

Compiler Project #2, 2022

The goal of the second term-project is to implement a bottom-up syntax analyzer (a.k.a., parser) as we've learned. More specifically, you will implement the syntax analyzer for a simplified Java programming language with the following context free grammar G ;

□ CFG G :

- 01: $\text{CODE} \rightarrow \text{VDECL CODE} \mid \text{FDECL CODE} \mid \text{CDECL CODE} \mid \epsilon$
- 02: $\text{VDECL} \rightarrow \text{vtype id semi} \mid \text{vtype ASSIGN semi}$
- 03: $\text{ASSIGN} \rightarrow \text{id assign RHS}$
- 04: $\text{RHS} \rightarrow \text{EXPR} \mid \text{literal} \mid \text{character} \mid \text{boolstr}$
- 05: $\text{EXPR} \rightarrow \text{EXPR addsub EXPR} \mid \text{EXPR multdiv EXPR}$
- 06: $\text{EXPR} \rightarrow \text{lparen EXPR rparen} \mid \text{id} \mid \text{num}$
- 07: $\text{FDECL} \rightarrow \text{vtype id lparen ARG rparen lbrace BLOCK RETURN rbrace}$
- 08: $\text{ARG} \rightarrow \text{vtype id MOREARGS} \mid \epsilon$
- 09: $\text{MOREARGS} \rightarrow \text{comma vtype id MOREARGS} \mid \epsilon$
- 10: $\text{BLOCK} \rightarrow \text{STMT BLOCK} \mid \epsilon$
- 11: $\text{STMT} \rightarrow \text{VDECL} \mid \text{ASSIGN semi}$
- 12: $\text{STMT} \rightarrow \text{if lparen COND rparen lbrace BLOCK rbrace ELSE}$
- 13: $\text{STMT} \rightarrow \text{while lparen COND rparen lbrace BLOCK rbrace}$
- 14: $\text{COND} \rightarrow \text{COND comp COND} \mid \text{boolstr}$
- 15: $\text{ELSE} \rightarrow \text{else lbrace BLOCK rbrace} \mid \epsilon$
- 16: $\text{RETURN} \rightarrow \text{return RHS semi}$
- 17: $\text{CDECL} \rightarrow \text{class id lbrace ODECL rbrace}$
- 18: $\text{ODECL} \rightarrow \text{VDECL ODECL} \mid \text{FDECL ODECL} \mid \epsilon$

□ Terminals (21):

1. **vtype** for the types of variables and functions
2. **num** for signed integers
3. **character** for a single character
4. **boolstr** for Boolean strings
5. **literal** for literal strings
6. **id** for the identifiers of variables and functions
7. **if**, **else**, **while**, and **return** for if, else, while and return statements, respectively
8. **class** for class declarations
9. **addsub** for + and - arithmetic operators

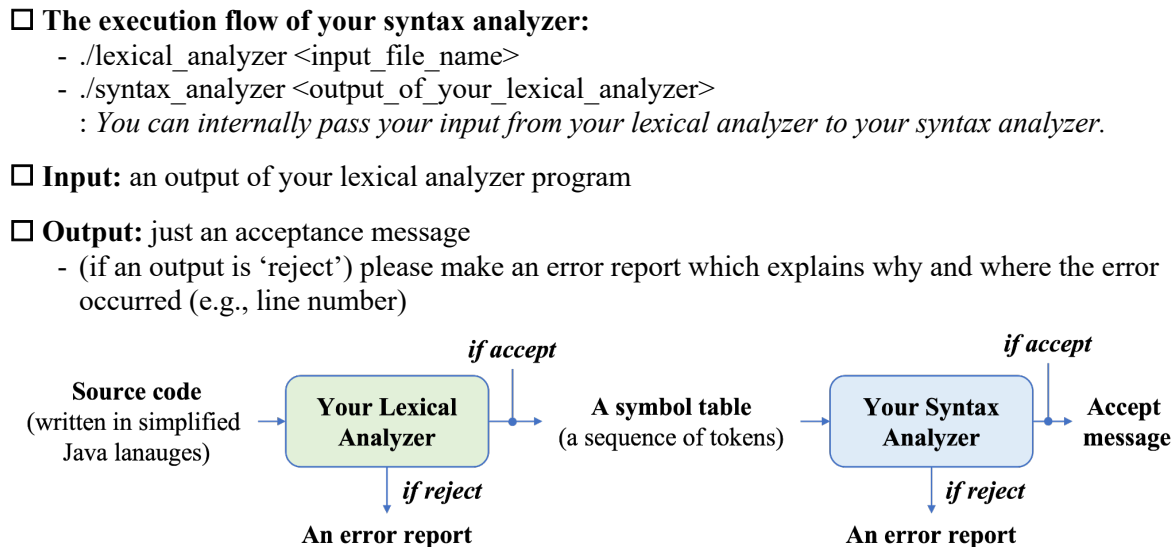
<p>10. multdiv for * and / arithmetic operators</p> <p>11. assign for assignment operators</p> <p>12. comp for comparison operators</p> <p>13. semi and comma for semicolons and commas, respectively</p> <p>14. lparen, rparen, lbrace, and rbrace for (,), {, and }, respectively</p>
<p>❑ Non-terminals (15)</p> <p>CODE, VDECL, ASSIGN, RHS, EXPR, FDECL, ARG, MOREARGS, BLOCK, STMT, COND, ELSE, RETURN, CDECL, ODECL</p>
<p>❑ Start symbol: CODE</p>
<p>❑ Descriptions</p> <ul style="list-style-type: none"> - The given CFG G is non-left recursive, but ambiguous - Codes include zero or more declarations of functions, variables, and classes (CFG 01) - Variables are declared with or without initialization (CFG 02, 03) - The right-hand side (RHS) of assignment operators can be classified into four types such as (1) arithmetic operators (expressions), (2) literal strings (CFG 13), (3) a single character, and (4) Boolean strings (CFG 04) - Arithmetic operations are the combinations of +, -, *, / operators (CFG 05,06,07) - Functions can have zero or more input arguments (CFG 08,09,10) - Function blocks include zero or more statements (CFG 11) - There are four types of statements: (1) variable declaration, (2) assignment operations, (3) if-else statements, and (4) while statements (CFG 12,13,14) - if and while statements include a conditional operation which consists of Boolean strings and an condition operator (CFG 13,14,15) - if statements can be used with or without an else statement (CFG 13 & 16) - return statements return (1) the computation result of arithmetic operations, (2) literal strings, (3) a single character, or (4) Boolean strings (CFG 17) - class is declared with zero or more declarations of functions and variables (CFG 18,19)

Based on this CFG, you should implement a bottom-up parser as follows:

- Discard an ambiguity in the CFG
- Construct a SLR parsing table for the non-ambiguous CFG through the following website:
<http://jsmachines.sourceforge.net/machines/slr.html>
- Implement a SLR parsing program for the simplified Java programming language by using the constructed table

For the implementation, you can use C, C++, JAVA, or Python as you want. However, your syntax analyzer must run on Linux or Unix-like OS without any error.

Your syntax analyzer should work as follows:



Term-project schedule and submission

- **Deadline:** 2022-06-21 (please use the e-Class)
- For a delayed submission, you will lose $0.1 \times$ your original project score per each delayed day
- **Submission file:** <your student ID>_<your_name>.tar.gz or .zip
- The compressed file should include
 - The source code of **your syntax and lexical analyzer** with detailed comments
 - The executable binary files of **your syntax analyzer + lexical analyzer**
 - Documentation (the most important thing!)
 - It must include (1) your non-ambiguous CFG G and (2) your SLR parsing table
 - It must also include any change in the CFG G and all about how your syntax analyzer works for validating token sequences (for example, overall procedures, implementation details like algorithms and data structures, working examples, and so on)
 - Test input files and outputs which you used in this project
 - The test input files are not given. You should make the test files, by yourself, which can examine all the syntax grammars.
- If there exist any error in the given CFG, please send an e-mail to kimjsung@cau.ac.kr