k-Complete Test Suite

What if k > 0?

- We should detect up to *k* extra states.
- $A \cdot I^{\leq k}$ reaches all implementation states!
- replace A in the 0-complete test suite by $A \cdot I^{\leq k}$

An *k*-complete test suite:

$$(A \cdot I^{\leq k}) \cdot I^{\leq 1} \cdot C$$

or simply

$$\textbf{T} = \textbf{A} \cdot \textbf{I}^{\leq k+1} \cdot \textbf{C}$$

Frits Vaandrager

Black Box Testing of Finite State Machines

4□ > 4回 > 4 = > 4 = > = 9 < 0</p>

Finite State Machine k-Complete Test Suites Characterization Sets Test Suites Without Resets

Special (Smaller) Characterisation Sets

- A sequence $c \in C$ is a Unique Input Output sequence (UIO) for some state q if:
 - for all other states q' of S: $\lambda^*(q,c) \neq \lambda^*(q',c)$
- ullet Hence, a characterisation set of UIOs needs only |S|-1 elements.
- A sequence $c \in C$ is a Distinguishing Sequence (DS) for S if:
 - For all states q, q' (with $q \neq q'$) of $S: \lambda^*(q, c) \neq \lambda^*(q', c)$
- Hence, a DS gives a singleton characterization set!
- Note:
 - A distinguishing sequence is for an entire specification
 - UIOs are per state
 - Separating sequences are per pair of states
- UIOs and DSs do not always exist...

Frits Vaandrager

Black Box Testing of Finite State Machines

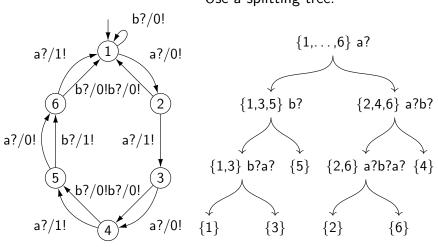
k-Complete Test Suites

Characterization Sets

Test Suites Without Resets

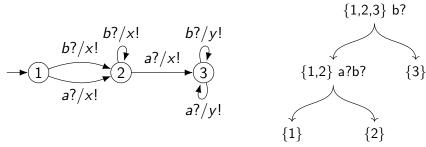
Algorithm for Finding Separating Sequences

Use a splitting tree:



 $C = \{a?, b?, a?b?, b?a?, a?b?a?\}$

Splitting node: Separate States by Input



 $\{1,2\}$ can be split based on a? and the split of $\{1,2,3\}$, because

- $\delta(1, a?) = 2$ and $\delta(2, a?) = 3$, and
- states 2 and 3 are already split in node $\{1,2,3\}$ (they are in different children of $\{1,2,3\}$)

$$C = \{b?, a?b?\}$$