

CS 512, Spring 2014

Assignment 2

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Due Monday January 27th

1 Problem 1

Algebraic Structure: $\tau = (T, Lt, \mathbb{N}, Rt, height)$. Define the size operation $|| : \tau \rightarrow \mathbb{N}$ and the height operation $height : \tau \rightarrow \mathbb{N}$ in the style of the definitions for the operations Node Lt, and Rt. If $t \in T$ is represented graphically by a finite binary tree, then $|t|$ should return should return the total number of leaf nodes in t + the total number of internal nodes in t , and the $height(t)$ should return the length of the longest path in t .

$$(1) \quad |t| = \begin{cases} |t_1| + |t_2|, & \text{if } t = \langle t_1 \ t_2 \rangle . \\ 1, & \text{otherwise} \end{cases}$$

$$(2) \quad height(x) = \begin{cases} \max(|t_1| + |t_2|) + 1, & \text{if } t = \langle t_1 \ t_2 \rangle . \\ 0, & \text{otherwise} \end{cases}$$

2 Problem 2

Algebraic Structure: $A = (\mathbb{N} - 0, lcm, gcd, \leq)$ where for all $m, n \in \mathbb{N} - 0$, we have that: $m \leq n$ iff "m divides n".

(a) Show that A is a lattice, where lcm plays the role of \vee and gcd plays the role of \wedge .

(b) Is A a distributive lattice? Justify your answer carefully, based on the distributivity axioms that A must satisfy.

(a)

Need to Show:

(1) Show (A, \leq) is a poset.

For all $a, b, c \in \mathbb{N} - 0$

reflexive: "a divides a" therefore $a \leq a$

anti-symmetric:

Assume $a \leq b$ and $b \leq a$. Then "a divides b" and "b divides a".

??

transitive:

Assume $a \leq b$ and $b \leq c$. Then "a divides b" and "b divides c".

If "a divides b" then b must be some multiple of a ($b = m \cdot a, m \in \mathbb{N} - 0$).

If "b divides c" then c must be some multiple of b ($c = n \cdot b, n \in \mathbb{N} - 0$).

So $c = n \cdot b = n \cdot (m \cdot a)$ where $m, n \in \mathbb{N} - 0$

But then c is some multiple of a. So "a divides c".

Therefore $a \leq c$.

(2) For all $m, n \in A$ the least upper bound of a and b in the ordering \leq exists, is unique m and is the result of the operation $m \vee n$

(3) For all $m, n \in A$ the greatest lower bound of m and n in the ordering \leq exists, is unique m and is the result of the operation $m \wedge n$

3 Problem 3: Exercise 1.5.3 parts (a),(b),(c)

(a) Show that $\{\neg, \wedge\}$, $\{\neg, \rightarrow\}$, $\{\rightarrow, \perp\}$ are adequate sets of connectives for propositional logic.

(a1) $\{\neg, \wedge\}$:

If I can show that there is an expression that is equivalent to $p \vee q$, then this will also show that there exists an equivalent for \rightarrow Since from the example we know $p \rightarrow q \equiv \neg p \vee q$.