SQL Overview

- Query capabilities
 - -SELECT-FROM-WHERE blocks,
 - -Basic features, ordering, duplicates
 - –Set operations (union, intersect, except)
 - –Aggregation & Grouping
 - Nested queries (correlation)
 - -Null values

Set operations

- UNION
- INTERSECTION
- EXCEPT (sometimes called MINUS)
- Recall: schemas must match for these operations.

UNION example

Find the names of sailors who have reserved a red or a green boat.

SELECT sname FROM Sailors S, Reserves R, Boats B WHERE S.sid = R.sid AND R.bid = B.bid AND B.color = 'red'

UNION

SELECT sname FROM Sailors S, Reserves R, Boats B WHERE S.sid = R.sid AND R.bid = B.bid AND B.color = 'green'

UNION

- Duplicates ARE NOT eliminated by default in basic SELECT-FROM-WHERE queries
- Duplicate ARE eliminated by default for UNION queries.
- To preserve duplicates in UNION, you must use UNION ALL

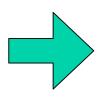
UNION example, alternative:

Find the names of sailors who have reserved a red or a green boat.

SELECT DISTINCT sname
FROM Sailors S, Reserves R, Boats B
WHERE S.sid = R.sid AND R.bid = B.bid
AND (B.color = 'red' OR B.color = 'green')

A small change in this query...

Find the names of sailors who have reserved a red or a green boat.



Find the names of sailors who have reserved a red and a green boat.

SELECT DISTINCT sname
FROM Sailors S, Reserves R, Boats B
WHERE S.sid = R.sid AND R.bid = B.bid
AND (B.color = 'red' **OR** B.color = 'green')

SELECT sname
FROM Sailors S, Reserves R, Boats B
WHERE S.sid = R.sid AND R.bid = B.bid
AND (B.color = 'red' **AND** B.color = 'green')

This doesn't work! What does this query return?

Find the names of sailors who have reserved a red and a green boat.

SELECT sname FROM Sailors S, Reserves R, Boats B WHERE S.sid = R.sid AND R.bid = B.bid AND B.color = 'red'

INTERSECT

SELECT sname FROM Sailors S, Reserves R, Boats B WHERE S.sid = R.sid AND R.bid = B.bid AND B.color = 'green'

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Aggregation

SELECT Avg(S.age)

FROM Sailors

WHERE S.rating = 10

SQL supports several aggregation operations:

COUNT (*)
COUNT ([DISTINCT] A)
SUM ([DISTINCT] A)
AVG ([DISTINCT] A)
MAX (A)
MIN (A)

Aggregation: Count

```
SELECT Count(*)
FROM Sailors
WHERE rating > 5
```

Except for COUNT, all aggregations apply to a single attribute

Aggregation: Count

COUNT applies to duplicates, unless otherwise stated:

```
SELECT Count(category)FROM ProductWHERE year > 1995
```

Better:

```
SELECT Count(DISTINCT category)
FROM Product
WHERE year > 1995
```

Simple Aggregation

Purchase(product, date, price, quantity)

Example 1: find total sales for the entire database

```
SELECT Sum(price * quantity)
FROM Purchase
```

Example 1': find total sales of bagels

```
SELECT Sum(price * quantity)
FROM Purchase
WHERE product = 'bagel'
```

GROUP BY and HAVING clauses

 We often want to apply aggregates to each of a number of groups of rows in a relation.

Find the age of the youngest sailor for each rating level.

SELECT MIN (S.age) FROM Sailors S WHERE S.rating = i

For
$$i = 1, 2, ... 10$$

Grouping

Sailors

sid	sname	rating	age		SI FF
 29	brutus	1	33		G
85	art	3	25.5		
95	bob	3	63.5		
96	frodo	3	25.5		
22	dustin	7	45		
64	horatio	7	35		
31	lubber	8	55.5		
32	andy	8	25.5		
74	horatio	9	35		
58	rusty	10	35		
71	zorba	10	16		

SELECT S.rating, MIN(S.age)
FROM Sailors S
GROUP BY S.rating

New Table



Queries With GROUP BY and HAVING

SELECT [DISTINCT] target-list

FROM relation-list

WHERE qualification

GROUP BY grouping-list

HAVING group-qualification

- The target-list contains (i) attribute names (ii) terms with aggregate operations (e.g., MIN (S.age)).
 - The <u>attribute list (i)</u> must be a subset of <u>grouping-list</u>. Intuitively, each answer tuple corresponds to a <u>group</u>, and these attributes must have a single value per group.

Conceptual Evaluation

- The cross-product of *relation-list* is computed, tuples that fail *qualification* are discarded, `*unnecessary*' fields are deleted, and the remaining tuples are partitioned into groups by the value of attributes in *grouping-list*.
- The group-qualification is then applied to eliminate some groups. Expressions in groupqualification must have a single value per group!
- One answer tuple is generated per qualifying group.

Find age of the youngest sailor with age ≥ 18 , for each rating with at least 2 <u>such</u> sailors

SELECT S.rating, MIN (S.age)
AS minage

WHERE S.age >= 18

FROM Sailors S

GROUP BY S.rating

HAVING COUNT (*) > 1

Answer relation:

rating	minage
3	25.5
7	35.0
8	25.5

Sailors instance:

sid	sname	rating	age
22	dustin	7	45.0
29	brutus	1	33.0
31	lubber	8	55.5
32	andy	8	25.5
58	rusty	10	35.0
64	horatio	7	35.0
71	zorba	10	16.0
74	horatio	9	35.0
85	art	3	25.5
95	bob	3	63.5
96	frodo	3	25.5

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Nested queries

- A nested query is a query with another query embedded within it.
- The embedded query is called the subquery.
- The subquery usually appears in the WHERE clause:

```
SELECT S.sname
FROM Sailors S
WHERE S.sid IN ( SELECT R.sid
FROM Reserves R
WHERE R.bid = 103 )
```

(Subqueries also possible in FROM or HAVING clause.)

Conceptual evaluation, extended

 For each row in cross product of outer query, evaluate the WHERE clause conditions, (re)computing the subquery.

```
SELECT S.sname
FROM Sailors S
WHERE S.sid IN ( SELECT R.sid
FROM Reserves R
WHERE R.bid = 103 )
```

equivalent to:

```
SELECT S.sname
FROM Sailors S, Reserves R
WHERE S.sid=R.sid AND R.bid=103
```

Correlated subquery

- If the inner subquery depends on tables mentioned in the outer query then it is a correlated subquery.
- In terms of conceptual evaluation, we must recompute subquery for each row of outer query.

```
SELECT S.sname
FROM Sailors S
WHERE EXISTS ( SELECT *
FROM Reserves R
WHERE R.bid = 103
AND R.sid = S.sid )
```

Set-comparison operators

- Optional NOT may precede these:
 - **-EXISTS** R -- true if R is non-empty
 - -attr IN R -- true if R contains attr
 - -UNIQUE R -- true if no duplicates in R
- For arithmetic operator op {<,<=,=,< >, >=,>}
 - -attr op ALL R -- all elements of R satisfy condition
 - -attr op ANY R -- some element of R satisfies condition

Example

Find the sailors with the highest rating

```
SELECT S.sid
FROM Sailors S
WHERE S.rating >= ALL (SELECT S2.rating
FROM Sailors S2)
```

Please write SQL

 Find sailors whose rating is higher than some sailor named Horatio.

```
SELECT S.sid
FROM Sailors S
WHERE S.rating > ANY (SELECT S2.rating
FROM Sailors S2
WHERE S2.name = 'Horatio')
```

 Find sailors whose rating is higher than all sailors named Horatio.

```
SELECT S.sid
FROM Sailors S
WHERE S.rating > ALL (SELECT S2.rating
FROM Sailors S2
WHERE S2.name = 'Horatio')
```

Find boats **not** reserved by sailor with sid = 100.

- B: all boats
- R: boats reserved by sailor with sid=100
- B R is what we want.

```
SELECT B.bid
FROM Boats B
WHERE B.bid NOT IN (SELECT R.bid
FROM Reserves R
WHERE R.sid = 100);
```

Existential conditions

- Find the names of sailors who have reserved some boat
- (i.e. *there exists* a boat they reserved)

```
SELECT S.sname
FROM Sailors S, Reserves R
WHERE S.sid=R.sid
```

Existential conditions are natural and easy.

Universal conditions

- Find the names of sailors who have reserved all boats.
- (i.e. *for each* boat, they have reserved it.)
- Universal conditions are harder.

```
SELECT S.sname
FROM Sailors S
WHERE NOT EXISTS (
Set of boats not reserved by S.sid
)
```

Universal conditions

 Find the names of sailors who have reserved all boats.

```
SELECT S.sname
FROM Sailors S
WHERE NOT EXISTS (

SELECT B.bid
FROM Boats B
WHERE B.bid NOT IN (SELECT R.bid
FROM Reserves R
WHERE R.sid = S.sid )
```

For each sailor, check that there is no boat s/he hasn't reserved.

Simulating INTERSECT

- Suppose we have tables R(a,b) and S(a,b)
- The following computes R ∩ S:

```
SELECT DISTINCT *
FROM R
WHERE (R.a, R.b) IN (SELECT *
FROM S );
```

This can be expressed without nesting:

Given R(a,b), S(a,b),
 what is R ⋈S?

Intersection!

SELECT DISTINCT R.a, R.b FROM R, S WHERE R.a = S.a AND R.b = S.b;

Find the names of sailors who reserved a red and a green boat.

using INTERSECT

SELECT sname FROM Sailors S, Reserves R, Boats B WHERE S.sid = R.sid AND R.bid = B.bid AND B.color = 'red'

INTERSECT

SELECT sname FROM Sailors S, Reserves R, Boats B WHERE S.sid = R.sid AND R.bid = B.bid AND B.color = 'green'

without INTERSECT

SELECT sname
FROM Sailors S, Reserves R, Boats B
WHERE S.sid = R.sid AND R.bid = B.bid AND B.color = 'red'
AND S.sid IN
(SELECT S2.sid
FROM Sailors S2, Reserves R2, Boats B2
WHERE S2.sid = R2.sid AND R2.bid = B2.bid AND B2.color = 'green')

"Find all sailors who have reserved a red boat and, further, have **sids** that are included in the set of **sids** of sailors who have reserved a green boat."

Simulating EXCEPT (set difference)

What does this query compute?

```
SELECT DISTINCT *
FROM R
WHERE (R.a, R.b) NOT IN (SELECT *
FROM S );
```

Can this be expressed without a nested query? No.

(But this fact is not obvious)