

lecture 3

Combinational logic 1

- truth tables
- Boolean algebra
- sum of products and product-of-sums
- logic gates

Let A, B be binary variables
("boolean")

$1 \equiv \text{true}, \quad 0 \equiv \text{false}$

Notation:

| | | |
|----------------|----------|--------------------|
| $A \bullet B$ | \equiv | $A \text{ and } B$ |
| $A + B$ | \equiv | $A \text{ or } B$ |
| \overline{A} | \equiv | not A |

One uses $+, \bullet$ instead of \vee, \wedge .
which you may have seen elsewhere.

Truth Tables

Notation: $A \cdot B \equiv A \text{ and } B$
 $A + B \equiv A \text{ or } B$
 $\overline{A} \equiv \text{not } A$

| A | B | $A \cdot B$ | $A + B$ |
|---|---|-------------|---------|
| 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 |

| A | \overline{A} |
|---|----------------|
| 0 | 1 |
| 1 | 0 |

(exclusive or)

NAND

NOR

XOR

| A | B | $\overline{A \cdot B}$ | $\overline{A + B}$ | $A \oplus B$ |
|---|---|------------------------|--------------------|--------------|
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 |

There are $2^4 = 16$ possible boolean functions.

$$f: \{0,1\} \times \{0,1\} \rightarrow \{0,1\}$$

| A | B | γ_1 | γ_2 | γ_3 | ... | γ_{16} |
|---|---|------------|------------|------------|-----|---------------|
| 0 | 0 | | | | | |
| 0 | 1 | | | | | |
| 1 | 0 | | | | | |
| 1 | 1 | | | | | |

We typically only work with AND, OR, NAND, NOR, XOR.

Laws of Boolean Algebra

identity

$$A + 0 = A$$

$$A \cdot 1 = A$$

inverse

$$A + \overline{A} = 1$$

$$A \cdot \overline{A} = 0$$

one and zero

$$A + 1 = 1$$

$$A \cdot 0 = 0$$

commutative

$$A + B = B + A$$

$$A \cdot B = B \cdot A$$

associative

$$(A + B) + C = A + (B + C)$$

$$(A \cdot B) \cdot C = A \cdot (B \cdot C)$$

distributive

$$\begin{aligned} A \cdot (B + C) \\ = (A \cdot B) + (A \cdot C) \end{aligned}$$

$$\begin{aligned} A + (B \cdot C) \\ = (A + B) \cdot (A + C) \end{aligned}$$

de Morgan

$$\overline{(A + B)} = \overline{A} \cdot \overline{B}$$

$$\overline{A \cdot B} = \overline{A} + \overline{B}$$

Laws of Boolean Algebra

distributive

$$A + (B \cdot C) = (A + B) \cdot (A + C)$$

Note this one behaves differently from integers or reals.

Example

$$Y = \overline{A \cdot B \cdot C} \cdot (A \cdot B + A \cdot C)$$

| A | B | C | Y |
|---|---|---|---|
| 0 | 0 | 0 | |
| 0 | 0 | 1 | |
| 0 | 1 | 0 | |
| 0 | 1 | 1 | |
| 1 | 0 | 0 | |
| 1 | 0 | 1 | |
| 1 | 1 | 0 | |
| 1 | 1 | 1 | |

$$Y = \overline{A \cdot B \cdot C} \cdot (A \cdot B + A \cdot C)$$

| A | B | C | $A \cdot B \cdot C$ | $\overline{A \cdot B \cdot C}$ | $A \cdot B$ | $A \cdot C$ | $A \cdot B + A \cdot C$ | Y |
|---|---|---|---------------------|--------------------------------|-------------|-------------|-------------------------|---|
| 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 1 |
| 1 | 1 | 1 | 1 | 0 | 1 | 1 | 1 | 0 |



Sum of Products

$$Y = \overline{A \cdot B \cdot C} \cdot (A \cdot B + A \cdot C)$$
$$= A \cdot \overline{B} \cdot C + A \cdot B \cdot \overline{C}$$

Q: For 3 variables A, B, C, how many terms can we have in a sum of products representation ?

A: $2^3 = 8$ i.e. previous slide

$$Y = A \cdot \overline{B} \cdot C + A \cdot B \cdot \overline{C}$$

$$\overline{Y} = A \cdot \overline{B} \cdot C + A \cdot B \cdot \overline{C}$$

$$= \overline{(A \cdot \overline{B} \cdot C)} \cdot \overline{(A \cdot B \cdot \overline{C})}$$

$$= (\overline{A} + B + \overline{C}) \cdot (\overline{A} + \overline{B} + C)$$

called a "product of sums"

How to write Y as a "product of sums" ?

First, write its complement \overline{Y} as a sum of products.

Because of time constraints, I decided to skip this example in the lecture.

You should go over it on your own.

| A | B | C | Y | \overline{Y} |
|---|---|---|---|----------------|
| 0 | 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 1 | 1 | 0 |
| 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 |

$$\overline{Y} = \overline{A} \cdot \overline{B} \cdot \overline{C} + \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + \overline{A} \cdot B \cdot C + A \cdot \overline{B} \cdot \overline{C} + A \cdot B \cdot C$$

Then write $Y = \overline{\overline{Y}}$ and apply de Morgan's Law.

$$\overline{Y} = \overline{A} \cdot \overline{B} \cdot \overline{C} + \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + \overline{A} \cdot B \cdot C + A \cdot \overline{B} \cdot \overline{C} + A \cdot B \cdot C$$

$$\overline{\overline{Y}} = (\overline{A} \cdot \overline{B} \cdot \overline{C} + \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + \overline{A} \cdot B \cdot C + A \cdot \overline{B} \cdot \overline{C} + A \cdot B \cdot C)$$

$$= (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C}) \cdot (\overline{A \cdot B \cdot C})$$

$$= (A + B + C) \cdot (A + B + \overline{C}) \cdot (A + \overline{B} + C) \cdot (A + \overline{B} + \overline{C}) \cdot (\overline{A} + B + C) \cdot (\overline{A} + \overline{B} + \overline{C})$$

Sometimes we have expressions where various combinations of input variables give the same output. In the example below, if A is false then any combination of B and C will give the same output (namely true).

$$Y = A \cdot \bar{B} \cdot C + \bar{A}$$

| A | B | C | $A \cdot \bar{B} \cdot C$ | \bar{A} | Y |
|---|---|---|---------------------------|-----------|---|
| 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 0 | 1 | 1 |
| 0 | 1 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 |

Don't Care

We can simplify the truth table in such situations.

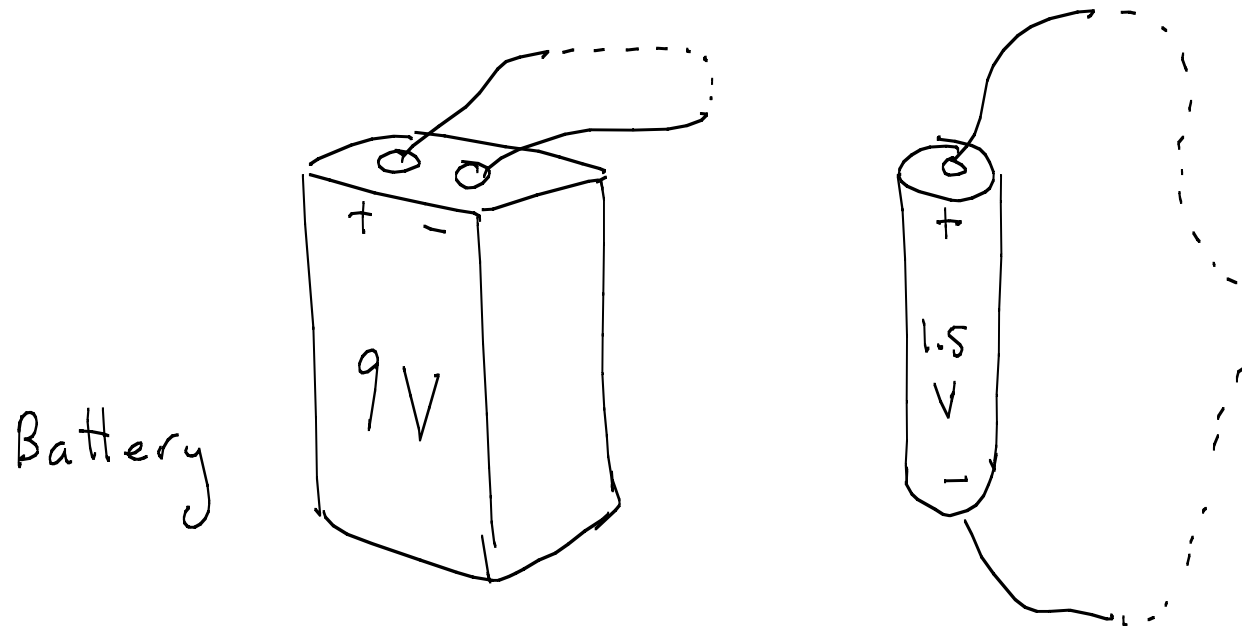
$$Y = A \cdot \bar{B} \cdot C + \bar{A}$$

| A | B | C | Y |
|---|---|---|---|
| 0 | X | X | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 0 |

X means we "don't care" what values are there.

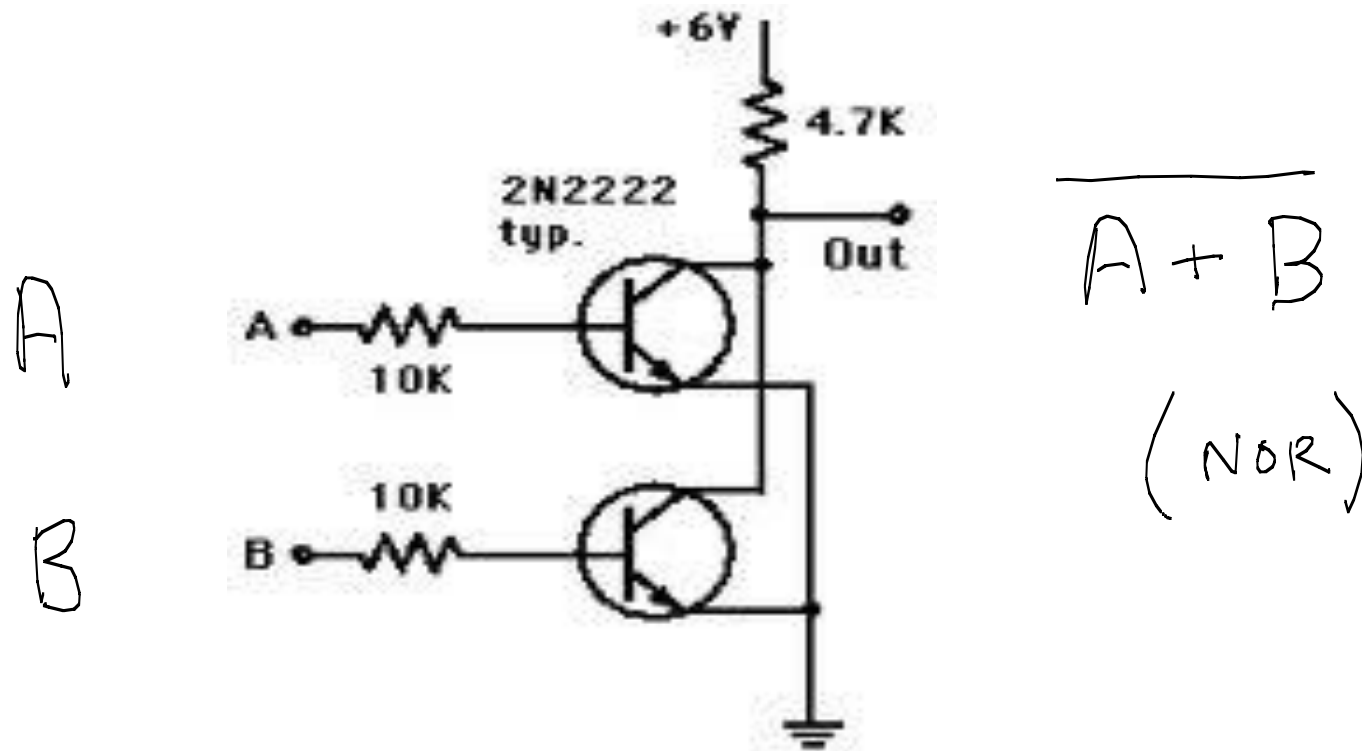
What are the 0's and 1's in a computer?

A wire can have a voltage difference between two terminals, which drives current.



In a computer, wires can have two voltages:
high (1, current ON) or low (0, current ~OFF)

Using circuit elements called "transistors" and "resistors", one can build circuits called "gates" that compute logical operations.



For each of the OR, AND, NAND, XOR gates, you would have a different circuit.

Moore's Law (Gordon Moore was founder of Intel)

The number of transistors per mm^2 approximately doubles every two years. (1965)

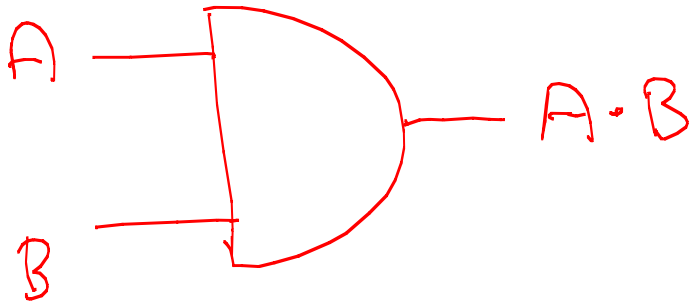
It is an observation, not a physical law.

It still holds true today, although people think that this cannot continue, because of limits on the size of atom and laws of quantum physics.

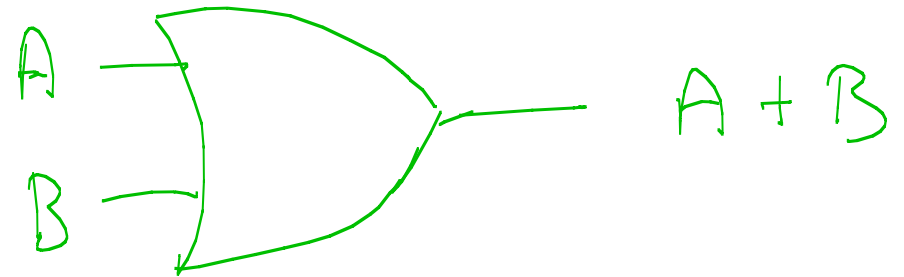
<http://phys.org/news/2015-07-law-years.html>

Logic Gates

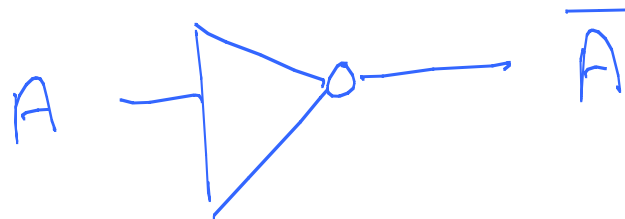
AND



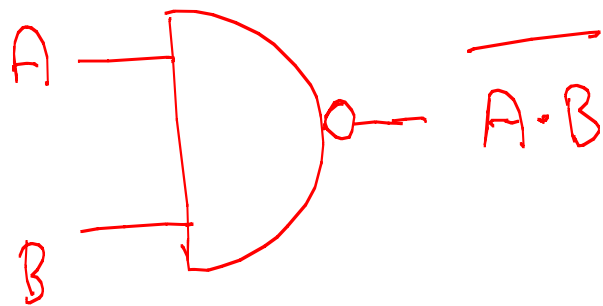
OR



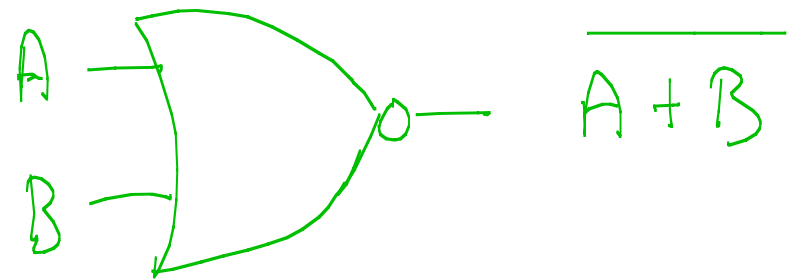
NOT



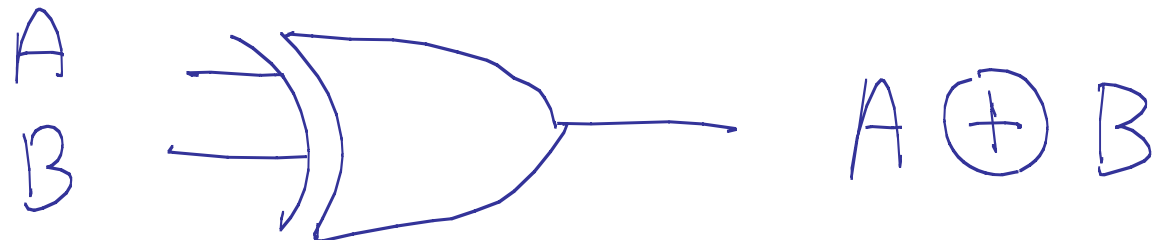
NAND



NOR



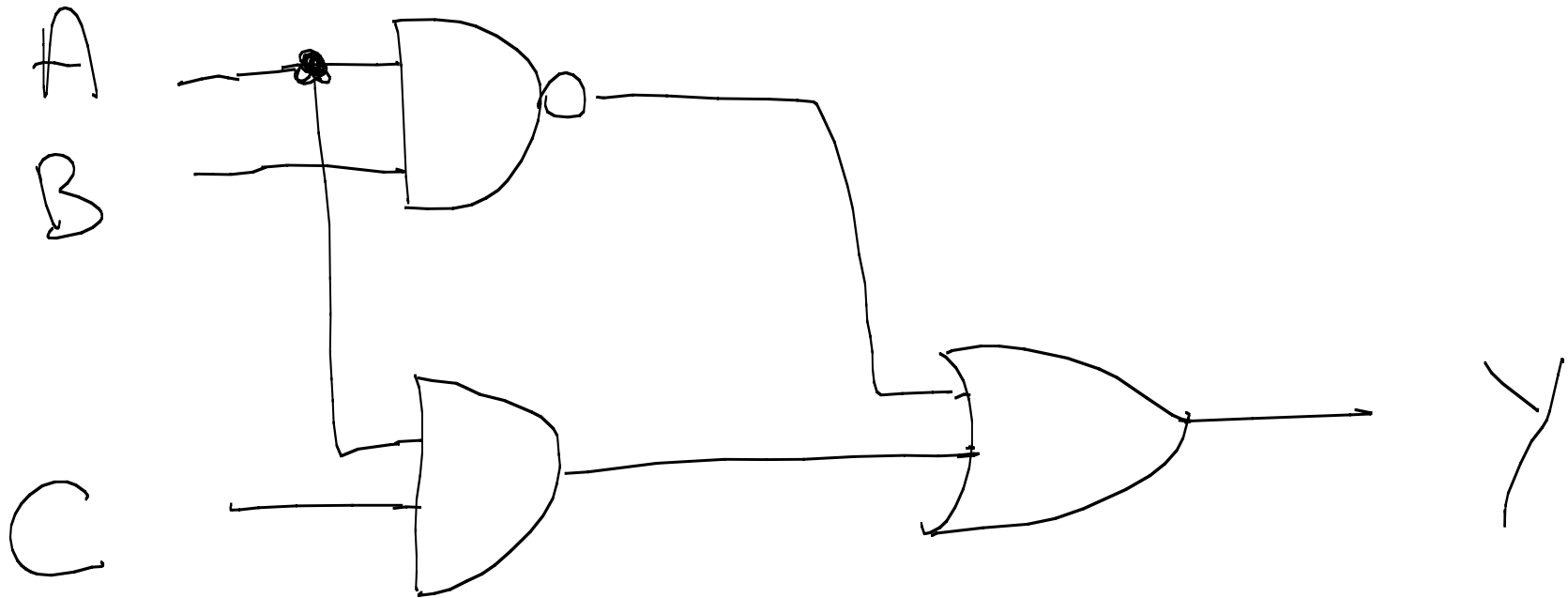
XOR



Logic Circuit

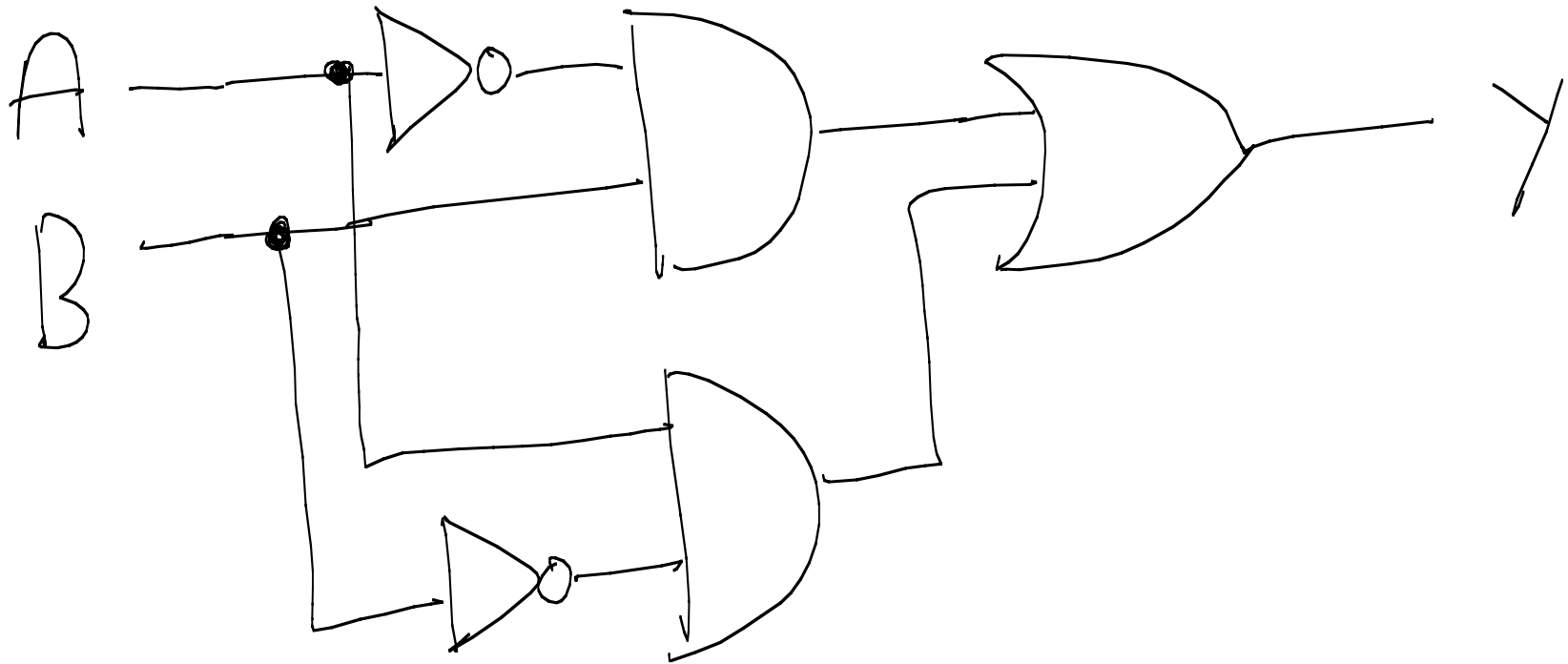
Example:

$$Y = \overline{A \cdot B} + A \cdot C$$



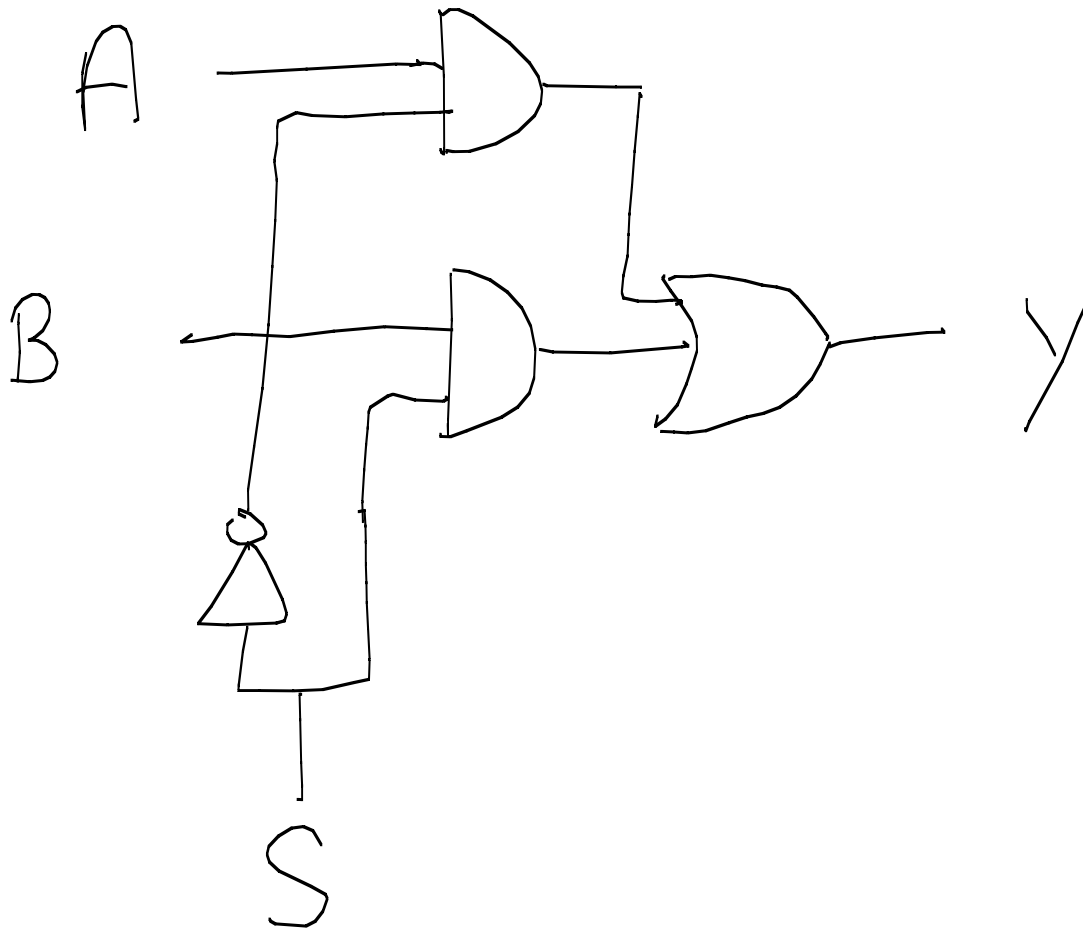
Example: XOR without using an XOR gate

$$Y = \overline{A} \cdot B + A \cdot \overline{B} = A \oplus B$$



Multiplexor (selector)

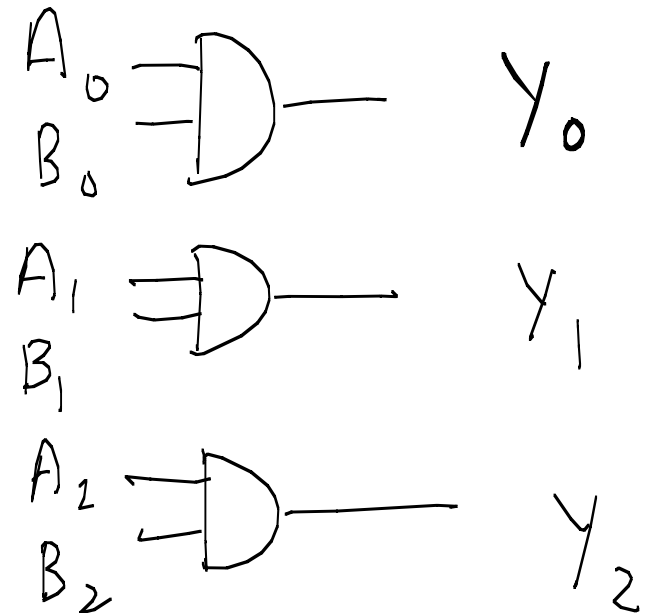
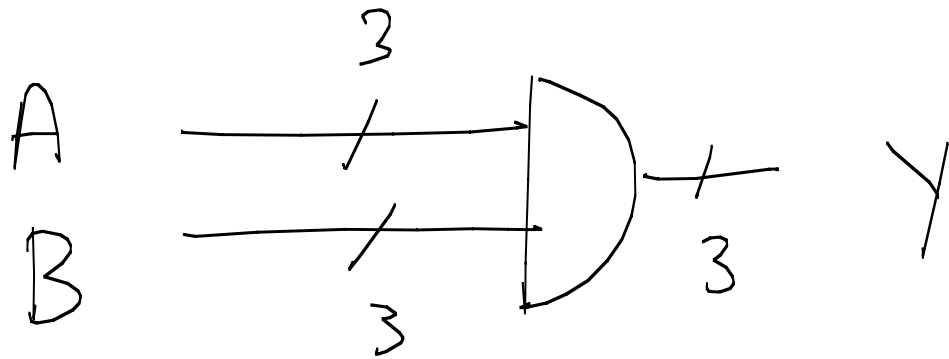
$$Y = \overline{S} \cdot A + S \cdot B$$



if *S*
 $Y = B$
else
 $Y = A$

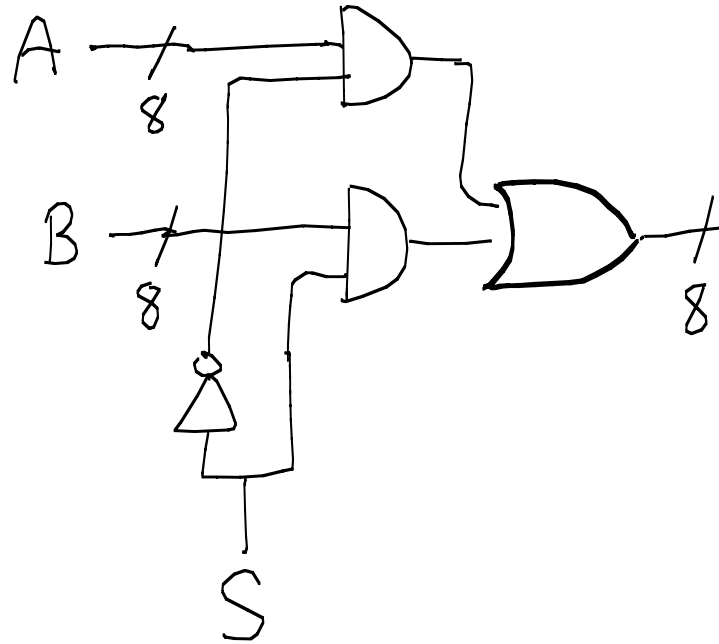
Notation

Suppose A and B are each 3 bits ($A_2 A_1 A_0$, $B_2 B_1 B_0$)



Suppose A and B are each 8 bits ($A_7 A_6 \dots A_0$, $B_7 B_6 \dots B_0$)
We can define an 8 bit multiplexor (selector).

Notation:



In fact we would build this from 8 separate one-bit multiplexors.

Note that the selector S is a single bit. We are selecting either all the A bits or all the B bits.