

Topic 6

Latches and Flip Flops

Introduction

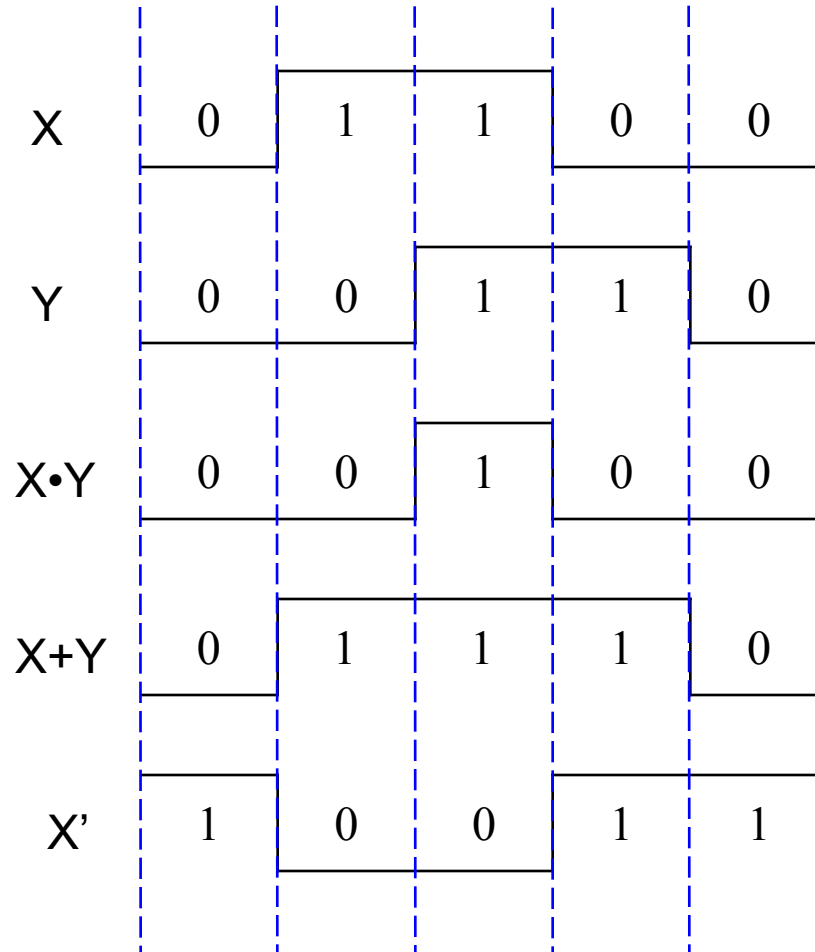
- **Combinational Circuit**
 - A digital circuit whose output depends only upon the **present combination** of its inputs
- **Sequential Circuit**
 - A digital circuit whose output depends not only upon the present input values, but also the history of input and output values
- **Beginning from this lecture, we will:**
 - Learn sequential circuits
 - Design a new type of building blocks, **latch & flip-flop**, that store value of a bit, a sequential circuit
 - Combine the blocks to build multi-bit storage – a **register**

Recall: Timing Diagrams for Gates

xy	F
0 0	0
0 1	0
1 0	0
1 1	1

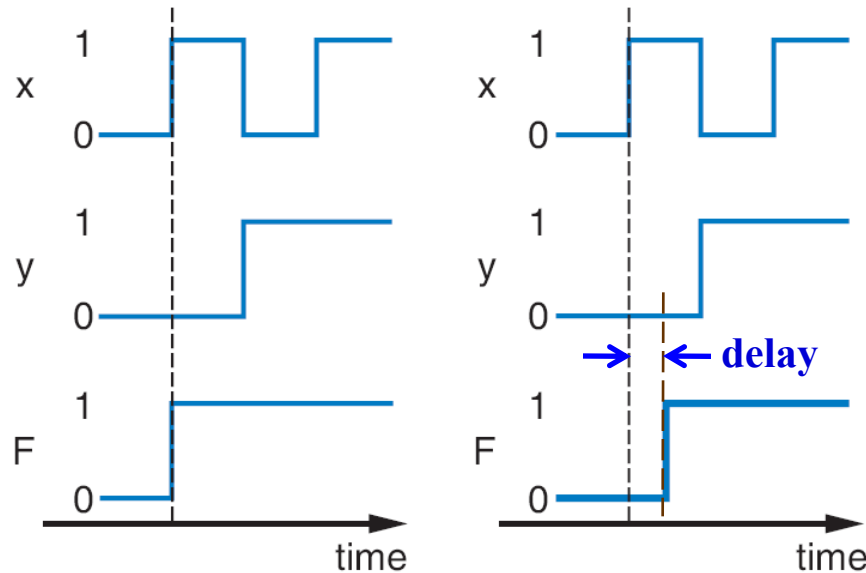
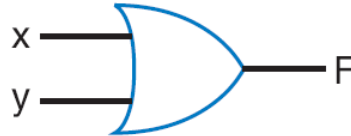
x+y	F
0 0	0
0 1	1
1 0	1
1 1	1

x	F
0	1
1	0



Reality of Combinational Circuit

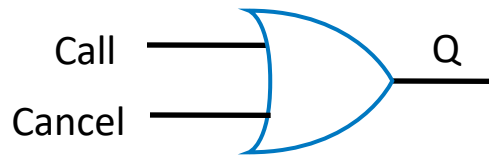
Non-Ideal Gate Behavior -- Delay



- Real gates have some delay
 - Outputs don't change immediately after inputs change

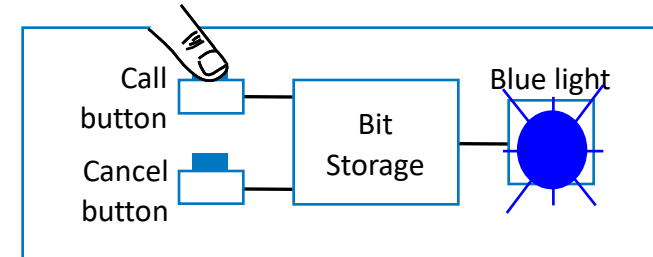
Example of Needing Bit Storage

- Flight attendant call button
 - Button pressed: provides a “1”
 - Press call: light turns on
 - **Stays on** after button released
 - Press cancel: light turns off
 - Logic gate circuit to implement this?

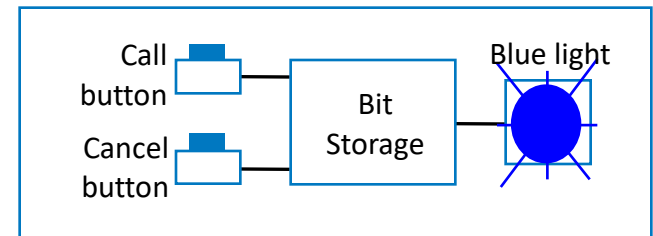


Doesn't work. $Q=1$ when $\text{Call}=1$, but doesn't stay 1 when Call returns to 0

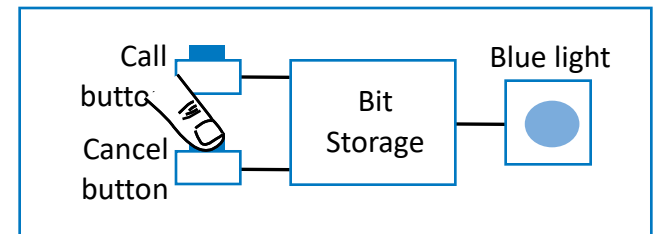
Need some form of “storage” in the circuit



1. Call button pressed – light turns on



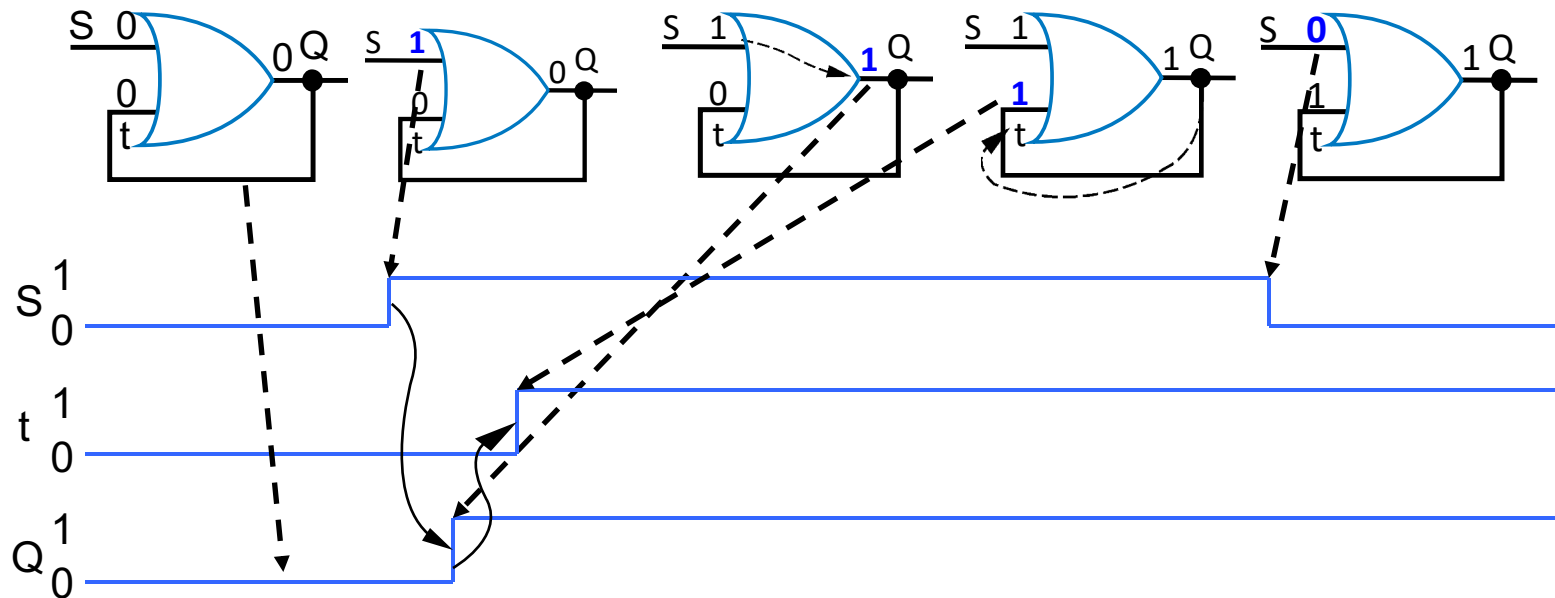
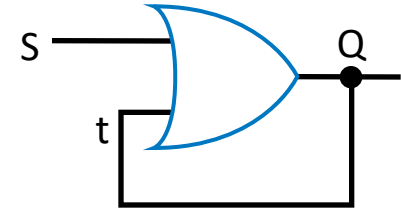
2. Call button released – light **stays on**



3. Cancel button pressed – light turns off

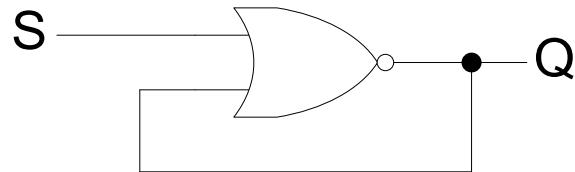
First Attempt at Implementation of Bit Storage

- We need some sort of feedback
 - Does circuit on the right do what we want?
 - Partially: Once Q becomes 1 (when S=1), Q stays 1 forever – no value of S can bring Q back to 0



Concepts of Sequential Circuit

- Sequential circuit
 - Combinational circuit with feedbacks

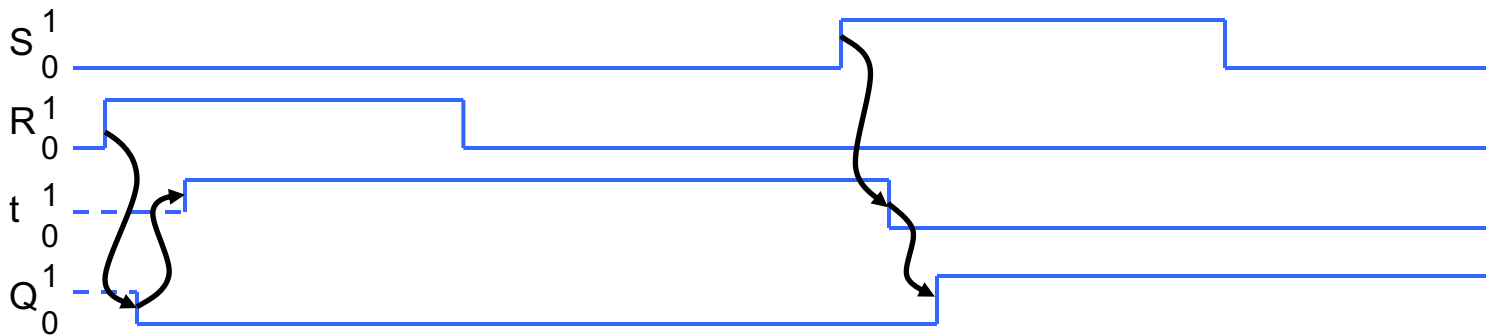
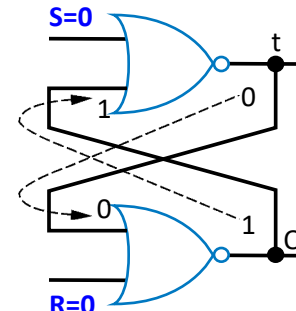
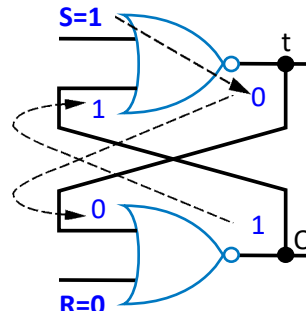
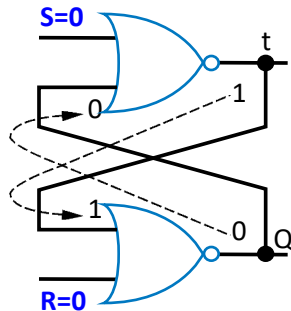
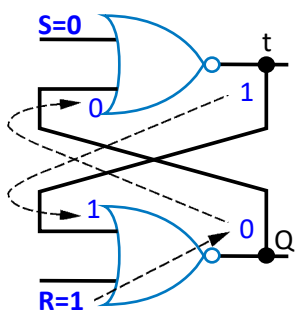
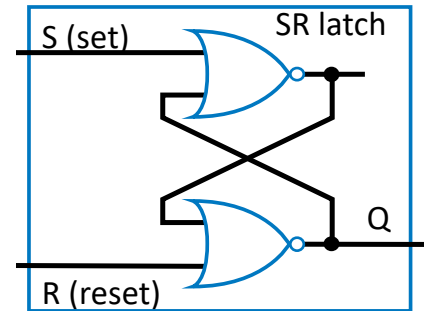
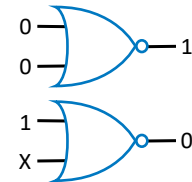


- Due to the feedback, output of a sequential circuit is decided by
 - Present inputs, and
 - Past input sequence and
 - Past outputs sequence
- Timing concepts
 - input-output propagation delay
 - clock
 - Other timing issues

Second Attempt at Bit Storage – SR Latch

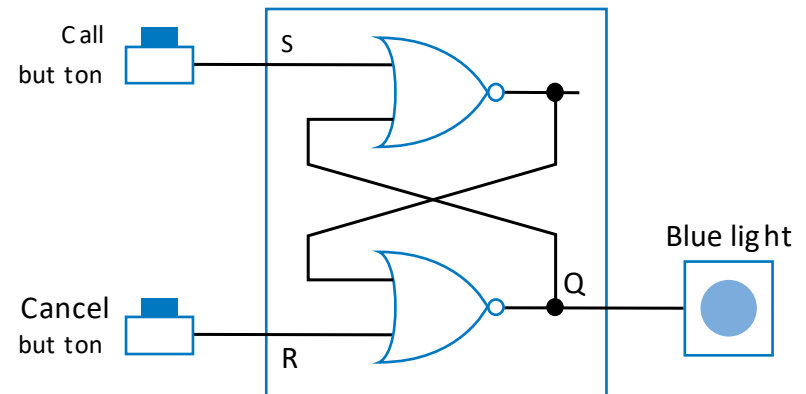
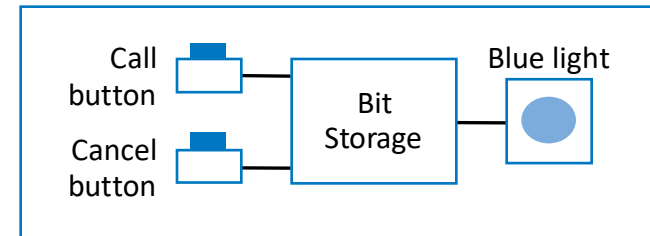
- Cross-coupled NOR gates
 - S: set (or preset) to 1
 - R: reset (or clear) to 0

Recall...



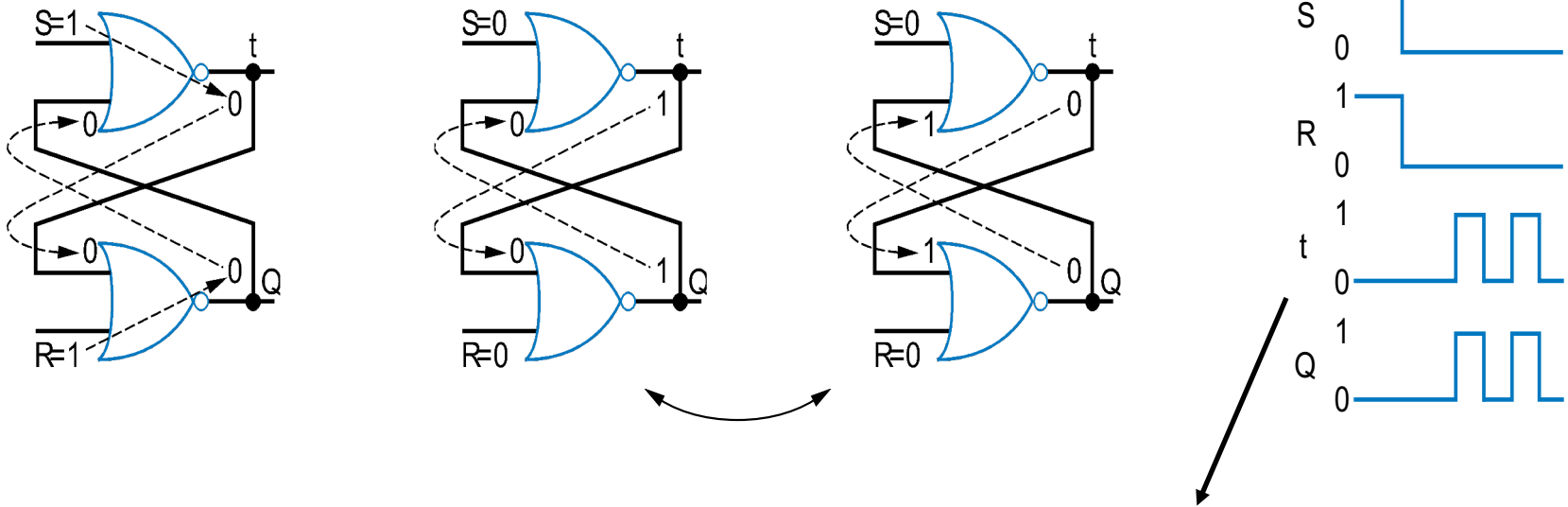
Example Using SR Latch for Bit Storage

- SR latch can serve as a bit storage, for example:
 - Call=1 : sets Q to 1
 - Q stays 1 even after Call=0
 - Cancel=1 : resets Q to 0
- But, there's a problem...

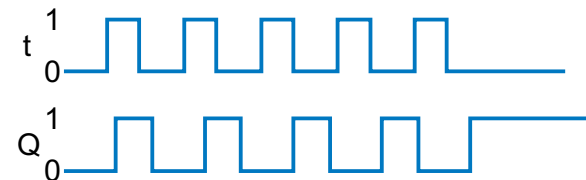


Problem with SR Latch

- Problem
 - If $S=1$ and $R=1$, we don't know what value Q will take when they both return to 0 *simultaneously*



Q may oscillate. Then, because one path will be slightly longer than the other, Q will eventually settle to 1 or 0 – but we don't know which.



Representation of SR Latch

- When discussing latches and flip-flops, we use
 - **present state** to represent current value of the Q output
 - **next state** to represent the new value of Q output responding to the current inputs and feedback of current output
- Characteristic table

S(t)	R(t)	Q(t)	Q(t+Δ) → Q ⁺	
0	0	0	0	hold
0	0	1	1	
0	1	0	0	reset
0	1	1	0	
1	0	0	1	set
1	0	1	1	
1	1	0	X	not allowed
1	1	1	X	

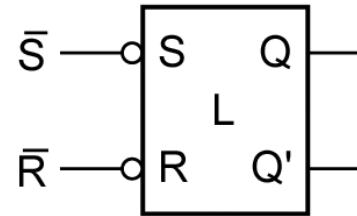
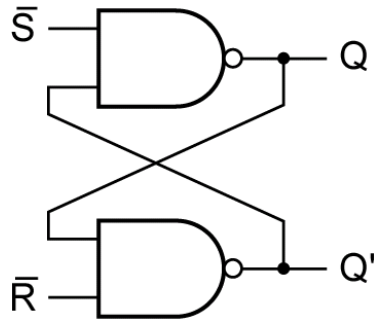
		Q ⁺	
		S(t)	R(t)
Q(t)	R(t)	0	1
		00	01
0	0	0	1
0	1	1	1
1	0	0	X
1	1	0	X

- Characteristic equation

$$Q^+ = S + R'Q$$

Alternative Implementation of SR Latch

- The cross-coupled RS latch can be implemented using NAND gates

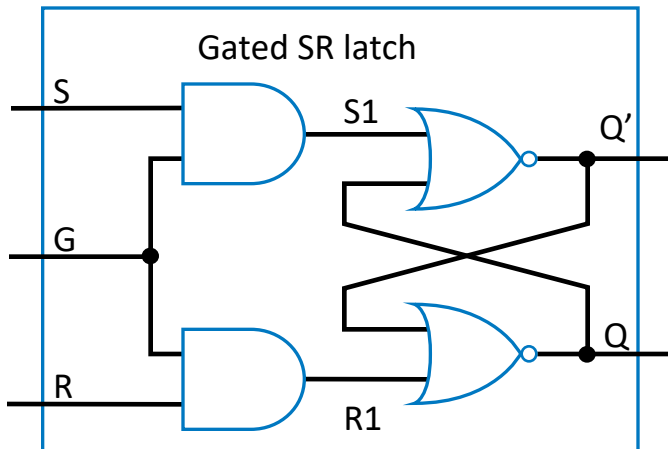


- Characteristic table

S	R	Q	Q ⁺	
0	0	0	X	not allowed
0	0	1	X	
0	1	0	1	set
0	1	1	1	
1	0	0	0	reset
1	0	1	0	
1	1	0	0	hold
1	1	1	1	

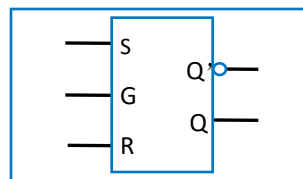
Gated SR Latch

- SR latch is enabled by a gate control signal G



Characteristic Table

G	S	R	Q ⁺
0	x	x	Q; Latch locked
1	0	0	Q; Hold state
1	0	1	0; Reset state
1	1	0	1; Set state
1	1	1	not allowed

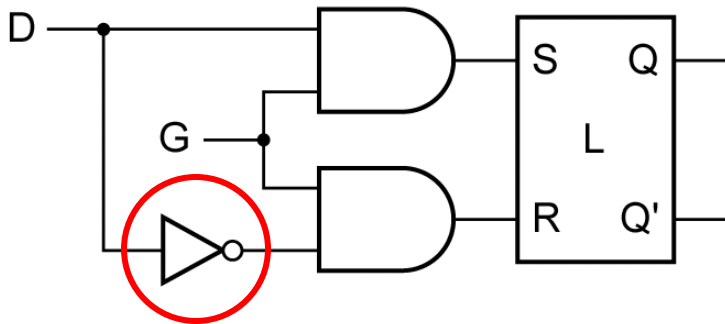


Gated SR latch
symbol

Solution to SR Latch Restriction

– Gated D Latch

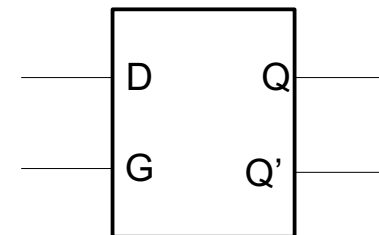
- Solution to the unstable state problem caused by $S = R = 1$ in SR latch



Characteristic Table

G	D	Q ⁺
1	0	0
1	1	1
0	X	Q

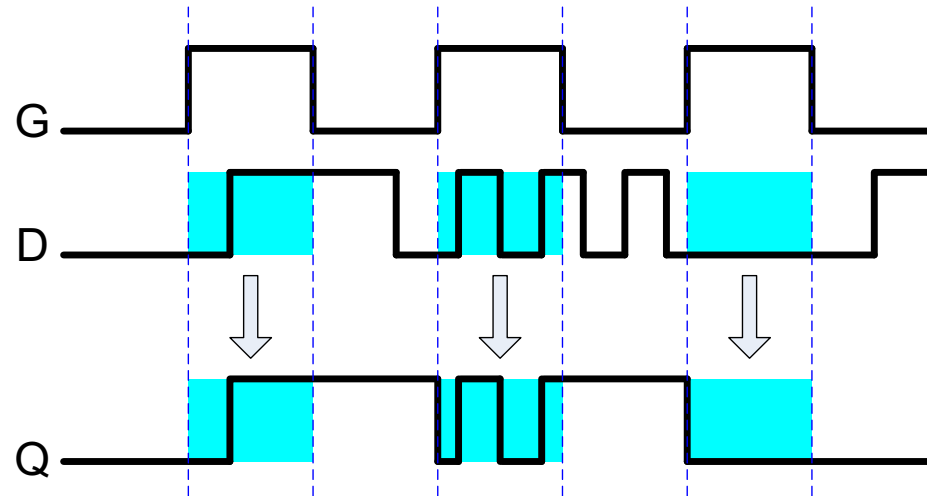
- The input value is stored into the latch only when gate control G has high level – **Level Sensitive**



D latch
symbol

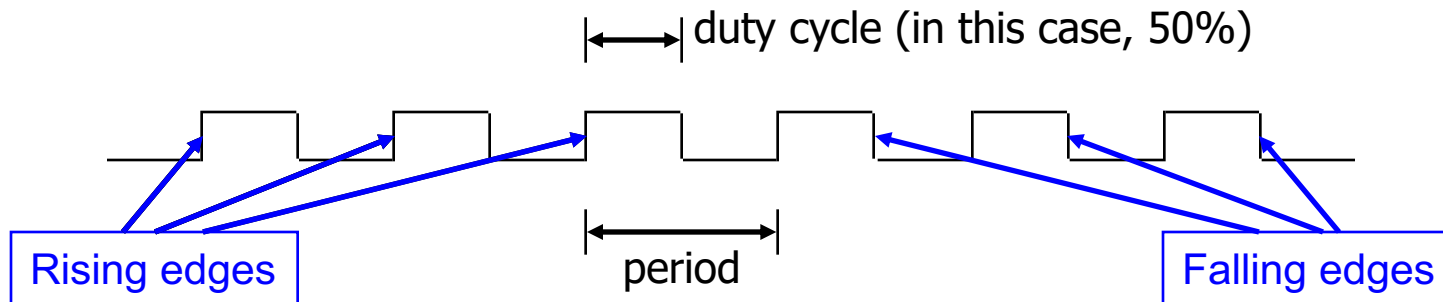
Gated D Latch – Transparent Latch

- Properties of the D latch
 - D latch is used as a **temporary storage** for a bit
 - The binary information at the data input of the D latch is **copied** to the Q output when the control input G is high (or enabled)
 - The output Q follows changes on the data input D as long as the control input G is enabled, so called a **transparent** latch



A Typical Control Input – Clock Signal

- Periodic pulse train used in sequential circuit to synchronize circuit behaviors

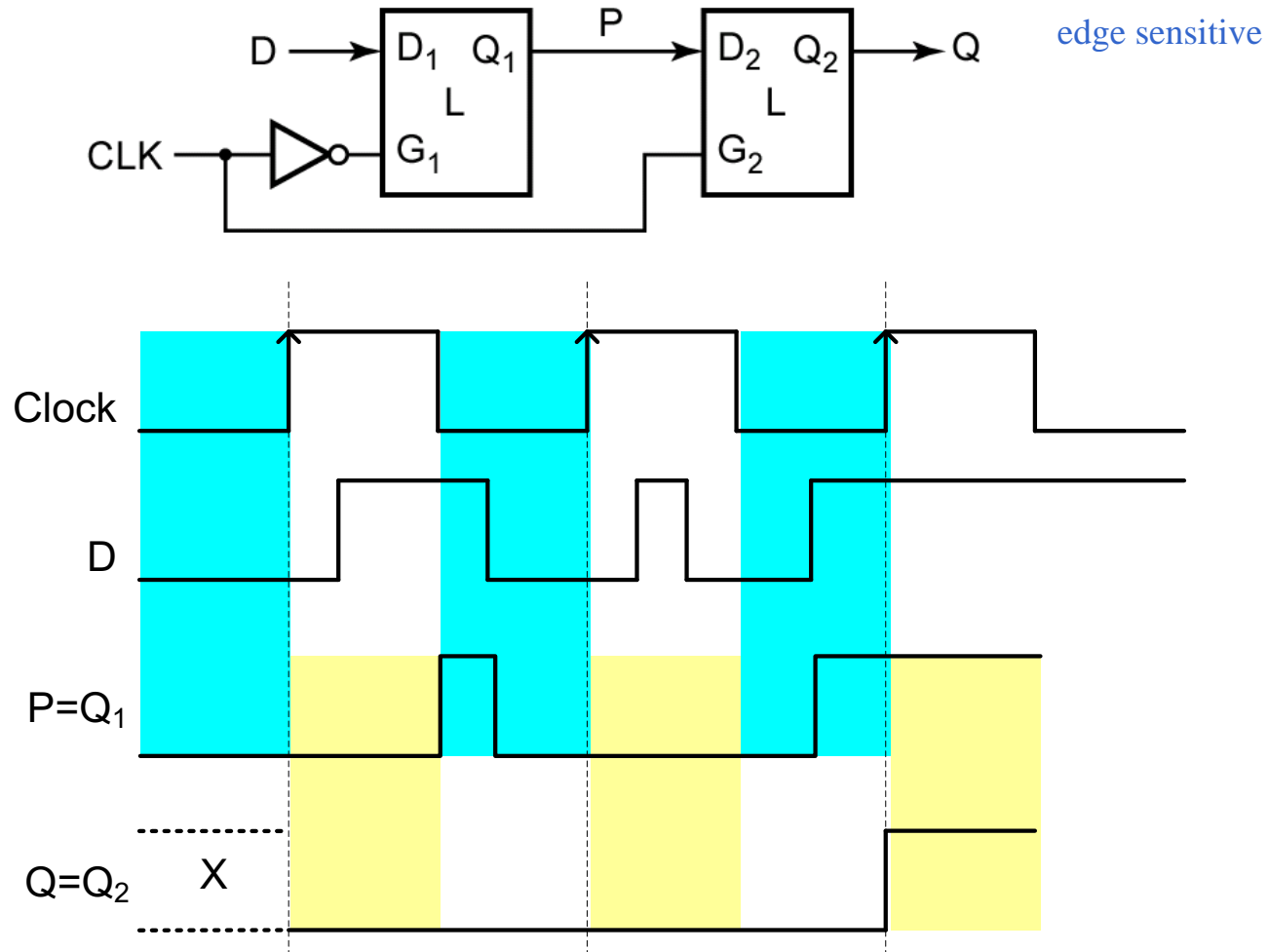


- ***Clock period***: time interval between pulses
- ***Clock cycle***: one such time interval
- ***Clock frequency***: $1/\text{period}$

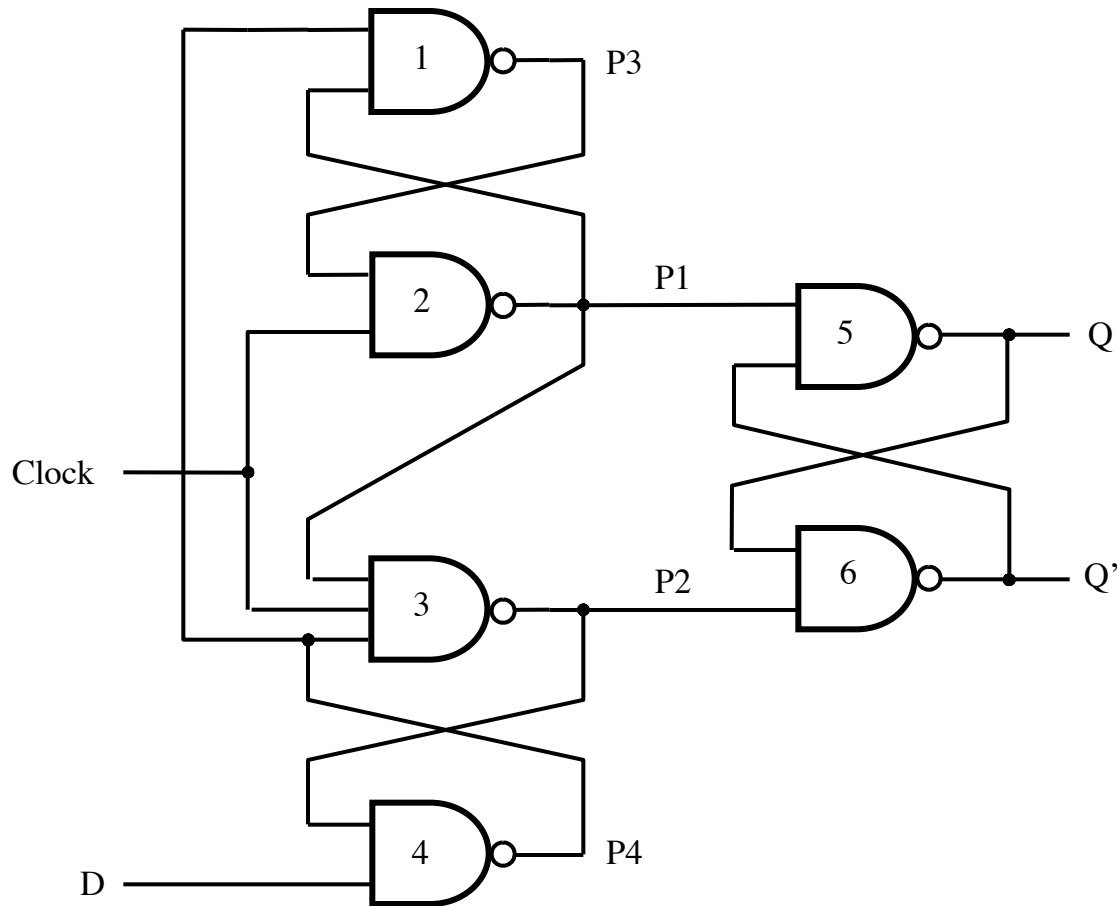
Freq	Period
100 GHz	0.01 ns
10 GHz	0.1 ns
1 GHz	1 ns
100 MHz	10 ns
10 MHz	100 ns

Rising-Edge Triggered D Flip Flop

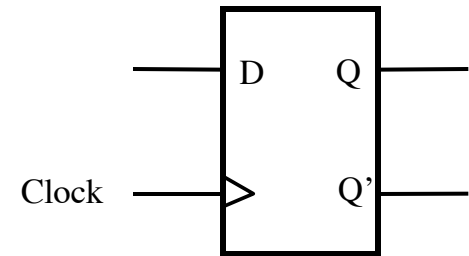
- Rising-edge triggered Master-Slave D flip flop



Alternative Implementation of Rising-Edge Triggered D Flip Flop



Alternative Implementation

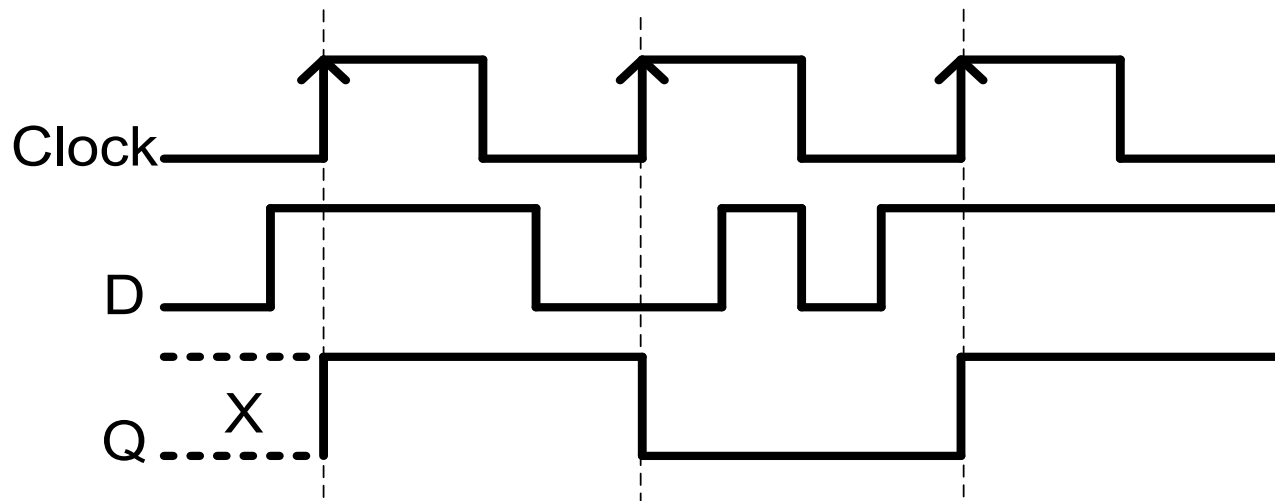


clock	D	Q ⁺
\uparrow	0	0
\uparrow	1	1
0	X	Q
1	X	Q

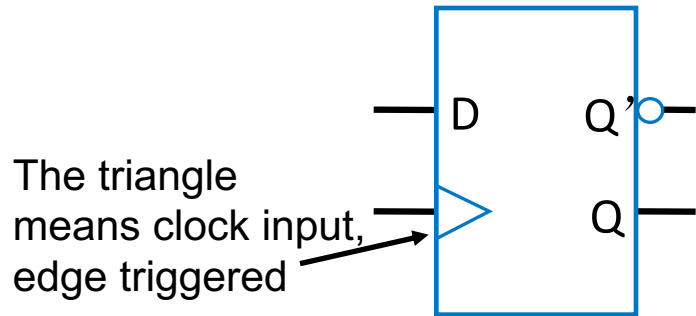
Characteristic equation:
 $Q^+ = D$ (at active clock edges)

Rising-Edge Triggered D Flip Flop

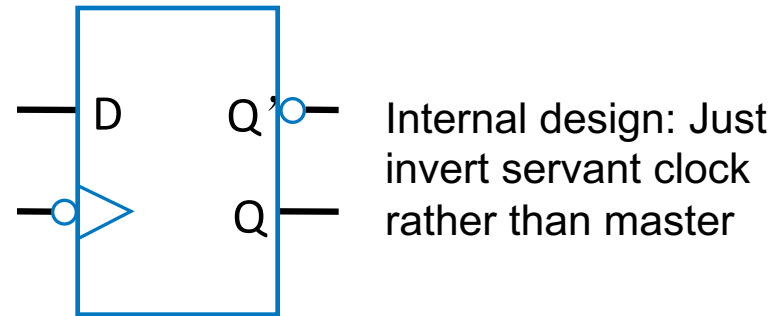
- Properties of the rising edge triggered D Flip Flop
 - The output changes only at the rising edges of the clock signal – **Edge Sensitive**
 - The output Q gets the value of input D at the time point of rising edge of clock



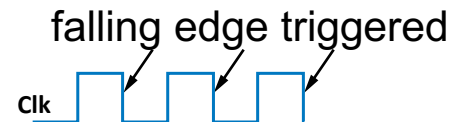
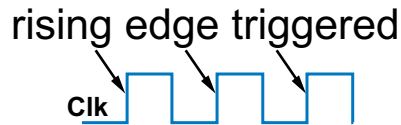
Symbols for D Flip-Flop



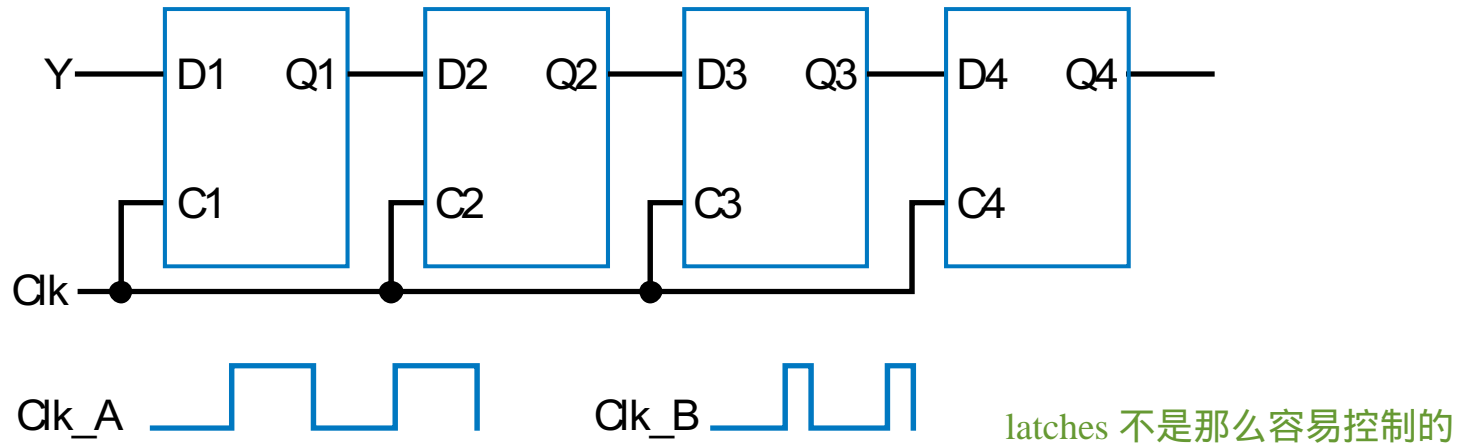
Symbol for rising-edge triggered D flip-flop



Symbol for falling-edge triggered D flip-flop



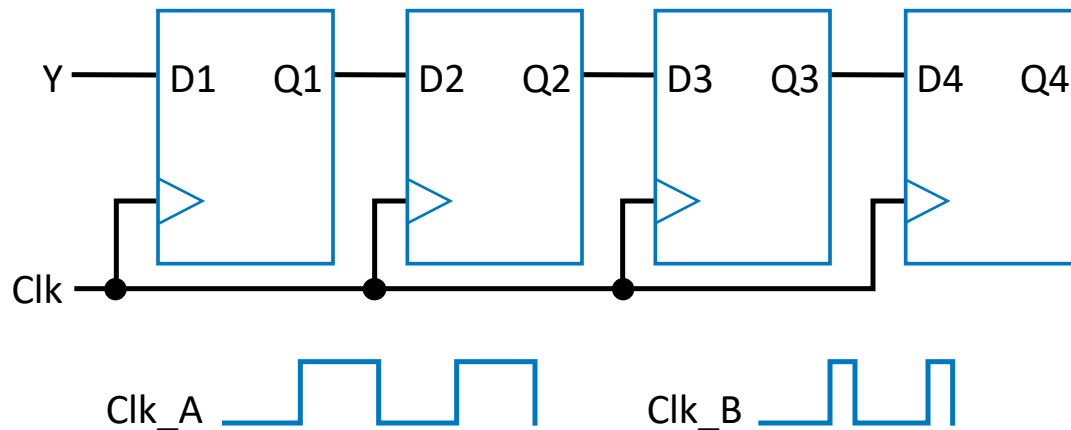
Problem with Level-Sensitive D Latch



- D latch still has problem
 - When Clk=1, through **how many latches will a signal travel?**
 - Depends on for how long Clk=1
 - Clk_A -- signal may travel through **multiple** latches
 - Clk_B -- signal may travel through **fewer** latches
 - Hard to control the storage elements

Application of D Flip-Flop

- Solves problem of concatenated D latches when $G=1$
 - In figure below, signal travels through exactly one flip-flop, for either Clk_A or Clk_B
 - On each rising edge of Clk, all four flip-flops are loaded simultaneously, doesn't matter how long Clk is 1.



而应用了FF了以后，
同时
可以控制这些latches

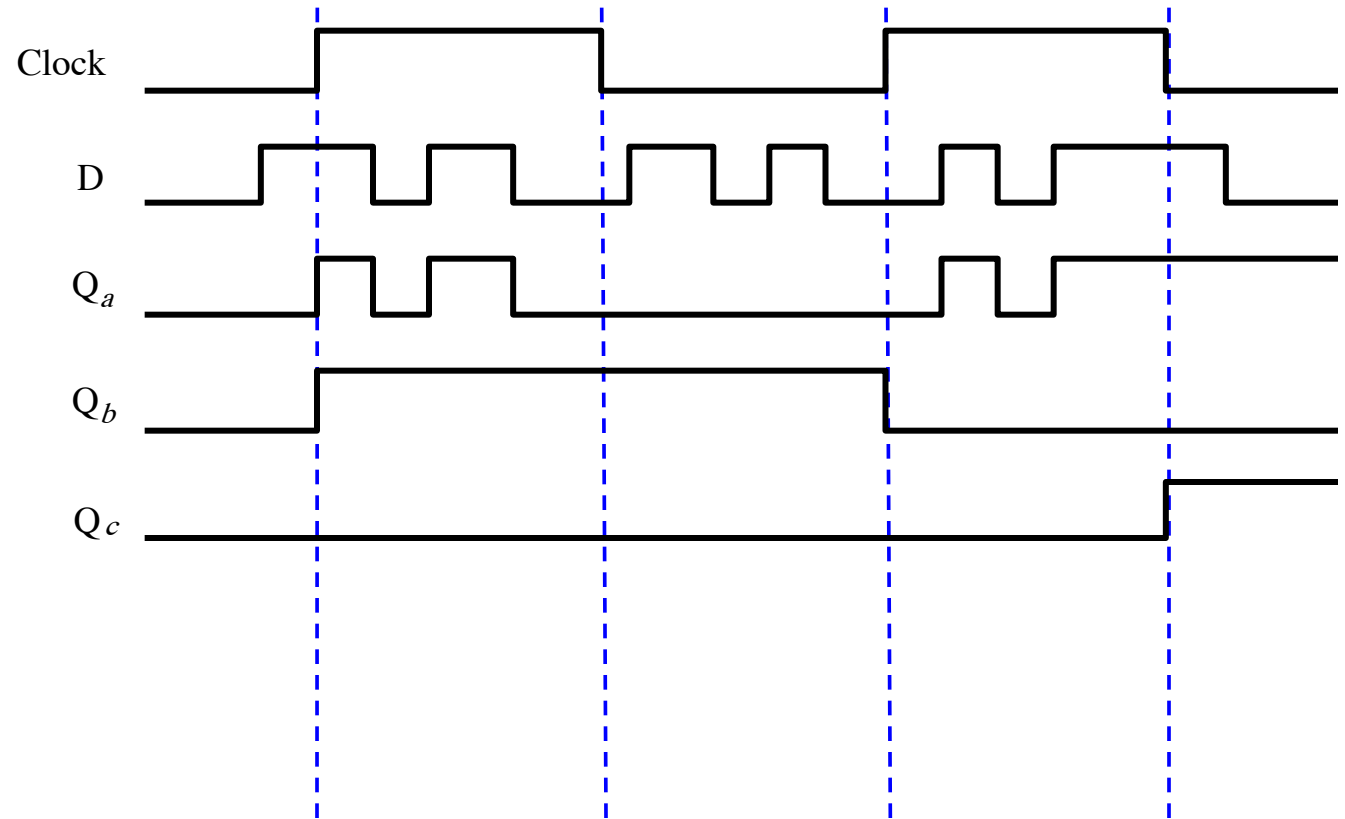
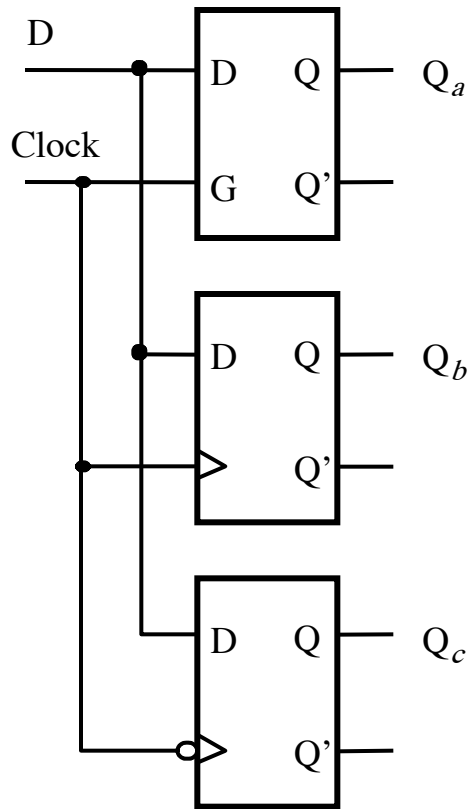
FF is easier to control!

Flip Flop vs. Latch

- **Both are storage elements in sequential circuits**
- **Flip flop**
 - **edge-sensitive**, the input matters only at active edges (rising or falling)
 - behaviors are **synchronous** to the clock signal
- **Latch**
 - **level-sensitive**, the input matters whenever control has active level (high or low)
 - behaviors are **asynchronous** to the clock signal

不同时的
不只被clock 控制，还会因input而改变

Flip Flop vs. Latch

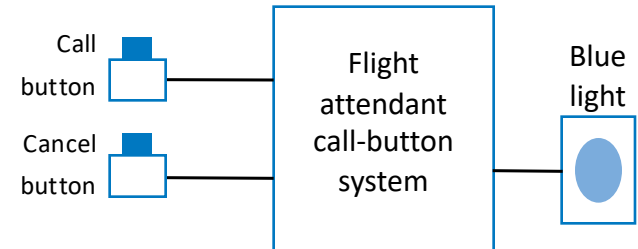


Flight-Attendant Call Button Using D Flip-Flop

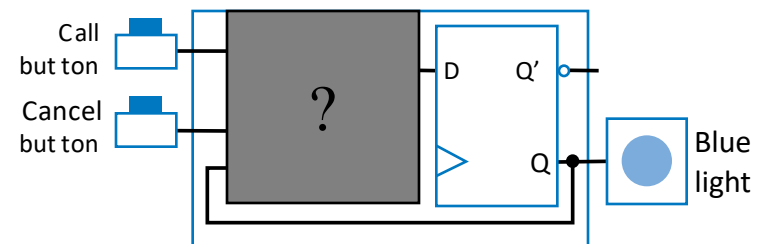
- D flip-flop will store bit
- Inputs are Call, Cancel, and present output Q of D flip-flop
- Truth table

Call	Cancel	Q	$Q^+ = D$
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

keep what it is used to be

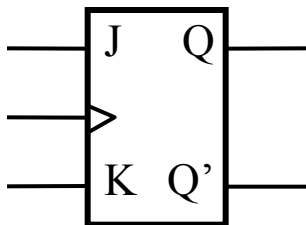
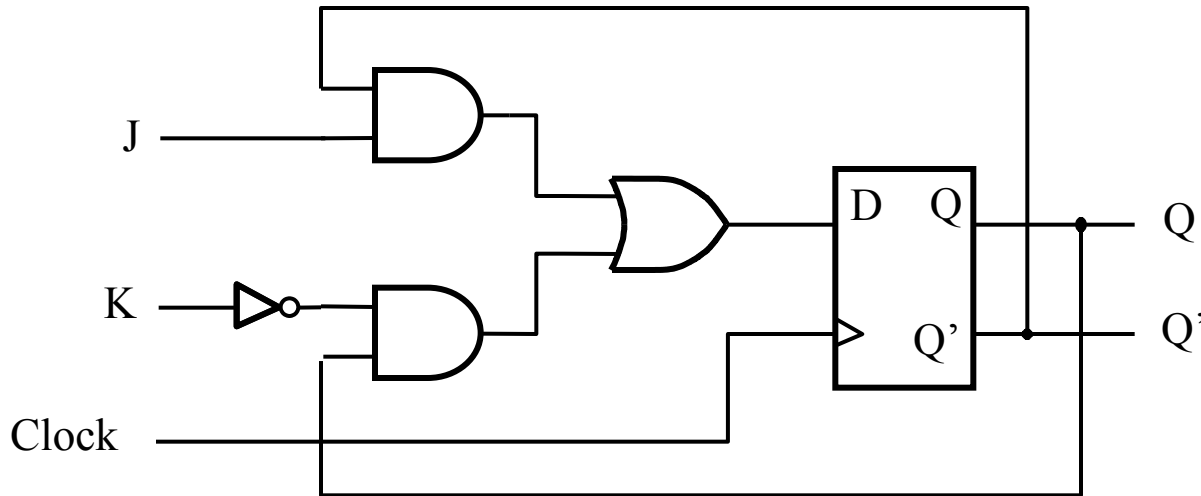


Circuit derived from truth table, using Chapter 2 combinational logic design process



虽然call 与cancel不能同时按，但是设计者可以决定当发生这种情况的时候应该发生什么

Rising Edge-triggered J-K Flip Flop



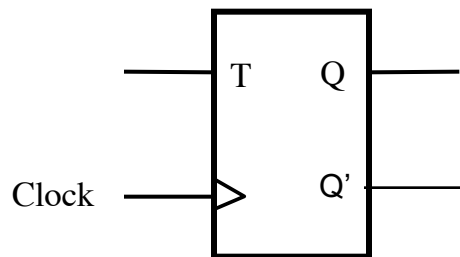
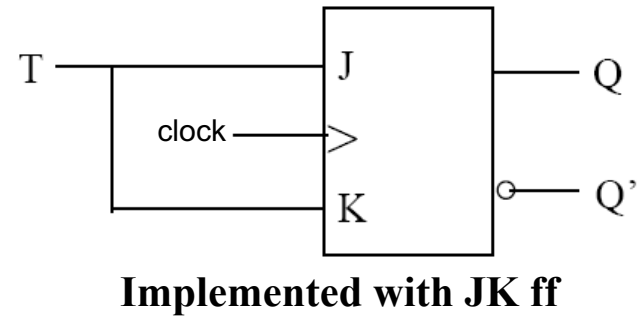
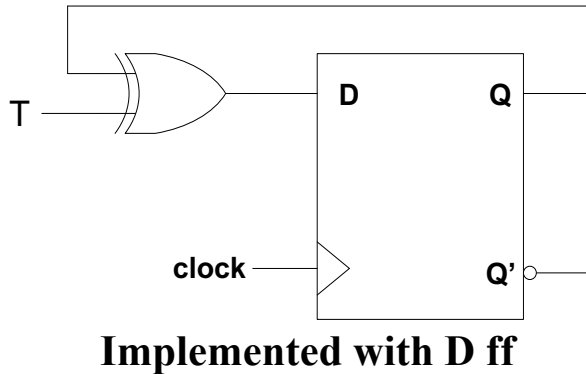
J	K	Q ⁺
0	0	Q
0	1	0
1	0	1
1	1	Q'

k for clear, J for set

Characteristic equation:
 $Q^+ = JQ' + K'Q$

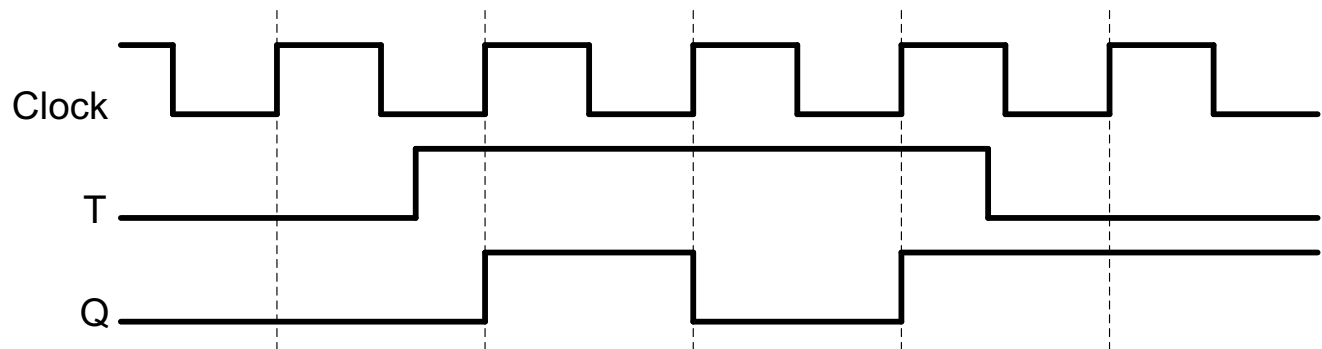
仍然是edge sensitive的

Rising Edge-triggered T Flip Flop



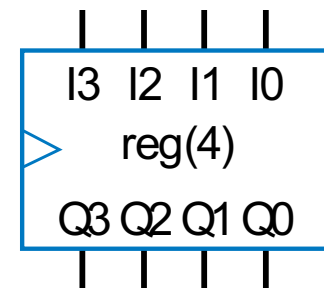
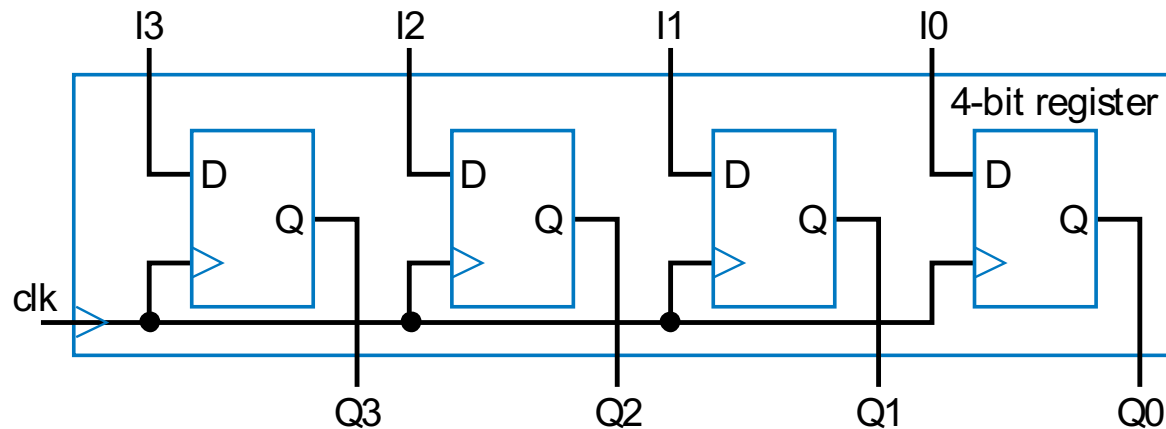
clock	T	Q ⁺
\uparrow	0	Q
\uparrow	1	Q'

Characteristic equation:
 $Q^+ = T'Q + TQ' = T \oplus Q$



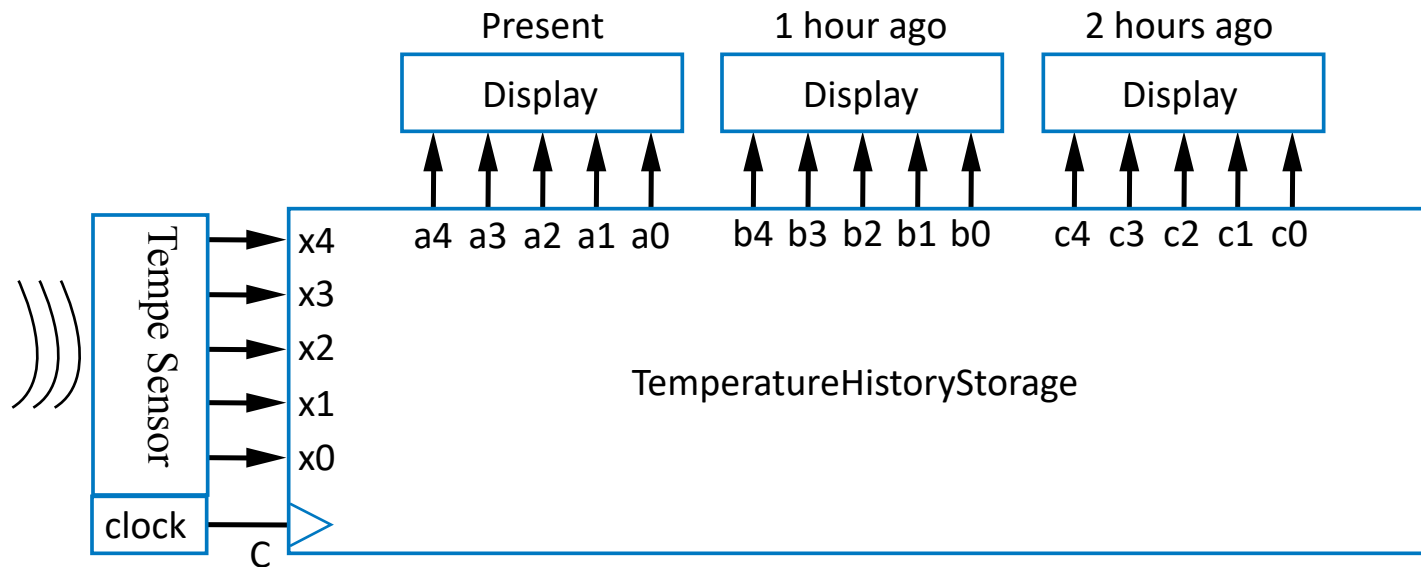
Basic Register

- Typically, we store multi-bit items
 - e.g., storing a 4-bit binary number
- **Register**: multiple flip-flops sharing clock signal



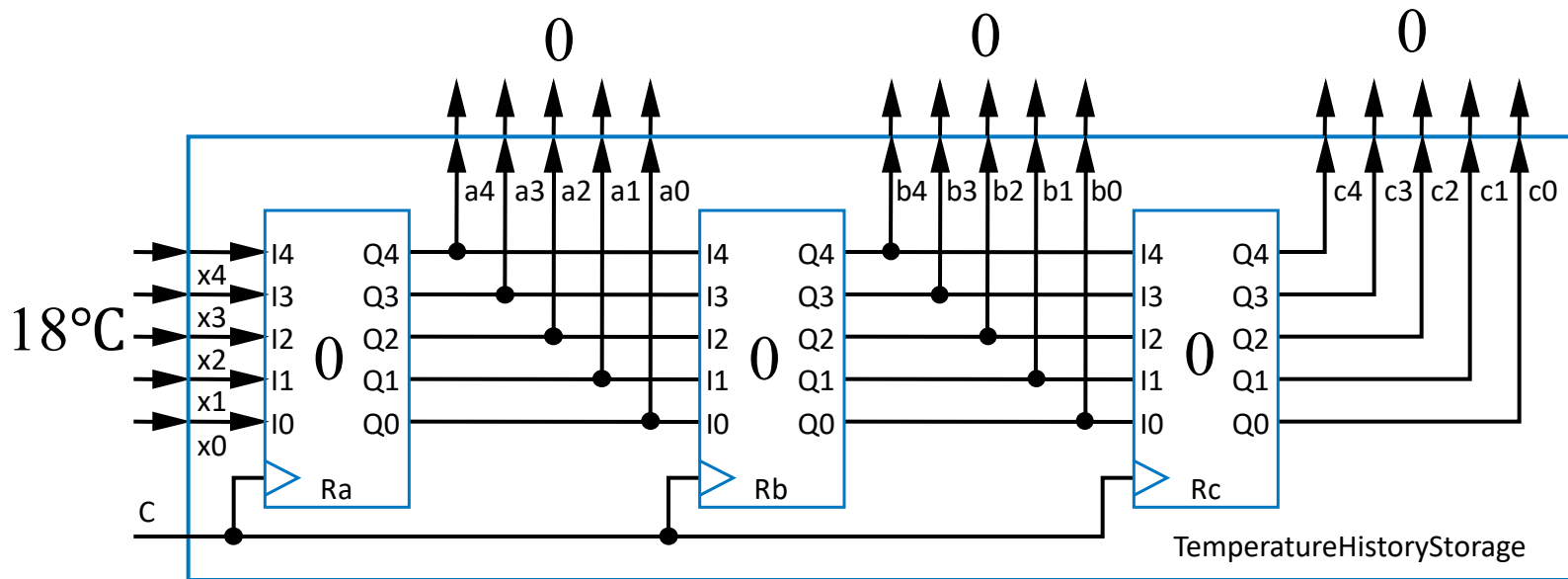
Example Using Registers: Temperature Display

- Temperature history display
 - Sensor outputs temperature as 5-bit binary number
 - Timer pulses C every hour
 - Record temperature on each pulse, display last three recorded values



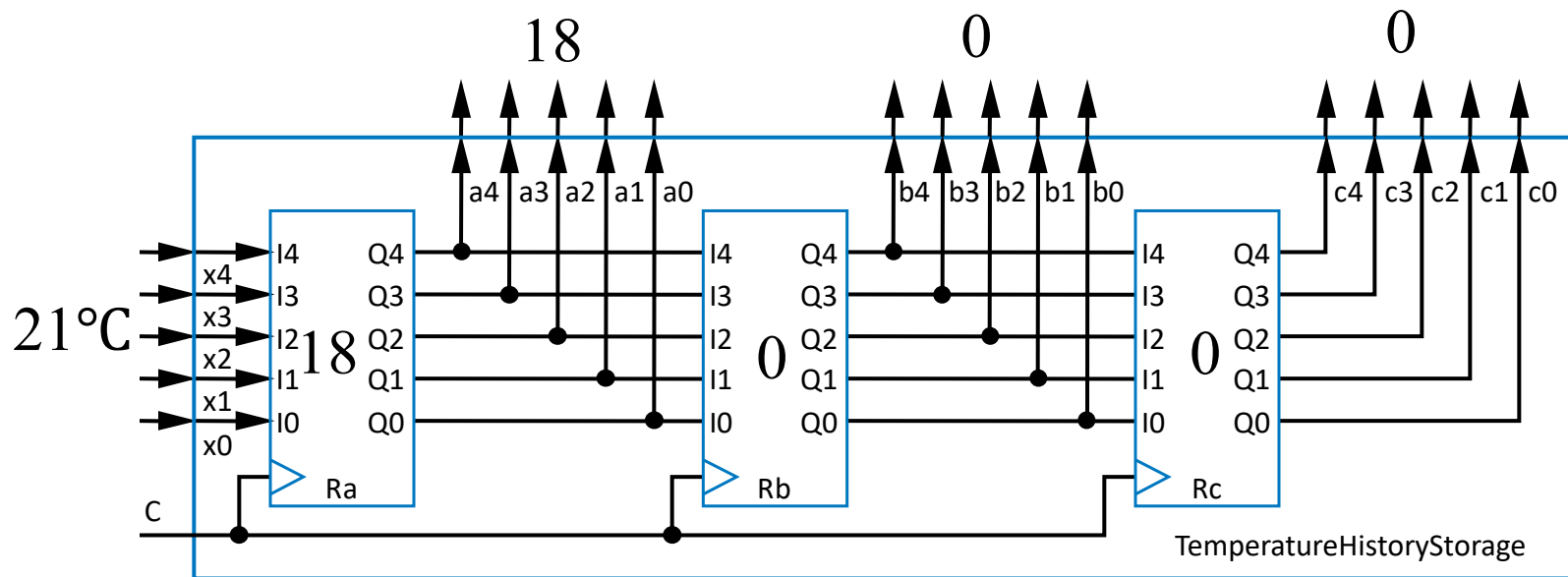
Example Using Registers: Temperature Display

- Use three 5-bit registers



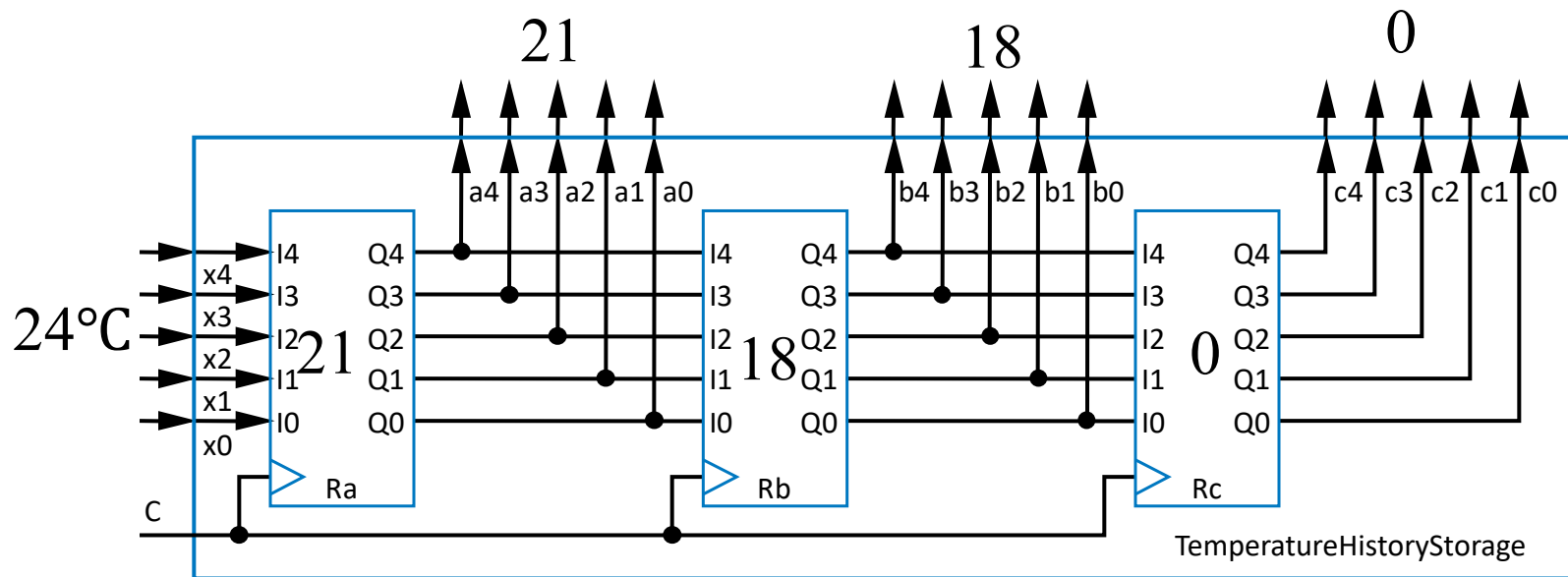
Example Using Registers: Temperature Display

- 1st rising edge of clock



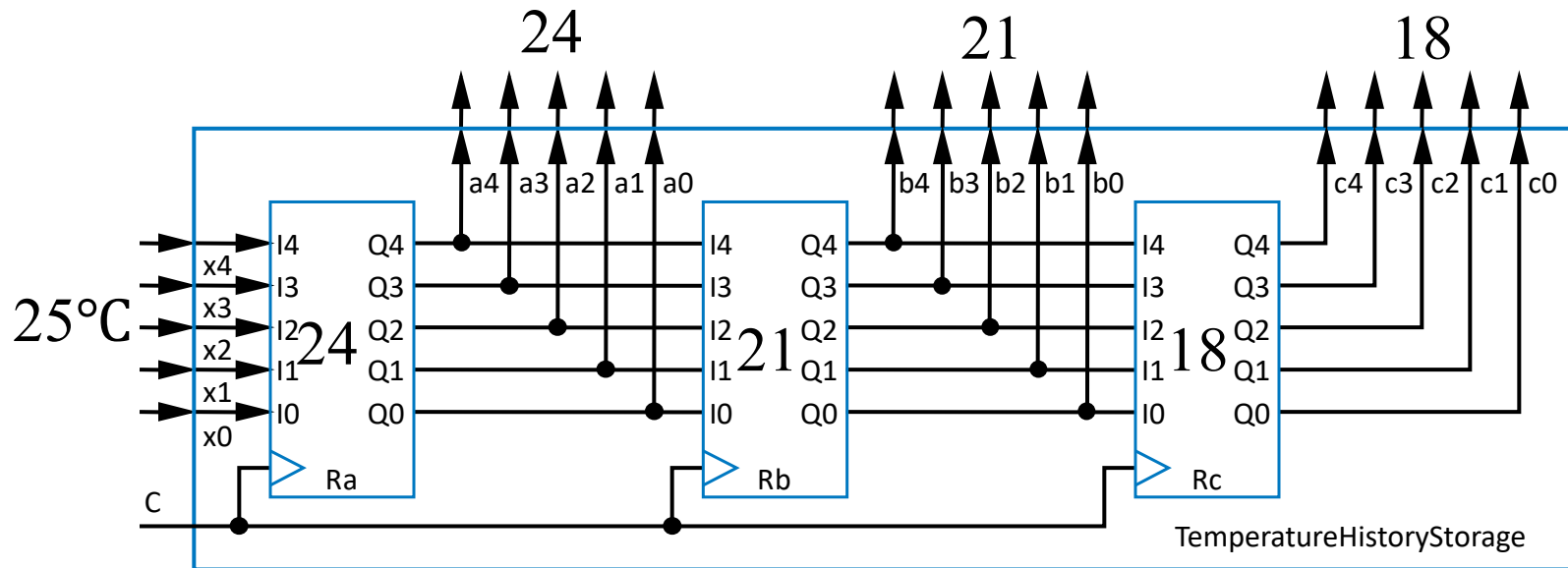
Example Using Registers: Temperature Display

- 2nd rising edge of clock



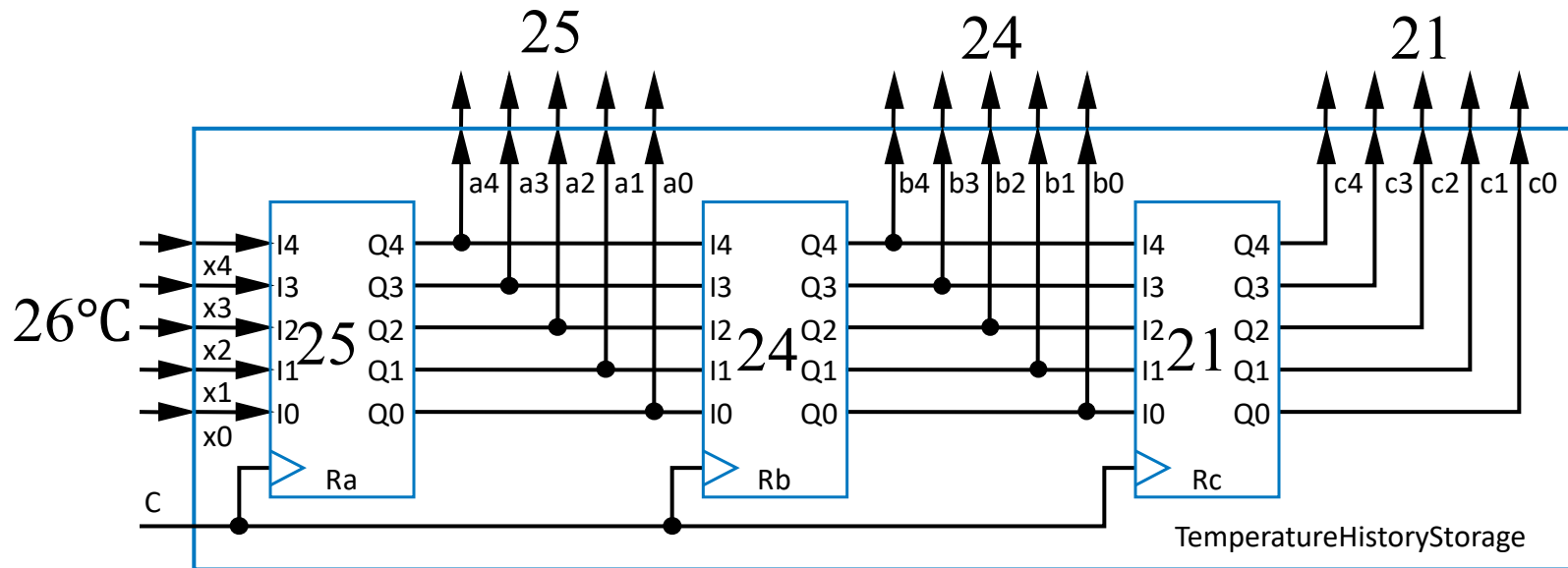
Example Using Registers: Temperature Display

- 3rd rising edge of clock



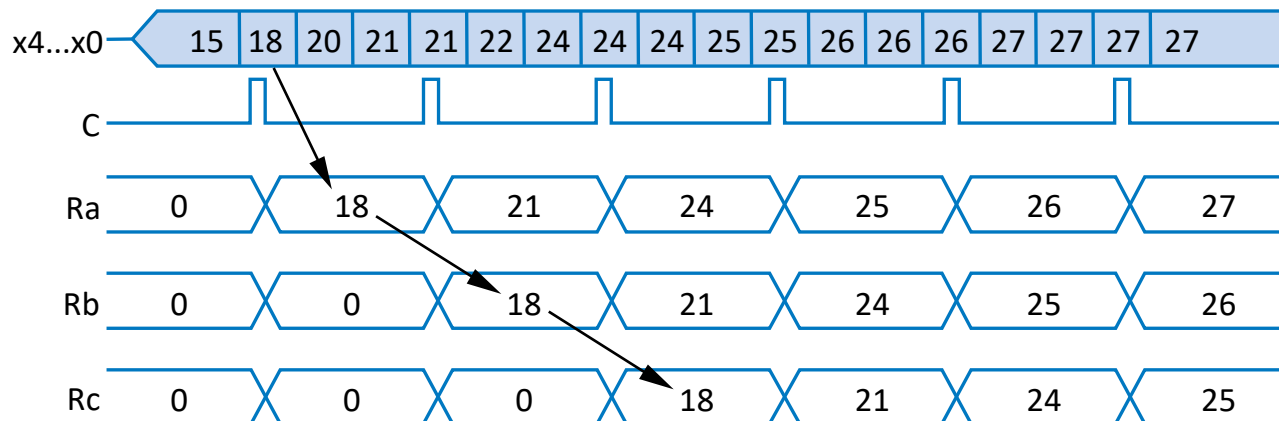
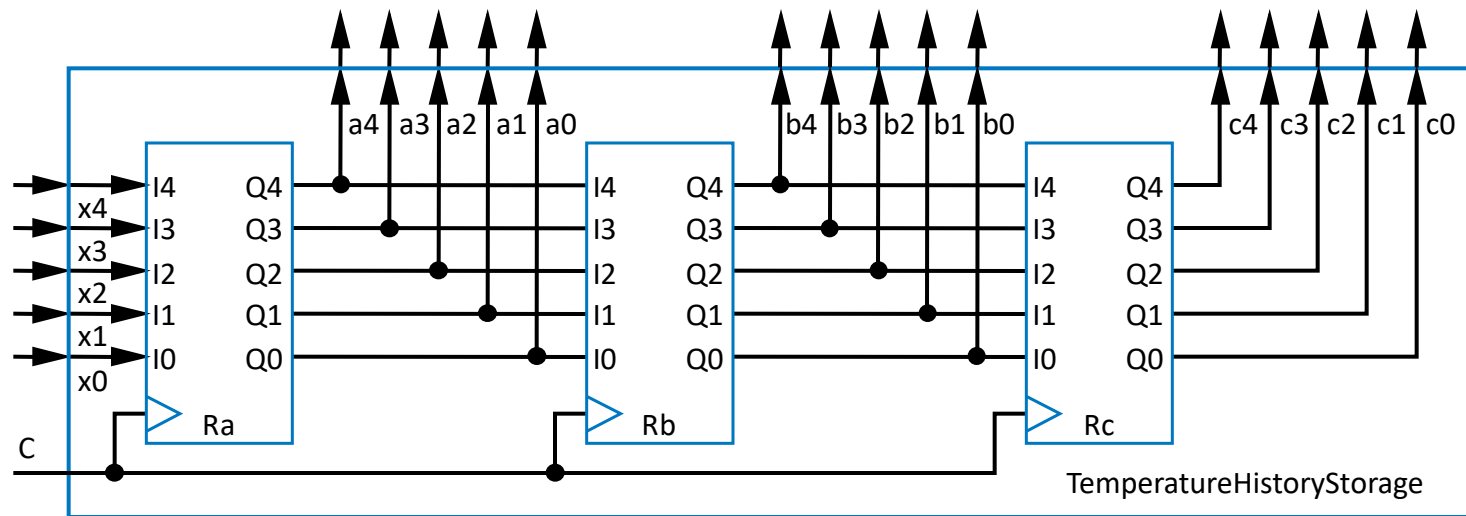
Example Using Registers: Temperature Display

- 4th rising edge of clock

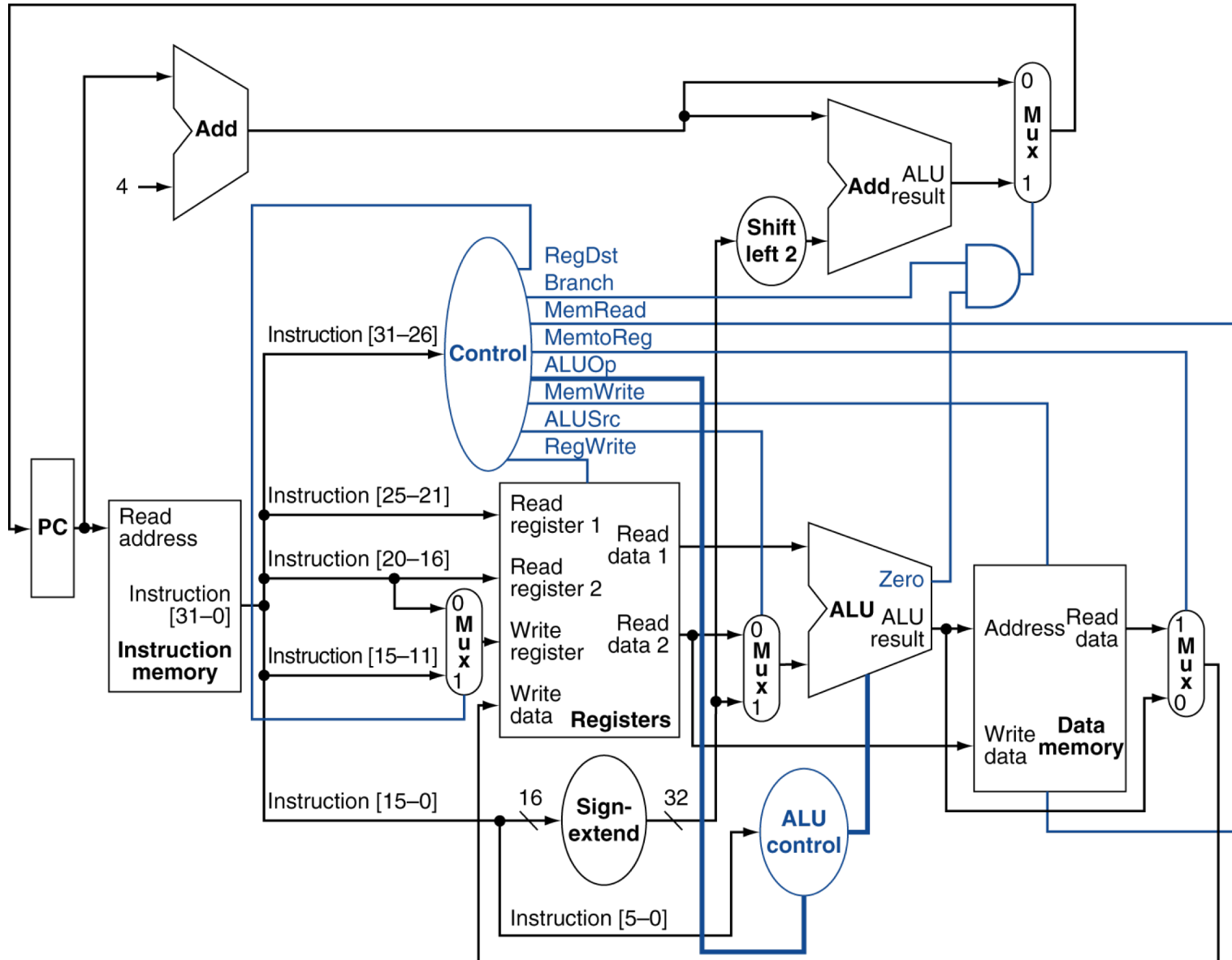


Example Using Registers: Temperature Display

- Use three 5-bit registers

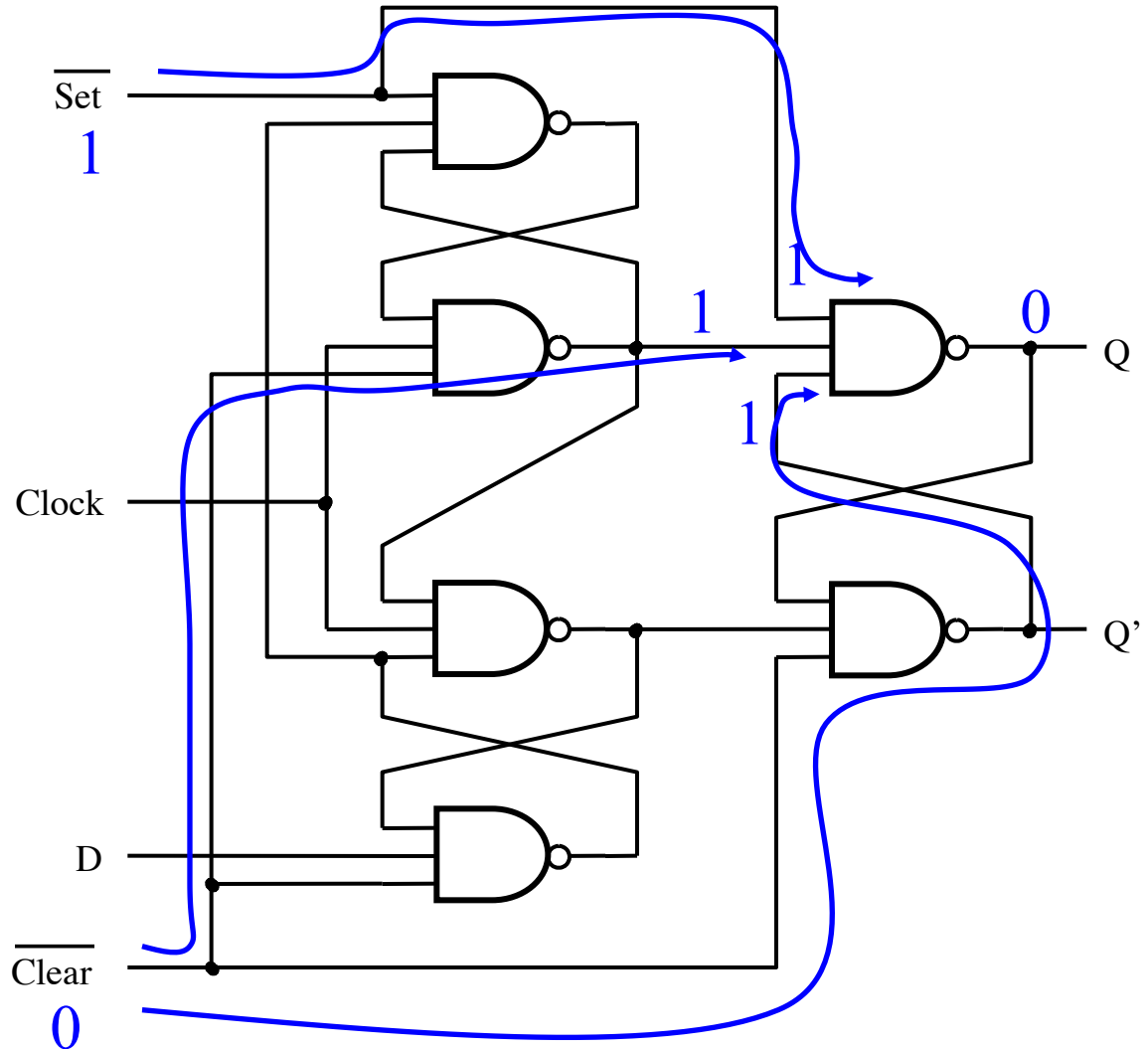


Big Picture – Simplified CPU



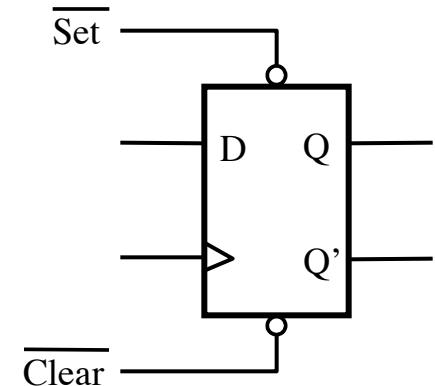
Implementation of Asynchronous Control Input

Control signals decides the output value directly independent of the clock signal



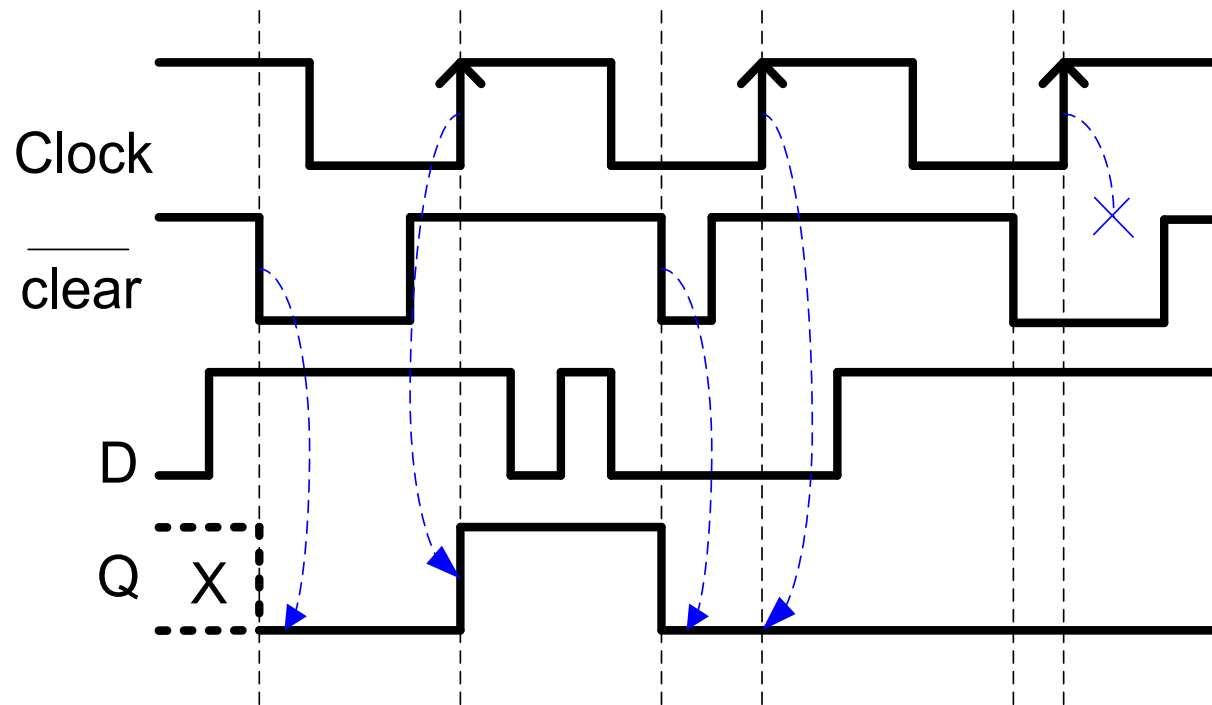
Control Inputs for Flip Flops

- Asynchronous:
 - control signals do not depend on the clock signal
- Synchronous:
 - control signals depend on the clock signal
- Active low:
 - It controls when it's low
- Active high:
 - It controls when it's high



Asynchronous Control Input

- D flip flop with active low asynchronous Clear

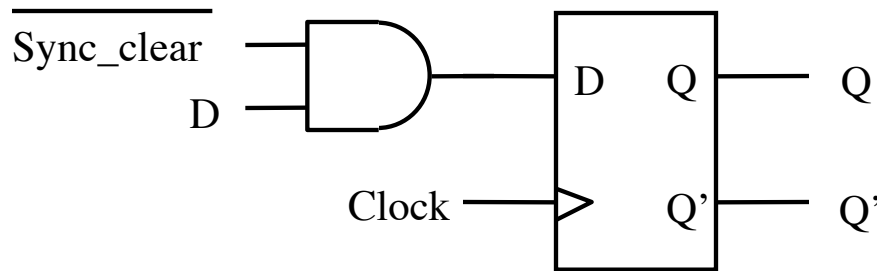


asynchronous clear has higher priority than the clock signal

Implementation of Synchronous Control Input

- Synchronous Clear
 - control signal depends on the active edge (either rising or falling) of the clock signal

Clock 说了算



Flip-Flops with Control Inputs

- D flip flop with active low synchronous Clear

