

## Report Section: UT Network

### a) Pseudocode of DP Optimal Run Time

```

For  $t = 0$  to  $n-1$ :
     $R[t][1] = \text{ceil} ((x[t] - x[0]) / 2)$ 
For  $j = 2$  to  $k$ :
    For  $t = 0$  to  $n-1$ :
        Best = positive infinity
        For  $l = 0$  to  $t$ :
             $\text{Prev} = (l > 0) ? R[l-1][j-1] : 0$ 
             $\text{Block} = \text{ceil} ((x[t] - x[l]) / 2)$ 
             $\text{Best} = \min(\text{best}, \text{cand})$ 
         $R[t][j] = \text{best}$ 
    Return  $R[n-1][k]$ 

```

### b) Pseudocode Set of Status

```

For  $t = 0$  to  $n-1$ :
     $R[t][1] = \text{ceil} ((x[t] - x[0]) / 2)$ 
     $\text{CUT}[t][1] = 0$ 

For  $j = 2$  to  $k$ :
    For  $t = 0$  to  $n-1$ :
        Best = positive infinity
        For  $l = 0$  to  $t$ :
             $\text{Prev} = (l > 0) ? R[l-1][j-1] : 0$ 
             $\text{Block} = \text{ceil} ((x[t] - x[l]) / 2)$ 
             $\text{Cand} = \max(\text{prev}, \text{block})$ 
            If  $\text{cand} < \text{best}$  or ( $\text{cand} == \text{best} \ \&\& \ l > \text{best index}$ ):
                Best =  $\text{cand}$ 
                Best Index =  $l$ 
         $R[t][j] = \text{best}$ 
         $\text{CUT}[t][j] = \text{best index}$ 

     $T = n - 1$ 
     $J = k$ 
     $C = \text{empty list}$ 
    While  $j \geq 1$ :
         $i = \text{CUT}[T][j]$ 
         $c = (x[i] + x[T]) / 2$ 
        Append  $c$  to  $C$ 
         $T = l - 1$ 
         $J = j - 1$ 

    Reverse  $C$  to ascending order
    Return  $C$ 

```

### c) Runtime and Space

**DP has 3 nested for loops and the split**

- **Each inner step does  $O(1)$  work (ceil and max), so the total time is  $O(n^2 * k)$**
- **Store two  $n * k$  tables, so the space is  $O(nk)$**

**For example:**

- **For houses 2,4,6,8 with  $k = 2$**
- **The optimal response time is 2 achieved by stations 3 and 10**