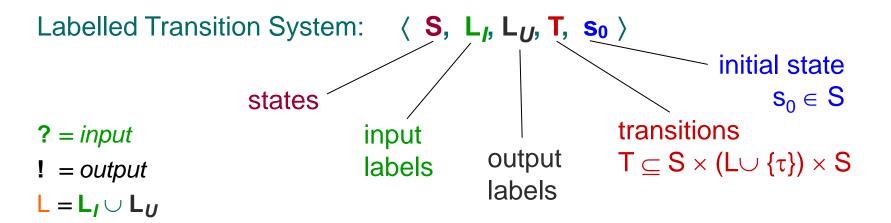
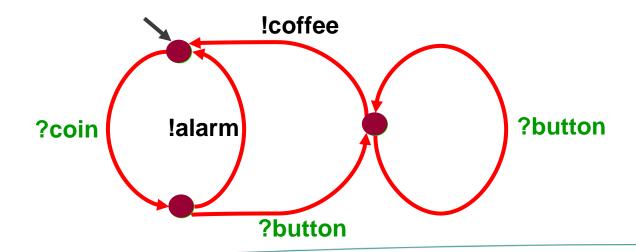
# Equivalences on Labelled Transition Systems

#### **Models: Labelled Transition Systems**

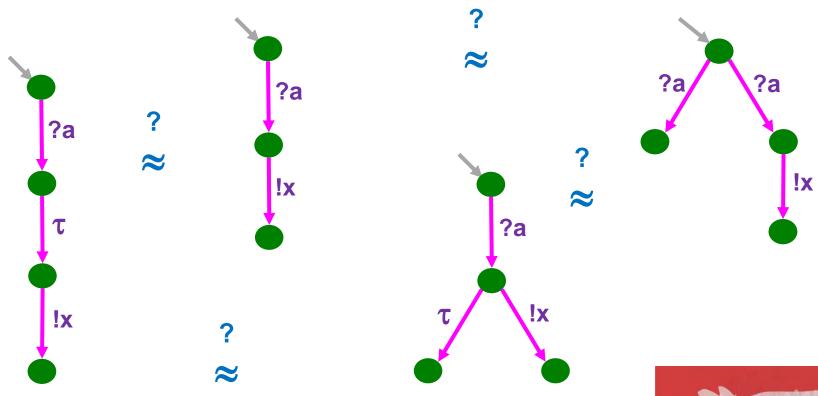




 $L_{I} \cap L_{U} = \emptyset$ 



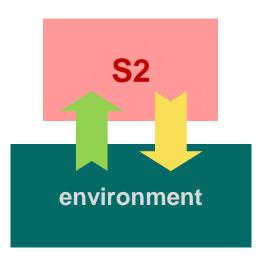
#### **Observable Behaviour**



"Some transition systems are more equal than others"







- Suppose an environment interacts with the systems:
  - the environment tests the system as black box by observing and actively controlling it;
  - the environment acts as a tester;
- Two systems are equivalent iff they pass the same tests.





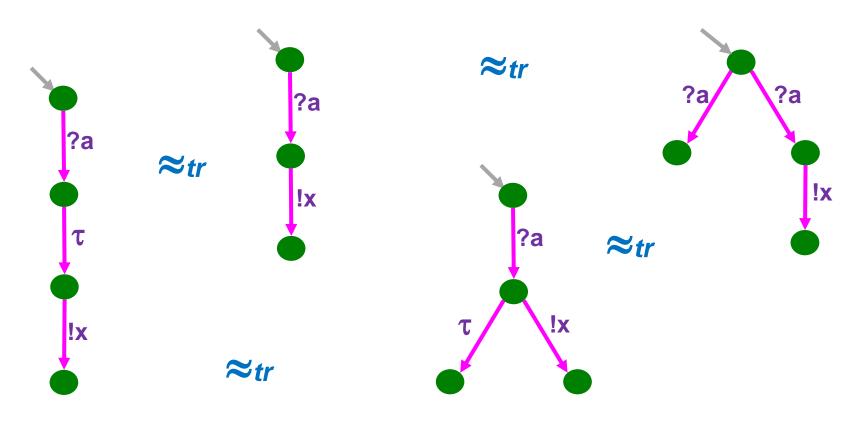




$$S1 \approx_{tr} S2 \Leftrightarrow traces(S1) = traces(S2)$$

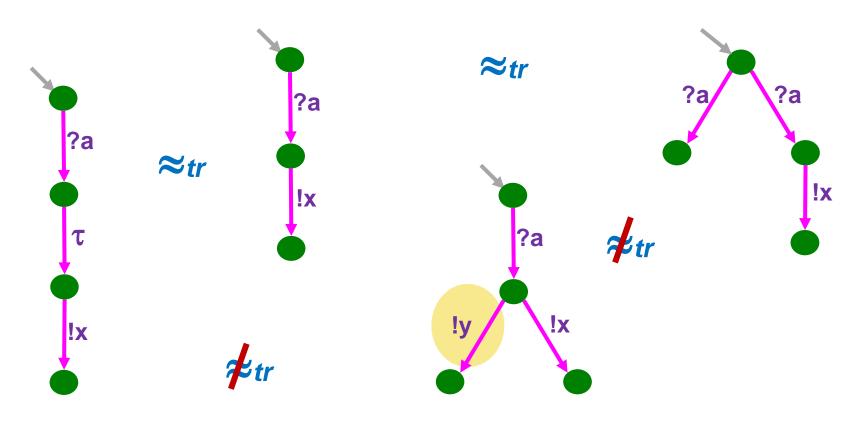
Traces: 
$$traces(s) = \{ \sigma \in L^* \mid s \Longrightarrow \}$$

#### **Trace Equivalence**



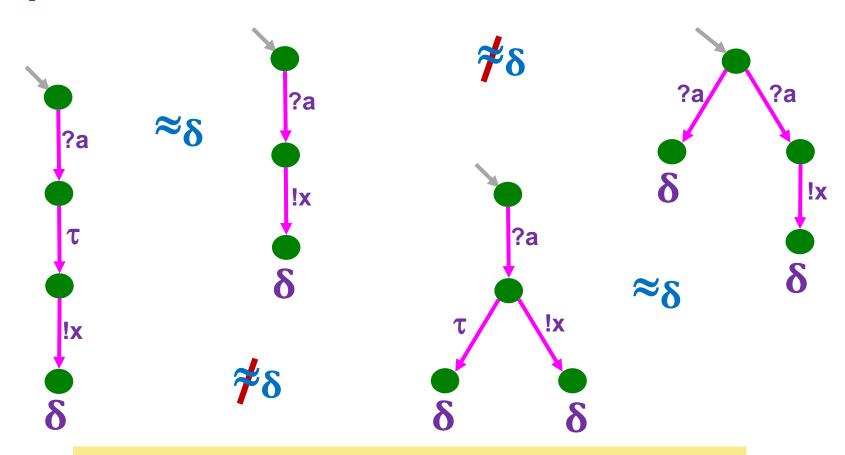
for all:  $traces(.) = \{ \epsilon, ?a, ?a.!x \}$ 

## **Trace Equivalence**



?a.  $!y \in traces(S3)$ 

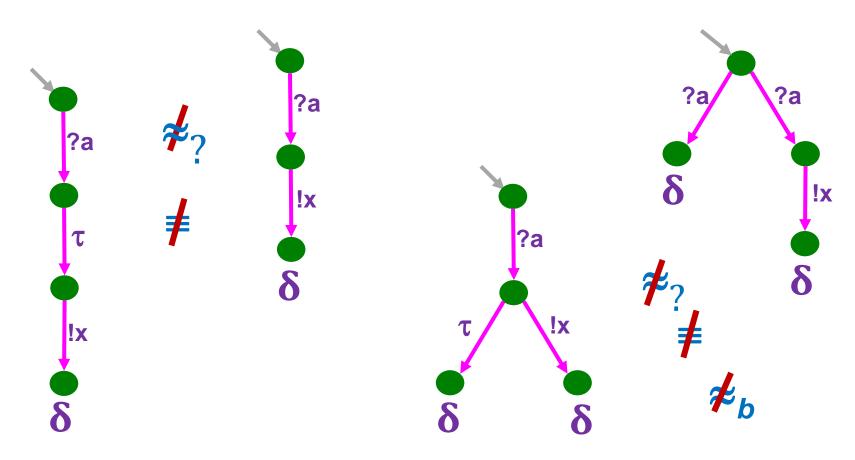
#### Equivalence with Quiescence δ



$$p \xrightarrow{\delta} p = \forall !x \in L_U \cup \{\tau\}. p \xrightarrow{!x}$$

$$Straces(s) = \{ \sigma \in (L \cup \{\delta\})^* \mid s \xrightarrow{\sigma} \}$$

### **Stronger Equivalences**



≡ isomorphism≈ bisimulation





#### **MBT**: Equivalences

1. equivalent if they have the same behaviours

```
S1 \approx S2 \Leftrightarrow traces(S1) = traces(S2)
```

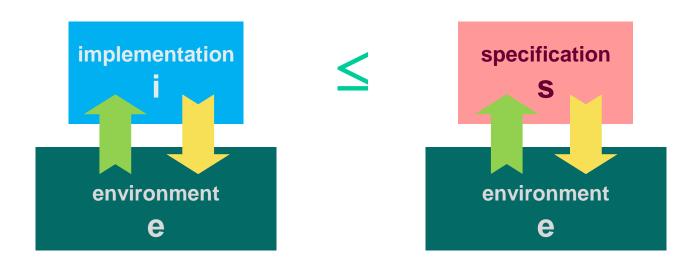
$$S1 \approx S2 \Leftrightarrow Straces(S1) = Straces(S2)$$

2. equivalent if they pass the same tests

$$S1 \approx S2 \Leftrightarrow \forall t \in T. obs(t, S1) = obs(t, S2)$$

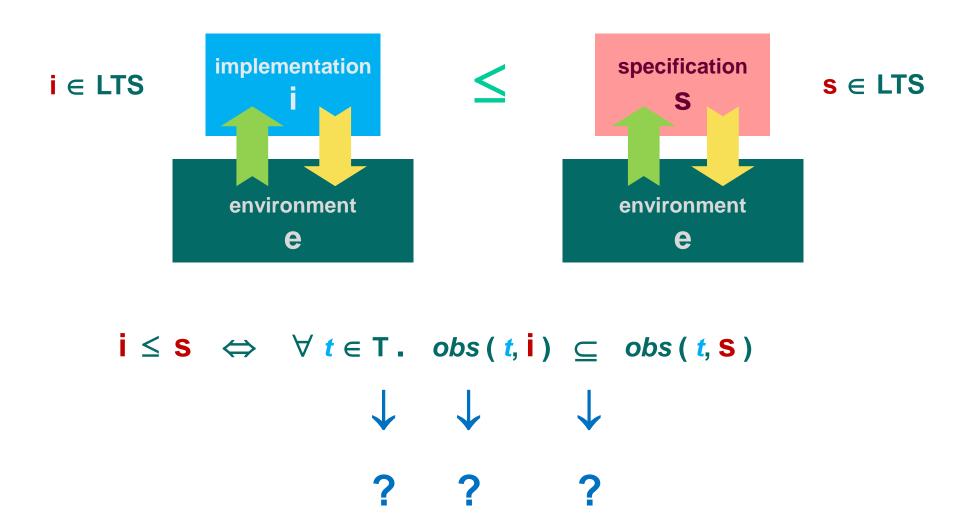
3. equivalent if they have the same implementations

$$S1 \approx S2 \Leftrightarrow Imp(S1) = Imp(S2)$$

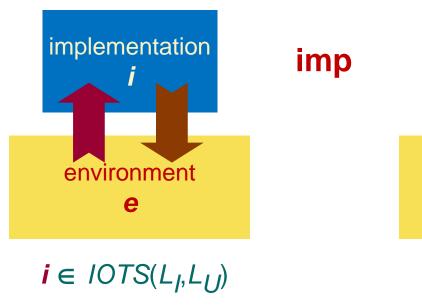


- Suppose environment e interacts with implementation i and with the specification s :
  - i correctly implements
     iff all observations of i can be related to observations of s

#### Implementation Relations on LTS



# **Implementation Relations for Input-Output Transition Systems**





$$s \in LTS(L_I, L_U)$$

imp 
$$\subseteq IOTS(L_{I},L_{U}) \times LTS(L_{I},L_{U})$$

i imp s

i is a conforming implementation of s

#### Implementation Relations on LTS

i imp s: implementation i implements specification s

```
imp reflexive? simps
imp symmetric? i imps ⇒ simpi
imp transitive? i imps, simpt ⇒ i impt
imp anti-symmetric? i imps, simpi ⇒ i = s
imp linear? i imps or simpi
```

```
equivalence : reflexive, symmetric, transitive preorder : reflexive, transitive anti-symmetric preorder linear/total order : linear partial order
```

$$Imp(s) = \{ i \mid i \text{ imp } s \}$$

#### MBT: our choices

```
s \in LTS(L_{I}, L_{U})
i \in IOTS(L_{I}, L_{U})
i \text{ imp } s \Leftrightarrow i \text{ iuioco } s
S1 \approx S2 \Leftrightarrow Ftraces(S1) = Ftraces(S2)
Ftraces: \text{ input-failure traces over } L_{I} \cup L_{U} \cup \{\delta\}
```

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#### **Test Cases**

#### model of a test case

- = labelled transition system
- labels in  $L_I \cup L_U \cup \{\theta\}$
- 'quiescence' / 'time-out' label  $\theta$
- tree-structured
- finite, deterministic
- sink states pass and fail
- from each state ≠ pass, fail :
  - either one input !a and all outputs ?x
  - or all outputs  $\mathbf{?x}$  and  $\mathbf{\theta}$



