Functional Programming

Lecture 14: Foldable and Traversable

Twan van Laarhoven

19 December 2022



Outline

- Foldable containers
- Traversable containers
- Ad hoc traversals
- Lenses

Foldable

Containers - again

What is a (finite) container?

- A container holds some number of 'values'
- It supports fmap (Functor)

Containers - again

What is a (finite) container?

- A container holds some number of 'values'
- It supports fmap (Functor)
- We can fold over the values in the container

Left and right folds for lists:

```
foldr :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
foldl :: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
```

Left and right folds for lists:

```
foldr :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
foldl :: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
```

For lists: foldr matches structure, foldl does not.

Left and right folds for lists:

```
foldr :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
foldl :: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
```

For lists: foldr matches structure, fold does not.

```
foldr (\diamondsuit) empty (1:(2:(3:[])) = (1 \Leftrightarrow (2 \Leftrightarrow (3 \Leftrightarrow empty))) foldl (\diamondsuit) empty (1:(2:(3:[])) = (((empty \Leftrightarrow 1) \Leftrightarrow 2) \Leftrightarrow 3)
```

Left and right folds for lists:

```
foldr :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
foldl :: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b
```

For lists: foldr matches structure, fold does not.

```
foldr (\diamondsuit) empty (1:(2:(3:[]))=(1 \diamondsuit (2 \diamondsuit (3 \diamondsuit empty))) foldl (\diamondsuit) empty (1:(2:(3:[]))=(((empty \diamondsuit 1) \diamondsuit 2) \diamondsuit 3)
```

When are these the same?

- <> is associative: a <>(b <> c) =(a <> b) <> c.
- empty is an identity element: z <>empty =empty =empty <>z

Associative operators with an identity elements form a monoid. In Haskell:

class Monoid m where

```
mempty :: m (\Leftrightarrow) :: m \to m \to m
```

Associative operators with an identity elements form a *monoid*. In Haskell:

```
class Monoid m where mempty :: m (\diamondsuit) :: m \to m \to m Actually class Semigroup m where (\diamondsuit) :: m \to m \to m class Semigroup m \Rightarrow Monoid m where mempty :: m
```

But we will pretend that it is a single class.

Associative operators with an identity elements form a *monoid*. In Haskell:

class Monoid m where

```
mempty :: m (\Leftrightarrow) :: m \to m \to m
```

With a monoid the way we fold does not matter, so there is just

```
fold :: (Monoid m) \Rightarrow [m] \rightarrow m
```

Associative operators with an identity elements form a *monoid*. In Haskell:

class Monoid m where

```
mempty :: m (\Leftrightarrow) :: m \to m \to m
```

With a monoid the way we fold does not matter, so there is just

```
fold :: (Monoid m) \Rightarrow [m] \rightarrow m
```

Often combined with a map

```
\mathsf{foldMap} \; :: \; (\mathsf{Monoid} \; \mathsf{m}) \; \Rightarrow \; (\mathsf{a} \; \rightarrow \; \mathsf{m}) \; \rightarrow \; [\mathsf{a}] \; \rightarrow \; \mathsf{m}
```

fold/foldMap for lists

Collapsing a list of values into a single value:

```
fold :: Monoid m \Rightarrow [m] \rightarrow m
fold [] = mempty
fold (x : xs) = x \Leftrightarrow fold xs
```

Is the same as

```
fold = foldr ( > ) mempty
```

Combining fold with map

```
foldMap :: Monoid m \Rightarrow (a \rightarrow m) \rightarrow [m] \rightarrow m foldMap f [] = mempty foldMap f (x : xs) = f x \Leftrightarrow foldMap f xs
```

```
class Monoid m where mempty :: m (\diamondsuit) :: m \to m \to m
```

foldMap examples

```
With the Sum monoid
  newtype Sum a = Sum { getSum :: a }
  instance Num a \Rightarrow Monoid (Sum a) where
    memptv = Sum 0
    Sum x \Leftrightarrow Sum y = Sum (x + y)
We can implement
  sum :: Num a \Rightarrow [a] \rightarrow a
  sum = getSum . foldMap Sum
and
  length :: [a] \rightarrow Int
  length = getSum . foldMap (\_ \rightarrow Sum 1)
```

Foldable containers

Generalize fold to any container type

```
class Foldable t where foldr :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow t \ a \rightarrow b foldl :: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow t \ a \rightarrow b fold :: Monoid m \Rightarrow t \ m \rightarrow m foldMap :: Monoid m \Rightarrow (a \rightarrow m) \rightarrow t \ a \rightarrow m
```

Foldable includes default implementations, defining either foldr or foldMap is sufficient.

Foldable containers

And is much easier to implement than foldr

```
data Tree a = Leaf a | Node (Tree a) (Tree a)
  instance Foldable Tree where
    foldMap :: Monoid m \Rightarrow (a \rightarrow m) \rightarrow Tree a \rightarrow m
    foldMap f (Leaf x) = f x
    foldMap f (Node | r) = foldMap f | coldMap f r
foldMap follows tree structure:
  foldMap f (Node (Node (Leaf 1) (Leaf 2)) (Node (Leaf 3) (Leaf 4)))
           = (f 1 \Leftrightarrow f 2) \Leftrightarrow (f 3 \Leftrightarrow f 4)
```

Other primitives

What primitives can we define for Foldable containers?

```
toList :: (Foldable t) \Rightarrow t a \rightarrow [a] null :: (Foldable t) \Rightarrow t a \rightarrow Bool length :: (Foldable t) \Rightarrow t a \rightarrow Int elem :: (Foldable t, Eq a) \Rightarrow a \rightarrow t a \rightarrow Bool maximum :: (Foldable t, Ord a) \Rightarrow t a \rightarrow a
```

Exercise: define these functions in terms of foldr or foldMap.

Defining foldMap using foldr and vice-versa

foldMap from foldr

```
foldMap :: Monoid m \Rightarrow (a \rightarrow m) \rightarrow t \ a \rightarrow m foldMap f = foldr (\a m \rightarrow f \ a \Leftrightarrow m) mempty
```

Defining foldMap using foldr and vice-versa

foldMap from foldr

```
foldMap :: Monoid m \Rightarrow (a \rightarrow m) \rightarrow t \ a \rightarrow m
foldMap f = foldr \ (\a m \rightarrow f \ a \colon m) mempty
foldr from foldMap
foldr :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow t \ a \rightarrow b
foldr f \in x = appEndo \ (foldMap \ (Endo . f) x) \ e
```

Endo explained

```
Functions in the same type (a \rightarrow a) are a monoid (endomorphisms)
  newtype Endo a = Endo \{ appEndo :: a \rightarrow a \}
  instance Monoid (Endo a) where
    mempty = Endo id
    Endo f \Leftrightarrow Endo g = Endo (f . g)
For example
    foldr f e [1,2.3]
  = (appEndo (foldMap (Endo . f) [1,2,3])) e
  = (appEndo (Endo (f 1) \Leftrightarrow Endo (f 2) \Leftrightarrow Endo (f 3) \Leftrightarrow Endo id)) e
  = (f 1 . f 2 . f 3 . id) e
  = f 1 (f 2 (f 3 e))
```

Traversable containers

Traversables

Mapping a function that may fail over a list

```
traverse :: (a \rightarrow Maybe \ b) \rightarrow [a] \rightarrow Maybe \ [b]
traverse f [] = pure []
traverse f (x : xs) = pure (:) <*> f x <*> traverse f xs
```

Generalize to

- any applicative f instead of Maybe (seen before as mapM).
- any container t instead of list

Gives

```
class (Functor t, Foldable t) \Rightarrow Traversable t where traverse :: Applicative f \Rightarrow (a \rightarrow f b) \rightarrow t a \rightarrow f (t b)
```

Traversable: examples

Instance for lists:

```
instance Traversable [] where
traverse :: Applicative f \Rightarrow (a \rightarrow f b) \rightarrow [a] \rightarrow f [b]
traverse g [] = pure []
traverse g (x : xs) = pure (:) <*> g x <*> traverse g xs
```

Instance for trees:

```
instance Traversable Tree where
```

```
traverse :: Applicative f \Rightarrow (a \rightarrow f b) \rightarrow Tree \ a \rightarrow f (Tree b)
traverse g (Leaf x) = pure Leaf <*> g x
traverse g (Node | r) = pure Node <*> traverse g | <*> traverse g r
```

Traverse generalizes mapM

Apply an action to every element of a list

```
askUser :: String \rightarrow IO Bool twoQuestions :: IO [Bool] twoQuestions = mapM askUser ["Do you like cake?", "Are you still here?"]
```

Traverse generalizes mapM

Apply an action to every element of a container

```
askUser :: String \rightarrow IO Bool twoQuestions :: IO [Bool] twoQuestions = traverse askUser ["Do you like cake?", "Are you still here?"] optQuestion :: Maybe String \rightarrow IO (Maybe Bool) optQuestion = traverse askUser
```

Traverse generalizes map

```
Recall: identity functor
  newtype Id a = Id { getId :: a }
  instance Functor Id where
    fmap f (Id x) = Id (f x)
  instance Applicative Id where
    pure = Id
    Id f < *> Id x = Id (f x)
Then
  amap :: Traversable f \Rightarrow (a \rightarrow b) \rightarrow f a \rightarrow f b
  amap f = getId . traverse (Id . f)
So Traversable containers are Functors
```

Traverse generalizes fold

```
The constant functor
```

```
newtype Const c a = Const { getConst :: c }
  instance Functor (Const c) where
    fmap (Const x) = Const x
  instance Monoid c \Rightarrow Applicative (Const c) where
    pure = Const mempty
    Const x \ll Const y = Const (x \Leftrightarrow y)
Gives
  afoldMap :: (Traversable f, Monoid m) \Rightarrow (a \rightarrow m) \rightarrow f a \rightarrow m
  afoldMap f = getConst . traverse (Const . f)
```

So Traversable containers are Foldable

Ad hoc traversals

Lenses



Ad hoc traversals

Traversable is about visiting all elements in a container

```
traverse :: (Traversable t, Applicative f) \Rightarrow (a \rightarrow f a) \rightarrow t a \rightarrow f (t a)
```

We can also traverse all literals in an expression

```
lits :: Applicative f \Rightarrow (Integer \rightarrow f \ Integer) \rightarrow Expr \rightarrow f \ Expr lits f (Lit \ x) = Lit < f \ x lits f (Add \ x \ y) = Add < lits <math>f \ x < lits \ f \ y lits f \ (Mul \ x \ y) = Mul < lits \ f \ x < lits \ f \ y
```

Or the first element of a pair

```
fst' :: Applicative f \Rightarrow (a \rightarrow f \ a) \rightarrow (a,b) \rightarrow f \ (a,b)
fst' f \ (x,y) = (\x' \rightarrow (x',y)) < f \ x
```

Using ad hoc traversals

```
Sum of all literals in an expression
  sumLits :: Expr \rightarrow Integer
  sumLits = getSum . getConst . lits (Const . Sum)
Count number of literals
  countLits :: Expr \rightarrow Int
  countLits = getSum . getConst . lits (\setminus \to Const (Sum 1))
  >>> sumLits (Add (Lit 1) (Mul (Lit 3) (Lit 3)))
  >>> countLits (Add (Lit 1) (Mul (Lit 3) (Lit 3)))
  3
```

Look familiar? Compare fold from traverse; sum and length from fold.

Using ad hoc traversals as folds

Let's give this a name

23

```
type Traversal s a = forall f. Applicative f \Rightarrow (a \rightarrow f a) \rightarrow (s \rightarrow f s)
  lits :: Traversal Expr Integer
  lits = -- as before
  foldMapOf :: Monoid m \Rightarrow Traversal s a \rightarrow (a \rightarrow m) \rightarrow (s \rightarrow m)
  foldMapOf trav f = getConst . trav (Const . f)
  lengthOf :: Traversal s a \rightarrow s \rightarrow Int
  lengthOf trav = getSum . foldMapOf trav (\rightarrow Sum 1)
Now
  length = lengthOf traverse
  countLits = lengthOf lits
```

Using ad hoc traversals as maps

Also

```
type Traversal s a = forall f. Applicative f \Rightarrow (a \rightarrow f \ a) \rightarrow (s \rightarrow f \ s) mapOf :: Traversal s a \rightarrow (a \rightarrow a) \rightarrow (s \rightarrow s) mapOf trav f = getId . trav (Id . f) — sometimes called over
```

Now

```
fmap = mapOf traverse
```

Example

```
>>> mapOf lits (*2) (Add (Lit 1) (Mul (Lit 3) (Lit 3))) (Add (Lit 2) (Mul (Lit 6) (Lit 6)))
```

More traversals

Traversing all items in a container

```
each :: Traversable t \Rightarrow Traversal (t a) a each = traverse
```

Traversing a single item in a list

```
at :: Applicative f \Rightarrow Int \rightarrow (a \rightarrow f \ a) \rightarrow [a] \rightarrow f \ [a]
at 0 f (x:xs) = (:xs) <$> f x
at n f (x:xs) = (x:) <$> at (n-1) f xs
at _ f _ = error "index out of range"
```

More traversals – examples

```
>>> foldMapOf (at 2) return "hello"
"l"
>>> mapOf each toUpper "hello"
"HELLO"
>>> mapOf (at 2) toUpper "hello"
"heLlo"
```

Odd and even

```
>>> mapOf evens toUpper "hello world"
"HeLlO WoRlD"
```

Where

```
— Traverse the even positions evens :: Applicative f \Rightarrow (a \rightarrow f \ a) \rightarrow [a] \rightarrow f \ [a] — Traversal [a] a evens \_[] = pure [] evens f(x:xs) = (:) < f(x) > f(x
```

Lenses

If the function is used exactly once, then Functor is enough

```
fst' :: Functor f \Rightarrow (a_1 \rightarrow f \ a_1) \rightarrow (a_1,b) \rightarrow f \ (a_2,b)
fst' f (x_1, y) = ((x_2 \rightarrow (x_2, y))) < $ f x_1
```

Such a traversal is called a *lens*.

type Lens s
$$a = forall f$$
. Functor $f \Rightarrow (a \rightarrow f a) \rightarrow (s \rightarrow f s)$

Compare

type Traversal s a = forall f. Applicative $f \Rightarrow (a \rightarrow f a) \rightarrow (s \rightarrow f s)$



Viewing through a lens

```
Const c is only an Applicative functor if c is a Monoid.
  foldOf :: Monoid m \Rightarrow Traversal s m \rightarrow s \rightarrow m
  foldOf = foldMapOf id
But Const c is always a functor even if c is not a Monoid, so
  view :: lens s a \rightarrow s \rightarrow a
  view lens = getConst. lens Const
Also
  modify :: Lens s a \rightarrow (a \rightarrow a) \rightarrow s \rightarrow s
  modify = mapOf
  set :: Lens s a \rightarrow a \rightarrow s \rightarrow s
  set lens x = modify (const x)
```

Lens examples

```
>>> view (at 2) [1,2,3,4]
3
>>> modify (at 2) negate [1,2,3,4]
[1,2,-3,4]
>>> view fst' (123,"hello")
123
>>> set fst' 456 (123,"hello")
(456,"hello")
```

Composing traversals

Visit the even elements, and for each of those visit the third element

```
evens . at 2 :: Traversal [[a]] a
So
  >>> foldMapOf (evens . at 2) return ["hello", "world", "where", "I", "live"]
  "lev"
  \implies mapOf (each . each) (+1) [Just 2, Nothing, Just 3]
  [Just 3, Nothing, Just 4]
Composing two lenses gives a lens
 >>> view (fst' . snd') (('a',3), "thing")
  3
```

Conditional traversal

Visit a value if a condition holds

```
when :: (a \rightarrow Bool) \rightarrow Traversal a a when pred f x | pred x = f x | otherwise = pure x | when (> 0)) succ [1,3,-4,2,0] [2,4,-4,3,0]
```

Traversing children

Traverse direct children of an expression

```
children :: Traversal Expr Expr
children f (Lit x) = pure (Lit x)
children f (Add x y) = Add <$> f x <*> f y
children f (Mul x y) = Mul <$> f x <*> f y

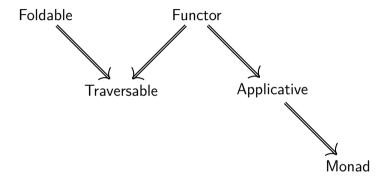
Children f (Mul x y) = Mul <$> f x <*> f y

Use as a default
flipMul :: Expr → Expr
flipMul (Mul x y) = Mul (flipMul y) (flipMul x)
flipMul expr = mapOf children flipMul expr
```

Flips arguments to all Mul constructors

Take away

The container-like class hierarchy



Summary

- Foldable generalizes folds to other containers
- Traversable generalizes mapM.
- Ad-hoc traversals can view and modify parts of a data structure