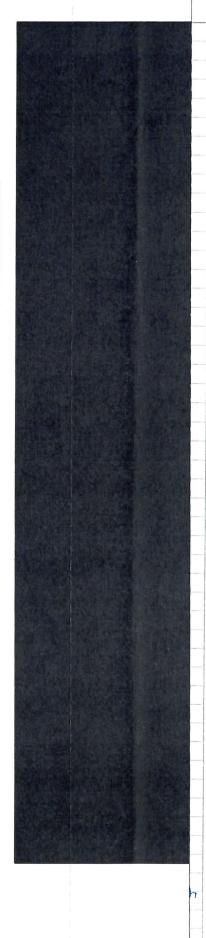
CHALMERS EXAMINATION/TENTAMEN

Course code/kurskod	Co	ourse name/kursnamn		
PIT 401	operating s	ystoms		
Anonymous code Anonym kod		Examination date Tentamensdatum	Number of pages Antal blad	Grade Betyg
PIT401-0007-LLZ		21 OCTOBEN 2023	07	G

^{*} I confirm that I've no mobile or other similar electronic equipment available during the examination. Jag intygar att jag inte har mobiltelefon eller annan liknande elektronisk utrustning tillgänglig under eximinationen.

eximination	onen.		in eller annan liknande elektronisk utrustning tillganglig under
Solved task Behandlade No/nr	e uppgifter	Points per task Poäng på uppgiften	Observe: Areas with bold contour are to completed by the teacher. Anmärkning: Rutor inom bred kontur ifylles av lärare.
1	✓	27	
2	√	12	
3	V	10	
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Bonus poäng			
Total exami points Summa poä på tentamer	ng	49	



or concurrently.

If he know how many cover the CPU has we are able to determine mether Processes are now in paramet. If the CPU only has one cover than the processes are now one after another. Otherwise they are now in paramet. To know if processes are no concurrently, we need so know if the there is no preemption of processes, then they are not now commencing. If each process is given a anantym, and preempted if they are do not complete in time, then there is commenced execution.

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