# SSY191 Individual Assignment 2

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## 1 Problem 1

#### 1.1 a

According to the given C code structure and the Petri Net figure, there is a risk of being in deadlock.

This situation can happen if Task A and Task B try to acquire the resources in a different order. For example, if Task B acquires resource A first and Task A acquires resource B first, a deadlock will occur because each task is waiting for the other task to release the resource it needs:

- 1. Task A:  $p_0 \to (\text{take } r_A)$  p<br/>1  $\to (\text{take } r_B)$  p2  $\to$  p3  $\to (\text{give } r_A)$  p4 Task A waits  $r_A$  to continue,  $r_B$  is taken by Task A
- 2. Task B:  $p_0 \to (\text{take } r_A)$  p<br/>1 Task B waits  $r_B$  to continue,  $r_A$  is taken by Task B
- 3. Task A and B wait for each other, deadlock happens.

#### 1.2 b

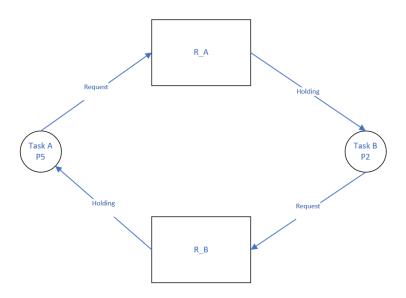


Figure 1: RAG graph

From figure 1, there is a loop. And from tuturioal, if there is a cycle in the Resource Allocation Graph and each resource in the cycle provides only one instance, then the processes will be in deadlock.

(Tuturioal: https://www.geeksforgeeks.org/resource-allocation-graph-rag-in-operating-system/)

### 1.3 c

From previous analysis, the key problem is to avoid the process  $p_0 \to p_1$  take away the  $r_A$  that process  $p_4 \to p_5$  need, then the deadlock can be sovled.

To achieve this, using the knowledge from discreate event system, I add a new resource to the Petri Net figure, that  $p_4$  requests it to happen the same time  $p_1$  requests it to happen.

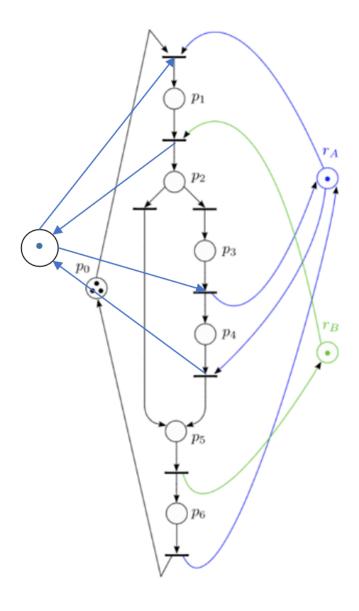


Figure 2: Modified Petri Net

The Modified code is listed below:

```
#include <stdio.h>
#include <stdlib.h>
#include "FreeRTOS.h"
```

```
#include "task.h"
5
       #include "semphr.h"
6
       xTaskHandle task_a_handle;
8
       xTaskHandle task_b_handle;
9
       xTaskHandle task_c_handle;
11
       SemaphoreHandle_t resource_a;
12
       SemaphoreHandle_t resource_b;
       SemaphoreHandle_t resource_c;
14
15
       void the_task(void *pvParameters)
17
           while (1)
18
            {
                xSemaphoreTake(resource_c, portMAX_DELAY);//Take
20
                   resouce c to go into p1
                xSemaphoreTake(resource_a, portMAX_DELAY);
21
                xSemaphoreTake(resource_b, portMAX_DELAY);
23
                xSemaphoreGive(resource_c, portMAX_DELAY);//Only
                   protect p1, after p1 give it out
25
                if (...)
26
                {
27
28
                    xSemaphoreTake(resource_c, portMAX_DELAY);//Take
29
                        resouce c to go into p4
                    xSemaphoreGive(resource_a);
30
31
                    xSemaphoreTake(resource_a, portMAX_DELAY);
32
                    xSemaphoreGive(resource_c, portMAX_DELAY);//
33
                        After p4 give it
                        out
                }
35
                xSemaphoreGive(resource_b);
36
37
                xSemaphoreGive(resource_a);
38
39
           }
40
```

```
}
41
42
       int main(int argc, char **argv)
43
       {
44
           resource_a = xSemaphoreCreateMutex();
45
           resource_b = xSemaphoreCreateMutex();
           resource_c = xSemaphoreCreateMutex();
47
           xTaskCreate(the_task, "Task 1",
48
              configMINIMAL_STACK_SIZE, NULL, 1, &task_a_handle);
           xTaskCreate(the_task, "Task 2",
49
              configMINIMAL_STACK_SIZE, NULL, 1, &task_b_handle);
           xTaskCreate(the_task, "Task 3",
50
              configMINIMAL_STACK_SIZE, NULL, 1, &task_c_handle);
51
           vTaskStartScheduler();
53
           for( ;; );
54
       }
55
```