c|ccccc} v & v_3 & v_4 & v_5 & v_6 \\ \hline d[v] & 6 & 7 & \infty & \infty \end{array}$$

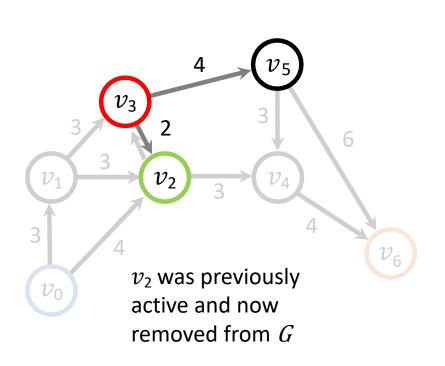
$$4+2 = 6$$



d = distance

P = predecessor

c = considered



#### Step 4:

v			$v_2$				$v_6$
d[v]	0	3	4	6	7	10	$\infty$
p[v]	-	$v_0$	$v_0$	$v_1$	$v_2$	$v_3$	-
c[ <i>v</i> ]	<b>/</b>	<b>/</b>	<b>/</b>	<b>/</b>	-		-

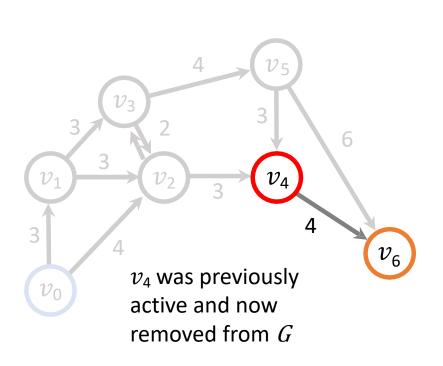
v	$v_4$	$v_5$	$v_6$				
d[v]	7	10	$\infty$				
	6+4 = 10						



d = distance

P = predecessor

c = considered



#### Step 5:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	3	4	6	7	10	11
p[ <i>v</i> ]	-	$v_0$	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$
c[v]	<b>/</b>	<b>/</b>	<b>/</b>	<b>/</b>	<b>/</b>	-	-

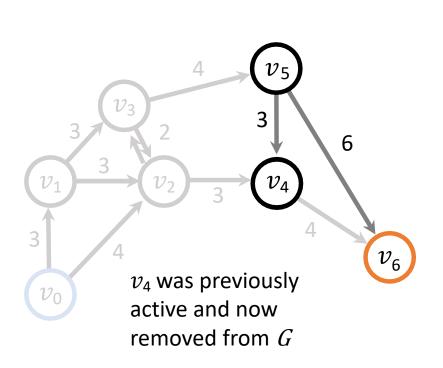
$$7+4 = 11$$



d = distance

P = predecessor

c = considered



#### Step 6:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	3	4	6	7	10	11
p[ <i>v</i> ]	-	$v_0$	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$
c[ <i>v</i> ]	<b>/</b>	<b>/</b>	<b>/</b>	<b>~</b>	<b>/</b>	<b>/</b>	-

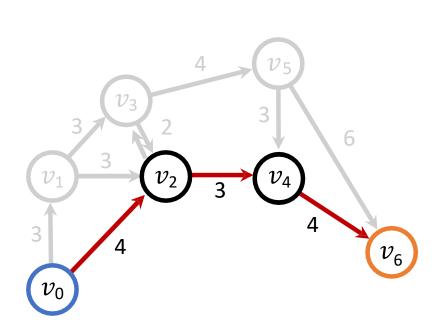
$$\begin{array}{c|c} v & v_6 \\ \hline d[v] & 11 \end{array}$$



d = distance

P = predecessor

c = considered

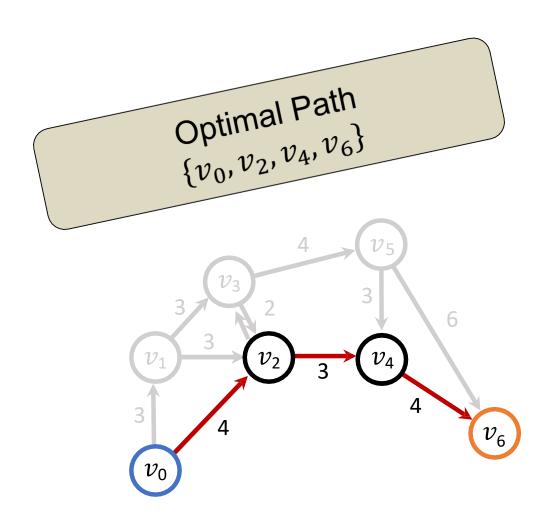


#### **Step 7:**

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	3	4	6	7	10	11
p[ <i>v</i> ]	-	$v_0$	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$
c[v]	<b>/</b>	<b>/</b>	<b>/</b>	<b>/</b>	<b>✓</b>	<b>/</b>	<b>✓</b>

$$\begin{array}{c|cccc} v & - \\ \hline d[v] & - \end{array} \qquad G = \emptyset$$





d = distance

P = predecessor

c = considered

#### **Step 7:**

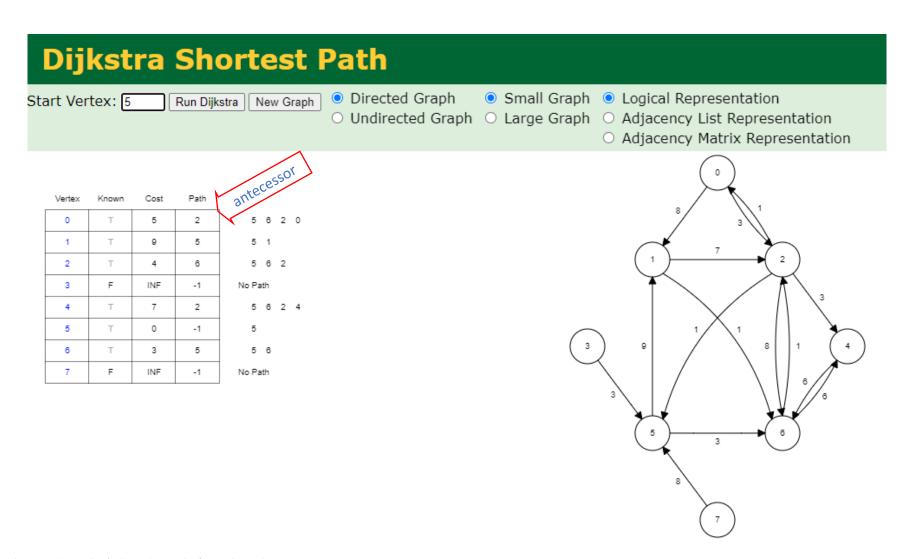
v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
d[v]	0	3	4	`6.	7	10	11
p[ <i>v</i> ]	-	$v_0$	$\dot{v}_0$	$v_1$	$v_2$	$v_3$	$v_4$
c[ <i>v</i> ]	<b>/</b>	<b>/</b>	<b>/</b>	<b>/</b>	<b>✓</b>	<b>/</b>	<b>/</b>

$$\begin{array}{c|c} v & - \\ \hline d[v] & - \end{array}$$

$$G = \emptyset$$

### Dijkstra Visualization, University of San Francisco

You can try: https://www.cs.usfca.edu/~galles/visualization/Dijkstra.html





- Invented by Peter Hart, Nils Nilsson and Bertram Raphael in 1968
- ❖ In literature frequently described as extension of Dijkstra's algorithm
- Informed, complete and optimal path search algorithm

#### Differences to Dijkstra's algorithm

	Inf	formed	heuristic	search	algorithm

	An estimation	n functior	n accelerates	the search	process
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Node costs =	distance to	o the start no	ode + the	estimated	distance to	the de	estination no	ode
INDUC COSES	aistailee ti	o the otal this	Jac - the	Cottillated	aistailee to	CIIC GC	,5(1110(1011 11)	Juc

☐ Node with lowest overa	Il costs has	priority
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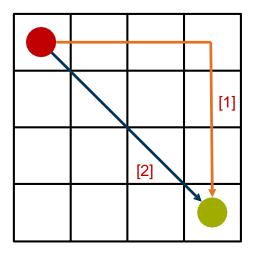
#### Estimation function

- ☐ Free choice of the estimation function
- ☐ Estimator must never overestimate the cost of a route



#### Cost estimation heuristics

- ☐ Exact heuristics → time consuming
- ☐ Approximation heuristics:
  - The Manhattan Distance Heuristic [1]
  - Euclidean Distance [2]



#### Cost estimation function

$$f(v) = g(v) + h(v)$$

f(v): total estimated cost of path through node v

g(v): cost so far to reach v

h(v): estimated cost from v to destination



#### Steps

Create open-list and closed list

Add starting node to open-list

Repeat algorithm until destination is found

- 1. Calculation of the temporary distances of the active node:
  - i. Consideration of all neighbor nodes
  - ii. Summing up its distance with the weights of the edges plus estimated distance to destination
- 2. If calculated distance is smaller than current:
  - i. Update the distance
  - ii. Set the current node as antecessor
  - iii. Add node to open-list
- 3. Add current node to closed-list
- 4. Proceed with minimal temporary distance node as active node

#### Pseudocode:

- 1: make open-list with starting node
- 2: make empty closed-list
- 3: while destination not reached:
- 4: consider node with min f[v]
- 5: for each child u of current node v:
- 6: set child costs f[u]
- 7: if child u is in open-list:
- 8: if g[u] < g[v]:
- 9: continue to line 17
- 10: else if child u is in closed-list:
- 11: if g[u] < g[v]:
- 12: add u to open-list
- 13: **else**:
- 15: add u to open-list
- 16: Set g[u] = successor current cost
- 17: Set prev u = v
- 18: add v to closed-list



#### Steps

Create open-list and closed list

Add starting node to open-list

Repeat algorithm until destination is found

- compare Lower Bound Estimate: 1. Calculation of the temporary distances of the active node:
  - i. Consideration of all neighbor nodes
  - ii. Summing up its distance with the edges plus estimated distar
- 2. If calculated distance
  - i. Update the distance
  - ii. Set the current node as
  - iii. Add node to open-list
- 3. Add current node to closed-list
- 4. Proceed with minimal temporary distance node as active node

#### Pseudocode:

- 1: make open-list with starting node
- 2: make empty closed-list
- ination not reached: 3: whi
  - node with min f[v]
  - whild u of current node v:
    - costs f[u]
    - is in open-list:
  - g[u] < g[v]:
  - continue to line 17
  - else if child u is in closed-list:
- 11: if g[u] < g[v]:
- 12: add u to open-list
- 13: else:
- 15: add u to open-list
- Set g[u] = successor current cost 16:
- 17: Set prev u = v
- add v to closed-list 18:



g=distance so far h=v to destination f=total estimated cost p=predecessor c=considered

#### Step 0:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
g[v]	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
h[v]	-	-	-	-	-	-	-
f[v]	-	-	-	-	-	-	-
p[v]	-	-	-	-	-	-	-
c[ <i>v</i> ]	-	-	-	-	-	-	-

#### **Priority Queue:**

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
f[v]	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$



g=distance so far h=v to destination f=total estimated cost p=predecessor c=considered

#### Step 0:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
g[v]	0	3	4	$\infty$	$\infty$	$\infty$	$\infty$
h[ <i>v</i> ]	10	10	7	-	-	-	-
f[v]	10	13	11	-	-	-	-
p[v]	-	$v_0$	$v_0$	-	-	-	-
c[ <i>v</i> ]	<b>/</b>	-	-	-	-	-	-

#### **Priority Queue:**

v	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
f[v]	13	11	$\infty$	$\infty$	$\infty$	$\infty$

Euclidean Airline Distance



g=distance so far
h=v to destination
f=total estimated cost
p=predecessor
c=considered

#### Step 0:

v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
g[v]	0	3	4	6	7	$\infty$	$\infty$
h[ <i>v</i> ]	10	10	7	9	4	-	-
f[v]	10	13	11	15	11	-	-
p[v]	-	$v_0$	$v_0$	$v_2$	$v_2$	-	-
c[v]	<b>/</b>	-	<b>✓</b>	-	-	-	-

#### **Priority Queue:**

v	$v_1$	$v_3$	$v_4$	$v_5$	$v_6$
f[v]	13	15	11	$\infty$	$\infty$

Euclidean Airline Distance



g=distance so far
h=v to destination
f=total estimated cost
p=predecessor
c=considered

#### Step 0:

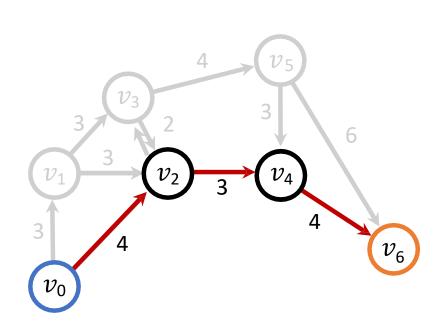
v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
g[v]	0	3	4	6	7	$\infty$	11
h[ <i>v</i> ]	10	10	7	9	4	-	0
f[v]	10	13	11	15	11	-	11
p[v]	-	$v_0$	$v_0$	$v_2$	$v_2$	-	$v_4$
c[v]	<b>/</b>	-	<b>✓</b>	-	<b>✓</b>	-	-

#### **Priority Queue:**

v	$v_1$	$v_3$	$v_5$	$v_6$
f[v]	13	15	$\infty$	11

Euclidean Airline Distance





g=distance so far
h=v to destination
f=total estimated cost
p=predecessor
c=considered

#### Step 0:

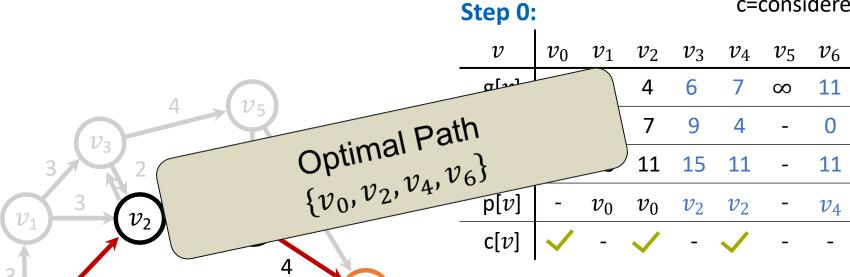
v	$v_0$	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
g[v]	0	3	4	6	7	$\infty$	11
h[ <i>v</i> ]	10	10	7	9	4	-	0
f[v]	10	13	11	15	11	-	11
p[v]	-	$v_0$	$v_0$	$v_2$	$v_2$	-	$v_4$
c[v]	<b>/</b>	-	<b>✓</b>	-	<b>✓</b>	-	-

#### **Priority Queue:**

v	$v_1$	$v_3$	$v_5$	$v_6$
f[v]	13	15	$\infty$	11



g=distance so far
h=v to destination
f=total estimated cost
p=predecessor
c=considered



#### **Priority Queue:**

In order to always obtain an optimal path, all nodes must be visited, or heuristic must not overestimate the cost to reach goal

 $v_0$ 

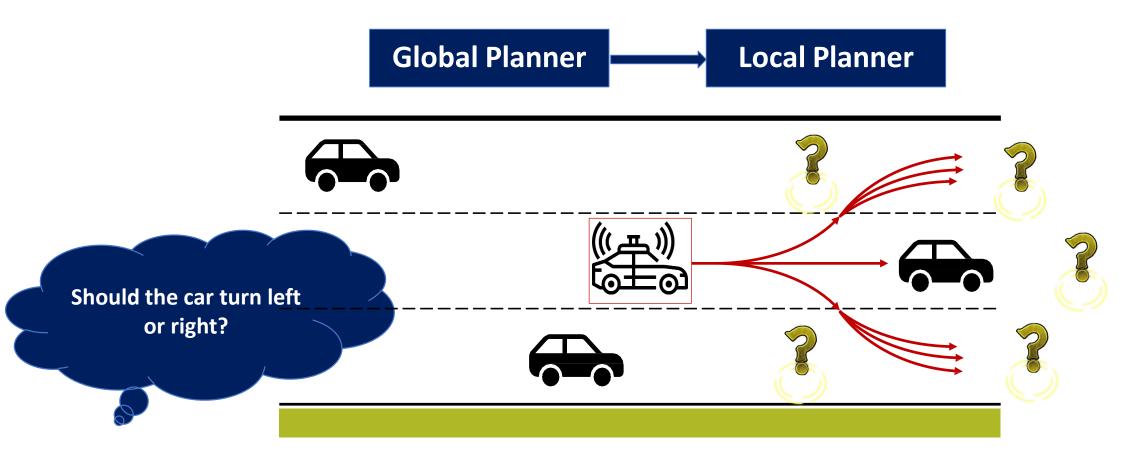


- ❖ A\* converges much **faster** toward the destination
- ❖ A\* explores all partial paths in the order of their potential to reach the destination with a minimum amount of steps
- ❖ A\* continues to explore the graph even **if a feasible path has been found**, because unlike Dijkstra it can not be sure that the first feasible path is an optimal path

### **Behavior Planning**



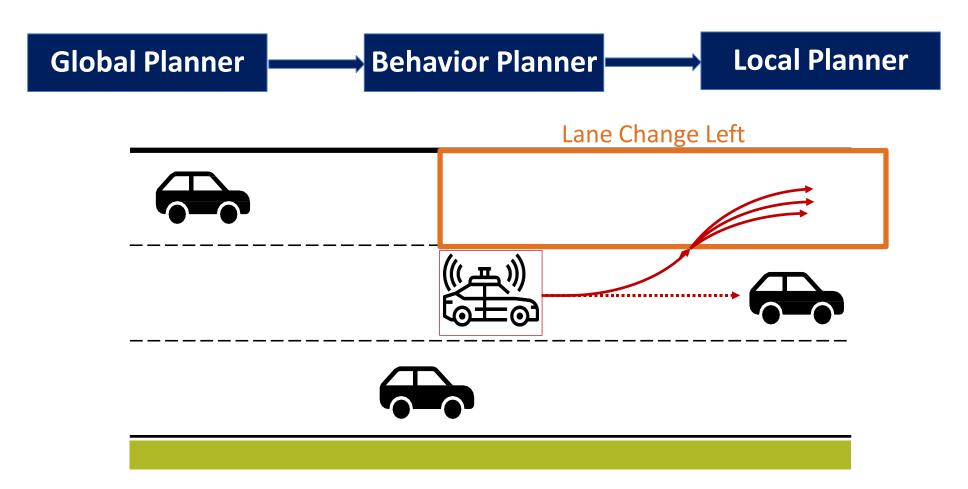
- Fills technical gap between Global Planner and Local Planner
- Considers rules of the road and static/dynamic objects around the vehicle



### **Behavior Planning**



Behavior planning plans high-level driving actions to safely achieve the driving mission under various driving situations



# **Behavior Planning**



- Decision of behavior can be determined by costs
  - ☐ Feasibility costs, security costs, legality costs, comfort costs, speed costs, etc.
- Possible constraints
  - ☐ Global objective, road speed limit, road lane boundaries, stop locations, set of interest vehicles, etc.

## Finite-state Machine (FSM)



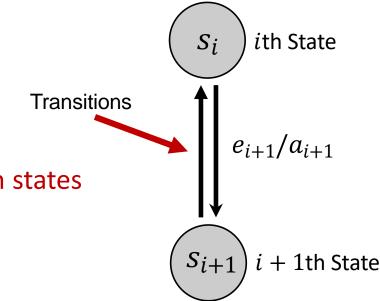
❖ Behavior of the vehicle can be modeled by **Finite-state Machines** 

### Finite-state Machine (FSM)

- Mathematical model of computation
- Consists of finite number of states
- ☐ Behavior of the vehicle can be modeled by transitions between states
- ☐ Transitions based on current state and given input
- Deterministic behavior

#### Predictions

- ☐ There is exactly one state per timestep
- ☐ The time delay at the state transition is irrelevant



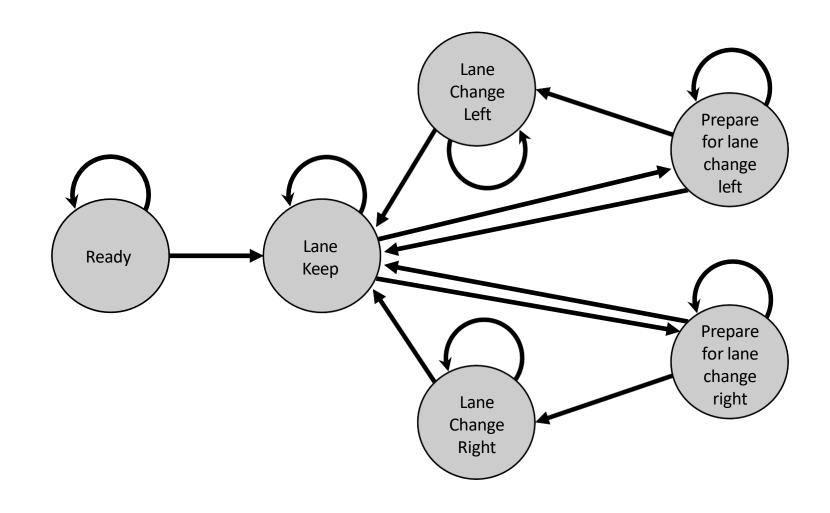
#### States $s_{i+1}$

represent a specific condition or configuration that the machine can be in

**Transition functions** defines the rules/actions  $a_i$  for machine to change from one state to another based on input events  $e_i$ 

# Finite-state Machine (FSM) - An Example





### Finite-state Machine (FSM)



#### Advantages

- ☐ Limiting number of rule checks
- ☐ Clear in structure
- ☐ Easy to calibrate / optimize
- ☐ Simple implementation
- ☐ High flexibility
- ☐ Easy determination of reachability of a state

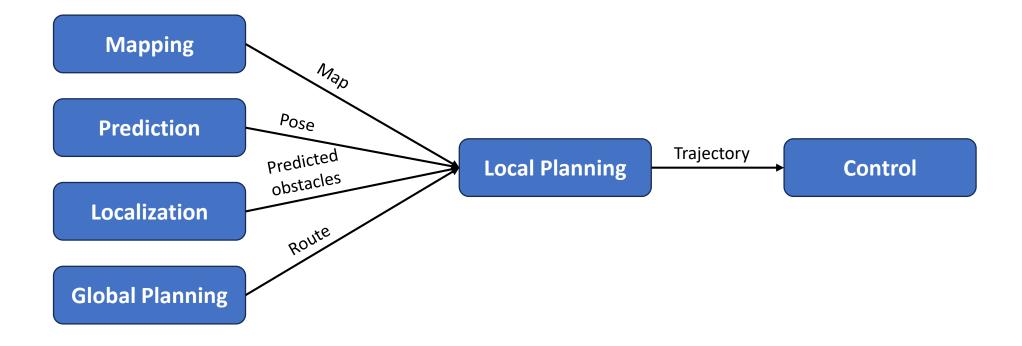
### Disadvantages

- ☐ High knowledge of system design required
- ☐ Large FSMs hard to visualize
- ☐ Difficult transferability to other projects

### **Local Planning**



❖The main task of the local planner is to generate a feasible and collision free trajectory that leads the vehicle in the current local environment and follow the global planning target as good as possible



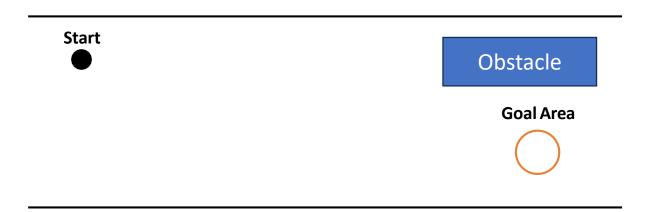


- Proposed by LaValle in 1998
- RRT a simple, iterative algorithm that quickly searches complicated, high-dimensional spaces for feasible paths
- The idea is to incrementally grow a space-filling tree by sampling the space at random and connecting the nearest point in the tree to the new random sample

Credit:
https://lavalle.pl/rrtpubs.ht

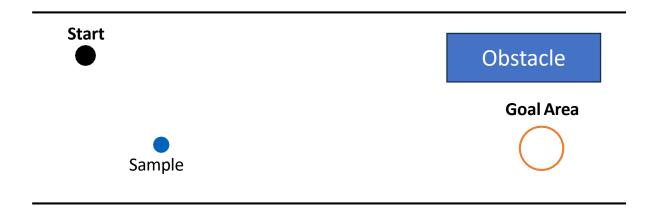
Rapidly-exploring random trees: A new tool for path planning. S. M. LaValle. TR 98-11, Computer Science Dept., Iowa State University, October 1998







- **Repeat until goal are is reached:** 
  - 1. Sample in configuration space





- 1. Sample in configuration space
- 2. Get the nearest node







- 1. Sample in configuration space
- 2. Get the nearest node
- 3. From the nearest node, extend tree in direction of sample point (if new node is collision free)





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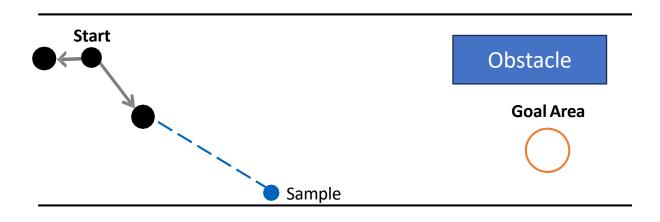
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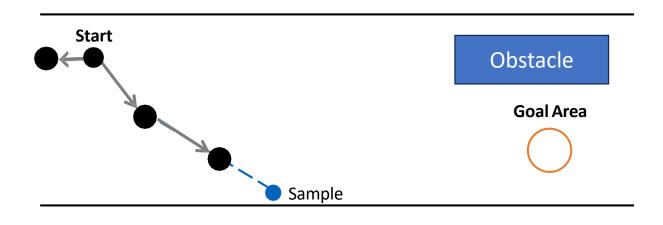
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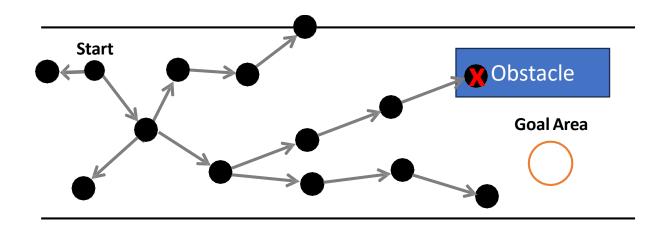
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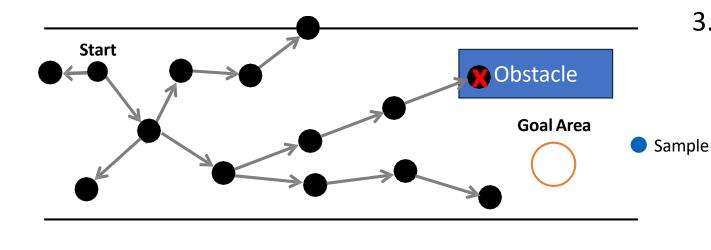
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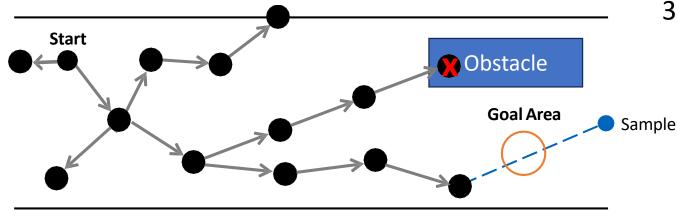
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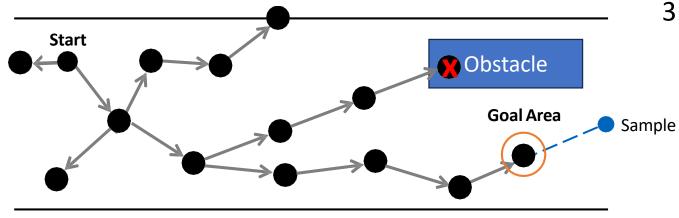
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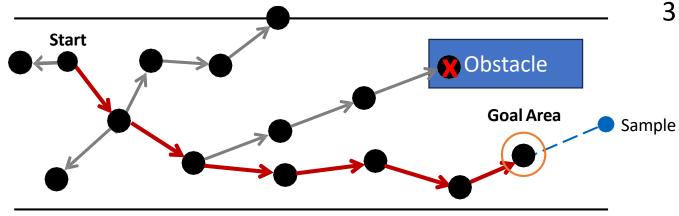
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### **Credit**



- Lectures from CS686: Robot Motion Planning and Applications (KAIST Fall 2013)
- Lectures from Autonomous Driving Software Engineering (TUM)



# Thank you