

# CHALMERS

## EXAMINATION / TENTAMEN

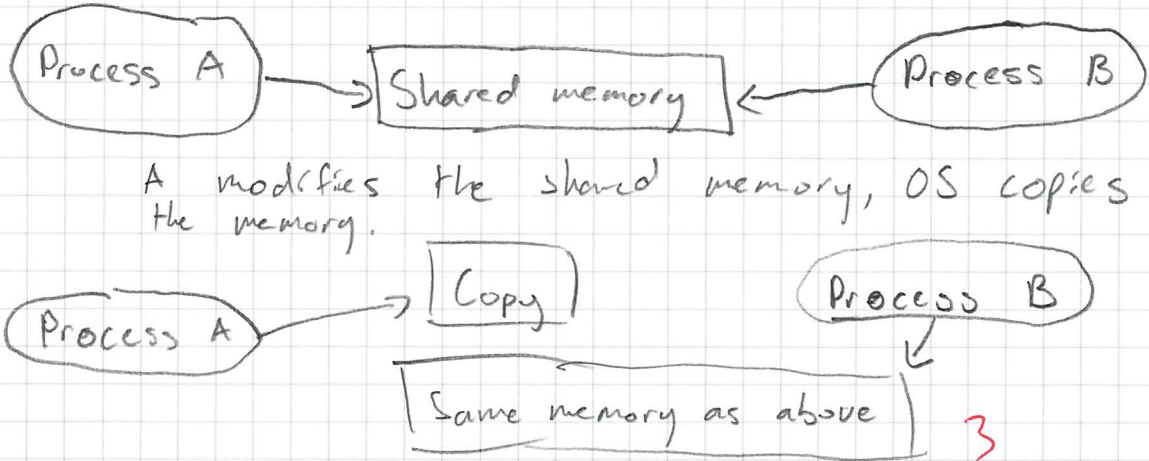
|                              |  |                                    |                               |                |
|------------------------------|--|------------------------------------|-------------------------------|----------------|
| Course code/kurskod          |  | Course name/kursnamn               |                               |                |
| EDA093                       |  | Operating Systems                  |                               |                |
| Anonymous code<br>Anonym kod |  | Examination date<br>Tentamensdatum | Number of pages<br>Antal blad | Grade<br>Betyg |
| EDA093-0015-GST              |  | 15-08-2023                         | 7                             | 3              |

\* I confirm that I've no mobile or other similar electronic equipment available during the examination.  
Jag intygar att jag inte har mobiltelefon eller annan liknande elektronisk utrustning tillgänglig under examinationen.

| Solved task<br>Behandlade uppgifter<br>No/nr              | Points per task<br>Poäng på<br>uppgiften | Observe: Areas with bold contour are to completed by the teacher.<br>Anmärkning: Rutor inom bred kontur ifylles av lärare. |
|---|--|--|
| 1   | X 7                                      |  |
| 2   | X 6                                      |  |
| 3   | X 8                                      |  |
| 4   | X 10                                     |  |
| 5   | X 4                                      |  |
| 6   |  |  |
| 7   |  |  |
| 8   |  |  |
| 9   |  |  |
| 10  |  |  |
| 11  |  |  |
| 12  |  |  |
| 13  |  |  |
| 14  |  |  |
| 15  |  |  |
| 16  |  |  |
| 17  |  |  |
| Bonus:<br>poäng   |  |  |
| Total examination<br>points<br>Summa poäng<br>på tentamen | 35                                       |  |

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- (COW)  
a) Copy-on-write enables different processes to share memory. Only when the shared memory will be modified will the memory be copied.



Benefits of COW are that memory will only be copied when necessary, meaning less computation which will be needed by the OS and more free memory.

b)

| Logical | Page table | Frame |
|---------|------------|-------|
| 0       | 0          | B 0   |
| A 1     | 1 5        | 1     |
| C 2     | 2 3        | 2     |
| 3       | 3          | C 3   |
| 4       | 4          | 4     |
| E 5     | 5 8        | A 5   |
| 6       | 6          | 6     |
| D 7     | 7 9        | 7     |
| 8       | 8          | E 8   |
| B 9     | 9 0        | D 9   |

bits?

- c)
- Handle page fault interrupt
    - Illegal memory access → Crash
    - Page not in frame → Continue with step 2.
  - Remove unwanted page (if there is any) to disk. Update the page table.
  - Find and put wanted page into the frame
  - Update page table.
  - Resume the process which caused the page fault.

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d) Yes, because then pages are only loaded into memory if and when they are needed.

and that decreases PF?



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- a) Amdahl's law describes the speedup of a program of which  $S \cdot 100\%$  runs sequentially and  $N$  cores/threads are used:

$$\text{speedup} \leq \frac{1}{S + \frac{(1-S)}{N}}$$

- b) 1. Scheduler interrupt: process A will be paused  
2. Store A's register values, PC, page table etc into memory.  
3. Retrieve B's register values, PC, page table etc from memory and load it into its corresponding unit.  
4. Resume process B.

The second context switch is the same as above, with the small change of A and B. This time B will be paused, the same operations done as before, only this time first on B, then A.

- c) User-level threads      Kernel-level threads  
+ Fast context switches



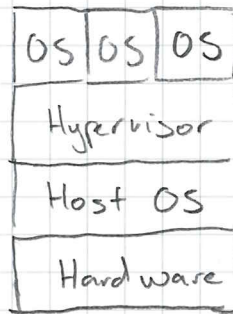
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a)

Type 1



Type 2



In type 1 the hypervisor runs directly on H/W. Type 2 runs on a host OS.

3

c) A hypervisor has a shadow page table for each VM. Each VM has a page table, and the shadow page table bridges the gap between a VM's page table and the actual frames in physical memory.

40

The two methods are guest-induced page faults and hypervisor-induced page faults.

2

d) Incremental snapshots are those where only the modifications since the last complete snapshot are stored. This reduces the memory footprint significantly.

3

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```

int thread-sem;

DataQueue queue;

main () {
    queue = data-queue-create();
    thread-sem = create-semaphore(2);

    pthreads A, B, C;

    A = p-threads-create(*a-method);
    B = p-threads-create(*b-method);
    C = p-threads-create(*c-method);
    shed-yield();
}

a-method() {
    while(true) {
        thread-sem.wait();

        files = check-files();

        if (files == null) {
            thread-sem.signal(); shed-yield();
        }
        else {
            for (file : files) {
                queue.add(file);
            }
            thread-sem.signal();
            shed-yield();
        }
    }
}

b-method() {
    while(true) {
        thread-sem.wait();
        element = queue.pop();
        if (element != null) {
            do-computation(element);
        }
        thread-sem.signal();
        shed-yield();
    }
}

```

*waiting for them?*

*10*

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```

c - method () {
    thread_sem.wait();
    int max_size = get_size_from_user();
    thread_sem.signal(); shed_yield();
    while (true) {
        thread_sem.wait();
        if (queue.size > max_size) {
            queue.remove(queue.size - max_size);
        }
        thread_sem.signal();
        shed_yield();
    }
}

```

Assumption: Since Data Queue is thread-safe, it is assumed that no races are possible in regards of the queue.

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Peterson's algorithm i (Same for thread B but vice versa)  
 boolean flagA, flagB // init false  
 int turn, idA, idB

repeat { // Thread A

1. flagA  $\leftarrow$  true

2. turn  $\leftarrow$  idA B

3. while (flagA and turn == idA) do nothing

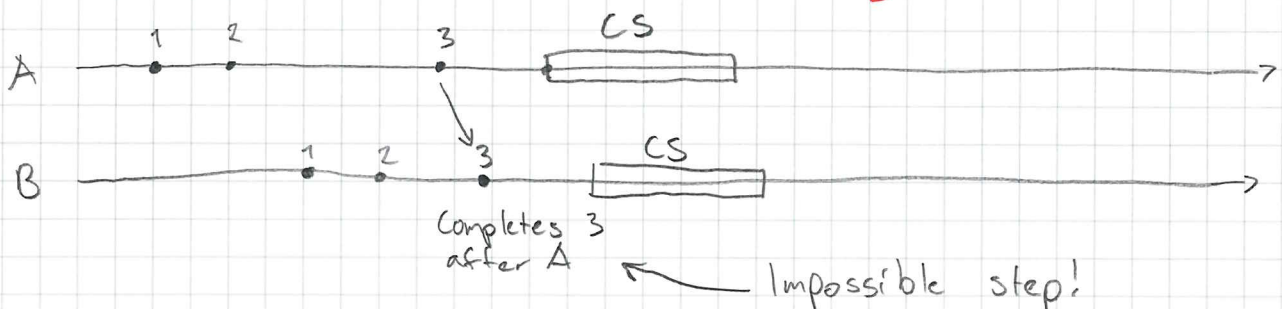
critical section

4. flagB  $\leftarrow$  false

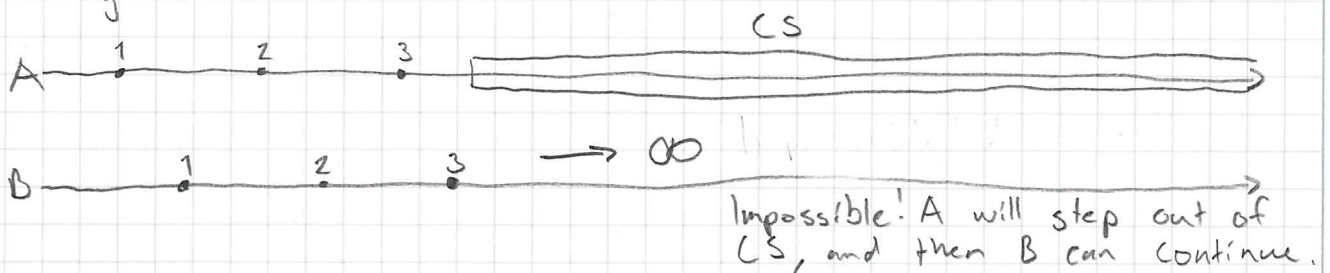
remainder section

}

Mutual exclusion:



Progress:



Fairness:

