

CS 512, Spring 2014

Assignment 5

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1 Problem 1

- (a) $\phi = \mathbf{G}a$
- (i) the path $q_3 \rightarrow q_4 \rightarrow q_3 \rightarrow q_4 \rightarrow q_3 \dots$ (q_3 and q_4 keep repeating) will satisfy $\mathbf{G}a$
 - (ii) No because there are other paths do not satisfy $\mathbf{G}a$
- (b) $\phi = a\mathbf{U}b$
- (i) the path $q_3 \rightarrow q_2 \dots$ (q_2 goes on forever) will satisfy $\phi = a\mathbf{U}b$
 - (ii) No because there are other paths that do not satisfy $\phi = a\mathbf{U}b$
- (c) $\phi = a\mathbf{U}X(a \wedge \neg b)$
- (i) the path $q_3 \rightarrow q_4 \rightarrow q_3 \rightarrow q_4 \rightarrow q_3 \dots$ (q_3 and q_4 keep repeating) will satisfy $\phi = a\mathbf{U}X(a \wedge \neg b)$
 - (ii) No because not all of the other paths satisfy $\phi = a\mathbf{U}X(a \wedge \neg b)$
- (d) $\phi = \mathbf{X}\neg b \wedge \mathbf{G}(\neg a \vee \neg b)$
- (i) the path $q_3 \rightarrow q_1 \rightarrow q_2 \dots$ will satisfy $\phi = \mathbf{X}\neg b \wedge \mathbf{G}(\neg a \vee \neg b)$
 - (ii) No because all the other paths do not satisfy $\phi = \mathbf{X}\neg b \wedge \mathbf{G}(\neg a \vee \neg b)$
- (e) $\mathbf{X}(a \wedge b) \wedge F(\neg a \wedge \neg b)$
- (i) the path $q_3 \rightarrow q_4 \rightarrow q_3 \rightarrow q_1 \rightarrow q_2 \dots$

(ii) No

(f) $a \wedge \mathbf{F}b$

(i) All possible paths satisfy $a \wedge \mathbf{F}b$.

(ii) Yes all paths satisfy $a \wedge \mathbf{F}b$

2 Problem 2

(a) $\mathbf{FG}\varphi$ and $\varphi \rightarrow \mathbf{X}\varphi$

The path $p \rightarrow \neg p \rightarrow \neg p \rightarrow p \rightarrow p \rightarrow p \rightarrow p \dots$ (p goes on forever) satisfies $\mathbf{FG}\varphi$ but not $\varphi \rightarrow \mathbf{X}\varphi$

(b) $\mathbf{FG}\varphi$ and $\neg\varphi\mathbf{UG}\varphi$

The path $\neg p \rightarrow \neg p \rightarrow p \rightarrow \neg p \rightarrow p \rightarrow p \rightarrow p \rightarrow p \dots$ (p go on forever) satisfies \mathbf{FG} but not $\neg\varphi\mathbf{UG}\varphi$

(c) $\mathbf{G}(\varphi \rightarrow \mathbf{X}\varphi)$ and $\neg\varphi\mathbf{UG}\varphi$ The path $p \rightarrow p \rightarrow p \rightarrow p \dots$ (p goes on forever), satisfies $\mathbf{G}(\varphi \rightarrow \mathbf{X}\varphi)$ but not $\neg\varphi\mathbf{UG}\varphi$.

(d) $\mathbf{F}(\varphi \wedge \psi)$ and $(\mathbf{F}\varphi \wedge \mathbf{F}\psi)$ The path $p \rightarrow q \rightarrow p \rightarrow q \rightarrow p \dots$ (have p, q alternate forever). This satisfies $(\mathbf{F}\varphi \wedge \mathbf{F}\psi)$ but not $\mathbf{F}(\varphi \wedge \psi)$

(e) (a) One implication has shown to be false in part a. To see why the implication is false going the other way, note that the path $\neg p \rightarrow \neg p \rightarrow \neg p \dots$ ($\neg p$ goes on forever) will satisfy $\varphi \rightarrow \mathbf{X}\varphi$ but not $\mathbf{FG}\varphi$.

(b) $\neg\varphi\mathbf{UG}\varphi$ implies $\mathbf{FG}\varphi$,

Because in the first case you will have $\mathbf{G}\varphi$ after some $\neg\varphi$ which implies sometime in the future you will have globally φ .

The implication does not exist the other way from the counter example shown in part b.

(c) $\neg\varphi\mathbf{UG}\varphi$ implies $\mathbf{G}(\varphi \rightarrow \mathbf{X}\varphi)$

Because once you have globally φ , then you know that if you see φ the next state must also have φ .

The implication does not exist the other way from the counter example shown in part c.

(d) $\mathbf{F}(\varphi \wedge \psi)$ implies $(\mathbf{F}\varphi \wedge \mathbf{F}\psi)$

Since if we know $\varphi \wedge \psi$ will be true at some point in the future, at that point both φ and ψ have to be true separately as well. (otherwise the \wedge could not be true. Then we know sometime in the future $(\mathbf{F}\varphi \wedge \mathbf{F}\psi)$

3 Problem 3

$\varphi \mathbf{U} \psi \equiv \varphi \mathbf{W} \psi \wedge \mathbf{F} \psi$
proof.

(1) $\pi \models \varphi \mathbf{U} \psi$ iff

By definition:

(2) $(\exists i \geq 1)[\pi^i \models \psi \wedge \pi^1 \models \varphi, \dots, \pi^{i-1} \models \varphi]$ iff

since we know $(\exists i \geq 1)\pi^i \models \psi$ is true, and using properties of \wedge :

(3) $((\exists i \geq 1)\pi^i \models \psi) \wedge (\exists i \geq 1)[\pi^i \models \psi \wedge \pi^1 \models \varphi, \dots, \pi^{i-1} \models \varphi]$ iff

Since we know that this statement is true, using properties of \vee :

(4) $((\exists i \geq 1)\pi^i \models \psi) \wedge (\exists i \geq 1)[\pi^i \models \psi \wedge \pi^1 \models \varphi, \dots, \pi^{i-1} \models \varphi] \vee \forall k \pi^k \models \varphi$ iff

by definitions of \mathbf{F} and \mathbf{W} :

(5) $\mathbf{F} \psi \wedge \varphi \mathbf{W} \psi$

(6) $\pi \models \mathbf{F} \psi \wedge \varphi \mathbf{W} \psi$

Since we have iff at every line we know that implications work both ways. So they are equivalent.

$\varphi \mathbf{W} \psi \equiv \varphi \mathbf{U} \psi \vee \mathbf{G} \varphi$
proof.

(1) $\pi \models \varphi \mathbf{W} \psi$ iff

By definition:

(2) $(\exists i \geq 1)[\pi^i \models \psi \wedge \pi^1 \models \varphi, \dots, \pi^{i-1} \models \varphi] \vee \forall k \pi^k \models \varphi$ iff

By definition of \mathbf{G} :

(3) $(\exists i \geq 1)[\pi^i \models \psi \wedge \pi^1 \models \varphi, \dots, \pi^{i-1} \models \varphi] \vee \mathbf{G} \varphi$ iff

by Definition of \mathbf{U} :

(4) $\varphi \mathbf{U} \psi \vee \mathbf{G} \varphi$

(5) $\pi \models \varphi \mathbf{U} \psi \vee \mathbf{G} \varphi$

Since we have iff at every line we know that implications work both ways. So they are equivalent.

4 Problem 4

function $NNF(\phi)$:

ϕ is a literal: return ϕ
 ϕ is $\neg\neg\phi_1$: return $NNF(\phi_1)$
 ϕ is $\phi_1 \wedge \phi_2$: return $NNF(\phi_1) \wedge NNF(\phi_2)$
 ϕ is $\phi_1 \vee \phi_2$: return $NNF(\phi_1) \vee NNF(\phi_2)$
 ϕ is $\neg(\phi_1 \wedge \phi_2)$: return $NNF(\neg\phi_1) \vee NNF(\neg\phi_2)$
 ϕ is $\neg(\phi_1 \vee \phi_2)$: return $NNF(\neg\phi_1) \wedge NNF(\neg\phi_2)$

 ϕ is $\mathbf{G}\phi_1$: return $\mathbf{G}NNF(\phi_1)$
 ϕ is $\neg\mathbf{G}\phi_1$: return $\mathbf{F}NNF(\neg\phi_1)$

 ϕ is $\mathbf{F}\phi_1$: return $\mathbf{F}NNF(\phi_1)$
 ϕ is $\neg\mathbf{F}\phi_1$: return $\mathbf{G}NNF(\neg\phi_1)$

 ϕ is $\mathbf{X}\phi_1$: return $\mathbf{X}NNF(\phi_1)$
 ϕ is $\neg\mathbf{X}\phi_1$: return $\mathbf{X}NNF(\neg\phi_1)$

 ϕ is $\phi_1 \mathbf{U} \phi_2$: return $NNF(\phi_1) \mathbf{U} NNF(\phi_2)$
 ϕ is $\neg(\phi_1 \mathbf{U} \phi_2)$: return $NNF(\neg\phi_1) \mathbf{R} NNF(\neg\psi)$

 ϕ is $\phi_1 \mathbf{R} \phi_2$: return $NNF(\phi_1) \mathbf{R} NNF(\phi_2)$
 ϕ is $\neg(\phi_1 \mathbf{R} \phi_2)$: return $NNF(\neg\phi_1) \mathbf{U} NNF(\neg\psi)$

 ϕ is $\phi_1 \mathbf{W} \phi_2$: return $NNF(\phi_1) \mathbf{U} NNF(\phi_2) \vee \mathbf{G} NNF(\phi)$
 ϕ is $\neg(\phi_1 \mathbf{W} \phi_2)$: return $NNF(\neg(\phi_1 \mathbf{U} \phi_2)) \wedge NNF(\neg\mathbf{G}\phi_1)$

5 Problem 5

- (a) No, take the path $s_2 \rightarrow s_4 \rightarrow s_2 \rightarrow s_4 \rightarrow s_2 \rightarrow s_4 \dots$. When you have s_2 , r is true, when you have s_4 r is false. So this cannot satisfy $\mathbf{FG}r$
- (b) Yes,
If you start from s_1 there are two ways to go s_3 and s_4 . s_3 has an r so that's okay. If we then follow a path to s_4 we see that from s_4 any state that we move to will have an r . If you start from s_2 , the next state is s_4 which like before any state you move to from there will have an r .
- (c) Yes,
If $\mathbf{X}\neg r$ then you know the next state has to be s_4 . Every state that you can move to from s_4 will have r as true. Then you know that in the next next state, r will be true. So $\mathbf{XX}r$
- (d) No, take the path $s_2 \rightarrow s_4 \rightarrow s_2 \rightarrow s_4 \rightarrow s_2 \rightarrow s_4 \dots$. When you have s_2 , q is false, when you have s_4 q is true. So this cannot satisfy $\mathbf{G}q$
- (e) Yes.
If you start from s_1 you have p and the next states contain a q or r . If you start from s_2 you start with r and then you have already satisfied the condition $q \vee r$ of the \mathbf{U} .
- (f) No.
If you take the path $s_1 \rightarrow s_4 \rightarrow s_2 \dots$ then this doesn't work.