

# Functional Programming

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Lists

**Lecture 3**

# Outline

- List notation
- Compositional programming
- Case study: DNA analysis
- List constructors
- List design pattern
- Some list operations
- List comprehensions



$(\textcolor{blue}{:}) :: a \rightarrow \textcolor{red}{[a]} \rightarrow \textcolor{red}{[a]}$

# List notation

- lists are central to functional programming (cf LISP!)
- enclosed in square brackets, comma-separated: `[1,2,3]`, `[ ]`
- variable-length sequences of elements of the same type
- the type of lists with elements of type `T` is `[T]`
- strings are just lists of characters: `['T','y','r','i','o','n']`

**type** `String` = `[Char]`

but with special syntax `"Tyrion"`

- list elements can be any type:

`[1,2,3] :: [Integer]`

`[[1,2],[ ],[3]] :: [[Integer]]`

`[(+),(*)] :: [Integer → Integer → Integer]`

`[(/2),(+ (-2))]` `:: [Integer → Integer]`

# Some library functions

- exploring the library `Data.List`

```
import Data.List
```

- see also <https://www.haskell.org/hoogle/>

```
length :: [a] → Int e.g.
```

```
length [1,2,3] = 3
```

```
reverse :: [a] → [a] e.g.
```

```
reverse "cersei" = "iesrec"
```

```
concat :: [[a]] → [a] e.g.
```

```
concat [[1,2],[ ],[3]] = [1,2,3]
```

```
map :: (a → b) → [a] → [b] e.g.
```

```
map (+ 1) [1,2,3] = [2,3,4]
```

```
filter :: (a → Bool) → [a] → [a] e.g.
```

```
filter (< 6) [2,7,6,5] = [2,5]
```

**remark:** some operations actually have more generic types, but here we pretend that you can only use them on lists.

# More library functions

- `lines :: String → [String]` e.g.  
`lines "a\nbc\nd" = ["a", "bc", "d"]`
- `unlines :: [String] → String` e.g.  
`unlines ["a", "bc", "d"] = "a\nbc\nd\n"`
- `tails :: [a] → [[a]]` e.g.  
`tails "arya" = ["arya", "rya", "ya", "a", ""]`
- `take :: Int → [a] → [a]` e.g.  
`take 3 "ramsay bolton" = "ram"`
- `sort :: (Ord a) ⇒ [a] → [a]` e.g.  
`sort "daenerys" = "adeenrsy"`
- `group :: (Eq a) ⇒ [a] → [[a]]` e.g.  
`group "joffrey" = ["j", "o", "ff", "r", "e", "y"]`
- and many more . . .

# How to solve it?

- write down the type (what's the input?, what's the output?)
- can you solve it using existing vocabulary?
- use function application
- some exercises: given a string (a list of characters)
  - remove newlines
  - count the number of lines
  - flip text upside down
  - flip text from left to right
  - determine the list of all segments

```
lines    :: String → [String]
unlines  :: [String] → String
tails    :: [a] → [[a]]
take     :: Int → [a] → [a]
sort     :: (Ord a) ⇒ [a] → [a]
group    :: (Eq a) ⇒ [a] → [[a]]

reverse  :: [a] → [a]
concat   :: [[a]] → [a]
map      :: (a → b) → ([a] → [b])
```

# Solutions

- remove newlines

```
unwrap :: String → String
```

```
unwrap s = concat (lines s)
```

- count the number of lines

```
countLines :: String → Int
```

```
countLines s = length (lines s)
```

- flip text upside down

```
upsideDown :: String → String
```

```
upsideDown s = unlines (reverse (lines s))
```

- flip text from left to right

```
leftRight :: String → String
```

```
leftRight s = unlines (map reverse (lines s))
```



# Solutions continued

- determine the list of all prefixes (actually, also defined in the library: `inits`)

```
suffixes, prefixes :: String → [String]
```

```
suffixes s = tails s
```

```
prefixes s = map reverse (tails (reverse s))
```

- determine the list of all segments

```
segments :: String → [String]
```

```
segments s = concat (map prefixes (suffixes s))
```

# Case study: DNA analysis

You are working on the evolutionary history of rodents. To this end you analyse the genome of various species, considering genome size, proportions of non-repetitive DNA, and repetitive DNA.

In particular, you are interested in finding repeated segments of length  $m$  in a DNA sequence of length  $n$ .

```
ATGTAAAGGGTCCAATGACTGGAATTACTTCACAGCCCTGACACTGTGGAGAGATGGATA
CCAGTTTAATTATGGACAGCTGAGAATAATCTGGGAGTCACCTGCTCTGGATATTTATCA
TCATTATGAACTTCTCACCAGGCTGTGGCCCGAACACTTTCATAACGTCCCATTTGTGTT
GGGCAGACATTATGATCTATACAGAACCTGCCTTGTCTGTCACCTTTGAATTTATGGATA
GACAATTTAATAT
GTGTTCTCTGGCAGCAAAGATAATCATGGAGAGTGGAGAGAACTAACCTTACCATTGATA
GGGAACTCTTGAAGTGTCAACTTCTCCATATTAAATCCAAGGACCAGCAGAGACGAGAA
AATGAAAAGAAGATGGTAGAAGAAAGAACAAAATCAGAAAAAGACAAAGGAAAAGGGAAG
TCTCCAAAGGAGAAGAAAGTTGCCAGTGCCAAGCCTGGGAAGGGAAAGAGCAAGGACCAG
...
```

# Think big

- representation of DNA data

```
data Base = A | C | G | T
    deriving (Eq,Ord,Show)
```

```
type DNA = [Base]
```

```
dna :: DNA
```

```
dna = [A,T,G,T,A,A,A,G,G,G,T,C,C,A,A,T,G,A]
```

- think in terms of entire lists, not individual list elements
- think in terms of transformations
- algorithmic idea:
  - step 1: generate all  $m$ -segments
  - step 2: identify repeated  $m$ -segments

# Step 1: generate all m-segments (m = 3)

ATGTAAAGGGTCCAATGA

↓ tails

[ATGTAAAGGGTCCAATGA, TGTAAAGGGTCCAATGA,  
...  
CCAATGA, CAATGA, AATGA, ATGA, TGA, GA, A, ]

↓ map (take 3)

[ATG, TGT, GTA, TAA, AAA, AAG, AGG, GGG, GGT, GTC,  
TCC, CCA, CAA, AAT, ATG, TGA, GA, A, ]

## Step 2: identify repeated m-segments

```
[ATG,TGT,GTA,TAA,AAA,AAG,AGG,GGG,GGT,GTC,  
TCC,CCA,CAA,AAT,ATG,TGA,GA,A,]
```

↓ sort

```
[,A,AAA,AAG,AAT,AGG,ATG,ATG,CAA,CCA,GA,  
GGG,GGT,GTA,GTC,TAA,TCC,TGA,TGT]
```

↓ group

```
[[ ],[A],[AAA],[AAG],[AAT],[AGG],[ATG,ATG],[CAA],  
[CCA],[GA],[GGG],[GGT],[GTA],[GTC],[TAA],  
[TCC],[TGA],[TGT]]
```

↓ filter (\s -> length s > 1)

```
[[ATG,ATG]]
```

# Solution

- chain transformations

`repeatedSegments :: Int → DNA → [[DNA]]`

`repeatedSegments m dna`

`= filter (\x → length x > 1) (group (sort (map (take m) (tails dna))))`

- using composition (more in week 5)

$$(g \circ f) \ x = g \ (f \ x)$$

`repeatedSegments :: Int → DNA → [[DNA]]`

`repeatedSegments m`

`= filter (\x → length x > 1) ∘ group ∘ sort ∘ map (take m) ∘ tails`

- transformational programming at work

# List constructors

- a list is either
  - empty, written `[]`
  - or consists of an element `x` followed by a list `xs`, written `x : xs`
- every (finite) list can be built up from `[]` using `:`
- e.g. `[1,2,3] = 1 : (2 : (3 : [])) = 1 : 2 : 3 : []`
- `[]` and `:` are called *constructors*

# Type of list constructors

- nil: the empty list

$[] :: [a]$

- cons: function for prefixing an element onto a list

$(:) :: a \rightarrow [a] \rightarrow [a]$

- $[]$  and  $:$  are polymorphic!
- puzzle: is  $[] : []$  well-typed? what's the type? what about  $[] : ([] : [])$  and  $([] : []) : []$ ?



# Pattern matching

- to define a function over lists, it suffices to consider the two cases `[]` and `:`

- e.g. to test if list is empty

```
null :: [a] → Bool
```

```
null [] = True
```

```
null (_:_) = False
```

- (subtle: why is this different from `(==[])`?)
- e.g. to return first element of non-empty list

```
head :: [a] → a
```

```
head (x:_) = x
```

# Case analysis

- cases can also be analysed using a *case-expression*

```
null :: [a] → Bool
```

```
null xs = case xs of
```

```
    []      → True
```

```
    (_:_)   → False
```

- declaration style: equations using patterns; expression style: case-expression using patterns

# Recursive definitions

- definitions by pattern-matching can be recursive too
- natural thing to do as the type is also recursively defined
- eg sum of a list of integers

`sum :: [Integer] → Integer`

`sum [] = 0`

`sum (x:xs) = x + sum xs`

- eg length of a list of elements

`length :: [a] → Int`

`length [] = 0`

`length (_:xs) = 1 + length xs`

# Computation = reduction

$[2,3,4] = 2 : 3 : 4 : []$

$\text{sum} :: [\text{Integer}] \rightarrow \text{Integer}$

$\text{sum} [] = 0$

$\text{sum} (x:xs) = x + \text{sum} xs$

$\text{sum} [2,3,4]$

$= 2 + \text{sum} [3,4]$

$= 2 + (3 + \text{sum} [4])$

$= 2 + (3 + (4 + \text{sum} []))$

$= 2 + (3 + (4 + 0))$

$= 2 + (3 + 4)$

$= \underline{2 + 7}$

$= 9$

# List design pattern

- remember: every type comes with a pattern of definition
- task: define a function  $f :: [P] \rightarrow S$
- step 1: solve the problem for the empty list

$f [] = \dots$

- step 2: solve the problem for non-empty lists;
- assume that you already have the solution for  $xs$  at hand; extend the intermediate solution to a solution for  $x:xs$

$f [] = \dots$

$f (x:xs) = \dots x \dots xs \dots f xs \dots$

- you have to program only a step
- put on your problem-solving glasses

# Some list operations

- append:  $[1,2,3] \mathbin{++} [4,5] = [1,2,3,4,5]$   
     $(++) :: [a] \rightarrow [a] \rightarrow [a]$   
     $[] \mathbin{++} ys = ys$   
     $(x:xs) \mathbin{++} ys = x:(xs \mathbin{++} ys)$
- concatenation:  $\text{concat } [[1,2],[],[3]] = [1,2,3]$   
     $\text{concat} :: [[a]] \rightarrow [a]$   
     $\text{concat } [] = []$   
     $\text{concat } (xs:xss) = xs \mathbin{++} \text{concat } xss$
- reverse:  $\text{reverse } [1,2,3] = [3,2,1]$   
     $\text{reverse} :: [a] \rightarrow [a]$   
     $\text{reverse } [] = []$   
     $\text{reverse } (x:xs) = \text{reverse } xs \mathbin{++} [x]$

solution for xs



# Computation = reduction

$[1,2,3] = 1 : 2 : 3 : []$

$(++) :: [a] \rightarrow [a] \rightarrow [a]$

$[] ++ ys = ys$

$(x:xs) ++ ys = x:(xs ++ ys)$

$[1,2,3] ++ [4,5]$

$= 1:([2,3] ++ [4,5])$

$= 1:(2:([3] ++ [4,5]))$

$= 1:(2:(3:([ ] ++ [4,5])))$

$= 1:(2:(3:([4,5])))$

$\equiv [1,2,3,4,5]$

# More list operations

- is a list ordered?

`ordered :: (Ord a) => [a] -> Bool`

`ordered [] = True`

`ordered [_] = True`

`ordered (x1:x2:xs) = x1 <= x2 && ordered (x2:xs)`

we distinguish three cases

- zip: eg `zip [1,2,3] "abc" = [(1,'a'),(2,'b'),(3,'c')]`

`zip :: [a] -> [b] -> [(a,b)]`

`zip [] [] = []`

`zip [] (_:_) = []`

`zip (_:_) [] = []`

`zip (x:xs) (y:ys) = (x,y):zip xs ys`

we pattern match on both arguments



# List comprehensions

- two useful operators on lists: `map` and `filter`
- list comprehensions provide a convenient syntax for expressions involving `map`, `filter`, and `concat`
- useful for constructing new lists from existing lists

# Map

- applies given function to every element of given list
- eg `map square [1,2,3] = [1,4,9]`
- eg `map succ "HAL" = "IBM"`
- definition

`map :: (a → b) → [a] → [b]`

`map f [] = []`

`map f (x:xs) = f x : map f xs`

- eg `sum (map square [1..10])`

# Computation = reduction

`even :: Int → Bool`

```
map :: (a → b) → [a] → [b]
map f [] = []
map f (x:xs) = f x : map f xs
```

```
map even [1,2,3,4]
= even 1 : map even [2,3,4]
= False : map even [2,3,4]
= False : even 2 : map even [3,4]
= False : True : map even [3,4]
= False : True : even 3 : map even [4]
= False : True : False : map even [4]
= False : True : False : even 4 : map even []
= False : True : False : True : map even []
= False : True : False : True : []
= [False, True, False, True]
```

# Filter

- returns sub-list of the argument whose elements satisfy given predicate
- eg `filter isDigit "more4u2say" = "42"`
- eg `(sum ◦ map square ◦ filter odd) [1..5] = 35`
- definition

```
filter :: (a → Bool) → [a] → [a]
filter _p []      = []
filter p (x:xs)
  | p x           = x:filter p xs
  | otherwise     = filter p xs
```

# List comprehensions

- special convenient syntax for list-generating expressions

- Map: alter every element **x** from **xs** (by applying **f**)

`[ f x | x ← xs ]`

- Filter: select element **x** from **xs** with predicate **pred**

`[ f x | x ← xs, pred x ]`

- Cartesian product: combine each element **x** from **xs** with each element **y** from **ys** (nested loops)

`[ g x y | x ← xs, y ← ys ]`

- Example: `sum [square x | x ← [1..5], odd x]`

# Comprehensions

- formally, a comprehension  $[e \mid Qs]$  for expression  $e$  and non-empty comma-separated sequence of qualifiers  $Qs$
- qualifier may be
  - *generator* of the form  $x \leftarrow xs$  or
  - *guard* i.e. a Boolean expression

# Examples of comprehensions

- eg primes up to a given bound

```
primes, divisors :: Integer → [Integer]
primes m      = [n | n ← [1..m], divisors n == [1,n]]
divisors n    = [d | d ← [1..n], n `mod` d == 0]
```

- eg database query

```
overdue =
  [(name,addr) | (key,name,addr) ← customers,
                 (key',date) ← invoices, key == key',
                 date < today]
```

- eg Quicksort

```
quicksort :: (Ord a) ⇒ [a] → [a]
quicksort [] = []
quicksort (x:xs) = quicksort [a | a ← xs, a < x]
                  ++ [x] ++
                  quicksort [a | a ← xs, x ≤ a]
```

# Another point of view

- list comprehension is 'really' a form of nested loop
- eg  $[f\ b \mid a \leftarrow x, b \leftarrow g\ a, p\ b]$  is related to

```
List<Integer> res = new LinkedList<>();
for ( int a : x ) {
    for ( int b: g(a) ) {
        if ( p(b) ) {
            res.add(f(b));
        }
    }
}
return res;
```



# Zip

- traversing two lists **xs**, **ys** simultaneously: combine each element **x** from **xs** pairwise with element **y** from **ys**
- example: define a function **odds** that returns all elements of the input lists that are located at odd positions.
- eg:
  - `odds [1,2,3,4,5] = [1,3,5]`
  - `odds "Stannis Baratheon" = "SansBrten"`
- Only basic/library functions and/or list comprehensions,
- Solution: `odds l = [ x | (i,x) ← zip [1..] l, odd i ]`
- In general: `[ ...x...y... | (x,y) ← zip xs ys ]`

patterns are allowed

# Summary: How to solve it?

- write down the type (what's the input?, what's the output?)
- can you solve the problem using existing vocabulary?
- if not, define new vocabulary
- use the list design pattern
  - remember: you only have to solve a step
- can you solve the step using existing vocabulary?
- if not, define new vocabulary (identify a sub-problem)
- solve the sub-problem in the same manner

