PBE-Based Selective Abstraction and Refinement for Efficient Property Falsification of Embedded Software

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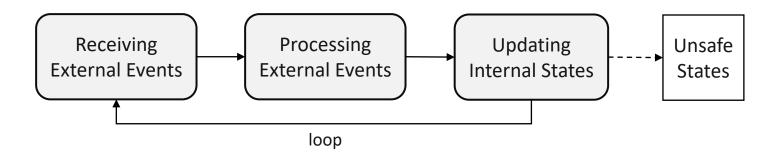
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Motivation

- Major characteristics of embedded software:
 - It constantly interacts with external events within an infinite loop
 - It maintains internal states whose updates must be traced
 - Safety properties may be violated only after many loop iterations





Motivation

- Verification/Falsification of embedded software:
 - e.g., (dis)proving "A car reduces its speed within a specified time"
 - We have seen decades of effort for this:
 - Most of them were model-based approaches applied to cyber-physical systems, automobiles, or medical devices, etc.
 - Code-based property verification/falsification is very expensive

Table 1: Code-based property falsification of an elevator controller (about 900 LoC)

Tools	# of refinements	Time usage	Memory usage
СВМС	N/A	8.3 hours	> 80 GB (OOM)
CPAchecker	418	> 72 hours (T/O)	7.1 GB



Motivating Example

int process(int *data) { int $max_area = -1$; int $max_index = -1$; for(int i = 0; i < data[0]; i++) { $if(data[5 * i + 1] != 0) {$ int xul = data[i * 5 + 2]; int xlr = data[i * 5 + 3];int yul = data[i * 5 + 4]; int ylr = data[i * 5 + 5];int width = abs(xul - xlr): int height = abs(yul - ylr); int area = width * height; if(area >= max area) { max area = area; $max_index = i;$ return max_index; }

Embedded Control Program

```
int timer = 0:
int speed = 0;
#define MAX_SPEED 100
int main() {
  int data[41];
  int i;
                                     void update(int i, int *data) {
                                       speed = speed + data[i * 5 + 1];
  while(1) {
                                       if(speed > MAX SPEED) {
    data = receive();
                                         speed = MAX SPEED; }}
    i = process(data);
    if(i >= 0) {
                                     int is safe() {
      update(i, data);
                                       if(speed >= MAX_SPEED){
                                         timer = timer + 1;
    assert(is safe());
                                       else timer = 0;
                                       return timer < 1000;}</pre>
```

- Checking this concrete program suffers from state space explosion (OOM)
- Abstracting this program too coarsely produces many false alarms



Motivating Example

Embedded Control Program

```
int timer = 0:
int process(int *data) {
                                  int speed = 0;
 int max_area = -1;
 int max_index = -1;
                                  #define MAX_SPEED 100
 for(int i = 0; i < data[0]; i++) {
   if(data[5 * i + 1] != 0) {
                                  int main()
    int
          We need to find a better abstraction approach
    int
    int
                                                                                                  + 1];
    int
                       for embedded software domain
    int
    int
    if(area >= max_area) {
                                      if(i >= 0) {
      max area = area;
                                                                      int is safe() {
                                        update(i, data);
      max index = i;
                                                                        if(speed >= MAX_SPEED){
                                                                         timer = timer + 1;
                                      assert(is safe());
                                                                       else timer = 0;
 return max_index; }
                                                                        return timer < 1000:}</pre>
```

- Checking this concrete program suffers from state space explosion (OOM)
- Abstracting this program too coarsely produces many false alarms



Main Idea 1: Selective Abstraction

Auxiliary Function

```
int process(int *data) {
 int max_area = -1;
 int max_index = -1;
 for(int i = 0; i < data[0]; i++) {
   if(data[5 * i + 1] != 0) {
      int xul = data[i * 5 + 2]:
     int xlr = data[i * 5 + 3]:
     int yul = data[i * 5 + 4];
     int ylr = data[i * 5 + 5];
     int width = abs(xul - xlr):
     int height = abs(yul - ylr);
     int area = width * height;
     if(area >= max area) {
       max area = area;
       max_index = i;
  return max_index; }
```

```
Main Control Logic
int timer = 0:
int speed = 0;
#define MAX_SPEED 100
int main() {
  int data[41];
  int i;
                                    void update(int i, int *data) {
  while(1) {
                                      speed = speed + data[i * 5 + 1];
                                      if(speed > MAX_SPEED) {
    data = receive();
                                        speed = MAX SPEED; }}
    i = process(data);
    if(i >= 0) {
                                    int is safe() {
      update(i, data);
                                      if(speed >= MAX_SPEED){
                                        timer = timer + 1;
    assert(is_safe());
```

else timer = 0;

return timer < 1000;}</pre>

 Auxiliary functions: functions that do not update any internal states (e.g., global variables)



Main Idea 1: Selective Abstraction

Auxiliary Function

```
int process(int *data) {
 int max_area = -1;
 int max_index = -1;
 for(int i = 0; i < data[0]; i++) {</pre>
   if(data[5 * i + 1] != 0) {
     int xul = data[i * 5 + 2];
     Abstraction!
     int height = abs(yul - ylr);
     int area = width * height;
     if(area >= max area) {
       max area = area;
       max_index = i;
  return max_index; }
```

```
Main Control Logic
int timer = 0:
int speed = 0;
#define MAX_SPEED 100
int main() {
                                  No
  int data[41];
  int i;
                                                   i, int *data) {
                                                   + data[i * 5 + 1];
  while(1) {
                          Abstraction
                                                  X_SPEED) {
    data = receive();
                                                   SPEED; }}
    i = process(data);
    if(i >= 0) {
                                   int is safe() {
      update(i, data);
                                     if(speed >= MAX_SPEED){
                                        timer = timer + 1;
    assert(is_safe());
                                     else timer = 0;
                                      return timer < 1000;}</pre>
```

- Main control logic has a significant impact on the program & property
- Abstracting the main control logic is likely to produce many false alarms



Main Idea 2: PBE-Based Abstraction

(1) Initial Abstraction

```
int process(int *data) {
  return nondet();
}
```

240 seconds, but false alarm

(2) Refined Abstraction

```
int process(int *data) {
  if(0 != data[36])
    if(data[0] >= data[15]) return 6;
    else return 7;
  else if(0 != data[1])
    if(data[12] < data[40]) return 5;
    else if(data[18] >= data[20]) return 1;
    else return 0;
  else if(0 != data[21]) return 4;
  else return -1; }
```

260 seconds, and true alarm!

(3) No Abstraction

```
int process(int *data) {
 int max area = -1;
 int max_index = -1;
 for(int i = 0; i < data[0]; i++) {</pre>
   if(data[5 * i + 1] != 0) {
      int xul = data[i * 5 + 2];
      int xlr = data[i * 5 + 3];
      int yul = data[i * 5 + 4];
      int ylr = data[i * 5 + 5];
      int width = abs(xul - xlr);
      int height = abs(yul - ylr);
      int area = width * height;
      if(area >= max_area) {
        max_area = area;
        max index = i;
 return max index; }
```

Out-Of-Memory

Precise & Expensive

Scalable & Imprecise

Precise & Scalable





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Main Idea 2: PBE-Based Abstraction

(1) Initial Abstraction

```
int process(int *data) {
  return nondet();
}
```

240 seconds, but false alarm

(2) Programming-By-Example (PBE) Based Abstraction

```
int process(int *data) {
   if(0 != data[36])
      if(data[0] >= data[15]) return 6;
      else return 7;
   else if(0 != data[1])
      if(data[12] < data[40]) return 5;
      else if(data[18] >= data[20]) return 1;
      else return 0;
   else if(0 != data[21]) return 4;
   else return -1; }
```

260 seconds, and true alarm!

We could synthesize this abstraction using only 33 input-output examples!

(3) No Abstraction

```
int process(int *data) {
 int max area = -1;
 int max_index = -1;
 for(int i = 0; i < data[0]; i++) {
   if(data[5 * i + 1] != 0) {
     int xul = data[i * 5 + 2];
     int xlr = data[i * 5 + 3];
      int yul = data[i * 5 + 4];
      int ylr = data[i * 5 + 5];
      int width = abs(xul - xlr);
      int height = abs(yul - ylr);
      int area = width * height;
      if(area >= max area) {
        max area = area;
        max_index = i;
     urn max index; }
```

Out-Of-Memory

Precise & Expensive



Scalable & Imprecise

Precise & Scalable

Concrete

KAIINGBUUK

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PBEAR: PBE-Based Selective Abstraction and Refinement

- ✓ PBEAR suggests selective abstraction for efficient property falsification of embedded software
- ✓ PBEAR is the first approach utilizing PBE for more precise & scalable abstraction
- ✓ Out of 15 safety properties from three embedded software,

 PBEAR finds 5 to 12 more true alarms than CBMC and CPAchecker

Imprecise: 13 T/Os with 2,556 refinements
Scalable: 30% less memory than PBEAR

Expensive: Out-Of-Memory in 6 cases

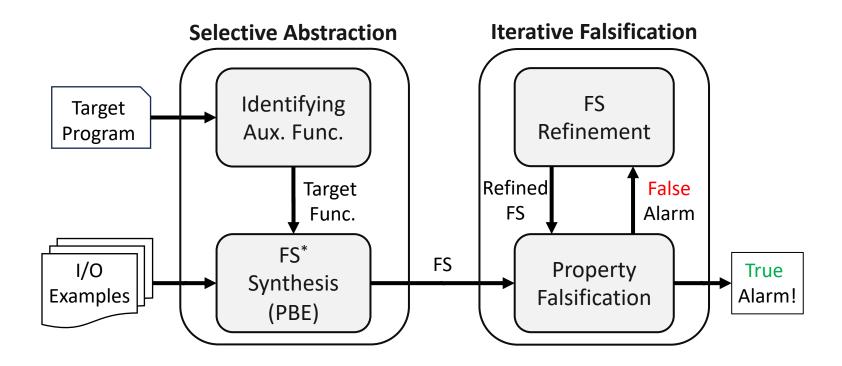
CPAchecker

PBEAR

CBMC



Overall Process of **PBEAR**



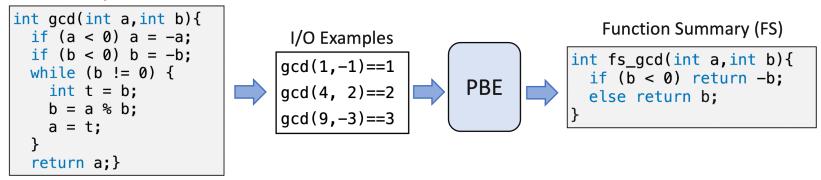
*FS: Function Summary



Problems & Solutions for Utilizing PBE

 PBE takes input/output (I/O) examples and produces a program that satisfies the given examples

Auxiliary Function



- 1. Overfitting problems in PBE:
- ✓ Iterative function summary synthesis
- 2. PBE-based refinement may reproduce infeasible paths:
- ✓ A novel PBE-based refinement with symbolic alarm filtering



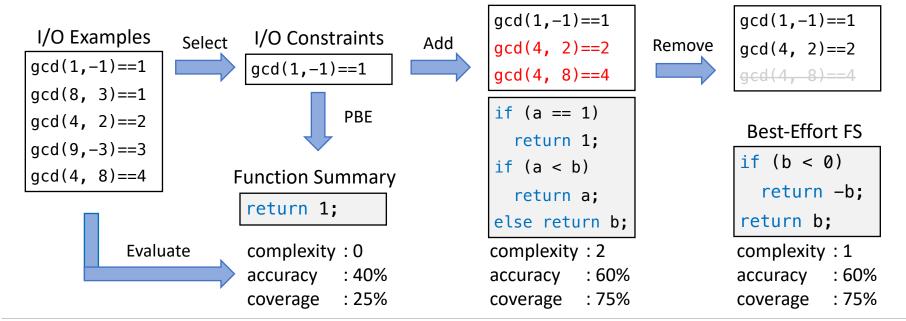
1. Overfitting → Iterative FS Synthesis

Measurement:

- Complexity: # of branch statements
- Similarity: accuracy and output coverage

Synthesize until a given time budget (1 hour)

10



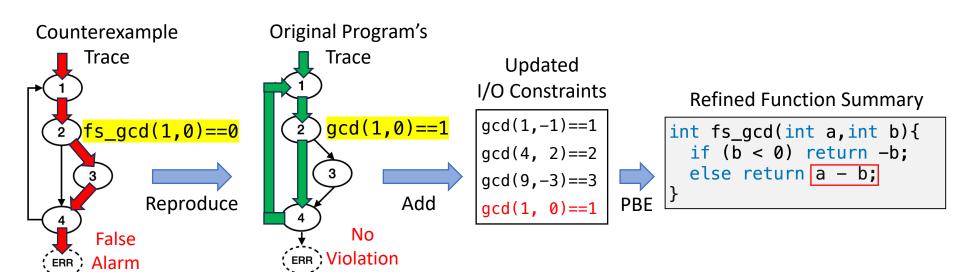


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2. PBE Produces Infeasible Paths

Baseline of refinement approach:

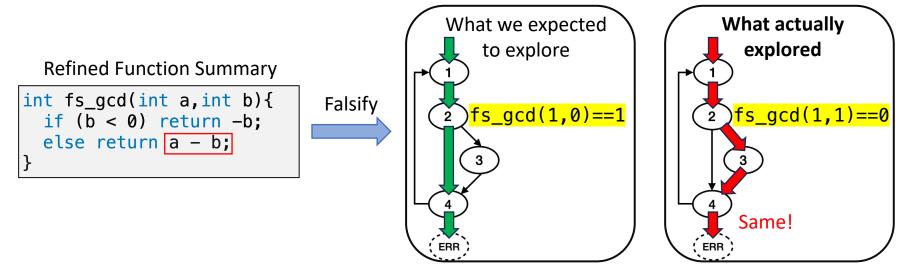
- 1. Add new I/O constraints from the false alarm
- 2. Re-synthesize FS using the updated I/O constraints





2. PBE Reproduces Infeasible Paths

Problem: No guarantee for avoiding re-exploration of **previously identified infeasible paths**



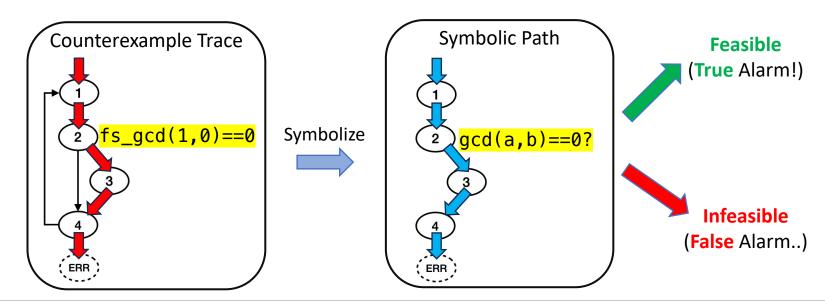
There is **another I/O example** leading to the same infeasible path



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2. Solution: PBE-Based Refinement with Symbolic Alarm Filtering

- 1. Convert the counterexample trace into a symbolic path
- Check the feasibility of the symbolic path using SAT solving

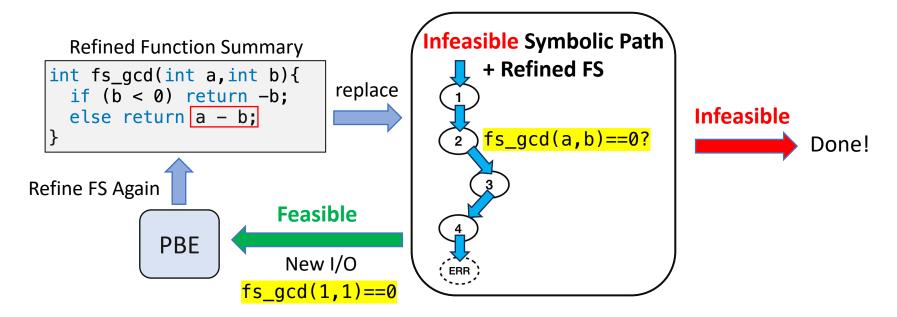




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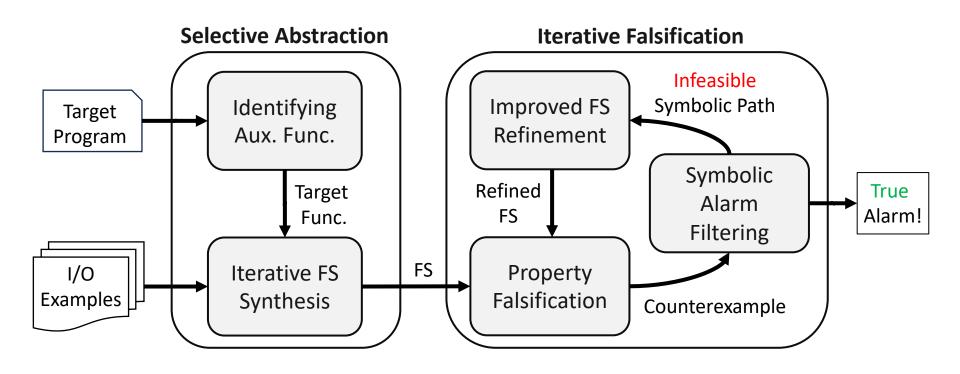
2. Solution: PBE-Based Refinement with Symbolic Alarm Filtering

3. Iteratively refine FS using PBE until it does not make infeasible symbolic path feasible





(Refined) Overall Process of PBEAR





Experiments

- Target
 - Checking assertion violations for 16 benchmark programs from SV-COMP
 - 15 safety properties from 3 embedded software (about 800~1,000 LoC)
 - Object following car, elevator controller, and clean-up robot

- RQ1: Performance of PBEAR vs. two state-of-the-art model checkers,
 - CBMC: bounded model checker
 - CPAchecker: predicate abstraction
- RQ2: Effectiveness of our improved FS refinement vs. PBEAR^{base}
 - PBEAR^{base}: baseline of refinement approach; no symbolic alarm filtering

* PBEAR uses DUET for function summary synthesis and CBMC for property falsification



RQ1: Performance of PBEAR (on SV-COMP)

Tools	# of refinements	Total time	Total memory	# of detected true alarms
PBEAR	9	485 s	3.1 GB	15/16
CBMC	N/A	177 s	3.9 GB	16/16
CPAchecker	57	4,334 s	10.3 GB	13/16

^{*} T/O: 900 s

- CBMC was the best: the overhead of refinement did not pay off on these programs
- PBEAR showed competitive performance on general-purpose programs



RQ1: Performance of PBEAR (on Embedded Software)

Tools	# of refinements	Total time	Total memory	# of detected true alarms
PBEAR	77	135.4 h	423.5 GB	14/15
CBMC	N/A	96.9 h	748.9 GB	9/15
CPAchecker	2,556	936.3 h	336.8 GB	2/15

^{*} T/O: 72 hours; Out-Of-Memory: 80 GB

Imprecise: 13 T/Os with 2,556 refinements

Scalable: 30% less memory than PBEAR

Precise: no need for refinement **Expensive: Out-Of-Memory in 6 cases**







CBMC



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RQ2: Effectiveness of Our Improved FS Refinement

Tools	# of refinements	Total time	# of detected true alarms
PBEAR	86	133.6 h	13/15
PBEAR ^{base} (without symbolic alarm filtering)	210	217.4 h	12/15

^{*} We omitted benchmarks where any refinements were not conducted by PBEARbase

PBEARbase took additional 124 refinements and 83.6 more hours



Conclusion

- Contributions
 - PBEAR is the first work utilizing PBE for selective abstraction
 - We suggested a novel PBE-based refinement approach with symbolic alarm filtering
 - PBEAR showed promising results on three embedded software and competitive performance on general-purpose programs
- Future Work
 - PBEAR is highly dependent on the performance & limitation of PBE
 - We may need to combine PBEAR with other abstraction techniques
 - We plan to further investigate the applicability of PBEAR to:
 - Programs in other domains with control-oriented structures



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