

LeetCode 10

<https://leetcode.com/problems/regular-expression-matching/description/>

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Description

10. Regular Expression Matching

Implement regular expression matching with support for '.' and '*'.

'.' Matches any single character.
'*' Matches zero or more of the preceding element.

The matching should cover the entire **input** string (not partial).

The function prototype should be:
bool isMatch(const char *s, const char *p)

Some examples:

isMatch("aa","a") → false
isMatch("aa","aa") → true
isMatch("aaa","aa") → false
isMatch("aa", "a*") → true
isMatch("aa", ".*") → true
isMatch("ab", ".*") → true
isMatch("aab", "c*a*b") → true

Idea Report

(Similar to edit distance) We use boolean[] f to represent the DP solution space, where f[i][j] means boolean representation to match for the first i characters of s and first j characters of p. For the given example "aab" and "c*a*b", we can have:

	0 ''	1 c	2 *	3 a	4 *	5 b

0 ''	T	F	T	F	T	F
1 a	F	F	F	T	T	F
2 a	F	F	F	F	T	F
3 b	F	F	F	F	F	T

The first row is the case when s is "", so as long as p has can match an empty string, the cell is true. The first column is the case when p is "", so as long as s is not "", the cell is false, these are the base cases we need to cover. For any $f[i][j]$ we need to discuss the situations:

- If $s.charAt(i - 1) == p.charAt(j - 1)$, then for $f[i][j]$, we need to see $f[i - 1][j - 1]$, e.g. in the case $f[1][3] = f[0][2] = \text{true}$.
- If $p.charAt(j - 1) == '.'$, it is the same situation as the above one, we check $f[i - 1][j - 1]$.
- If $p.charAt(j - 1) == '*'$, there may be a few different cases:
 - We need to check the character before '*' in p ($p[j - 2]$) and see if it is the same with the character in $s[i - 1]$. If $p[j - 2] != '.'$ && $p[j - 2] != s[i - 1]$, that means we cannot use this '*', which means '*' = 0, so we need to check $f[i][j - 2]$.
 - Otherwise, we are able to use '*', but we don't necessary have to, so there may be three different situations:
 - We do not use '*' at all, then we check $f[i][j - 2]$, at $f[2][4]$ we check $f[2][2]$
 - We use '*' as 1, so we check $f[i][j - 1]$, at $f[2][4]$ we check $f[2][3]$
 - We use '*' multiple times, we check $f[i - 1][j]$, at $f[2][4]$ we check $f[1][4]$.
- Other cases $f[i][j]$ are all false.

Code:

```
class Solution {
    public boolean isMatch(String s, String p) {
        int len1 = s.length();
        int len2 = p.length();
        boolean[][] f = new boolean[len1 + 1][len2 + 1];
        f[0][0] = true;
        for (int i = 1; i <= len2; i++) {
            f[0][i] = p.charAt(i - 1) == '*' && f[0][i - 2];
        }

        for (int i = 1; i <= len1; i++) {
            for (int j = 1; j <= len2; j++) {
                if (s.charAt(i - 1) == p.charAt(j - 1)
                    || p.charAt(j - 1) == '.') {

```

```

        f[i][j] = f[i - 1][j - 1];
    } else if (p.charAt(j - 1) == '*') {
        if (s.charAt(i - 1) != p.charAt(j - 2)
            && p.charAt(j - 2) != '.') {
            f[i][j] = f[i][j - 2];
        } else {
            f[i][j] = f[i - 1][j] || f[i][j - 1] || f[i][j - 2];
        }
    }
}

return f[len1][len2];
}

```

Summary

- List the transition matrix really helps analyzing the state transfer.