

SAFE 2.0 User Manual

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Chapter 1

Foreword

1.1 Audience

This document is for users of SAFE (Scalable Analysis Framework for ECMAScript) 2.0, a scalable and pluggable analysis framework for JavaScript web applications. General information on the SAFE project is available at an invited talk at ICFP 2016 [26]:

<https://www.youtube.com/watch?v=gEU9utf0sxE>

and the source code and publications are available at:

<https://github.com/sukyoung/safe>

For more information, please contact the main developers of SAFE at `safe [at] plrg.kaist.ac.kr`.

SAFE has been used by:

- JSAI [11] @ UCSB
- ROSAEC [1] @ Seoul National University
- K framework [21] @ UIUC
- Ken Cheung [4] @ HKUST
- Web-based vulnerability detection [13] @ Oracle
- Tizen [8] @ Linux Foundation

1.2 Contributors

The current developers of SAFE 2.0 are as follows:

- Jihyeok Park
- Youngseo Choi
- Jaemin Hong
- Joonyoung Park
- Sukyoung Ryu

and the following have contributed to the source code:

- Yeonhee Ryou (SAFE 2.0 core)
- Minsoo Kim (Built-in function modeling)
- PLRG @ KAIST and our colleagues in S-Core and Samsung Electronics (SAFE 1.0)

1.3 License

The SAFE source code is released under the BSD license:

github.com/sukyoung/safe/blob/master/LICENSE

1.4 Installation

We assume you are using an operating system with a Unix-style shell (for example, Mac OS X, Linux, or Cygwin on Windows). Assuming `SAFE_HOME` points to the SAFE directory, you will need to have access to the following:

- J2SDK 1.8. See <http://java.sun.com/javase/downloads/index.jsp>
- Scala 2.12. See <http://scala-lang.org/download>
- sbt version 0.13. See <http://www.scala-sbt.org>
- Bash version 2.5, installed at `/bin/bash`. See <http://www.gnu.org/software/bash/>

In your shell startup script, add `$SAFE_HOME/bin` to your path. The shell scripts in this directory are Bash scripts. To run them, you must have Bash accessible in `/bin/bash`.

Type `sbt compile` and then `sbt test` in the SAFE directory to make sure that your installation successfully finishes the tests. Two regression test suites are provided with SAFE and can be analyzed automatically:

```
$ sbt test
$ sbt test262Test
```

In addition to the SAFE-specific test suite, SAFE 2.0 has been tested using Test262, the official ECMAScript (ECMA-262) conformance suite:

<https://github.com/tc39/test262>

Not a single test should end in a failure.

Once you have built the framework, you can call it from any directory, on any JavaScript file, simply by typing one of available commands at a command line as explained in Chapter 3.

1.4.1 IntelliJ configuration

[Thanks to Alexander Jordan @ Oracle Labs in Australia]

IntelliJ users can use IntelliJ 2016.2.4 with the latest Scala plugin as follows:

1. Create a new project from existing sources (aka. **Import project**).
2. Choose **build.sbt** in the SAFE 2.0 root to import.
3. Choose JDK 1.8 as the project JDK.
4. Manually download **xtc.jar** in to **lib/**.
5. Goto **Project Settings** → **Modules** → **root (module)** → **Dependencies**.
6. Open **SBT:unmanaged-jars** dependencies.
7. Remove broken entries for **spray-json** and **xtc**.
8. Add (+) **.jars** for the two libraries above.
9. Run the **buildParsers** task in SBT.

[Thanks to Zhen Zhang @ USTC in China]

For debugging purpose, one can first add a configuration as follows:

1. Go to **Run** → **Edit Configurations**.
2. Click the + button on the left to add a new configuration.
3. Choose **Remote** in the drop-down menu.
4. Name the configuration something like **remote-debugging**. (It should work out of box and the port should be 5005.)
5. Save the configuration by clicking OK.

In order to debug **something.js** for example, launch SAFE first with the **-debug** flag as the first argument: **safe -debug analyze something.js**. The application waits for a connection from the debugger once started, and users can start the debugging configuration created in IntelliJ.

Chapter 2

SAFE

2.1 Introduction to SAFE 1.0

Analyzing real-world JavaScript web applications is a challenging task. On top of understanding the semantics of JavaScript [2], it requires modeling of web documents [27], platform objects [8], and interactions between them. Not only JavaScript itself but also its usage patterns are extremely dynamic [25, 24]. Most of web applications load JavaScript code dynamically, which makes pure static analysis approaches inapplicable.

To analyze JavaScript web applications in the wild mostly statically, we have developed SAFE and extended it with various approaches. We first described quirky language features and semantics of JavaScript that make static analysis difficult and designed SAFE to analyze pure JavaScript benchmarks [14]. It provides a default static analyzer based on the abstract interpretation framework [6], and it supports flow-sensitive and context-sensitive analyses of stand-alone JavaScript programs. It performs several preprocessing steps on JavaScript code to address some quirky semantics of JavaScript such as the `with` statement [17]. The pluggable and scalable design of the framework allowed experiments with JavaScript variants like adding a module system [10, 5] and detecting code clones [4].

We then extended SAFE to model web application execution environments of various browsers [20] and platform-specific library functions [3, 22]. To provide a faithful (partial) model of browsers, we support the configurability of HTML/DOM tree abstraction levels so that users can adjust a trade-off between analysis performance and precision depending on their applications. To analyze interactions between applications and platform-specific libraries specified in Web APIs written in Web IDLs, we developed automatic modeling of library functions from Web APIs and detect possible misuses of Web APIs by web applications. The same technique can support analysis of libraries specified in TypeScript [16]. Analyzing real-world web applications requires more scalable analysis than analyzing stand-alone JavaScript programs [12, 19].

The baseline analysis is designed to be sound, which means that the properties it computes should over-approximate the concrete behaviors of the analyzed program. However, SAFE may contain implementation bugs leading to unsound analysis results. Moreover, some components of SAFE may be intentionally unsound, or soundy [15]. To lessen the burden of analyzing the entire concrete behaviors of programs, we may use approximate call graphs [7] from WALA [9] to analyze a fraction of them, or utilize dynamic information statically [23] to prune relatively unrelated code.

2.2 Introduction to SAFE 2.0

Based on our experiments and experiences with SAFE 1.0, we now release SAFE 2.0, which is aimed to be a playground for advanced research in JavaScript web applications. Thus, we intentionally designed it to be lightweight, highly parametric, and modular.

The important changes from SAFE 1.0 include the following:

- SAFE 2.0 has been tested using Test262, the official ECMAScript (ECMA-262) conformance suite.
- SAFE 2.0 now uses sbt instead of ant to build the framework.
- SAFE 2.0 provides a library of abstract domains that supports parameterization and high-level specification of abstract semantics on them.
- Most Java source files are replaced by Scala code and the only Java source code remained is the generated parser code.
- Several components from SAFE 1.0 may not be integrated into SAFE 2.0. Such components include interpreter, concolic testing, clone detector, clone refactoring, TypeScript support, Web API misuse detector, and several abstract domains like the string automata domain.

We have the following roadmap for SAFE 2.0:

- SAFE 2.0 will make monthly updates.
- The next update will include a SAFE document, browser benchmarks, and more Test262 tests.
- We plan to support some missing features from SAFE 1.0 incrementally such as a bug detector, DOM modeling, and jQuery analysis.
- Future versions of SAFE 2.0 will address various analysis techniques, dynamic features of web applications, event handling, modeling framework, compositional analysis, and selective sensitivity among others.

2.3 A sample use of SAFE

Let us consider a very simple JavaScript program stored in a file name “sample.js” located in the current directory:

```
with({a: 1}) {a = 2;}
```

Then, one can see how SAFE desugars the `with` statement by the command below:

```
safe astRewrite sample.js
```

which shows an output like the following:

```
{
  <>alpha<>1 = @ToObject({
    a : 1
  });
  ("a" in <>alpha<>1 ? <>alpha<>1.a = 2 : a = 2);
}
```

The command ‘astRewrite’ took 100 ms.

where the names prefixed by `<>` are generated by SAFE. SAFE translates the rewritten JavaScript source code to its intermediate representation format, and one can see the result by the command below:

```
safe compile sample.js
```

which shows an output like the following:

```
{
  {
    <>new0<>1 = {
      a : 1
    }
    <>Global<>ignore1 = @ToObject(<>new0<>1)
    <>alpha<>2 = <>Global<>ignore1
  }
  if("a" in <>alpha<>2)
  {
    <>obj<>3 = @ToObject(<>alpha<>2)
    <>obj<>3["a"] = 2
    <>Global<>ignore2 = <>obj<>3["a"]
  }
  else
  {
    a = 2
    <>Global<>ignore2 = 2
  }
}
```

The command ‘compile’ took 123 ms.

The SAFE analysis is performed on control flow graphs of programs, which can be built by the command below:

```
safe cfgBuild sample.js
```

resulting an output as follows:

```
function[0] top-level {
  Entry[-1] -> [0]

  Block[0] -> [2], [1], ExitExc
```

```
[0] noop(StartOfFile)
[1] <>new0<>1 := alloc() @ #1
[2] <>new0<>1["a"] := 1
[3] <>Global<>ignore1 := @ToObject(<>new0<>1) @ #2
[4] <>alpha<>2 := <>Global<>ignore1
```

```
Block[1] -> [3], ExitExc
[0] assert("a" in <>alpha<>2)
[1] <>obj<>3 := @ToObject(<>alpha<>2) @ #3
[2] <>obj<>3["a"] := 2
[3] <>Global<>ignore2 := <>obj<>3["a"]
```

```
Block[2] -> [3], ExitExc
[0] assert(! "a" in <>alpha<>2)
[1] a := 2
[2] <>Global<>ignore2 := 2
```

```
Block[3] -> Exit
[0] noop(EndOfFile)
```

```
Exit[-2]
```

```
ExitExc[-3]
```

```
}
```

The command ‘cfgBuild’ took 160 ms.

Finally, the following command:

```
safe analyze sample.js -analyzer:exitDump
```

analyzes the JavaScript program in the file and shows the analysis results:

```
** heap **
R#1 -> Internal Properties:
  [[Class]] -!> "Object"
  [[Extensible]] -!> true
  [[Prototype]] -!> R#0Object.prototype
  DEFAULT: (âĽĚElem, Top(absent))
Normal Properties:
  a -!> [ttt] 2
  DEFAULT: (âĽĚDataProp, Top(absent))
```

```
...
```

```
** context **
```

```
...
```

```
#PureLocal -> ...
  <>obj<>3 |-> [t]R#1
```

```
...
```

```
this: R#Global
```

```
** old allocation site set **
```

```
mayOld: (#1)
```

```
mustOld: (#1)
```

```
- # of iteration: 6
```

```
- # of user functions: 1
```

```
- # of touched blocks: 6
```

```
  user blocks: 6
```

```
  modeling blocks: 0
```

```
- # of instructions: 13
```

The command ‘analyze’ took 406 ms.

Chapter 3

Reference manual

We describe SAFE commands and their basic usage.

3.1 SAFE commands

Each command consists of phases and one can run a SAFE command as follows:

```
safe {command} [-{option}]*  
    [-{phase}:{option} [{input}]]* {filename}+
```

For example, The command `compile` consists of six phases; `parser`, `astRewriter`, and `compiler`. The following command compiles JavaScript code stored in a file name “sample.js” located in the current directory without showing detailed information but printing the result of the `cfgBuilder` phase into a file name “out”:

```
safe analyze -silent  
    -cfgBuilder:out=out sample.js
```

Each phase has its own available options. The most common options are as follows:

- `-{phase}:silent`
SAFE does not show messages during the phase.
- `-{phase}:out={string}`
SAFE writes the result of the phase to a file `{string}`.

and the global options are `-silent` and `-testMode`.

The currently supported commands and their phases are as follows:

<code>parse</code>	<code>parser</code>
<code>astRewrite</code>	<code>parse >> astRewriter</code>
<code>compile</code>	<code>astRewrite >> compiler</code>
<code>cfgBuild</code>	<code>compile >> cfgBuilder</code>
<code>heapBuild</code>	<code>cfgBuild >> heapBuilder</code>
<code>analyze</code>	<code>heapBuild >> analyzer</code>
<code>bugDetect</code>	<code>analyze >> bugDetector</code>
<code>web</code>	<code>web</code>
<code>help</code>	<code>helper</code>

and phases and their options are as follows:

- `parser -out={string}`
parses the JavaScript code and generates an Abstract Syntax Tree (AST) in given files. If multiple files are given, they are concatenated in the given order before being parsed.
- `astRewriter -silent
 -out={string}`
rewrites a given AST into a simplified AST.
- `compiler -silent
 -out={string}`
generates an Intermediate Representation (IR) from a given AST.
- `cfgBuilder -silent
 -out={string}`
generates a Control Flow Graph (CFG) from a given IR.
- `heapBuilder -silent
 -maxStrSetSize={number}
 -aaddrType={string}
 -callsiteSensitivity={number}
 -loopIter={number}
 -loopDepth={number}
 -snapshot={string}
 -number={string}
 -jsModel`

builds initial abstract heap used in the `analyzer` phase and attaches it into a given CFG.

If `-maxStrSetSize={number}` is given, the analyzer will use the string set domain with given size limit `n`.

If `-aaddrType={string}` is given, the analyzer will use a given address abstraction type; “normal”: allocation-site abstraction and “recency”: recency abstraction with allocation-site abstraction.

If `-callsiteSensitivity={number}` is given, number-depth callsite-sensitive analysis will be executed.

If `-loopIter={number}` is given, number-iteration loop-sensitive analysis will be executed.

If `-loopDepth={number}` is given, number-depth loop-sensitive analysis will be executed.

If `-snapshot={string}` is given, analysis with an initial heap generated from a dynamic snapshot (*.json) will be executed.

If `-jsModel` is given, analysis with JavaScript models will be executed.

- **analyzer** `-silent`
`-console`
`-exitDump`
`-out={string}`

analyzes the given CFG with the given initial abstract heap.

If `-console` is given, a user can debug analysis results by investigating the intermediate status of the analysis.

If `-exitDump` is given, it shows the abstract state of the top-level exit node.

- **bugDetector** `-silent`
detects possible bugs in a given analysis result.
- **web** `-port={number}`
allows a user to investigate the analysis status from a browser based on network connection. It starts web server with the given port number (default: 8080). We describe this in Section 3.2.2
- **help**
shows the usage of SAFE commands to the standard output.

The **parser** phase parses the JavaScript code in a given file and rewrites obvious dynamic code generation into other statements without using dynamic code generation but with the same semantics. For example, the following JavaScript code

```
function f() { return 3; }
eval("f()")
```

is rewritten as follows:

```
function f() { return 3; }
f();
```

The **astRewriter** phase rewrites its AST representation into a simpler AST using three kinds of AST transformations:

- **Hoister** lifts the declarations of functions and variables inside programs and functions up to the beginning of them.
- **Disambiguator** checks some static restrictions and renames identifiers to unique names.
- **WithRewriter** rewrites the `with` statements that do not include any dynamic code generation such as `eval` into other statements without using the `with` statement but with the same semantics [17, 18].

3.2 SAFE analyzer debugging

When the `-analyzer:console` option is given to the **analyze** command, SAFE provides a REPL-style console debugger. For example, the following command:

```
safe analyze -analyzer:console test.js
```

shows the list of available commands for debugging and the starting point of the analysis:

Command list:

- help	
- next	jump to the next iteration. (same as "")
- jump	Continue to analyze until the given iteration.
- print	Print out various information.
- result	Print out various information.
- run_insts	Run instruction by instruction.
- move	Change a current position.
- home	Reset the current position.
- run	Run until meet some break point.
- break	Add a break point.
- break-list	Show the list of break points.
- break-rm	Remove a break point.
- find-bot	Find the instruction whose result is bottom.
- restart	Restart the analysis.

For more information, see 'help <command>'.

```
<top-level[0]: Entry[-1], NoLoop> @sample.js:1:1
Iter[0] >
```

The current status is denoted as follows:

```
<{fun-name}[{fid}]: {block-kind}[{bid}],
  {sensitivity}> @{filename}:{span}
Iter[{#iteration}] >
```

where fun-name and fid are the name and the id of the current function, respectively, block-kind and bid are the kind and the id of the current block, respectively, sensitivity is the sensitivity of the current analysis, filename is the name of the file being analyzed, span is the location of the current analysis, and #iteration is the iteration number of the current analysis. We support call-context sensitivity based on call-sites and loop sensitivity based on loop iterations and inner-loop depths as analysis sensitivities.

A block is one of the following kinds:

- **Entry**: the entry block of a function
- **Exit**: the exit block of a function
- **ExitExc**: a block denoting uncaught exceptions in a function
- Call: a block denoting a function call
- **AfterCall**: a block receiving a return value of a function call
- **AfterCatch**: a block receiving uncaught exceptions after a function call

- **Block**: a normal block with instructions
- **LoopHead**: a block denoting a loop head
- **LBreak**: a block right after the **break** statement
- **LCont**: a block right after the **continue** statement
- **Branch**: a block denoting a branch
- **Switch**: a block denoting a head of **switch** statement
- **Case**: a block denoting a **case** clause in a **switch** statement
- **Default**: a block denoting the **default** clause in a **switch** statement
- **Try**: a block denoting a head of **try** block
- ~~Finally: a block denoting a finally block~~
当前没有实现，但是不重要
- **Catch**: a block denoting a **catch** block
- **Label(string)**: a block denoting a user label {string}
- **ModelBlock**: a block denoting a modeled function

The **help** command displays a list of available commands and the **help <command>** command displays the usage of the <command>. For example:

```
<function[0] top-level: Entry[-1], ()> @test.js:1:1
Iter[0] > help print
usage: print state(-all) ({keyword})
       print heap(-all) ({keyword})
       print context ({keyword})
       print block ({fid}:{bid})
       print loc {LocName} ({keyword})
       print func ({fid})
       print function ({fid})
       print worklist
       print ipsucc
       print trace
```

shows the usage of the **print** command.

The **next** command proceeds the analysis of the current block, which is the default command. For example:

```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] >
```

```
<top-level[0]: Block[0], NoLoop> @test.js:2:1-8:20
```

The **jump #iteration** command proceeds the analysis until the given number of iterations. For example:

```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] > jump 10
```

```
<f[1]: Block[4], NoLoop> @test.js:5:8-19
Iter[10] >
```

The **print** command displays the status just before analyzing the current block. We describe it in Section 3.2.1.

The **result** command displays the status after analyzing the current block:

- **result (exc-)state(-all) ({keyword})**

It displays the state in the same way as the **print** command does, and it can additionally show the exception state generated after the analysis.

- **result (exc-)loc {LocName}**

It finds and displays the location in the same way as the **print** command does, and it can additionally find and display the location from the exception state generated after the analysis.

The **run_insts** command shows the list of instructions in the current block, and it enables to analyze each instruction. It opens a sub-console, which provides 3 kinds of commands:

- **s** shows the state
- **q** quits the analysis
- **n** analyzes the next instruction; the default command

For example:

```
<top-level[0]: Block[0], NoLoop> @test.js:2:1-8:20
Iter[1] > run_insts
Block[0] -> [1], ExitExc
[0] f := function (1) @ #6, #7
[1] noop(StartOfFile)
[2] x := 1
[3] <>obj<>9 := @ToObject(f) @ #8
[4] <>arguments<>10 := allocArg(0) @ #9
[5] <>fun<>11 := @GetBase(f)
[6] <>this<> := enterCode(<>fun<>11)

inst: [0] f := function (1) @ #6, #7
('s': state / 'q': stop / 'n',': next)
>
```

```
inst: [1] noop(StartOfFile)
('s': state / 'q': stop / 'n',': next)
>
```

The **move {fid}:{bid}|entry|exit|exitExc** command moves the current block to the given block denoted by the id of a function, the id of a block, and the kind of the block. For example:

```
<top-level[0]: Block[0], NoLoop> @test.js:2:1-8:20
Iter[1] > move 0:entry
* current control point changed.
```

```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[1] >
```

The **home** command moves the current block back to the original block to be analyzed. For example:


```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[1] > home
* reset the current control point.

<top-level[0]: Block[0], NoLoop> @test.js:2:1-8:20
Iter[1] >
```

The `run` command proceeds the analysis until encountering a break point. A short-key for this command is `Ctrl-d`. For example:

```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] > break 0:exit

<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] > run

<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] >
```

The `break` command sets up a break point at the given block. For example:

```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] > break 0:exit

<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] > run

<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] >
```

The `break-list` command shows a list of blocks with break points. For example:

```
top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] > break-list
* 1 break point(s).
[0] function[0] Exit[-2]

<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] >
```

The `break-rm {break-order}` command removes the break point of a given block denoted by the order in the result of `break-list`. For example:

```
<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] > break-list
* 1 break point(s).
[0] function[0] Exit[-2]

<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] > break-rm 0
* break-point[0] removed.
[0] function[0] Exit[-2]

<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] > break-list
* no break point.

<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] >
```

The `find-bot` command finds the instruction that generates the bottom abstract state. For example:

```
<top-level[0]: Block[2], NoLoop> @branch.js:2:5-12
Iter[3] > find-bot
The result of the following instruction is bottom:
[0] assert(x !== 0)

<top-level[0]: Block[2], NoLoop> @branch.js:2:5-12
Iter[3] >
```

The `restart` command restarts the analysis. For example:

```
<top-level[0]: Exit[-2], NoLoop> @test.js:10:1
Iter[14] > restart

<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] >
```

3.2.1 Analyzer debugging with printing

The `print` command displays the status just before analyzing the current block.

```
print state(-all) ({keyword})
```

The `print state` command displays the current state, and the `print state-all` command displays the current state including all system addresses. When a keyword is given, it displays only the parts that include the keyword. For example:

```
<top-level[0]: Exit[-2], NoLoop> @test.js:31:1
Iter[15] > print state result
__result1 -!> [ttf] UInt
__result2 -!> [ttf] UInt
__result3 -!> [ttf] UInt
```

```
print heap(-all) ({keyword})
```

The `print heap` command displays the heap of the current state in the same way as `print state`.

```
print context(-all) ({keyword})
```

The `print heap` command displays the context of the current state in the same way as `print state`.

```
print block ({fid}{bid})
```

The `print block` command displays the information of the current block. For example:

```
<top-level[0]: Block[0], NoLoop> @test.js:2:1-10:20
Iter[1] > print block
span: test.js:2:1-10:20
Block[0] -> [1], ExitExc
[0] f := function (1) @ #6, #7
[1] noop(StartOfFile)
[2] x := 1
[3] <>obj<>9 := @ToObject(f) @ #8
[4] <>arguments<>10 := allocArg(0) @ #9
[5] <>fun<>11 := @GetBase(f)
[6] <>this<> := enterCode(<>fun<>11)
```

Moreover, it displays the information of the given block denoted by the id of a function, and the id or the kind of the block. For example:

```
<top-level[0]: Block[0], NoLoop> @test.js:2:1-10:20
Iter[1] > print block 1:0
span: test.js:4:8-5:19
Block[0] -> [1], ExitExc
[0] <>ff<>1 := function (2) @ #1, #2
[1] <>x<>2 := 2
[2] <>obj<>5 := @ToObject(<>ff<>1) @ #3
[3] <>arguments<>6 := allocArg(0) @ #4
[4] <>fun<>7 := @GetBase(<>ff<>1)
[5] <>this<> := enterCode(<>fun<>7)
```

print loc {LocName} ({keyword})

The `print loc {LocName}` command shows the object bound at a given location in the current state. When a keyword is given, it displays only the parts that include the keyword in the object. For example:

```
<top-level[0]: Exit[-2], NoLoop> @test.js:2:1
Iter[2] > print loc R#1
R#1 -> Internal Properties:
[[Class]] -!> "Object"
[[Extensible]] -!> true
[[Prototype]] -!> R#Object.prototype
DEFAULT: (â&#x2013;Elem, Top(absent))
Normal Properties:
DEFAULT: (â&#x2013;DataProp, Top(absent))
```

print func ({fid})

It displays the list of functions. If a function id is given, it displays the name and the span of it. For example:

```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] > print func 1
* function name: f
* span info. : test.js:1:1-15
```

print function ({fid})

It displays the detailed control flows in the current function. If a function id is given, it displays that of the corresponding function. For example:

```
<top-level[0]: Entry[-1], NoLoop> @test.js:1:1
Iter[0] > print function 1
function[1] f {
  Entry[-1] -> [0]

  Block[0] -> Exit, ExitExc
    [0] x := 1

  Exit[-2]

  ExitExc[-3]
}
```

print worklist

It shows the work in the current worklist. For example:

```
<top-level[0]: Block[1], NoCall x NoLoop> @test.js:0:0
Iter[2] > print worklist
* Worklist set
(0:Block[1], NoCall x NoLoop)
(0:Block[2], NoCall x NoLoop)
```

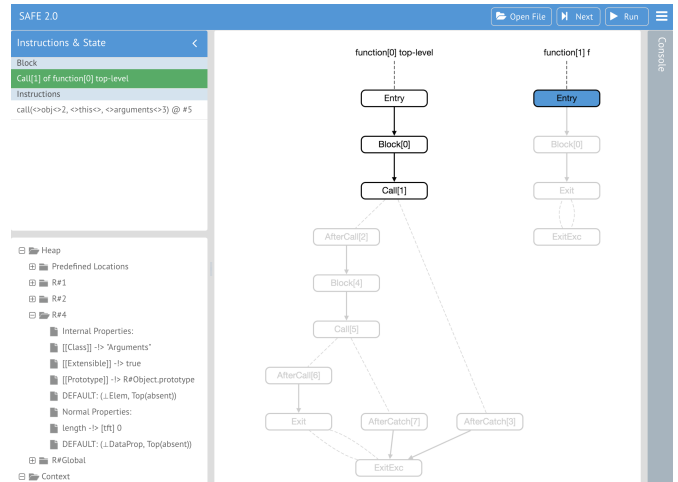
print ipsucc

It displays the information of the current inter-procedural successor. For example:

```
<f[1]: Exit[-2], NoLoop> @test.js:1:15
Iter[17] > print ipsucc
* successor map
- src: (1:Exit[-2], NoLoop)
- dst:
  (0:AfterCall[2], NoLoop), mayOld: (#4, #1, #5, #2),
  mustOld: (#4, #1, #5, #2)
  (0:AfterCall[6], NoLoop), mayOld: (#4, #7, #1, #5,
  #8, #2), mustOld: (#4, #7, #1, #5, #8, #2)
```

3.2.2 Analyzer debugging with browsing

The `safe web` command starts a web server to allow users to investigate the analysis status from a browser. Consider the following snapshot:



which shows the current CFG in the middle. Nodes in black lines denote the blocks that are analyzed, those in gray lines denote the blocks not yet being analyzed, and colored nodes denote the blocks that are currently in the worklist of the analyzer. One can toggle whether to show the nodes in the worklist by the menu button on the top right. When a user selects a block from the CFG, the list of the instructions in the block and the state just before analyzing the block are displayed on the left. Moreover, users can use all functionalities of the console debugger through the console screen in the right side.

Appendix A

Appendix

This appendix provides the definitions of AST, IR, and translation rules from AST to IR in SAFE.

A.1 AST

This section describes each construct of the JavaScript language in both the BNF notation and its corresponding implementation. The implementation of AST nodes is available at:

`$SAFE_HOME/src/main/scala/kr/ac/kaist/safe/nodes/ast/`

```

p ::= fd* vd* s*
      Program(body: TopLevel)
      TopLevel(fds: List[FunDecl], vds: List[VarDecl],
               stmts: List[SourceElements])

fd ::= function f((x,)* {fd* vd* s*}
      FunDecl(ftn: Functional, strict: Boolean)
      Functional(fds: List[FunDecl],
                vds: List[VarDecl],
                stmts: SourceElements, name: Id,
                params: List[Id], body: String)

vd ::= x (= e)?
      VarDecl(name: Id, expr: Option[Expr],
              strict: Boolean)

s ::= {s*}
      ABlock(stmts: List[Stmt], internal: Boolean)
      | var vd(, vd)*;
      | VarStmt(vds: List[VarDecl])
      | ;
      | EmptyStmt()
      | e;
      | ExprStmt(expr: Expr, internal: Boolean)
      | if (e) s (else s)?
      | If(cond: Expr, trueBranch: Stmt,
          falseBranch: Option[Stmt])

```

```

switch (e) {cc* (default:s*)? cc*}
Switch(cond: Expr, frontCases: List[Case],
       defopt: Option[List[Stmt]],
       backCases: List[Case])

do s while (e);
DoWhile(body: Stmt, cond: Expr)

while (e) s
While(cond: Expr, body: Stmt)

for (e?; e?; e?) s
For(init: Option[Expr], cond: Option[Expr],
    action: Option[Expr], body: Stmt)

for (lhs in e) s
ForIn(lhs: LHS, expr: Expr, body: Stmt)

for (var vd(, vd)*; e?; e?) s
ForVar(vars: List[VarDecl], cond: Option[Expr],
       action: Option[Expr], body: Stmt)

for (var vd in e) s
ForVarIn(vd: VarDecl, expr: Expr, body: Stmt)

continue x?;
Continue(target: Option[Label])

break x?;
Break(target: Option[Label])

return e?;
Return(expr: Option[Expr])

with (e) s
With(expr: Expr, stmt: Stmt)

l : s
LabelStmt(label: Label, stmt: Stmt)

throw e;
Throw(expr: Expr)

try {s*} (catch(x) {s*})? (finally {s*})?
Try(body: List[Stmt], catchBlock: Option[Catch],
    fin: Option[List[Stmt]])

catch (id: Id, body: List[Stmt])
Catch(id: Id, body: List[Stmt])

debugger;
Debugger()

cc ::= case e : s*
      Case(cond: Expr, body: List[Stmt])

e ::= e, e
      ExprList(exprs: List[Expr])
      | e ? e : e
      | Cond(cond: Expr, trueBranch: Expr,
              falseBranch: Expr)
      | e ⊗ e
      | InfixOpApp(left: Expr, op: Op, right: Expr)
      | ⊖ e
      | PrefixOpApp(op: Op, right: Expr)
      | lhs ⊙
      | UnaryAssignOpApp(lhs: LHS, op: Op)
      | lhs ⊙ e
      | AssignOpApp(lhs: LHS, op: Op, right: Expr)
      | lhs
      | LHS()

```

```

lhs ::= lit
      Literal()
      | x
      VarRef(id: Id)
      | [(e?,)*]
      ArrayExpr(elements: List[Option[Expr]])
      | {(m,)*}
      ObjectExpr(members: List[Member]
      (e)
      Parenthesized(expr: Expr)
      | function x?((x,)* {fd* vd* s*}
      FunExpr(ftn: Functional)
      | lhs[e]
      Bracket(obj: LHS, index: Expr)
      | lhs.x
      Dot(obj: LHS, member: Id)
      | new lhs
      New(lhs: LHS)
      | lhs((e,)*
      FunApp(fun: LHS, args: List[Expr])

lit ::= this
      This()
      | null
      Null()
      | true
      Bool(bool: Boolean)
      | false
      Bool(bool: Boolean)
      | num
      DoubleLiteral(text: String, num: Double)
      | str
      IntLiteral(intVal: BigInteger, radix: Integer)
      | str
      StringLiteral(quote: String, escaped: String)
      | reg
      RegularExpression(body: String, flag: String)

m ::= pr : e
    Field(prop: Property, expr: Expr)
    | get pr() {fd* vd* s*}
    GetProp(prop: Property, ftn: Functional)
    | set pr(x) {fd* vd* s*}
    SetProp(prop: Property, ftn: Functional)

pr ::= x
      PropId(id: Id)
      | str
      PropStr(str: String)
      | num
      PropNum(num: NumberLiteral)

⊙ ::= = | * | / | % | + | - | [ | > | >> | & | ^ | = |
⊗ ::= && | || | & | ^ | [[ | >> | >>> | + | - | * | /
      | % | == | != | === | !== | [ | > | [= | >=
      | instanceof | in
⊖ ::= ++ | -- | ~ | ! | + | - | delete | void | typeof
⊙ ::= ++ | --

```

Note that after the `Hoister` transformation of the `astRewriter` phase, ASTs should have the following invariants:

- The `expr` field of `VarDecl` should be `None`.
- `VarStmt`, `ForVar`, and `ForVarIn` should be removed.
- `StmtUnit` is an internally generated statement unit.

A.2 IR

This section describes each construct of the SAFE IR in both the BNF notation and its corresponding implementation. The implementation of IR nodes is available at:

`$SAFE_HOME/src/main/scala/kr/ac/kaist/safe/nodes/ir/`

```

p ::= s*
    IRoot(val fds: List[IRFunDecl],
          val vds: List[IRVarStmt],
          val irs: List[IRStmt])

s ::= x = e
    IRExprStmt(val lhs: IRId, val right: IRExpr,
               val ref: Boolean)
    | x = delete x
    IRDelete(val lhs: IRId, val id: IRId)
    | x = delete x[x]
    IRDeleteProp(val lhs: IRId, val obj: IRId,
                 val index: IRExpr)
    | x[x] = e
    IRStore(val obj: IRId, val index: IRExpr,
            val rhs: IRExpr)
    | x = {(m,)*}
    IRObject(val lhs: IRId, val members: List[IRMember],
             val proto: Option[IRId])
    | x = [(e,)*]
    IRArray(val lhs: IRId,
            val elements: List[Option[IRExpr]])
    | IRArgs(val lhs: IRId,
            val elements: List[Option[IRExpr]])
    | x = x(x, x)
    IRCall(val lhs: IRId, val fun: IRId,
           val thisB: IRId, val args: IRId)
    | x = x(x, x)?
    IRInternalCall(val lhs: IRId, val fun: IRId,
                   val first: IRExpr,
                   val second: Option[IRId])
    | toObject, toNumber, isObject, getBase, iteratorInit,
    iteratorHasNext, iteratorKey
    | x = new x((x,)*
    IRNew(val lhs: IRId, val fun: IRId,
          val args: List[IRId])
    | x = function f(x, x) {s*}
    IRFunExpr(val lhs: IRId, val fun: IRFunctional)
    | IRFunctional(val fromSource: Boolean,
                  val name: IRId,
                  val params: List[IRId],
                  val args: List[IRStmt],
                  val fds: List[IRFunDecl],
                  val vds: List[IRVarStmt],
                  val body: List[IRStmt])
    | function f(x, x) {s*}
    IRFunDecl(val ftn: IRFunctional)

```

```

|    $\underline{x}$  = eval( $\underline{e}$ )
|   IREval(val lhs: IRId, val arg: IRExpr)
|   break  $\underline{x}$ 
|   IRBreak(val label: IRId)
|   return  $\underline{e}$ ?
|   IRReturn(val expr: Option[IRExpr])
|   with ( $\underline{x}$ )  $\underline{s}$ 
|   IRWith(val id: IRId, val stmt: IRStmt)
|    $\underline{l}$  : {  $\underline{s}$  }
|   IRLabelStmt(val label: IRId, val stmt: IRStmt)
|   var  $\underline{x}$ 
|   IRVarStmt(val lhs: IRId, val fromParam: Boolean)
|   throw  $\underline{e}$ 
|   IRThrow(val expr: IRExpr)
|    $\underline{s}^*$ 
|   IRSeq(val stmts: List[IRStmt])
|   if ( $\underline{e}$ ) then  $\underline{s}$  (else  $\underline{s}$ )?
|   IRIIf(val expr: IRExpr, val trueB: IRStmt,
|         val falseB: Option[IRStmt])
|   while ( $\underline{e}$ )  $\underline{s}$ 
|   IRWhile(val cond: IRExpr, val body: IRStmt)
|   try { $\underline{s}$ } (catch ( $\underline{x}$ ){ $\underline{s}$ }?) (finally { $\underline{s}$ }?)?
|   IRTry(val body: IRStmt, val name: Option[IRId],
|         val catchB: Option[IRStmt],
|         val finallyB: Option[IRStmt])
|    $\langle \underline{s}^* \rangle$ 
|   IRStmtUnit(List[IRStmt] stmts)

 $\underline{e}$  ::=  $\underline{e} \otimes \underline{e}$ 
|   IRBin(val first: IRExpr, val op: IROp,
|         val second: IRExpr)
|    $\ominus \underline{e}$ 
|   IRUn(val op: IROp, val expr: IRExpr)
|    $\underline{x}[\underline{e}]$ 
|   IRLoad(val obj: IRId, val index: IRExpr)
|    $\underline{x}$ 
|   IRUserId(val global: Boolean, val isWith: Boolean)
|    $\diamond \underline{x}$ 
|   IRTmpId(val global: Boolean)
|    $\underline{num}$ 
|   IRNumber(val text: String, val num: Double)
|    $\underline{str}$ 
|   IRString(val str: String)
|   true
|   IRBool(val bool: Boolean)
|   false
|   IRBool(val bool: Boolean)
|   undefined
|   IRUndef()
|   null
|   IRNull()
|   this
|   IRThis()

 $\underline{m}$  ::=  $\underline{x} : \underline{e}$ 
|   IRField(val prop: IRId, val expr: IRExpr)
|   get  $\underline{f}(\underline{x}, \underline{x})$  { $\underline{s}^*$ }
|   IRGetProp(val ftn: IRFunctional)
|   set  $\underline{f}(\underline{x}, \underline{x})$  { $\underline{s}^*$ }
|   IRSetProp(val ftn: IRFunctional)

```

The SAFE IR has the following assumptions and notations:

- We denote a list as a possibly empty, semicolon-separated sequence, enclosed by \langle and \rangle .
- We denote a series of list appends as superscripted $*$ such as \underline{s}^* .
- We abuse our notations by interchanging semicolon-separated sequences and lists.
- To denote an AST-level statement granularity in the translated IR statements, we use `IRStmtUnit` which is represented as green angle brackets $\langle \rangle$ in this document.
- Declarations of functions and variables are hoisted to their closest enclosing functions or the top level via the `Hoister` transformation of the `astRewriter` phase.
- Identifiers and labels that exist in the source program, except when they appear at top level or within the `with` statement, are already disambiguated via the `Disambiguator` transformation of the `astRewriter` phase so that they have unique names.

A.3 AST to IR

This section describes the SAFE translation rules from AST to IR, whose implementation is available at:

`SAFE_HOME/src/main/scala/kr/ac/kaist/safe/compiler/Translator.scala`

- We use Σ to disambiguate the generated labels and temporary variables in the AST to IR translation. For the presentation brevity, we simply add the newly generated names to Σ .
 - In the actual implementation, we need to create a unique id for each generated name and add the binding information from the general name to the unique id to Σ . For example, when we say “ $\Sigma; \diamond \text{break}$ ”, we actually create a unique id for $\diamond \text{break}$, say $\diamond \text{break}_{42}$, and add it to Σ as $\Sigma; \diamond \text{break} \mapsto \diamond \text{break}_{42}$. When we look up the environment by $\Sigma(\diamond \text{break})$, the unique $\diamond \text{break}_{42}$ is returned.
 - In the scope when the generated name is created, we don’t add it to the environment but use the unique id instead of the general name. For example, when we say “ $\diamond \text{eq} = \Sigma(\diamond \text{val}) === \diamond \text{break};$ ”, we create a unique id for $\diamond \text{eq}$, say $\diamond \text{eq}_{910157}$, and it is actually “ $\diamond \text{eq}_{910157} = \Sigma(\diamond \text{val}) === \diamond \text{break}_{42};$ ”.
 - To be clear, we use blue for the binding sites of such names and red for the use sites of such names.
- We denote a fresh variable name as \diamond and its variants.
- We use the following:
 - `===`, `toObject`, `toNumber`, `isObject`, `iteratorInit`, `iteratorHasNext`, `iteratorNext`, `global`, `getBase`
- To reduce the number of temporary variables, we use global variables to denote constants such as 1 and `true` which is represented in green `1` and `true` in this document.
- We wrap a possibly identical assignment with a box so that the actual implementation, `Translator`, can eliminate identical assignments.

Here are the types of the environment used to disambiguate the generated labels and temporary variables and the translation functions for different language constructs:

```

Σ : Env
ast2irp :
Program -> IRRoot
ast2irfd :
FunDecl -> Env -> IRFunDecl
ast2irvd :
VarDecl -> Env -> IRVarStmt
ast2irs :
Stmt -> Env -> IRStmt
ast2ircase :
List[Case] * Option[List[Stmt]] * List[Case]
-> Env -> List[Option[Expr] * IRId] -> IRStmt
ast2irscond :
List[Option[Expr] * IRId] -> Env -> IRStmt
ast2irlval :
Expr -> Env -> List[IRStmt] -> IRExpr -> boolean
-> List[IRStmt] * IRExpr
ast2ire :
Expr -> Env -> IRId -> List[IRStmt] * IRExpr
ast2irlhs :
LHS -> Env -> IRId -> List[IRStmt] * IRExpr
ast2irlit :
LIT -> Env -> IRId -> List[IRStmt] * IRExpr
ast2irm :
Member -> Env -> IRId -> List[IRStmt] * IRMember
ast2irpr :
Property -> IRId

```

```

⊙ ::= * | / | % | + | - | [[ | >> | >>> | & | ^ | |
⊖ ::= ~ | ! | + | - | delete | void | typeof
⊗ ::= | | & | ^ | [[ | >> | >>> | + | - | * | / | % | ==
| != | === | !== | [ | > | [= | >= | instanceof | in

```

$$ast2ir_p[f d^* v d^* s^*] = \langle (ast2ir_{fd}[f d](\langle \rangle))^* (ast2ir_{vd}[v d](\langle \rangle))^* (ast2ir_s[s](\langle \rangle))^* \rangle$$

$$ast2ir_{fd}[\text{function } f((x,)*) \{f d^* v d^* s^*\}](\Sigma) = \text{function } \underline{f}(\underline{\diamond this}, \underline{\diamond arguments})\{$$

$$\begin{aligned}
& (ast2ir_{fd}[f d](\Sigma))^* \\
& (\text{var } x_i)^* \\
& (ast2ir_{vd}[v d](\Sigma))^* \\
& (x_i = \underline{\diamond arguments}["i"])^*
\end{aligned}$$

where x_i is not the name of any of fd

$$(ast2ir_s[s](\Sigma; \underline{\diamond this}; \underline{\diamond arguments}))^*$$

A function always receives explicit “this” and “arguments” arguments so that the desugaring of this and arguments is correct. Currently, “arguments” denotes copies of the arguments instead of their aliases.

$$ast2ir_{vd}[\text{var } x](\Sigma) =$$

$$\text{var } \underline{x}$$

$$ast2ir_s[\{s^*\}](\Sigma) = \langle (ast2ir_s[s](\Sigma))^* \rangle$$

$$ast2ir_s[\langle \rangle](\Sigma) = \langle \rangle$$

$$\begin{aligned}
ast2ir_s[e;](\Sigma) &= \\
\text{LET } (\underline{s^*}, \underline{e}) &= ast2ir_e[e](\Sigma)(\underline{\diamond new}) \\
\text{IN } \langle \underline{s^*}; \boxed{\underline{\diamond new} = \underline{e}} \rangle
\end{aligned}$$

$$\begin{aligned}
ast2ir_s[\text{if } (e_1 \&\& \dots \&\& e_n) s_1 (\text{else } s_2)^?](\Sigma) &= \\
\text{LET } (\underline{s_i^*}, \underline{e_i}) &= ast2ir_e[e_i](\Sigma)(\underline{\diamond new_i}) \quad \text{where } 1 \leq i \leq n \\
\text{IN } \langle \underline{s_1^*}; \\
& \underline{\diamond label} : \{ \\
& \quad \text{if } (e_1) \\
& \quad \text{then } \langle \underline{s_2^*}; \dots; \\
& \quad \quad \text{if } (e_{n-1}) \\
& \quad \quad \text{then } \langle \underline{s_n^*}; \\
& \quad \quad \quad \text{if } (e_n) \\
& \quad \quad \quad \text{then } \{ast2ir_s[s_1](\Sigma); \text{break } \underline{\diamond label}\} \dots \}; \\
& \quad (ast2ir_s[s_2](\Sigma))^? \rangle
\end{aligned}$$

Candidate for optimization

$$\begin{aligned}
ast2ir_s[\text{if } (e_1 | e_2) s_1 (\text{else } s_2)^?](\Sigma) &= \\
\text{LET } (\underline{s_1^*}, \underline{e_1}) &= ast2ir_e[e_1](\Sigma)(\underline{\diamond new_1}) \\
& (\underline{s_2^*}, \underline{e_2}) = ast2ir_e[e_2](\Sigma)(\underline{\diamond new_2}) \\
\text{IN } \langle \underline{s_1^*}; \\
& \underline{\diamond label_2} : \{ \\
& \quad \underline{\diamond label_1} : \{ \\
& \quad \quad \text{if } (\underline{e_1}) \\
& \quad \quad \text{then break } \underline{\diamond label_1}; \underline{s_2^*}; \\
& \quad \quad \text{if } (\underline{e_2}) \text{ then break } \underline{\diamond label_1}; \\
& \quad \quad (ast2ir_s[s_2](\Sigma); \text{break})^? \underline{\diamond label_2} \\
& \quad \}; ast2ir_s[s_1](\Sigma) \} \rangle
\end{aligned}$$

$$\begin{aligned}
ast2ir_s[\text{if } (e) s_1 (\text{else } s_2)^?](\Sigma) &= \\
\text{LET } (\underline{s^*}, \underline{e}) &= ast2ir_e[e](\Sigma)(\underline{\diamond new}) \\
\text{IN } \langle \underline{s^*}; \text{if } (e) \text{ then } ast2ir_s[s_1](\Sigma) (\text{else } ast2ir_s[s_2](\Sigma))^? \rangle \\
ast2ir_s[\text{switch } (e) \{cc_1^* (\text{default}: s^*)^? cc_2^*\}](\Sigma) &= \\
\text{LET } (\underline{s^*}, \underline{e}) &= ast2ir_e[e](\Sigma)(\underline{\diamond val}) \\
\text{IN } \langle \underline{\diamond break} : \{ \\
& \quad \underline{s^*}; \boxed{\underline{\diamond val} = \underline{e}}; \\
& \quad ast2ir_{case}[(\text{rev } cc_2^*)(s^*)^? (\text{rev } cc_1^*)](\Sigma; \underline{\diamond break}; \underline{\diamond val}) \} \rangle
\end{aligned}$$

$$\begin{aligned}
ast2ir_{case}[(\text{case } e : s_1^*) :: cc_2^* (s_2^*)^? cc_1^*](\Sigma)(c^*) &= \\
\langle \underline{\diamond label} : \{ast2ir_{case}[cc_2^* (s_2^*)^? cc_1^*](\Sigma)((e, \underline{\diamond label}) :: c^*)\}; \\
& (ast2ir_s[s_1](\Sigma))^* \rangle
\end{aligned}$$

$$\begin{aligned}
ast2ir_{case}[(\text{ } (s^*)^? cc_1^*](\Sigma)(c^*) &= \\
\langle \underline{\diamond label} : \{ast2ir_{case}[(\text{ } () cc_1^*](\Sigma)(c^* @ ((\text{ }, \underline{\diamond label})))\}; \\
& ((ast2ir_s[s](\Sigma))^*)^? \rangle
\end{aligned}$$

$$\begin{aligned}
ast2ir_{case}[(\text{ } () (\text{case } e : s^*) :: cc_1^*](\Sigma)(c^*) &= \\
\langle \underline{\diamond label} : \{ast2ir_{case}[(\text{ } () cc_1^*](\Sigma)((e, \underline{\diamond label}) :: c^*)\}; \\
& (ast2ir_s[s](\Sigma))^* \rangle
\end{aligned}$$

$$\begin{aligned}
ast2ir_{case}[(\text{ } () ()](\Sigma)((e, l)^*) &= \\
\langle ast2ir_{scond}[(e, l)^*](\Sigma); \\
& \text{break } \Sigma(\underline{\diamond break}) \rangle
\end{aligned}$$

$$\begin{aligned}
ast2ir_{scond}[(e, l) :: (c^*)](\Sigma) &= \\
\text{LET } (\underline{s^*}, \underline{e}) &= ast2ir_e[e](\Sigma)(\underline{\diamond cond}) \\
\text{IN } \langle \underline{s^*}; \\
& \text{if } (\Sigma(\underline{\diamond val}) == \underline{e}) \text{ then break } \underline{l} \text{ else } ast2ir_{scond}[(c^*)](\Sigma) \rangle
\end{aligned}$$

$$ast2ir_{scond}[\llbracket (((), l) \rrbracket](\Sigma) =$$

$$\langle \text{break } l \rangle$$

$$ast2ir_{scond}[\llbracket () \rrbracket](\Sigma) =$$

$$\langle \rangle$$

Where c is either (e, l) or $((), l)$.

$$ast2ir_s[\llbracket \text{do } s \text{ while } (e); \rrbracket](\Sigma) =$$

$$\text{LET } (s^*, e) = ast2ir_e[\llbracket e \rrbracket](\Sigma)(\diamond new_1)$$

$$\text{IN } \langle \diamond break : \{$$

$$\quad \diamond continue : \{ ast2ir_s[\llbracket s \rrbracket](\Sigma; \diamond break; \diamond continue) \};$$

$$\quad s^*;$$

$$\quad \text{while } (e) \{$$

$$\quad \quad \diamond continue : \{ ast2ir_s[\llbracket s \rrbracket](\Sigma; \diamond break; \diamond continue) \};$$

$$\quad \quad s^*;$$

$$\quad \}$$

$$\rangle$$

$$ast2ir_s[\llbracket \text{while } (e) s \rrbracket](\Sigma) =$$

$$\text{LET } (s^*, e) = ast2ir_e[\llbracket e \rrbracket](\Sigma)(\diamond new_1)$$

$$\text{IN } \langle \diamond break : \{$$

$$\quad s^*;$$

$$\quad \text{while } (e) \{$$

$$\quad \quad \diamond continue : \{ ast2ir_s[\llbracket s \rrbracket](\Sigma; \diamond break; \diamond continue) \};$$

$$\quad \quad s^*;$$

$$\quad \}$$

$$\rangle$$

$$ast2ir_s[\llbracket \text{for } (e_1^?; e_2^?; e_3^?) s \rrbracket](\Sigma) =$$

$$\text{LET } ((s_1^*, e_1) = ast2ir_e[\llbracket e_1 \rrbracket](\Sigma)(\diamond \underline{\hspace{0.5cm}}))^?$$

$$((s_3^*, e_3) = ast2ir_e[\llbracket e_3 \rrbracket](\Sigma)(\diamond \underline{\hspace{0.5cm}}))^?$$

$$\text{IN } \langle \diamond break : \{$$

$$\quad (s_1^*; \diamond \underline{\hspace{0.5cm}} = e_1)^?$$

$$\quad \text{while } (\text{true}) \{$$

$$\quad \quad \diamond continue : \{ ast2ir_s[\llbracket s \rrbracket](\Sigma; \diamond break; \diamond continue) \};$$

$$\quad \quad (s_3^*; \diamond \underline{\hspace{0.5cm}} = e_3)^?$$

$$\quad \}$$

$$\rangle$$

$$ast2ir_s[\llbracket \text{for } (e_1^?; e_2; e_3^?) s \rrbracket](\Sigma) =$$

$$\text{LET } ((s_1^*, e_1) = ast2ir_e[\llbracket e_1 \rrbracket](\Sigma)(\diamond \underline{\hspace{0.5cm}}))^?$$

$$(s_2^*, e_2) = ast2ir_e[\llbracket e_2 \rrbracket](\Sigma)(\diamond new_2)$$

$$((s_3^*, e_3) = ast2ir_e[\llbracket e_3 \rrbracket](\Sigma)(\diamond \underline{\hspace{0.5cm}}))^?$$

$$\text{IN } \langle \diamond break : \{$$

$$\quad (s_1^*; \diamond \underline{\hspace{0.5cm}} = e_1)^?$$

$$\quad s_2^*;$$

$$\quad \text{while } (e_2) \{$$

$$\quad \quad \diamond continue : \{ ast2ir_s[\llbracket s \rrbracket](\Sigma; \diamond break; \diamond continue) \};$$

$$\quad \quad (s_3^*; \diamond \underline{\hspace{0.5cm}} = e_3)^?$$

$$\quad \quad s_2^*;$$

$$\quad \}$$

$$\rangle$$

$$ast2ir_s[\llbracket \text{for } (lhs \text{ in } e) s \rrbracket](\Sigma) =$$

$$\text{LET } (s^*, e) = ast2ir_e[\llbracket e \rrbracket](\Sigma)(\diamond new_1)$$

$$\text{IN } \langle \diamond break : \{$$

$$\quad s^*;$$

$$\quad \diamond obj = \diamond toObject(e);$$

$$\quad \diamond iterator = \diamond iteratorInit(\diamond obj);$$

$$\quad \diamond cond_1 = \diamond iteratorHasNext(\diamond obj, \diamond iterator);$$

$$\quad \text{while } (\diamond cond_1) \{$$

$$\quad \quad \diamond key = \diamond iteratorNext(\diamond obj, \diamond iterator);$$

$$\quad \quad ast2ir_{lval}[\llbracket lhs \rrbracket](\Sigma)(; \diamond key)(\text{false})._1;$$

$$\quad \quad \diamond continue : \{ ast2ir_s[\llbracket s \rrbracket](\Sigma; \diamond break; \diamond continue) \};$$

$$\quad \quad \diamond cond_1 = \diamond iteratorHasNext(\diamond obj, \diamond iterator);$$

$$\quad \}$$

$$\rangle$$

$$ast2ir_s[\llbracket \text{continue}; \rrbracket](\Sigma) =$$

$$\langle \text{break } \Sigma(\diamond continue) \rangle$$

$$ast2ir_s[\llbracket \text{continue } l; \rrbracket](\Sigma) =$$

$$\langle \text{break } \Sigma(l) \rangle$$

$$ast2ir_s[\llbracket \text{break}; \rrbracket](\Sigma) =$$

$$\langle \text{break } \Sigma(\diamond break) \rangle$$

$$ast2ir_s[\llbracket \text{break } l; \rrbracket](\Sigma) =$$

$$\langle \text{break } l \rangle$$

$$ast2ir_s[\llbracket \text{return}; \rrbracket](\Sigma) =$$

$$\langle \text{return} \rangle$$

$$ast2ir_s[\llbracket \text{return } e; \rrbracket](\Sigma) =$$

$$\text{LET } (s^*, e) = ast2ir_e[\llbracket e \rrbracket](\Sigma)(\diamond new_1)$$

$$\text{IN } \langle s^*; \text{return } e \rangle$$

$$ast2ir_s[\llbracket \text{with } (e) s \rrbracket](\Sigma) =$$

$$\text{LET } (s^*, e) = ast2ir_e[\llbracket e \rrbracket](\Sigma)(\diamond new_1)$$

$$\text{IN } \langle s^*;$$

$$\quad \diamond new_2 = \diamond toObject(e);$$

$$\quad \text{with } (\diamond new_2) ast2ir_s[\llbracket s \rrbracket](\Sigma) \rangle$$

$$ast2ir_s[\llbracket l : s \rrbracket](\Sigma) =$$

$$\langle l : \{ ast2ir_s[\llbracket s \rrbracket](\Sigma; l) \} \rangle$$

$$ast2ir_s[\llbracket \text{throw } e; \rrbracket](\Sigma) =$$

$$\text{LET } (s^*, e) = ast2ir_e[\llbracket e \rrbracket](\Sigma)(\diamond new_1)$$

$$\text{IN } \langle s^*; \text{throw } e \rangle$$

$$ast2ir_s[\llbracket \text{try } \{s_1^*\} (\text{catch}(x) \{s_2^*\})^? (\text{finally } \{s_3^*\})^? \rrbracket](\Sigma) =$$

$$\langle \text{try } \{ (ast2ir_s[\llbracket s_1 \rrbracket](\Sigma))^* \}$$

$$\quad (\text{catch}(x) \{ (ast2ir_s[\llbracket s_2 \rrbracket](\Sigma))^* \})^?$$

$$\quad (\text{finally } \{ (ast2ir_s[\llbracket s_3 \rrbracket](\Sigma))^* \})^? \rangle$$

$$ast2ir_s[\llbracket \text{debugger}; \rrbracket](\Sigma) =$$

$$\langle \rangle$$

$$\text{ast2ir}_{\text{lval}}[\langle e \rangle](\Sigma)(\underline{s}^*; \underline{e}')(\text{keepOld}) =$$

$$\text{ast2ir}_{\text{lval}}[e](\Sigma)(\underline{s}^*; \underline{e}')(\text{keepOld})$$

$$\text{ast2ir}_{\text{lval}}[x](\Sigma)(\underline{s}^*; \underline{e})(\text{keepOld}) =$$

$$\text{IF keepOld THEN } (\langle \text{old} = x; \underline{s}^*; x = e \rangle, x)$$

$$\text{ELSE } \langle \underline{s}^*; x = e \rangle$$

$$\text{ast2ir}_{\text{lval}}[lhs.x](\Sigma)(\underline{s}^*; \underline{e})(\text{keepOld}) =$$

$$\text{ast2ir}_{\text{lval}}[lhs["x"]](\Sigma)(\underline{s}^*; \underline{e})(\text{keepOld})$$

$$\text{ast2ir}_{\text{lval}}[lhs[e]](\Sigma)(\underline{s}^*; \underline{e}')(\text{keepOld}) =$$

$$\text{LET } (\underline{s}_1^*, e_1) = \text{ast2ir}_{\text{lhs}}[lhs](\Sigma)(\text{obj}_1)$$

$$(\underline{s}_2^*, e_2) = \text{ast2ir}_e[e](\Sigma)(\text{field}_1)$$

$$\text{IN IF keepOld}$$

$$\text{THEN } (\langle \underline{s}_1^*; \text{obj} = \text{toObject}(e_1); \underline{s}_2^*;$$

$$\text{old} = \text{obj}[e_2]; \underline{s}^*; \text{obj}[e_2] = e' \rangle,$$

$$\text{obj}[e_2])$$

$$\text{ELSE } (\langle \underline{s}_1^*; \text{obj} = \text{toObject}(e_1); \underline{s}_2^*;$$

$$\underline{s}^*; \text{obj}[e_2] = e' \rangle, \text{obj}[e_2])$$

$$\text{ast2ir}_{\text{lval}}[e](\Sigma)(\underline{s}^*; \underline{e})(\text{keepOld}) =$$

Warning: ReferenceError!

$$\text{ast2ir}_e[e_1, e_2](\Sigma)(x) =$$

$$\text{LET } (\underline{s}_1^*, e_1) = \text{ast2ir}_e[e_1](\Sigma)(\text{oy})$$

$$(\underline{s}_2^*, e_2) = \text{ast2ir}_e[e_2](\Sigma)(x)$$

$$\text{IN } (\underline{s}_1^*; \text{oy} = e_1; \underline{s}_2^*, e_2)$$

Candidate for optimization

$$\text{ast2ir}_e[e_a \& \& e_b ? e_2 : e_3](\Sigma)(x) =$$

$$\text{LET } (\underline{s}_a^*, e_a) = \text{ast2ir}_e[e_a](\Sigma)(\text{new}_a)$$

$$(\underline{s}_b^*, e_b) = \text{ast2ir}_e[e_b](\Sigma)(\text{new}_b)$$

$$(\underline{s}_2^*, e_2) = \text{ast2ir}_e[e_2](\Sigma)(x)$$

$$(\underline{s}_3^*, e_3) = \text{ast2ir}_e[e_3](\Sigma)(x)$$

$$\text{IN } (\underline{s}_a^*;$$

$$\text{label} : \{$$

$$\text{if } (e_a)$$

$$\text{then } \langle \underline{s}_b^*; \text{if } (e_b) \text{ then } \{ \underline{s}_2^*; x = e_2 \}; \text{break label} \rangle;$$

$$\underline{s}_3^*; x = e_3 \}, x)$$

Candidate for optimization

$$\text{ast2ir}_e[e_a || e_b ? e_2 : e_3](\Sigma)(x) =$$

$$\text{LET } (\underline{s}_a^*, e_a) = \text{ast2ir}_e[e_a](\Sigma)(\text{new}_a)$$

$$(\underline{s}_b^*, e_b) = \text{ast2ir}_e[e_b](\Sigma)(\text{new}_b)$$

$$(\underline{s}_2^*, e_2) = \text{ast2ir}_e[e_2](\Sigma)(x)$$

$$(\underline{s}_3^*, e_3) = \text{ast2ir}_e[e_3](\Sigma)(x)$$

$$\text{IN } (\underline{s}_a^*;$$

$$\text{label}_2 : \{$$

$$\text{label}_1 : \{$$

$$\text{if } (e_a)$$

$$\text{then break label}_1; \underline{s}_b^*;$$

$$\text{if } (e_b) \text{ then break label}_1;$$

$$\underline{s}_3^*; x = e_3 \}; \text{break label}_2$$

$$\}; \underline{s}_2^*; x = e_2 \}, x)$$

$$\text{ast2ir}_e[e_1 ? e_2 : e_3](\Sigma)(x) =$$

$$\text{LET } (\underline{s}_1^*, e_1) = \text{ast2ir}_e[e_1](\Sigma)(\text{new}_1)$$

$$(\underline{s}_2^*, e_2) = \text{ast2ir}_e[e_2](\Sigma)(x)$$

$$(\underline{s}_3^*, e_3) = \text{ast2ir}_e[e_3](\Sigma)(x)$$

$$\text{IN } (\underline{s}_1^*; \text{if } (e_1) \text{ then } \{ \underline{s}_2^*; x = e_2 \} \text{ else } \{ \underline{s}_3^*; x = e_3 \} \}, x)$$

$$\text{ast2ir}_e[lhs = e](\Sigma)(x) =$$

$$\text{LET } (\underline{s}^*, \underline{e}) = \text{ast2ir}_e[e](\Sigma)(x)$$

$$\text{IN IF } \underline{e} \text{ contains lhs}$$

$$\text{THEN } \text{ast2ir}_{\text{lval}}[lhs](\Sigma)(\underline{s}^*; \underline{e})(\text{false})$$

$$\text{ELSE } (\text{ast2ir}_{\text{lval}}[lhs](\Sigma)(\underline{s}^*; \underline{e})(\text{false})._1, \underline{e})$$

$$\text{ast2ir}_e[lhs \odot e](\Sigma)(x) =$$

$$\text{LET } (\underline{s}^*, \underline{e}) = \text{ast2ir}_e[e](\Sigma)(\text{oy})$$

$$\text{IN } (\text{ast2ir}_{\text{lval}}[lhs](\Sigma)(\underline{s}^*; \text{old} \odot \underline{e})(\text{true})._1, \text{old} \odot \underline{e})$$

$$\text{ast2ir}_e[++e](\Sigma)(x) =$$

$$(\text{ast2ir}_{\text{lval}}[e](\Sigma)(\text{new} = \text{toNumber}(\text{old});$$

$$\text{new} + 1)(\text{true})._1, \text{new} + 1)$$

$$\text{ast2ir}_e[--e](\Sigma)(x) =$$

$$(\text{ast2ir}_{\text{lval}}[e](\Sigma)(\text{new} = \text{toNumber}(\text{old});$$

$$\text{new} - 1)(\text{true})._1, \text{new} - 1)$$

$$\text{ast2ir}_e[\text{delete } x](\Sigma)(y) =$$

$$(\langle y = \text{delete } x \rangle, y)$$

$$\text{ast2ir}_e[\text{delete } (x)](\Sigma)(y) =$$

$$(\langle y = \text{delete } x \rangle, y)$$

$$\text{ast2ir}_e[\text{delete } lhs.x](\Sigma)(y) =$$

$$\text{ast2ir}_e[\text{delete } lhs["x"]](\Sigma)(y)$$

$$\text{ast2ir}_e[\text{delete } lhs[e]](\Sigma)(x) =$$

$$\text{LET } (\underline{s}_1^*, e_1) = \text{ast2ir}_{\text{lhs}}[lhs](\Sigma)(\text{obj}_1)$$

$$(\underline{s}_2^*, e_2) = \text{ast2ir}_e[e](\Sigma)(\text{field}_1)$$

$$\text{IN } (\underline{s}_1^*; \text{obj} = \text{toObject}(e_1); \underline{s}_2^*;$$

$$x = \text{delete } \text{obj}[e_2], x)$$

$$\text{ast2ir}_e[\text{delete } e](\Sigma)(x) =$$

$$\text{LET } (\underline{s}^*, \underline{e}) = \text{ast2ir}_e[e](\Sigma)(\text{oy})$$

$$\text{IN } (\underline{s}^*; \underline{e} = \underline{e}, \text{true})$$

$$\text{ast2ir}_e[\ominus e](\Sigma)(x) =$$

$$\text{LET } (\underline{s}^*, \underline{e}) = \text{ast2ir}_e[e](\Sigma)(\text{oy})$$

$$\text{IN } (\underline{s}^*, \ominus \underline{e})$$

$$\text{ast2ir}_e[lhs++](\Sigma)(x) =$$

$$(\text{ast2ir}_{\text{lval}}[lhs](\Sigma)(\text{new} = \text{toNumber}(\text{old});$$

$$\text{new} + 1)(\text{true})._1, \text{new})$$

$$\text{ast2ir}_e[lhs--](\Sigma)(x) =$$

$$(\text{ast2ir}_{\text{lval}}[lhs](\Sigma)(\text{new} = \text{toNumber}(\text{old});$$

$$\text{new} - 1)(\text{true})._1, \text{new})$$

Candidate for optimization

$ast2ir_e[e_1 \& \& e_2](\Sigma)(x) =$
 LET $(s_1^*, e_1) = ast2ir_e[e_1](\Sigma)(\diamond y)$
 $(s_2^*, e_2) = ast2ir_e[e_2](\Sigma)(\diamond z)$
 IN $(s_1^*; \text{if } (e_1) \text{ then } s_2^*; x = e_2 \text{ else } x = e_1, x)$

Candidate for optimization

$ast2ir_e[e_1 \parallel e_2](\Sigma)(x) =$
 LET $(s_1^*, e_1) = ast2ir_e[e_1](\Sigma)(\diamond y)$
 $(s_2^*, e_2) = ast2ir_e[e_2](\Sigma)(\diamond z)$
 IN $(s_1^*; \text{if } (e_1) \text{ then } x = e_1 \text{ else } s_2^*; x = e_2, x)$

In order to preserve the semantics when the evaluation of e_1 throws an exception, we force to evaluate e_1 before evaluating s_2^* by introducing an assignment " $\diamond new_1 = e_1$ " to avoid any side effects by s_2^* . Note that we add the assignment only when s_2^* is not empty for a simple optimization.

Candidate for optimization

$ast2ir_e[e_1 \otimes e_2](\Sigma)(x) =$
 LET $(s_1^*, e_1) = ast2ir_e[e_1](\Sigma)(\diamond y)$
 $(s_2^*, e_2) = ast2ir_e[e_2](\Sigma)(\diamond z)$
 IN IF s_2^* is empty
 THEN $(s_1^*, e_1 \otimes e_2)$
 ELSE $(s_1^*; \diamond y = e_1; s_2^*, \diamond y \otimes e_2)$

$ast2ir_e[lhs](\Sigma)(x) =$
 $ast2ir_{lhs}[lhs](\Sigma)(x)$

$ast2ir_{lhs}[lit](\Sigma)(x) =$
 $ast2ir_{lit}[lit](\Sigma)(x)$

$ast2ir_{lhs}[arguments](\Sigma)(x) =$
 $(\langle \rangle, \Sigma(\diamond arguments))$

$ast2ir_{lhs}[x](\Sigma)(y) =$
 $(\langle \rangle, x)$

Candidate for optimization

$ast2ir_{lhs}[(e^?, *)](\Sigma)(x) =$
 LET $((s^*, e) = ast2ir_e[e](\Sigma)(\diamond elem))^*$
 IN $((s^*; \diamond elem = e)^*; x = [(\diamond elem)^*,], x)$

$ast2ir_{lhs}[\{f(m, *)\}](\Sigma)(x) =$
 LET $((s^*, mem) = ast2ir_m[m](\Sigma)(\diamond member))^*$
 IN $((s^*)^*; x = \{f(mem, *)\}, x)$

$ast2ir_{lhs}[(e)](\Sigma)(x) =$
 $ast2ir_e[e](\Sigma)(x)$

$ast2ir_{lhs}[\text{function } f^?((x, *)^*) \{fd^* vd^* s^*\}](\Sigma)(y) =$
 $(\langle y = \text{function } f^?(\diamond this, \diamond arguments) \{$
 $(ast2ir_{fd}[fd](\Sigma))^*$
 $(\text{var } x_i)^*$
 $(ast2ir_{vd}[vd](\Sigma))^*$
 $(x_i = \diamond arguments["i"])^*$
 $\text{where } x_i \text{ is not the name of any of } fd$
 $(ast2ir_s[s](\Sigma; \diamond this; \diamond arguments))^*\} \rangle, y)$

$ast2ir_{lhs}[lhs.x](\Sigma)(y) =$
 $ast2ir_{lhs}[lhs["x"]](\Sigma)(y)$

$ast2ir_{lhs}[lhs["x"]](\Sigma)(y) =$
 LET $(s_1^*, e_1) = ast2ir_{lhs}[lhs](\Sigma)(\diamond obj_1)$
 IN $(s_1^*; \diamond obj = \diamond toObject(e_1), \diamond obj["x"])$

$ast2ir_{lhs}[lhs[e]](\Sigma)(x) =$
 LET $(s_1^*, e_1) = ast2ir_{lhs}[lhs](\Sigma)(\diamond obj_1)$
 $(s_2^*, e_2) = ast2ir_e[e](\Sigma)(\diamond field_1)$
 IN $(s_1^*; \diamond obj = \diamond toObject(e_1); s_2^*, \diamond obj[e_2])$

Candidate for optimization

$ast2ir_{lhs}[\text{new } lhs((e, *)^*)](\Sigma)(x) =$
 LET $(s_l^*, e_l) = ast2ir_{lhs}[lhs](\Sigma)(\diamond fun_1)$
 $((s^*, e) = ast2ir_e[e](\Sigma)(\diamond y))^*$
 IN $(s_l^*; \diamond fun = \diamond toObject(e_l); (s^*; \diamond y = e)^*;$
 $\diamond arguments = [(\diamond y_i)^*];$
 $\diamond proto = \diamond fun["prototype"];$
 $\diamond obj = \{[[Prototype]] = \diamond proto\};$
 $\diamond newObj = \text{new } \diamond fun(\diamond obj, \diamond arguments);$
 $\diamond cond = \diamond isObject(\diamond newObj);$
 if $(\diamond cond)$ then $x = \diamond newObj$ else $x = \diamond obj, x)$

$ast2ir_{lhs}[\text{new } lhs](\Sigma)(x) =$
 LET $(s^*, e) = ast2ir_{lhs}[lhs](\Sigma)(\diamond fun_1)$
 IN $(s^*; \diamond fun = \diamond toObject(e);$
 $\diamond arguments = [];$
 $\diamond proto = \diamond fun["prototype"];$
 $\diamond obj = \{[[Prototype]] = \diamond proto\};$
 $\diamond newObj = \text{new } \diamond fun(\diamond obj, \diamond arguments);$
 $\diamond cond = \diamond isObject(\diamond newObj);$
 if $(\diamond cond)$ then $x = \diamond newObj$ else $x = \diamond obj, x)$

$ast2ir_{lhs}[\text{eval}(e)](\Sigma)(x) =$
 LET $(s^*, e) = ast2ir_e[e](\Sigma)(\diamond new_1)$
 IN $(s^*; x = \text{eval}(e), x)$

$ast2ir_{lhs}[(f)((e, *)^*)](\Sigma)(x) =$
 $ast2ir_{lhs}[f((e, *)^*)](\Sigma)(x)$

Candidate for optimization

$ast2ir_{lhs}[f((e, *)^*)](\Sigma)(x) =$
 LET $((s^*, e) = ast2ir_e[e](\Sigma)(\diamond y))^*$
 IN $(\diamond obj = \diamond toObject(f); (s^*; \diamond y = e)^*;$
 $\diamond arguments = [(\diamond y_i)^*];$
 $\diamond fun = \diamond getBase(f);$
 $x = \diamond obj(\diamond fun, \diamond arguments), x)$

$ast2ir_{lhs}[(lhs.x)((e, *)^*)](\Sigma)(y) =$
 $ast2ir_{lhs}[lhs["x"]((e, *)^*)](\Sigma)(y)$

$ast2ir_{lhs}[lhs.x((e, *)^*)](\Sigma)(y) =$
 $ast2ir_{lhs}[lhs["x"]((e, *)^*)](\Sigma)(y)$

$ast2ir_{lhs}[(lhs[e'])((e, *)^*)](\Sigma)(x) =$
 $ast2ir_{lhs}[lhs[e']((e, *)^*)](\Sigma)(x)$

Candidate for optimization

$$\begin{aligned}
& ast2ir_{lhs}[\llbracket lhs[e'] \rrbracket((e, *)^*)](\Sigma)(x) = \\
& \text{LET } (s^*, e_l) = ast2ir_{lhs}[\llbracket lhs \rrbracket(\Sigma)(\diamond obj_1)] \\
& \quad (s'^*, e') = ast2ir_e[\llbracket e' \rrbracket(\Sigma)(\diamond field_1)] \\
& \quad ((s^*, e) = ast2ir_e[\llbracket e \rrbracket(\Sigma)(\diamond y)]^*) \\
& \text{IN } (s^*; \diamond obj = \diamond toObject(e_l); s'^*; \\
& \quad (s^*; \diamond y = e)^*; \\
& \quad \diamond arguments = [(\diamond y_i,)^*]; \\
& \quad \diamond fun = \diamond toObject(\diamond obj[\llbracket e' \rrbracket]); \\
& \quad x = \diamond fun(\diamond obj, \diamond arguments), x)
\end{aligned}$$
Candidate for optimization

$$\begin{aligned}
& ast2ir_{lhs}[\llbracket lhs((e, *)^*) \rrbracket(\Sigma)(x) = \\
& \text{LET } (s^*, e_l) = ast2ir_{lhs}[\llbracket lhs \rrbracket(\Sigma)(\diamond obj_1)] \\
& \quad ((s^*, e) = ast2ir_e[\llbracket e \rrbracket(\Sigma)(\diamond y)]^*) \\
& \text{IN } (s^*; \diamond obj = \diamond toObject(e_l); (s^*; \diamond y = e)^*; \\
& \quad \diamond arguments = [(\diamond y_i,)^*]; \\
& \quad x = \diamond obj(\diamond global, \diamond arguments), x)
\end{aligned}$$

$$\begin{aligned}
& ast2ir_{lit}[\llbracket this \rrbracket(\Sigma)(x) = \\
& (\langle \rangle, \Sigma(\diamond this))
\end{aligned}$$

$$\begin{aligned}
& ast2ir_{lit}[\llbracket null \rrbracket(\Sigma)(x) = \\
& (\langle \rangle, null)
\end{aligned}$$

$$\begin{aligned}
& ast2ir_{lit}[\llbracket true \rrbracket(\Sigma)(x) = \\
& (\langle \rangle, true)
\end{aligned}$$

$$\begin{aligned}
& ast2ir_{lit}[\llbracket false \rrbracket(\Sigma)(x) = \\
& (\langle \rangle, false)
\end{aligned}$$

$$\begin{aligned}
& ast2ir_{lit}[\llbracket num \rrbracket(\Sigma)(x) = \\
& (\langle \rangle, num)
\end{aligned}$$

$$\begin{aligned}
& ast2ir_{lit}[\llbracket str \rrbracket(\Sigma)(x) = \\
& (\langle \rangle, str)
\end{aligned}$$
ast2ir_{lit}[reg](Σ)

Regular expressions are desugared into construction of **RegExp** objects by **Disambiguator**

$$\begin{aligned}
& ast2ir_m[\llbracket pr : e \rrbracket(\Sigma)(y) = \\
& \text{LET } (s^*, e) = ast2ir_e[\llbracket e \rrbracket(\Sigma)(y)] \\
& \text{IN } (s^*, ast2ir_{pr}[\llbracket pr \rrbracket : e])
\end{aligned}$$

$$\begin{aligned}
& ast2ir_m[\llbracket get\ pr() \{fd^* vd^* s^*\} \rrbracket(\Sigma)(x) = \\
& (\langle \rangle, get\ ast2ir_{pr}[\llbracket pr \rrbracket(\diamond this, \diamond arguments)\{ \\
& \quad (ast2ir_{fd}[\llbracket fd \rrbracket(\Sigma))^* \\
& \quad (ast2ir_{vd}[\llbracket vd \rrbracket(\Sigma))^* \\
& \quad (ast2ir_s[\llbracket s \rrbracket(\Sigma; \diamond this; \diamond arguments))^* \})
\end{aligned}$$

$$\begin{aligned}
& ast2ir_m[\llbracket set\ pr(x) \{fd^* vd^* s^*\} \rrbracket(\Sigma)(y) = \\
& (\langle \rangle, set\ ast2ir_{pr}[\llbracket pr \rrbracket(\diamond this, \diamond arguments)\{ \\
& \quad (ast2ir_{fd}[\llbracket fd \rrbracket(\Sigma))^* \\
& \quad \text{var } x \\
& \quad (ast2ir_{vd}[\llbracket vd \rrbracket(\Sigma))^* \\
& \quad x = \diamond arguments["0"]; \\
& \quad \text{where } x \text{ is not the name of any of fd} \\
& \quad (ast2ir_s[\llbracket s \rrbracket(\Sigma; \diamond this; \diamond arguments))^* \})
\end{aligned}$$

A.4 CFG

This section describes each construct of the SAFE CFG in both the BNF notation and its corresponding implementation. The implementation of CFG nodes is available at:

`$SAFE_HOME/src/main/scala/kr/ac/kaist/safe/nodes/cfg/`

A.4.1 Syntax of CFG

$$\begin{aligned}
& cfg \in \text{CFG} = \wp(\text{CFGFunction}) \\
& f, \langle B, \hookrightarrow, C \rangle \in \text{CFGFunction} = \\
& \quad \wp(\text{CFGBlock}) \times \wp(\text{CFGEdge}) \times \wp(\text{CFGCallTriple}) \\
& b \in \text{CFGBlock} = \text{Entry} \mid \text{Exit} \mid \text{ExitExc} \mid \text{Call} \mid \text{AfterCall} \\
& \quad \mid \text{AfterCatch} \mid \text{NormalBlock} \mid \text{ModelBlock} \\
& \text{NormalBlock} = \text{CFGInst}^+ \\
& \text{CFGEdge} = \text{CFGBlock} \times \text{CFGBlock} \times \text{EdgeType} \\
& \text{EdgeType} = \text{EdgeNormal} \mid \text{EdgeExc} \\
& C \in \text{CFGCallTriple} = \text{Call} \times \text{AfterCall} \times \text{AfterCatch} \\
& i \in \text{CFGInst} = \text{CFGNormalInst} \mid \text{CFGCallInst} \\
& x \in \text{CFGId} = \text{String} \times \text{VarKind} \\
& s \in \text{String} \\
& \text{VarKind} ::= \text{GlobalVar} \mid \text{PureLocalVar} \\
& \quad \mid \text{CapturedVar} \mid \text{CapturedCatchVar} \\
& a \in \text{Address} = \text{Integer} \\
& v \in \text{EJSVal} = \text{Number} \mid \text{String} \mid \text{Boolean} \mid \text{Null} \\
& \quad \mid \text{Undefined}
\end{aligned}$$

CFGNormalInst

$$\begin{aligned}
& ::= x := \text{alloc}(e^?)@a && \text{CFGAlloc} \\
& \mid x := \text{allocArray}(n)@a && \text{CFGAllocArray} \\
& \mid x := \text{allocArg}(n)@a && \text{CFGAllocArg} \\
& \mid x := \text{enterCode}(e) && \text{CFGEnterCode} \\
& \mid x := e && \text{CFGExprStmt} \\
& \mid x := \text{delete}(e) && \text{CFGDelete} \\
& \mid x := \text{delete}(e, e) && \text{CFGDeleteProp} \\
& \mid e[e] := e && \text{CFGStore} \\
& \mid e[s] := e && \text{CFGStoreStringIdx} \\
& \mid x := \text{function } x^?(f)@(a, a(, a)^?) && \text{CFGFunExpr} \\
& \mid \text{assert}(e) && \text{CFGAssert} \\
& \mid \text{catch}(x) && \text{CFGCatch} \\
& \mid \text{return}(e^?) && \text{CFGReturn} \\
& \mid \text{throw}(e) && \text{CFGThrow} \\
& \mid \text{noop} && \text{CFGNoOp} \\
& \mid x := x(e^*)(@a)^? && \text{CFGInternalCall}
\end{aligned}$$

CFGCallInst

$$\begin{aligned}
& ::= \text{call}(e, e, e)@(a, a) && \text{CFGCall} \\
& \mid \text{construct}(e, e, e)@(a, a) && \text{CFGConstruct}
\end{aligned}$$
 $e \in \text{CFGExpr}$

$$\begin{aligned}
& ::= x && \text{CFGVarRef} \\
& \mid e \otimes e && \text{CFGBin} \\
& \mid \ominus e && \text{CFGUn} \\
& \mid e[e] && \text{CFGLoad} \\
& \mid \text{this} && \text{CFGThis} \\
& \mid v && \text{CFGVal}
\end{aligned}$$

$$\ominus ::= \text{void} \mid \text{typeof} \mid + \mid - \mid \sim \mid !$$

$$\begin{aligned}
& \otimes ::= \text{instanceof} \mid \text{in} \mid | \mid \& \mid ^ \mid \ll \mid \gg \mid + \mid - \mid * \mid / \mid \% \\
& \mid == \mid != \mid === \mid !== \mid < \mid > \mid <= \mid >=
\end{aligned}$$

- Entry, Exit, and ExitExc blocks always uniquely exist for each CFGFunction, and an Entry block has no predecessor, Exit and ExitExc blocks have no successor.

- $\forall \langle B, _, _ \rangle \in \text{CFGFunction}.$
 $(\nexists \text{entry} \in \text{Entry} \cap B),$
 $(\nexists \text{exit} \in \text{Exit} \cap B),$
 $(\nexists \text{exitExc} \in \text{ExitExc} \cap B)$
- $\forall \langle B, \hookrightarrow, _ \rangle \in \text{CFGFunction}.$ $\nexists b \in B.$ s.t.
 $(b \hookrightarrow \text{entry}) \vee (\text{exit} \hookrightarrow b) \vee (\text{exitExc} \hookrightarrow b)$ where
 $\text{entry} \in \text{Entry}, \text{exit} \in \text{Exit}, \text{exitExc} \in \text{ExitExc}$

- Each edge connects 2 blocks in the same CFGFunction.

- $\forall \langle B, \hookrightarrow, _ \rangle \in \text{CFGFunction}.$ $\hookrightarrow \subseteq B \times B$

- Each CFGFunction records Call, AfterCall, and AfterCatch without any edges between them, and the triple consisting of them is unique.

- $\forall \langle B, _, C \rangle \in \text{CFGFunction}.$
 $\forall \text{call} \in \text{Call} \cap B.$
 $\forall \text{acall} \in \text{AfterCall} \cap B.$
 $\forall \text{acatch} \in \text{AfterCatch} \cap B.$
 $(\text{call}, _, _), (_, \text{acall}, _), (_, _, \text{acatch}) \in C$
- $\forall \langle _, _, C \rangle \in \text{CFGFunction}.$
 $\forall (\text{call}, \text{acall}, \text{acatch}), (\text{call}', \text{acall}', \text{acatch}') \in C.$
 $\text{call} = \text{call}' \Leftrightarrow \text{acall} = \text{acall}'$
 $\Leftrightarrow \text{acatch} = \text{acatch}'$

- There is no instruction after **return** in a block. There is no instruction before **catch** in a block.

- $\forall b \in \text{NormalBlock}.$
 $(i_k \in b \wedge (i_k = \text{return})) \rightarrow (\nexists i_{k'} \in b. k < k')$
- $\forall b \in \text{NormalBlock}.$
 $(i_k \in b \wedge (i_k = \text{catch})) \rightarrow (\nexists i_{k'} \in b. k > k')$

A.4.2 CFG Implementation

When a function f expects a collection B and we pass a singleton set $\{b\}$ as its argument, we abuse the notation and simply write $f(b)$ instead of $f(\{b\})$.

Helper Functions

$\text{toSeq}(S)$: $\wp(\text{Any}) \rightarrow \text{Any}^*$
= convert a set of any elements into a sequence of them by an arbitrary order.

$\text{fold}(A)(b)(f)$: $\text{Any}^* \times \text{Any}' \times (\text{Any} \times \text{Any}' \rightarrow \text{Any}') \rightarrow \text{Any}'$
= if $(\text{length}(A) = 0)$ then b
else $\text{fold}(\text{tailOf}(A))(f(\text{headOf}(A), b))(f)$

$\text{iter}(A)(f)$: $\text{Any}^* \times (\text{Any} \rightarrow \text{Unit}) \rightarrow \text{Unit}$
= if $(\text{length}(A) = 0)$ then **Unit**
else $f(\text{headOf}(A))$
 $\text{iter}(\text{tailOf}(A))(f)$

$\text{iter}(A)(\text{idx})(f)$: $\text{Any}^* \times \text{Integer} \times (\text{Any} \times \text{Integer} \rightarrow \text{Unit}) \rightarrow \text{Unit}$
= if $(\text{length}(A) = 0)$ then **Unit**
else $f(\text{headOf}(A), \text{idx})$
 $\text{iter}(\text{tailOf}(A))(\text{idx} + 1)(f)$

$\text{getTail}(B, f)$: $\text{CFGBlock}^* \times \text{CFGFunction} \rightarrow \text{CFGBlock}$
= if $(\text{length}(B) = 0)$ then
 $f.\text{createBlock}$
else if $(\text{length}(B) = 1)$ then
 $\text{headOf}(B)$
else $b \stackrel{\text{let}}{=} f.\text{createBlock}$
 $\text{cfg.addEdge}(B, b)$
 b

CFG

new CFG : $\text{CFGId}^* \rightarrow \text{CFG}$
new CFG(globalVars)
= let cfg be a new CFG
 cfg.globalFunc
 \leftarrow **new** CFGFunction($\text{cfg}, \text{"", Nil, globalVars, "top-level"}$)
 cfg

CFG: $\left\{ \begin{array}{ll} \text{globalFunc} & : \text{CFGFunction} \\ \text{createFunction} & : \text{String} \times \text{CFGId}^* \times \text{CFGId}^* \\ & \times \text{String} \rightarrow \text{CFGFunction} \\ \text{addEdge} & : \wp(\text{CFGBlock}) \times \wp(\text{CFGBlock}) \\ & \times \text{EdgeType} \rightarrow \text{Unit} \end{array} \right\}$

CFG. $\text{createFunction}(\text{argsName}, \text{argVars}, \text{localVars}, \text{name})$
= **new** CFGFunction(**this**, $\text{argsName}, \text{argVars}, \text{localVars}, \text{name}$)

CFG. $\text{addEdge}(B_1, B_2, \text{ty} = \text{EdgeNormal})$
= $\text{iter}(\text{toSeq}(B_1))(\lambda(b_1) \Rightarrow \text{iter}(\text{toSeq}(B_2))(\lambda(b_2) \Rightarrow$
 $b_1.\text{succs}(\text{ty}) \leftarrow b_1.\text{succs}(\text{ty}) \cup \{b_2\}$
 $b_2.\text{preds}(\text{ty}) \leftarrow b_2.\text{preds}(\text{ty}) \cup \{b_1\}$
 $))$

CFGFunction

new CFGFunction : $\text{CFG} \times \text{String} \times \text{CFGId}^* \times \text{CFGId}^* \times \text{String} \rightarrow \text{CFGFunction}$
new CFGFunction($\text{cfg}, \text{argsName}, \text{argVars}, \text{localVars}, \text{name}$)
= let func be a new CFGFunction
 $\text{entry} \leftarrow$ **new** Entry
 $\text{exitExc} \leftarrow$ **new** ExitExc
 $\text{exit} \leftarrow$ **new** Exit
 $\text{func.cfg} \leftarrow \text{cfg}$
 $\text{func.argumentsName} \leftarrow \text{argsName}$
 $\text{func.argVars} \leftarrow \text{argVars}$
 $\text{func.localVars} \leftarrow \text{localVars}$
 $\text{func.entry} \leftarrow \text{entry}$
 $\text{func.exit} \leftarrow \text{exit}$
 $\text{func.exitExc} \leftarrow \text{exitExc}$
 $\text{func.blocks} \leftarrow \{\text{entry}, \text{exit}, \text{exitExc}\}$
 $\text{func.captured} \leftarrow \emptyset$
 func

$$\text{CFGFunction: } \left\{ \begin{array}{ll} \text{cfg} & : \text{CFG} \\ \text{argumentsName} & : \text{String} \\ \text{argVars} & : \text{CFGId}^* \\ \text{localVars} & : \text{CFGId}^* \\ \text{entry} & : \text{Entry} \\ \text{exit} & : \text{Exit} \\ \text{exitExc} & : \text{ExitExc} \\ \text{blocks} & : \wp(\text{CFGBlock}) \\ \text{captured} & : \wp(\text{CFGId}) \\ \text{createCall} & : \text{CFGCallInst} \times \text{CFGId} \rightarrow \text{Call} \\ \text{createBlock} & : \text{Unit} \rightarrow \text{CFGBlock} \\ \text{createModelBlock} & : \text{Any} \rightarrow \text{ModelBlock} \end{array} \right\}$$

$\text{CFGFunction.createCall}(\text{callInst}, \text{retVar})$
 $= \text{new Call}(\text{this}, \text{callInst}, \text{retVar})$

$\text{CFGFunction.createBlock}$
 $= \text{new NormalBlock}(\text{this})$

$\text{CFGFunction.createModelBlock}(\text{data})$
 $= \text{new ModelBlock}(\text{this}, \text{data})$

CFGBlock

$\text{new CFGBlock} : \text{CFGFunction} \rightarrow \text{CFGBlock}$
 $\text{new CFGBlock}(\text{func}) = \text{let } b \text{ be a new CFGBlock}$
 $\quad b.\text{func} \leftarrow \text{func}$
 $\quad b.\text{succs} \leftarrow \emptyset$
 $\quad b.\text{preds} \leftarrow \emptyset$
 $\quad b$

$\text{CFGBlock: } \left\{ \begin{array}{ll} \text{func} & : \text{CFGFunction} \\ \text{succs} & : \text{EdgeType} \mapsto \wp(\text{CFGBlock}) \\ \text{preds} & : \text{EdgeType} \mapsto \wp(\text{CFGBlock}) \end{array} \right\}$

Call

$\text{new Call} : \text{CFGFunction} \times \text{CFGCallInst} \times \text{CFGId} \rightarrow \text{Call}$
 $\text{new Call}(\text{func}, \text{callInst}, \text{retVar})$
 $= \text{let } \text{call} \text{ be a new Call initialized to } \text{new CFGBlock}(\text{func})$
 $\quad \text{call.afterCall} \leftarrow \text{new AfterCall}(\text{func}, \text{call}, \text{retVar})$
 $\quad \text{call.afterCatch} \leftarrow \text{new AfterCatch}(\text{func}, \text{call})$
 $\quad \text{call.callInst} \leftarrow \text{callInst}$
 $\quad \text{call}$

$\text{Call} : \left\{ \begin{array}{ll} \text{afterCall} & : \text{AfterCall} \\ \text{afterCatch} & : \text{AfterCatch} \\ \text{callInst} & : \text{CFGCallInst} \end{array} \right\}$

AfterCall

$\text{new AfterCall} : \text{CFGFunction} \times \text{Call} \times \text{CFGId} \rightarrow \text{AfterCall}$
 $\text{new AfterCall}(\text{func}, \text{call}, \text{retVar})$
 $= \text{let } \text{acall} \text{ be a new AfterCall initialized to}$
 $\quad \text{new CFGBlock}(\text{func})$
 $\quad \text{acall.call} \leftarrow \text{call}$
 $\quad \text{acall.retVar} \leftarrow \text{retVar}$
 $\quad \text{acall}$

$\text{AfterCall} : \left\{ \begin{array}{ll} \text{call} & : \text{Call} \\ \text{retVar} & : \text{CFGId} \end{array} \right\}$

AfterCatch

$\text{new AfterCatch} : \text{CFGFunction} \times \text{Call} \rightarrow \text{AfterCatch}$
 $\text{new AfterCatch}(\text{func}, \text{call})$
 $= \text{let } \text{acatch} \text{ be a new AfterCatch initialized to}$
 $\quad \text{new CFGBlock}(\text{func})$
 $\quad \text{acatch.call} \leftarrow \text{call}$
 $\quad \text{acatch}$

$\text{AfterCatch: } \{ \text{call} : \text{Call} \}$

NormalBlock

$\text{new NormalBlock} : \text{CFGFunction} \rightarrow \text{NormalBlock}$
 $\text{new NormalBlock}(\text{func})$
 $= \text{let } b \text{ be a new NormalBlock initialized to}$
 $\quad \text{new CFGBlock}(\text{func})$
 $\quad b.\text{insts} \leftarrow \text{Nil}$
 $\quad b$

$\text{NormalBlock: } \left\{ \begin{array}{ll} \text{insts} & : \text{CFGInst}^* \\ \text{createInst} & : \text{CFGInst} \rightarrow \text{Unit} \end{array} \right\}$

$\text{NormalBlock.createInst}(\text{inst})$
 $= \text{this.insts} \leftarrow \text{this.insts} :: \text{inst}$

ModelBlock

$\text{new ModelBlock} : \text{CFGFunction} \times \text{Any} \rightarrow \text{ModelBlock}$
 $\text{new ModelBlock}(\text{func}, \text{data})$
 $= \text{let } \text{modelB} \text{ be a new ModelBlock initialized to}$
 $\quad \text{new CFGBlock}(\text{func})$
 $\quad \text{modelB.data} \leftarrow \text{data}$
 $\quad \text{modelB}$

$\text{ModelBlock: } \{ \text{data} : \text{Any} \}$

CFGId

$\text{new CFGId} : \text{String} \times \text{VarKind} \rightarrow \text{CFGId}$
 $\text{new CFGId}(\text{name}, \text{kind})$
 $= \text{let } \text{id} \text{ be a new CFGId}$
 $\quad \text{id.name} \leftarrow \text{name}$
 $\quad \text{id.kind} \leftarrow \text{kind}$
 $\quad \text{id}$

$\text{CFGId: } \left\{ \begin{array}{ll} \text{name} & : \text{String} \\ \text{kind} & : \text{VarKind} \end{array} \right\}$

A.5 IR to CFG

This section describes the SAFE translation rules from IR to CFG, whose implementation is available at:

`$SAFE_HOME/src/main/scala/kr/ac/kaist/safe/cfg_builder/`
`CFGBuilder.scala`

$\text{LabelMap} : \text{JSLabel} \mapsto \wp(\text{CFGBlock})$
 $\text{JSLabel} = \text{RetLabel} \mid \text{ThrowLabel} \mid \text{ThrowEndLabel}$
 $\quad \mid \text{AfterCatchLabel} \mid \text{UserLabel}$
 $\text{UserLabel} = \text{String}$

$\llbracket - \rrbracket_{\text{root}} : \text{IRRoot} \rightarrow \text{CFG}$
 $\llbracket - \rrbracket_{\text{fdvars}} : \text{IRFunDecl}^* \rightarrow \text{CFGId}^*$
 $\llbracket - \rrbracket_{\text{vds}} : \text{IRVarStmt}^* \rightarrow \text{CFGId}^*$
 $\llbracket - \rrbracket_{\text{args}} : \text{IRStmt}^* \rightarrow \text{CFGId}^*$
 $\llbracket - \rrbracket_{\text{functional}} : \text{IRFunctional} \rightarrow \text{CFGFunction}$
 $\llbracket - \rrbracket_{\text{fd}} : \text{IRFunDecl} \times \text{CFGFunction} \times \text{NormalBlock} \rightarrow \text{Unit}$
 $\llbracket - \rrbracket_{\text{fd}^*} : \text{IRFunDecl}^* \times \text{CFGFunction} \times \text{NormalBlock} \rightarrow \text{Unit}$
 $\llbracket - \rrbracket_{\text{stmt}} : \text{IRStmt} \times \text{CFGFunction} \times \text{CFGBlock}^* \times \text{LabelMap}$
 $\rightarrow \text{CFGBlock}^* \times \text{LabelMap}$
 $\llbracket - \rrbracket_{\text{stmt}^*} : \text{IRStmt}^* \times \text{CFGFunction} \times \text{CFGBlock}^* \times \text{LabelMap}$
 $\rightarrow \text{CFGBlock}^* \times \text{LabelMap}$
 $\llbracket - \rrbracket_{\text{mem}} : \text{IRField} \times \text{NormalBlock} \times \text{IRId} \rightarrow \text{Unit}$
 $\llbracket - \rrbracket_{\text{elem}} : \text{IRExpr} \times \text{NormalBlock} \times \text{IRId} \times \text{Integer} \rightarrow \text{Unit}$
 $\llbracket - \rrbracket_{\text{expr}} : \text{IRExpr} \rightarrow \text{CFGExpr}$
 $\llbracket - \rrbracket_{\text{id}} : \text{IRId} \rightarrow \text{CFGId}$

We use the following global variables for describing the translation from IR to CFG:

$* \text{catchVars} : \wp(\text{String})$ initialized to \emptyset
 $* \text{cfg} : \text{CFG}$
 $* \text{cfgIdMap} : \text{String} \mapsto \text{CFGId}$ initialized to $[\]$
 $* \text{currentFunc} : \text{CFGFunction}$

$\llbracket \text{IRRoot}(fds, vds, stmts) \rrbracket_{\text{root}}$
 $= \text{globalVars} \stackrel{\text{let}}{=} \llbracket fds \rrbracket_{\text{fdvars}} ++ \llbracket vds \rrbracket_{\text{vds}}$
 $\text{cfg} \leftarrow \text{new CFG}(\text{globalVars})$
 $\text{globalFunc} \stackrel{\text{let}}{=} \text{cfg.globalFunc}$
 $\text{currentFunc} \leftarrow \text{globalFunc}$
 $\text{startBlock} \stackrel{\text{let}}{=} \text{globalFunc.createBlock}$
 $\text{cfg.addEdge}(\text{globalFunc.entry}, \text{startBlock})$
 $\llbracket fds \rrbracket_{\text{fd}^*}(\text{globalFunc}, \text{startBlock})$
 $(B, L) \stackrel{\text{let}}{=} \llbracket stmts \rrbracket_{\text{stmt}^*}(\text{globalFunc}, \text{startBlock}, [\])$
 $\text{cfg.addEdge}(B, \text{globalFunc.exit})$
 $\text{cfg.addEdge}(L(\text{ThrowLabel}), \text{globalFunc.exitExc}, \text{EdgeExc})$
 $\text{cfg.addEdge}(L(\text{ThrowEndLabel}), \text{globalFunc.exitExc})$
 $\text{cfg.addEdge}(L(\text{AfterCatchLabel}), \text{globalFunc.exitExc})$
 cfg

$\llbracket fds \rrbracket_{\text{fdvars}}$
 $= \text{fold}(fds)(\text{Nil})(\lambda(\text{vars}, \text{function } f(_, _) \Rightarrow \text{vars} :: \llbracket f \rrbracket_{\text{id}})$

$\llbracket vds \rrbracket_{\text{vds}}$
 $= \text{fold}(vds)(\text{Nil})(\lambda(\text{vars}, \text{var } x \Rightarrow \text{vars} :: \llbracket x \rrbracket_{\text{id}})$

$\llbracket args \rrbracket_{\text{args}}$
 $= \text{fold}(args)(\text{Nil})(\lambda(\text{args}, x = \text{arguments}[i]) \Rightarrow \text{args} :: \llbracket x \rrbracket_{\text{id}})$

$\llbracket \text{IRFunctional}(_, \text{name}, \text{params}, \text{args}, \text{fds}, \text{vds}, \text{body}) \rrbracket_{\text{functional}}$
 $= \text{argVars} \stackrel{\text{let}}{=} \llbracket \text{args} \rrbracket_{\text{args}}$
 $\text{localVars} \stackrel{\text{let}}{=} \llbracket fds \rrbracket_{\text{fdvars}} ++ \llbracket vds \rrbracket_{\text{vds}}$
 $\text{argumentsName} \stackrel{\text{let}}{=} \text{headOf}(\text{params}).\text{name}$
 $\text{nameStr} \stackrel{\text{let}}{=} \text{name.name}$
 $\text{newFunc} \stackrel{\text{let}}{=} \text{cfg.createFunction}(\text{argumentsName}, \text{argVars}, \text{localVars}, \text{nameStr})$
 $\text{oldFunc} \stackrel{\text{let}}{=} \text{currentFunc}$
 $\text{currentFunc} \leftarrow \text{newFunc}$
 $\text{startBlock} \stackrel{\text{let}}{=} \text{newFunc.createBlock}$
 $\text{cfg.addEdge}(\text{newFunc.entry}, \text{startBlock})$
 $\llbracket fds \rrbracket_{\text{fd}^*}(\text{newFunc}, \text{startBlock})$
 $(B, L) \stackrel{\text{let}}{=} \llbracket \text{body} \rrbracket_{\text{stmt}^*}(\text{newFunc}, \text{startBlock}, [\])$
 $\text{cfg.addEdge}(B, \text{newFunc.exit})$
 $\text{cfg.addEdge}(L(\text{RetLabel}), \text{newFunc.exit})$
 $\text{cfg.addEdge}(L(\text{ThrowLabel}), \text{newFunc.exitExc}, \text{EdgeExc})$
 $\text{cfg.addEdge}(L(\text{ThrowEndLabel}), \text{newFunc.exitExc})$
 $\text{cfg.addEdge}(L(\text{AfterCatchLabel}), \text{newFunc.exitExc})$
 $\text{currentFunc} \leftarrow \text{oldFunc}$
 newFunc

$\llbracket \text{IRFunDecl}(\text{functional}) \rrbracket_{\text{fd}}(\text{func}, \text{block})$
 $= \text{lhs} \stackrel{\text{let}}{=} \llbracket \text{functional.name} \rrbracket_{\text{id}}$
 $\text{newFunc} \stackrel{\text{let}}{=} \llbracket \text{functional} \rrbracket_{\text{functional}}$
 $(\text{addr1}, \text{addr2}) \stackrel{\text{let}}{=} (\text{new Address}, \text{new Address})$
 $\text{block.createInst}(\text{lhs} := \text{function } (\text{newFunc}) @ (\text{addr1}, \text{addr2}))$

$\llbracket fds \rrbracket_{\text{fd}^*}(\text{func}, \text{block})$
 $= \text{iter}(fds)(\lambda(f) \Rightarrow \llbracket f \rrbracket_{\text{fd}}(\text{func}, \text{block}))$

$\llbracket \text{IRStmtUnit}(\text{stmts}) \rrbracket_{\text{stmt}}(f, B, L)$
 $= \llbracket \text{stmts} \rrbracket_{\text{stmt}^*}(f, B, L)$

$\llbracket \text{IRSeq}(\text{stmts}) \rrbracket_{\text{stmt}}(f, B, L)$
 $= \llbracket \text{stmts} \rrbracket_{\text{stmt}^*}(f, B, L)$

$\llbracket \text{IRFunExpr}(\text{lhs}, \text{functional}) \rrbracket_{\text{stmt}}(f, B, L)$
 $= \text{newFunc} \stackrel{\text{let}}{=} \llbracket \text{functional} \rrbracket_{\text{functional}}$
 $(\text{addr1}, \text{addr2}, \text{addr3}) \stackrel{\text{let}}{=} (\text{new Address}, \text{new Address}, \text{new Address})$

$\text{name} \stackrel{\text{let}}{=} \llbracket \text{functional.name} \rrbracket_{\text{id}}$
 $\text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$
 $\text{if } (\text{name.kind} = \text{CapturedVar})$
 $\text{tailBlock.createInst}(\text{lhs} := \text{function } \text{name} (\text{newFunc}) @ (\text{addr1}, \text{addr2}, \text{addr3}))$

else
 $\text{tailBlock.createInst}(\text{lhs} := \text{function } (\text{newFunc}) @ (\text{addr1}, \text{addr2}))$
 $(\text{tailBlock}, L)$

$$\begin{aligned}
& \llbracket x = \{members\} \rrbracket_{stmt}(f, B, L) \\
&= tailBlock \stackrel{let}{=} getTail(B, f) \\
&\quad addr1 \stackrel{let}{=} \text{new Address} \\
&\quad tailBlock.createInst(x := alloc() @ addr1) \\
&\quad iter(members)(\lambda(m) \Rightarrow \llbracket m \rrbracket_{mem}(tailBlock, x)) \\
&\quad (tailBlock, L[ThrowLabel \mapsto tailBlock \cup L(ThrowLabel)])
\end{aligned}$$

$$\begin{aligned}
& \llbracket x = \{members, proto\} \rrbracket_{stmt}(f, B, L) \\
&= tailBlock \stackrel{let}{=} getTail(B, f) \\
&\quad protoId \stackrel{let}{=} \llbracket proto \rrbracket_{expr} \\
&\quad addr1 \stackrel{let}{=} \text{new Address} \\
&\quad tailBlock.createInst(x := alloc(protoId) @ addr1) \\
&\quad iter(members)(\lambda(m) \Rightarrow \llbracket m \rrbracket_{mem}(tailBlock, x)) \\
&\quad (tailBlock, L[ThrowLabel \mapsto tailBlock \cup L(ThrowLabel)])
\end{aligned}$$

$$\begin{aligned}
& \llbracket \text{try}\{body\} \text{ catch}(x) \{catIR\} \rrbracket_{stmt}(f, B, L) \\
&= catchVars \leftarrow catchVars \cup \{x.name\} \\
&\quad tryB \stackrel{let}{=} f.createBlock \\
&\quad cfg.addEdge(B, tryB) \\
&\quad catB \stackrel{let}{=} f.createBlock \\
&\quad catB.createInst(catch(\llbracket x \rrbracket_{id})) \\
&\quad (trybs, trylmap) \stackrel{let}{=} \llbracket body \rrbracket_{stmt}(f, tryB, []) \\
&\quad cfg.addEdge(trylmap(ThrowLabel), catB, EdgeExc) \\
&\quad cfg.addEdge(trylmap(ThrowEndLabel), catB) \\
&\quad cfg.addEdge(trylmap(AfterCatchLabel), catB) \\
&\quad tmplmap \stackrel{let}{=} trylmap - ThrowLabel - ThrowEndLabel \\
&\quad \quad - AfterCatchLabel \\
&\quad (catchbs, catchlmap) \stackrel{let}{=} \llbracket catIR \rrbracket_{stmt}(f, catB, tmplmap) \\
&\quad (trybs ++ catchbs, L ++ catchlmap)
\end{aligned}$$

$$\begin{aligned}
& \llbracket \text{try}\{body\} \text{ finally } \{finIR\} \rrbracket_{stmt}(f, B, L) \\
&= tryB \stackrel{let}{=} f.createBlock \\
&\quad cfg.addEdge(B, tryB) \\
&\quad finB \stackrel{let}{=} f.createBlock \\
&\quad (trybs, trylmap) \stackrel{let}{=} \llbracket body \rrbracket_{stmt}(f, tryB, []) \\
&\quad (finbs, finlmap) \stackrel{let}{=} \llbracket finIR \rrbracket_{stmt}(f, finB, L) \\
&\quad cfg.addEdge(trybs, finB) \\
&\quad reslmap \stackrel{let}{=} fold(trylmap - AfterCatchLabel) \\
&\quad \quad (finlmap)(\lambda((l, B'), L') \Rightarrow \\
&\quad \quad \quad dupB \stackrel{let}{=} f.createBlock \\
&\quad \quad \quad (B'', L'') \stackrel{let}{=} \llbracket finIR \rrbracket_{stmt}(f, dupB, L') \\
&\quad \quad \quad \text{if } (l = \text{ThrowLabel}) \\
&\quad \quad \quad \quad cfg.addEdge(trylmap(AfterCatchLabel), dupB) \\
&\quad \quad \quad \quad cfg.addEdge(B', dupB, EdgeExc) \\
&\quad \quad \quad \quad L''[ThrowEndLabel \mapsto (L''(ThrowEndLabel) ++ B'')] \\
&\quad \quad \quad \text{else} \\
&\quad \quad \quad \quad cfg.addEdge(B', dupB, EdgeExc) \\
&\quad \quad \quad \quad L''[l \mapsto (L''(l) ++ B'')] \\
&\quad \quad \quad) \\
&\quad \quad (finbs, reslmap)
\end{aligned}$$

$$\begin{aligned}
& \llbracket \text{try}\{body\} \text{ catch}(x) \{catIR\} \text{ finally } \{finIR\} \rrbracket_{stmt}(f, B, L) \\
&= catchVars \leftarrow catchVars \cup \{x.name\} \\
&\quad tryB \stackrel{let}{=} f.createBlock \\
&\quad cfg.addEdge(B, tryB) \\
&\quad catB \stackrel{let}{=} f.createBlock \\
&\quad catB.createInst(catch(\llbracket x \rrbracket_{id})) \\
&\quad finB \stackrel{let}{=} f.createBlock \\
&\quad (trybs, trylmap) \stackrel{let}{=} \llbracket body \rrbracket_{stmt}(f, tryB, []) \\
&\quad cfg.addEdge(trylmap(ThrowLabel), catB, EdgeExc) \\
&\quad cfg.addEdge(trylmap(ThrowEndLabel), catB) \\
&\quad cfg.addEdge(trylmap(AfterCatchLabel), catB) \\
&\quad tmplmap \stackrel{let}{=} trylmap - ThrowLabel - ThrowEndLabel \\
&\quad \quad - AfterCatchLabel \\
&\quad (catchbs, catchlmap) \stackrel{let}{=} \llbracket catIR \rrbracket_{stmt}(f, catB, tmplmap) \\
&\quad (finbs, finlmap) \stackrel{let}{=} \llbracket finIR \rrbracket_{stmt}(f, finB, L) \\
&\quad cfg.addEdge(trybs ++ catchbs, finB) \\
&\quad reslmap \stackrel{let}{=} fold(catchlmap - AfterCatchLabel) \\
&\quad \quad (finlmap)(\lambda((l, B'), L') \Rightarrow \\
&\quad \quad \quad dupB \stackrel{let}{=} f.createBlock \\
&\quad \quad \quad (B'', L'') \stackrel{let}{=} \llbracket finIR \rrbracket_{stmt}(f, dupB, L') \\
&\quad \quad \quad \text{if } (l = \text{ThrowLabel}) \\
&\quad \quad \quad \quad cfg.addEdge(catchlmap(AfterCatchLabel), dupB) \\
&\quad \quad \quad \quad cfg.addEdge(B', dupB, EdgeExc) \\
&\quad \quad \quad \quad L''[ThrowEndLabel \mapsto (L''(ThrowEndLabel) ++ B'')] \\
&\quad \quad \quad \text{else} \\
&\quad \quad \quad \quad cfg.addEdge(B', dupB, EdgeExc) \\
&\quad \quad \quad \quad L''[l \mapsto (L''(l) ++ B'')] \\
&\quad \quad \quad) \\
&\quad (finbs, reslmap)
\end{aligned}$$

$$\begin{aligned}
& \llbracket x = [elems] \rrbracket_{stmt}(f, B, L) \quad \text{for IRArgs} \\
&= tailBlock \stackrel{let}{=} getTail(B, f) \\
&\quad addr \stackrel{let}{=} \text{new Address} \\
&\quad tailBlock.createInst(x := allocArg(length(elems)) @ addr) \\
&\quad iter(elems)(0)(\lambda(e, k) \Rightarrow \llbracket e \rrbracket_{elem}(tailBlock, x, k)) \\
&\quad (tailBlock, L[ThrowLabel \mapsto L(ThrowLabel) \cup \{tailBlock\}])
\end{aligned}$$

$$\begin{aligned}
& \llbracket x = [elems] \rrbracket_{stmt}(f, B, L) \quad \text{for IRRarray} \\
&= tailBlock \stackrel{let}{=} getTail(B, f) \\
&\quad addr \stackrel{let}{=} \text{new Address} \\
&\quad tailBlock.createInst(x := allocArray(length(elems)) @ addr) \\
&\quad iter(elems)(0)(\lambda(e, k) \Rightarrow \llbracket e \rrbracket_{elem}(tailBlock, x, k)) \\
&\quad (tailBlock, L[ThrowLabel \mapsto L(ThrowLabel) \cup \{tailBlock\}])
\end{aligned}$$

$$\begin{aligned}
& \llbracket x = \text{fun}(thisB, args) \rrbracket_{stmt}(f, B, L) \\
&= tailBlock \stackrel{let}{=} getTail(B, f) \\
&\quad this \stackrel{let}{=} \text{a CFG local variable of name "<>this<>"} \\
&\quad tailBlock.createInst(this := enterCode(\llbracket thisB \rrbracket_{expr})) \\
&\quad (addr1, addr2) \stackrel{let}{=} (\text{new Address}, \text{new Address}) \\
&\quad call \stackrel{let}{=} tailBlock.func.createCall(call(\llbracket fun \rrbracket_{expr}, this, \llbracket args \rrbracket_{expr}) \\
&\quad \quad @ (addr1, addr2), \llbracket x \rrbracket_{id}) \\
&\quad cfg.addEdge(tailBlock, call) \\
&\quad (call.afterCall, L[\\
&\quad \quad ThrowLabel \mapsto L(ThrowLabel) \cup \{call, tailBlock\}, \\
&\quad \quad AfterCatchLabel \mapsto L(AfterCatchLabel) \\
&\quad \quad \quad \cup \{call.afterCatch\} \\
&\quad \quad])
\end{aligned}$$

$$\llbracket \text{break } \text{label} \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{key} \stackrel{\text{let}}{=} \text{label.name}$$

$$\text{bs} \stackrel{\text{let}}{=} L(\text{key}) \cup B$$

$$(\emptyset, L[\text{key} \mapsto \text{bs}])$$

$$\llbracket x_1 = \text{new } x_2(\text{args}) \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$(\text{addr1}, \text{addr2}) \stackrel{\text{let}}{=} (\text{new Address}, \text{new Address})$$

$$\text{call} \stackrel{\text{let}}{=} \text{tailBlock.func.createCall}(\text{construct}(\llbracket x_2 \rrbracket_{\text{expr}}, \llbracket \text{args}(0) \rrbracket_{\text{expr}}, \llbracket \text{args}(1) \rrbracket_{\text{expr}})$$

$$\quad @ (\text{addr1}, \text{addr2}), \llbracket x_1 \rrbracket_{\text{id}})$$

$$)$$

$$\text{cfg.addEdge}(\text{tailBlock}, \text{call})$$

$$(\text{call.afterCall}, L[\text{ThrowLabel} \mapsto L(\text{ThrowLabel}) \cup \{\text{call}, \text{tailBlock}\},$$

$$\quad \text{AfterCatchLabel} \mapsto L(\text{AfterCatchLabel})$$

$$\quad \cup \{\text{call.afterCatch}\}]$$

$$))$$

$$\llbracket x = \text{delete } e \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$\text{tailBlock.createInst}(\llbracket x \rrbracket_{\text{id}} := \text{delete}(\llbracket e \rrbracket_{\text{expr}}))$$

$$(\text{tailBlock}, L[\text{ThrowLabel} \mapsto L(\text{ThrowLabel}) \cup \{\text{tailBlock}\}])$$

$$\llbracket x = \text{delete } e_1[e_2] \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$\text{tailBlock.createInst}(\llbracket x \rrbracket_{\text{id}} := \text{delete}(\llbracket e_1 \rrbracket_{\text{expr}}, \llbracket e_2 \rrbracket_{\text{expr}}))$$

$$(\text{tailBlock}, L[\text{ThrowLabel} \mapsto L(\text{ThrowLabel}) \cup \{\text{tailBlock}\}])$$

$$\llbracket x = e \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$\text{tailBlock.createInst}(\llbracket x \rrbracket_{\text{id}} := \llbracket e \rrbracket_{\text{expr}})$$

$$(\text{tailBlock}, L[\text{ThrowLabel} \mapsto L(\text{ThrowLabel}) \cup \{\text{tailBlock}\}])$$

$$\llbracket \text{if}(\text{cond}) \text{trueIR} \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{trueB} \stackrel{\text{let}}{=} f.\text{createBlock}$$

$$\text{cfg.addEdge}(B, \text{trueB})$$

$$\text{falseB} \stackrel{\text{let}}{=} f.\text{createBlock}$$

$$\text{cfg.addEdge}(B, \text{falseB})$$

$$\text{cfgCond} \stackrel{\text{let}}{=} \llbracket \text{cond} \rrbracket_{\text{expr}}$$

$$\text{trueB.createInst}(\text{assert}(\text{cfgCond}))$$

$$\text{falseB.createInst}(\text{assert}(\neg \text{cfgCond}))$$

$$(B', L') \stackrel{\text{let}}{=} \llbracket \text{trueIR} \rrbracket_{\text{stmt}}(f, \text{trueB}, L)$$

$$\text{endB} \stackrel{\text{let}}{=} f.\text{createBlock}$$

$$\text{cfg.addEdge}(\{\text{falseB}, B'\}, \text{endB})$$

$$(\text{endB}, L'[\text{ThrowLabel} \mapsto L'(\text{ThrowLabel})$$

$$\quad \cup \{\text{trueB}, \text{falseB}\}])$$

$$\llbracket \text{if}(\text{cond}) \text{trueIR} \text{ else } \text{falseIR} \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{trueB} \stackrel{\text{let}}{=} f.\text{createBlock}$$

$$\text{cfg.addEdge}(B, \text{trueB})$$

$$\text{falseB} \stackrel{\text{let}}{=} f.\text{createBlock}$$

$$\text{cfg.addEdge}(B, \text{falseB})$$

$$\text{cfgCond} \stackrel{\text{let}}{=} \llbracket \text{cond} \rrbracket_{\text{expr}}$$

$$\text{trueB.createInst}(\text{assert}(\text{cfgCond}))$$

$$\text{falseB.createInst}(\text{assert}(\neg \text{cfgCond}))$$

$$(B', L') \stackrel{\text{let}}{=} \llbracket \text{trueIR} \rrbracket_{\text{stmt}}(f, \text{trueB}, L)$$

$$\text{endB} \stackrel{\text{let}}{=} f.\text{createBlock}$$

$$(B'', L'') \stackrel{\text{let}}{=} \llbracket \text{falseIR} \rrbracket_{\text{stmt}}(f, \text{falseB}, L')$$

$$\text{cfg.addEdge}(B' ++ B'', \text{endB})$$

$$(\text{endB}, L''[\text{ThrowLabel} \mapsto L''(\text{ThrowLabel})$$

$$\quad \cup \{\text{trueB}, \text{falseB}\}])$$

$$\llbracket l : \{s\} \rrbracket_{\text{stmt}}(f, B, L)$$

$$= b \stackrel{\text{let}}{=} f.\text{createBlock}$$

$$(B', L') \stackrel{\text{let}}{=} \llbracket s \rrbracket_{\text{stmt}}(f, B, L)$$

$$\text{label} \stackrel{\text{let}}{=} l.\text{name}$$

$$\text{cfg.addEdge}(B', b)$$

$$\text{cfg.addEdge}(L'(\text{label}), b)$$

$$(b, L' - l)$$

$$\llbracket \text{return} \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$\text{tailBlock.createInst}(\text{return}())$$

$$(\emptyset, L[\text{RetLabel} \mapsto L(\text{RetLabel}) \cup \{\text{tailBlock}\}])$$

$$\llbracket \text{return } e \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$\text{tailBlock.createInst}(\text{return}(\llbracket e \rrbracket_{\text{expr}}))$$

$$(\emptyset, L[\text{RetLabel} \mapsto L(\text{RetLabel}) \cup \{\text{tailBlock}\}])$$

$$\llbracket x_1[x_2] = e \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$\text{tailBlock.createInst}(\llbracket x_1 \rrbracket_{\text{expr}}[\llbracket x_2 \rrbracket_{\text{expr}}] := \llbracket e \rrbracket_{\text{expr}})$$

$$(\text{tailBlock}, L[\text{ThrowLabel} \mapsto L(\text{ThrowLabel}) \cup \{\text{tailBlock}\}])$$

$$\llbracket \text{throw } e \rrbracket_{\text{stmt}}(f, B, L)$$

$$= \text{tailBlock} \stackrel{\text{let}}{=} \text{getTail}(B, f)$$

$$\text{tailBlock.createInst}(\text{throw}(\llbracket e \rrbracket_{\text{expr}}))$$

$$(\emptyset, L[\text{ThrowLabel} \mapsto L(\text{ThrowLabel}) \cup \{\text{tailBlock}\}])$$

$$\begin{aligned}
\llbracket \text{while}(e) \ s \rrbracket_{stmt}(f, B, L) &= \text{tailBlock} \stackrel{let}{=} \text{getTail}(B, f) \\
&\quad \text{head} \stackrel{let}{=} f.createBlock \\
&\quad \text{body} \stackrel{let}{=} f.createBlock \\
&\quad \text{out} \stackrel{let}{=} f.createBlock \\
&\quad \text{cfgCond} \stackrel{let}{=} \llbracket e \rrbracket_{expr} \\
&\quad \text{body.createInst}(\text{assert}(\text{cfgCond})) \\
&\quad \text{out.createInst}(\text{assert}(\neg \text{cfgCond})) \\
&\quad \text{cfg.addEdge}(\text{tailBlock}, \text{head}) \\
&\quad \text{cfg.addEdge}(\text{head}, \text{body}) \\
&\quad \text{cfg.addEdge}(\text{head}, \text{out}) \\
&\quad (B', L') \stackrel{let}{=} \llbracket s \rrbracket_{stmt}(f, [\text{body}], L) \\
&\quad \text{cfg.addEdge}(B', \text{head}) \\
&\quad (\text{out}, L'[\text{ThrowLabel} \mapsto L'(\text{ThrowLabel}) \\
&\quad \quad \cup \{\text{body}, \text{out}\}])
\end{aligned}$$

$$\begin{aligned}
\llbracket stmts \rrbracket_{stmt^*}(f, B, L) &= \text{fold}(stmts)((B, L))(\lambda((B', L'), stmt) \Rightarrow \\
&\quad \llbracket stmt \rrbracket_{stmt}(f, B', L'))
\end{aligned}$$

$$\begin{aligned}
\llbracket mem \rrbracket_{mem}(b, x) &= \text{lhs} \stackrel{let}{=} \llbracket x \rrbracket_{expr} \\
&\quad \text{str} \stackrel{let}{=} \text{mem.prop.name} \\
&\quad \text{expr} \stackrel{let}{=} \llbracket mem.expr \rrbracket_{expr} \\
&\quad b.createInst(\text{lhs}[\text{str}] := \text{expr})
\end{aligned}$$

$$\begin{aligned}
\llbracket e \rrbracket_{elem}(b, x, k) &= \text{lhs} \stackrel{let}{=} \llbracket x \rrbracket_{expr} \\
&\quad \text{expr} \stackrel{let}{=} \llbracket e \rrbracket_{expr} \\
&\quad b.createInst(\text{lhs}["k"] := \text{expr})
\end{aligned}$$

$$\llbracket x \rrbracket_{expr} = x$$

$$\llbracket e_1 \otimes e_2 \rrbracket_{expr} = \llbracket e_1 \rrbracket_{expr} \otimes \llbracket e_2 \rrbracket_{expr}$$

$$\llbracket \ominus e \rrbracket_{expr} = \ominus \llbracket e \rrbracket_{expr}$$

$$\llbracket x[e] \rrbracket_{expr} = \llbracket x \rrbracket_{expr}[\llbracket e \rrbracket_{expr}]$$

$$\llbracket \text{this} \rrbracket_{expr} = \text{this}$$

$$\llbracket v \rrbracket_{expr} = v$$

$$\begin{aligned}
\llbracket x \rrbracket_{id} &= \text{let } kind \text{ be a VarKind based on the information} \\
&\quad \text{from CapturedVariableCollector} \\
&\quad \text{new CFGId}(x.name, kind)
\end{aligned}$$

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