



# ***ZERO-KNOWLEDGE PROOFS (ZKP)***

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## DISCUSSION

# Food for Thought



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# AGENDA

- 01** ZKPs: Cryptographic Intro
- 02** Applications
- 03** ZK-Rollups and Layer 2
- 04** Developer Ecosystems
- 05** Discussion

# AGENDA

**01** ZKPs: Cryptographic Intro

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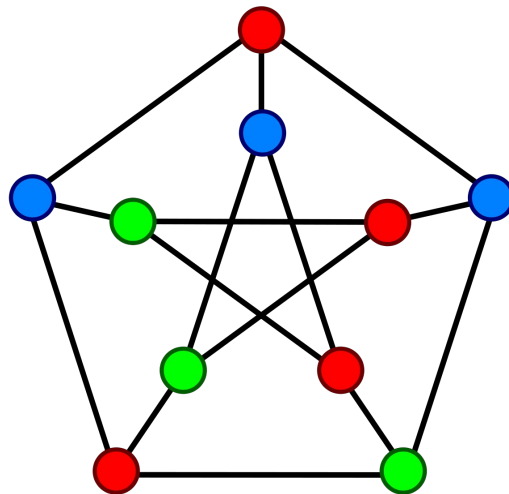
# ZKPs: Definition and Origin

- ZKPs: Zero-Knowledge Proofs
- Allows one party (the prover) to prove to another (the verifier) that a statement is true, without revealing any information beyond the validity of the statement itself.
- Coined by MIT researchers Shafi Goldwasser, Silvio Micali, Charles Rackoff

## EXAMPLE

# 3-Coloring Problem

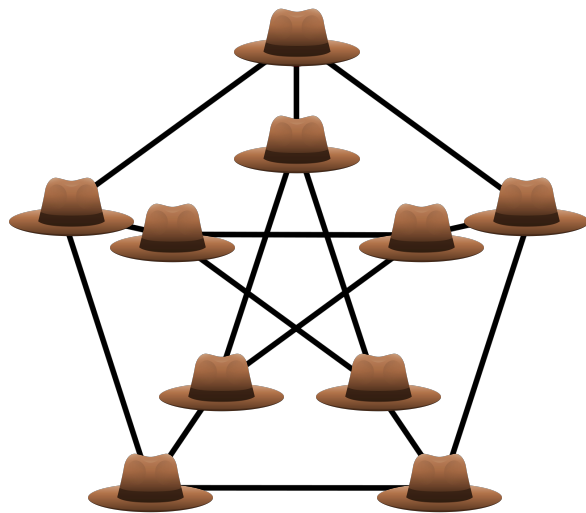
- Given a graph  $G$ , can all vertices be colored with 3 colors so that no edge has vertices of the same color?
- This is an NP-complete problem
  - i.e. very hard in the general case
- Andy wishes to prove existence of solution to Conan
  - Conan starts a new uncolored version of  $G$ .
  - Privately, Andy colors all vertices with a **permutation** of the 3 colors, and then hides all vertices under "hat"s.
  - Conan picks an edge, of which Andy reveals both vertices.
  - If the vertices are of the same color, Andy loses.
  - If the vertices are of different colors, then it remains possible that Andy has a solution. To start another round of *challenge-response*, they repeat from the first step.



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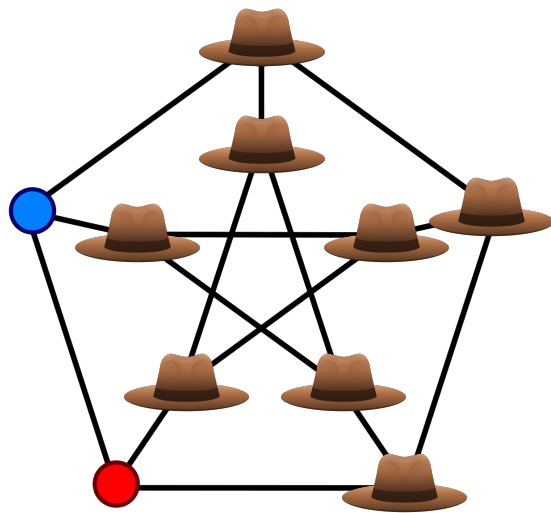




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# Three Important Properties

Valid ZKPs must satisfy the following conditions:

## **Completeness**

If B is telling the truth,  
then B will eventually  
convince A.

## **Soundness**

B can only convince A if  
they're telling the truth.

## **Zero-Knowledgeness**

A does not learn  
anything else about B's  
proof.

# ZK-SNARKs

- ZK-SNARK: Zero-Knowledge Succinct Non-Interactive ARGument of Knowledge
  - Succinct: short and well-defined
  - Non-interactive: once for all; no challenge-response dynamic
  - Sublinear verification time with respect to input
- Example written in Circom on the right

```

51 lines (42 sloc) | 1.43 KB
Raw Blame
1 include "../node_modules/circomlib/circuits/mimc/mimc.circom";
2
3 // Computes MIMC([left, right])
4 template HashLeftRight() {
5   signal input left;
6   signal input right;
7   signal output hash;
8
9   component hasher = MIMCSponge(2, 1);
10  hasher.ins[0] <== left;
11  hasher.ins[1] <== right;
12  hasher.k <== 0;
13  hash <== hasher.outs[0];
14 }
15
16 // If s == 0 returns [in[0], in[1]]
17 // If s == 1 returns [in[1], in[0]]
18 template DualMux() {
19   signal input in[2];
20   signal input s;
21   signal output out[2];
22
23   s * (1 - s) == 0
24   out[0] <== (in[1] - in[0]) * s + in[0];
25   out[1] <== (in[0] - in[1]) * s + in[1];
26 }
27
28 // Verifies that merkle proof is correct for given merkle root and a leaf
29 // pathIndices input is an array of 0/1 selectors telling whether given pathElement is on the left or right side of merkle path
30 template MerkleTreeChecker(levels) {
31   signal input leaf;
32   signal input root;
33   signal input pathElements[levels];
34   signal input pathIndices[levels];
35
36   component selectors[levels];
37   component hashers[levels];
38
39   for (var i = 0; i < levels; i++) {
40     selectors[i] = DualMux();
41     selectors[i].in[0] <== i == 0 ? leaf : hashers[i - 1].hash;
42     selectors[i].in[1] <== pathElements[i];
43     selectors[i].s <== pathIndices[i];
44
45     hashers[i] = HashLeftRight();
46     hashers[i].left <== selectors[i].out[0];
47     hashers[i].right <== selectors[i].out[1];
48   }
49
50   root == hashers[levels - 1].hash;
51 }

```

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# Blockchains and ZK-SNARKs

- Due to the technical design of blockchains, everything on-chain is **fully transparent** to everyone.
  - This includes one's entire transaction history (including income stream), NFT possessions, gameplay stats, etc.
  - As blockchain use cases expand, the list could soon include much more sensitive information such as one's health data.
  - A land of zero privacy...
- Founding value of blockchains: zero-trust/permissionlessness
- Dilemma: permissionlessness or privacy?
- Another problem: incomplete-information for fully-on-chain game
  - Imagine a game of poker
  - Fact: Current crypto games like Axie are not fully on-chain

# Applications in DeFi



## Zcash

- A Bitcoin fork with privacy features, created in 2016
- Transactions can be "transparent" and similar to bitcoin transactions in which case they are controlled by a t-addr, or they can be "shielded" and are controlled by a z-addr.
- Allows "selective disclosure", allowing a user to prove payment for auditing purposes. One such reason is to allow private transactors the choice to comply with anti-money laundering or tax regulations.



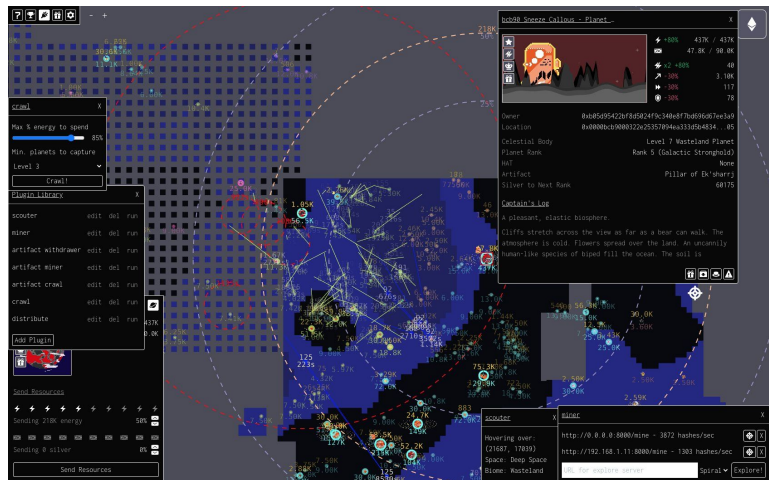
## Tornado

- A protocol for private transactions for ERC-20 tokens

$\mathcal{S}[R, h, A, f, t] = \{ \text{I KNOW } k, r \in \mathbb{B}^{248}, l \in \mathbb{B}^{16}, O \in Z_p^{16} \text{ SUCH THAT } h = H_1(k) \\ \text{AND } O \text{ is the opening of } H_2(k||r) \text{ at position } l \text{ to } R \}$

# Applications in Gaming

- In traditional games, nothing is on-chain.
- In current crypto games (Axie, Cryptomines, etc.), the ownables are on-chain in the form of NFTs, and thus tradable.
- In Dark Forest, the ownables as well as game logic is on-chain.
- Metaverse, composability, and interoperability



Dark Forest

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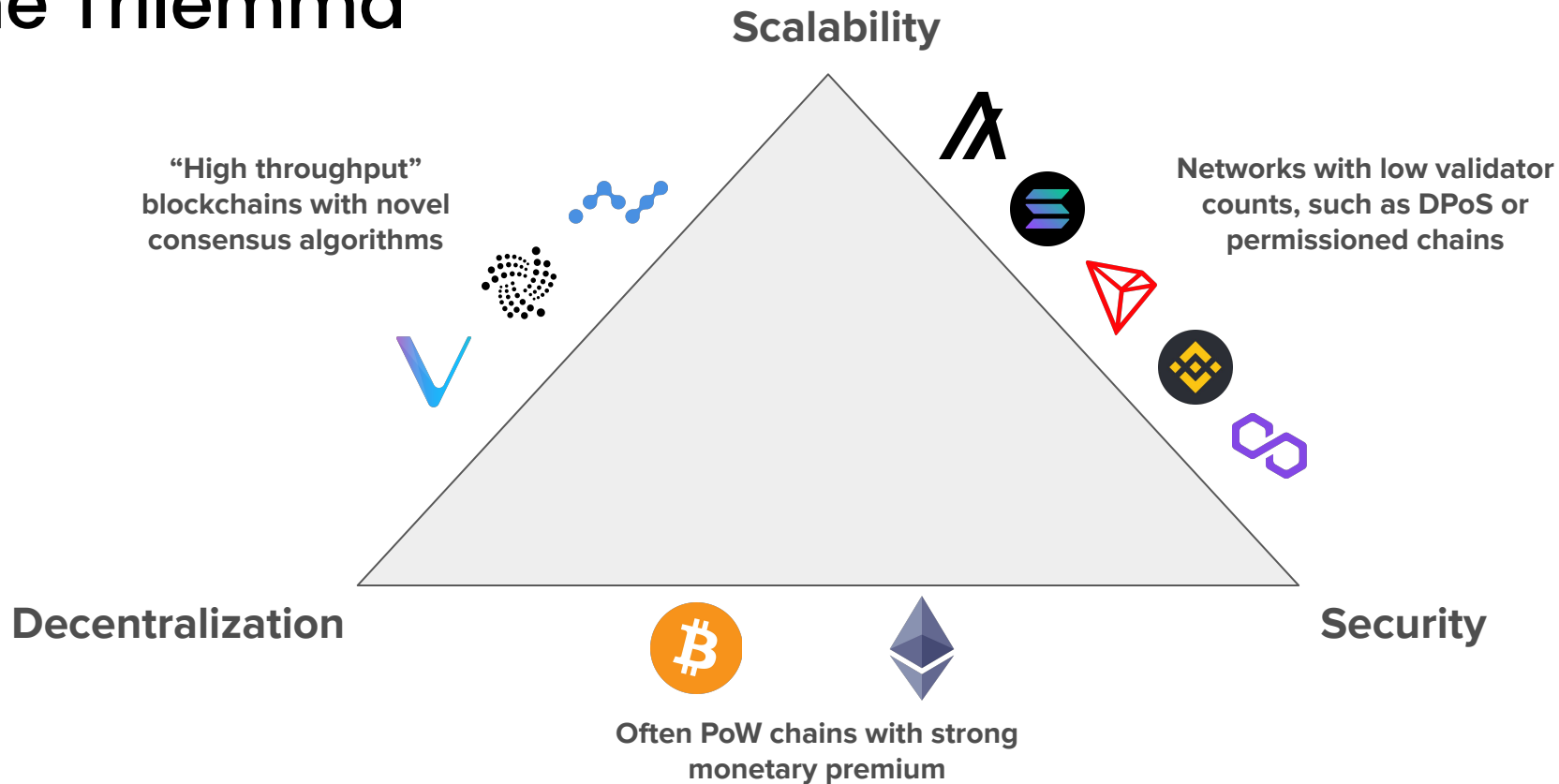
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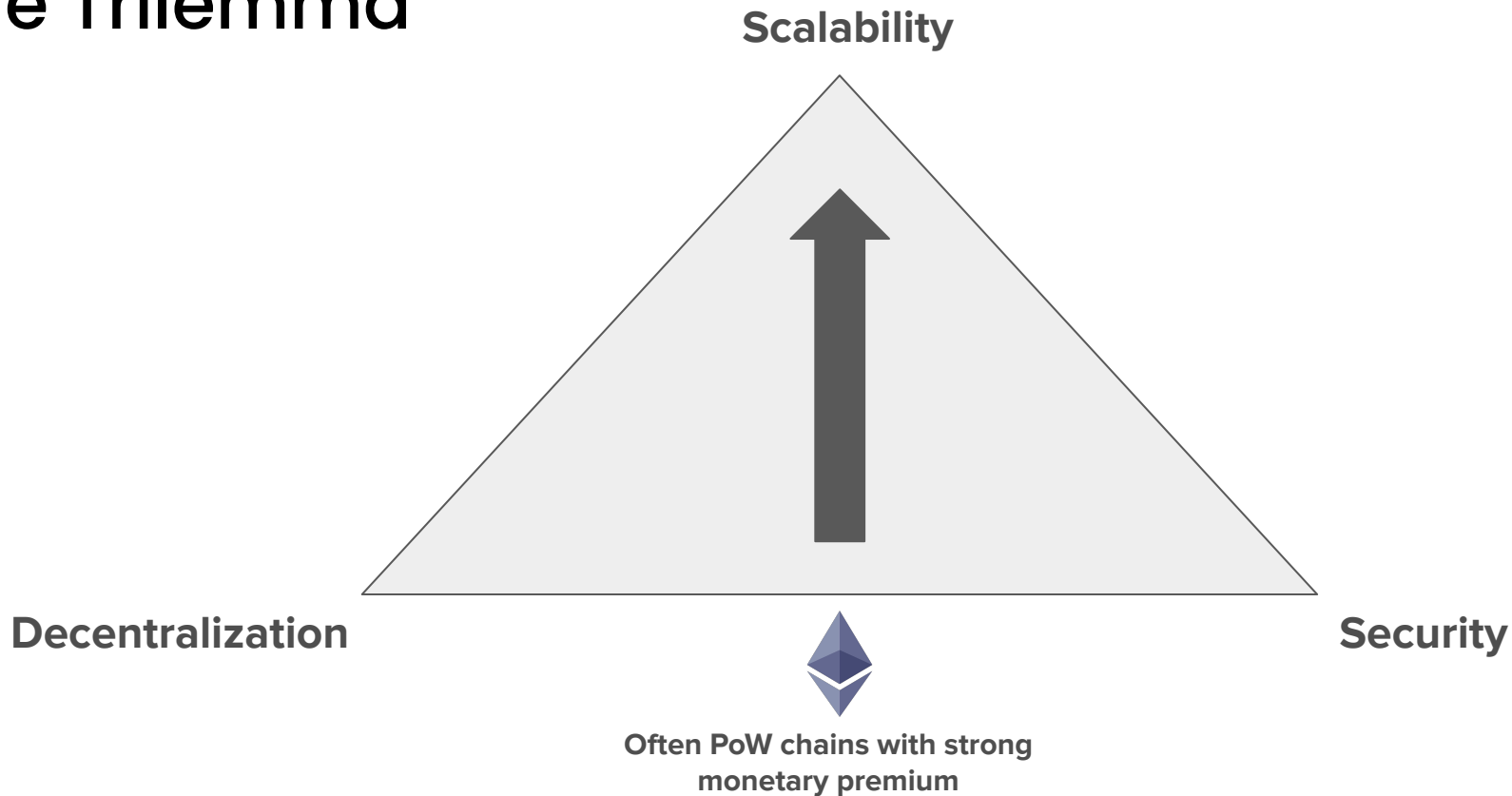


# The Trilemma



TOPIC

# The Trilemma



# Why L2 Rollup + Sharding

- Layer 1 scaling: increase throughput -> harder to verify -> higher threshold for verifier -> more centralized -> less sustainable
- Technical sustainability
  - Whatever the most centralized of L1s can do, zkR can do much better, with significantly higher TPS. Further, multiple zk-rollups can effectively attain global scale in aggregate.
- Economic sustainability
  - Zk-rollups have a miniscule fraction of the cost overhead of a centralized L1, allowing it to offer orders of magnitude greater throughput with similar fees; or similar throughput with a fraction of the fees.
  - e.g. Solana yearly \$36.5M revenue, \$4B inflationary reward
- Conclusion
  - The blockchain industry does not yet possess the technology to achieve global scale.
  - Zk-rollups are the only solution that can scale to millions of TPS, attaining global scale, while remaining technically and economically sustainable. They can do this while remaining highly secure, decentralized, permissionless, trustless, and credibly neutral.

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# Developer Ecosystems

- Layer 2
  - Cairo by StarkWare
  - Zinc by Matter Labs (zkSync)
- ZK-SNARK
  - Circom by iden3, written in Rust
- ZK-EVMs
  - C, Rust, etc.
- ZK hardware acceleration

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# Some Interesting Directions for ZKP

- ZK-NFT
  - Yi Protocol
- ZK games
  - Interoperability
  - Test ground for under-collateralized lending?
  - Challenges
- ZK in pure social settings, such as a DAO
  - Whistleblower, manifested by Semaphore (used by Worldcoin)
    - Allows any Ethereum user to signal their endorsement of an arbitrary string, revealing only that they have been previously approved to do so, and not their specific identity
- etc.