

Wk 2 Lecture Note



Abstract Data Types

- defined in terms of its data items & operations
- x on implementation (more like computational ϕ)
- supported by python (+ many other langs)

Q Abstract?

ex) Driving a car

Interface

- wheel 
- pedals 

implementation

- engine
- electronic chips

Q Why approach with ADT?

- easier code to understand (or experimenting ^{performance} trade-offs)
- When making choices of implementation, the cost is much reduced.
- code is reused (from standard library) \rightarrow faster building

ex) movie theatre reservation system

Q operations needed?

- capacity - available(c) ^{\swarrow all seats}
- capacity - sold(c) ^{\swarrow seats sold}
- customer(x): check who reserved seat x
- ~~add(x)~~ ~~reserve(x)~~ ^{\swarrow seats left}
- reserve(x, customer)
- release(x)
- get-capacity

Q Think about ~~at~~ edge cases!

- reserve(x, customer): if x is reserved already?
- release(x): x isn't reserved?

ADT = specification of the desired behavior (from user perspective)

input \rightarrow output

Data structure = concrete representation of data (perspective of implementor)

Comp Problem vs Algorithm

ADT vs Data Structure

implementation \leftrightarrow implementation

Q Index-based ADT

· size()

· isEmpty()

· get(i) $\checkmark O(1)$

· add(i, e) $\checkmark O(n)$: existing elements with idx $\geq i$ are shifted up.

· set(i, e) $\checkmark O(1)$ $\checkmark O(n)$

· remove(i) : remove & return the element at i

existing elements with idx $> i$ are shifted down.

Q how the operations would work?

if $ls = [A, B]$

get(1) = B

get(2) = "error"

add(2, C) = — $ls = [A, B, C]$

add(5, D) = "error"

set(3, D) = "error"

add(3, D) = — $ls = [A, B, C, D]$

memory

D-B-B-D

Q Positional Lists

· store elements at "positions"

· positions not altered even after addition/deletion

- size()

- isEmpty()

- first() : pos of 1st elem (null if empty)

- ~~last~~ after() : pos of last elem (null if empty)

- before(p) : pos of elem before p

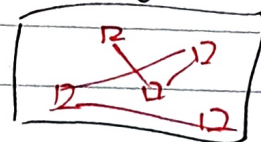
- ~~at~~ after(p) : pos of elem after p

- insertBefore(p, e) : insert e in front of p

- insertAfter(p, e) : insert e following the elem at p

- remove(p) : remove and return the removed elem

memory



ex) Singly Linked List

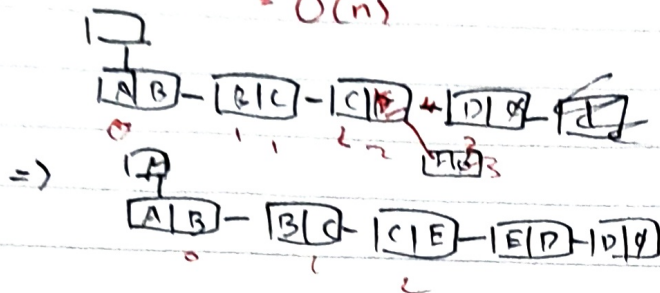
head



insertBefore()

[A, B, C, D] → [A, B, C, E, D]

$O(n)$



(3, E)

• Start at head and go to pos 3.

• Create a new Node with e: E and linking to p.

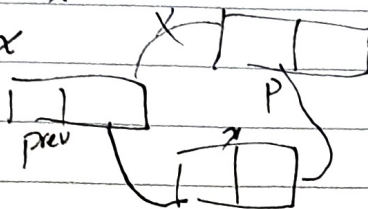
• Change prev node to link to the created Node and link the new node to where prev Node originally pointed to.

EX) Doubly Linked List

insertBefore(p, e) $O(1)$

p. prev. next ← x

p. next ← x
prev



Q. Iterators

in python: `__iter__ (self)`

`__next__` method

for x in ls:

// process x

`it = iter(ls)`

try:

while True:

x = it.next()

// process

except:

pass

Stacks & Queues

= Restricted forms of List

Stack: last-in-first out (LIFO)

queue: first in-first out (FIFO)

Stack ADT

↳ push(s): inserts an elem (puts plate on top of stack)

↳ pop(): remove + returns the top elem.

push(s) = [s]

push(3) = [s, 3]

pop() = 3 [s]

push(7) = [s, 7]

top() = 7

ex) browsing history (only backwards)
editing doc undo (no redos)

Queue ADT

↳ enqueue(s): inserts s at the end

↳ dequeue: removes the first elem

enqueue(s) = [s]

enqueue(3) = [s, 3]

dequeue: s [3]

ex) waiting in a line

buffering packets in stream
(video, audio)