Description

57. Insert Interval

Given a set of non-overlapping intervals, insert a new interval into the intervals (merge if necessary).

You may assume that the intervals were initially sorted according to their start times.

Example 1:

Given intervals [1,3],[6,9], insert and merge [2,5] in as [1,5],[6,9].

Example 2:

Given [1,2],[3,5],[6,7],[8,10],[12,16], insert and merge [4,9] in as [1,2],[3,10],[12,16].

This is because the new interval [4,9] overlaps with [3,5],[6,7],[8,10].

Idea

Μ

Java

```
class Solution {
  public List<Interval> insert(List<Interval> intervals, Interval newInterval) {
    List<Interval> res = new ArrayList<>();

  int i = 0;
  while (i < intervals.size() && intervals.get(i).end < newInterval.start) {
    res.add(intervals.get(i++));
  }

  while (i < intervals.size() && newInterval.end >= intervals.get(i).start) {
    newInterval.start = Math.min(newInterval.start, intervals.get(i).start);
    newInterval.end = Math.max(newInterval.end, intervals.get(i).end);
    i++;
```

```
res.add(newInterval);

while (i < intervals.size()) {
    res.add(intervals.get(i++));
  }
  return res;
}</pre>
```

C++

```
class Solution {
public:
    vector<Interval> insert(vector<Interval>& intervals, Interval newInterval) {
        vector<Interval> res;
        int i = 0;
        while (i < intervals.size() && intervals[i].end < newInterval.start) {</pre>
            res.push_back(intervals[i++]);
        }
        while (i < intervals.size() && intervals[i].start <= newInterval.end) {</pre>
            newInterval.start = min(newInterval.start, intervals[i].start);
            newInterval.end = max(newInterval.end, intervals[i].end);
            i++;
        }
        res.push_back(newInterval);
        while (i < intervals.size()) {</pre>
            res.push_back(intervals[i++]);
        }
        return res;
   }
};
```

Summary

- $O(n^3)$ brute force may be speed up by $O(n\log n)$ sorting then achieve $O(n^2)$ time complexity.
- Fix outter loop and optimize inner loop.