

Counterexample Handling

f produces a 1 if number of a 's and b 's are both even, otherwise 0.

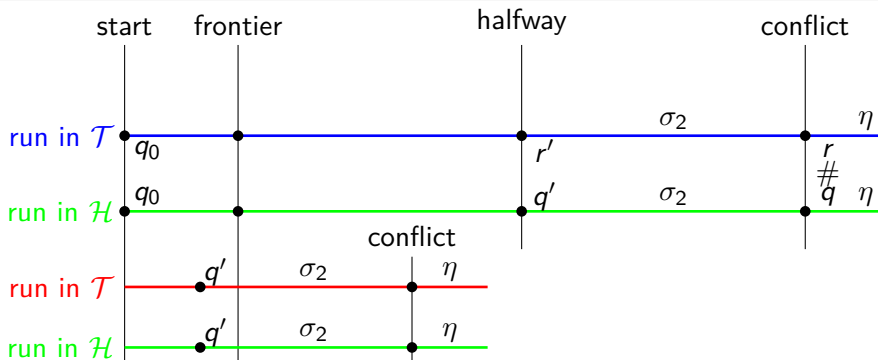
Two a -transition to states which differ on suffix a : add suffix aa

	a	b	aa	"colour"
$\rightarrow \epsilon$	0	0		
$\rightarrow a$	1	0		a/0 ↓
$\rightarrow b$	0	1		
$\rightarrow ab$	0	0		
$\rightarrow aba$	0	1		a/0 ↓
aa	0	0		
ba	0	0		
bb	0	0		

Angluin's L^* Algorithm

- 1 Maintain a set \mathcal{U} of (marked) prefixes, initially $\mathcal{U} = \{\epsilon\}$
- 2 Maintain a set \mathcal{V} of suffixes, initially $\mathcal{V} = I$
- 3 Maintain an observation table with rows $\mathcal{U} \cup \mathcal{U}I$ and columns \mathcal{V}
- 4 Fill the table using output queries
- 5 Table is **closed** when every row from $\mathcal{U}I$ is also a row from \mathcal{U} ;
if table is not closed extend \mathcal{U} and go to step 4
- 6 Table is **consistent** if whenever rows $u, v \in \mathcal{U}$ are the same,
rows ui and vi are also the same, for all $i \in I$;
if table is not consistent extend \mathcal{V} and go to step 4
- 7 When table is both **closed** and **consistent** construct hypothesis
and perform equivalence query
- 8 If reply is “no” add all prefixes of counterexample to \mathcal{U} and go
to step 4

Counterexample Processing



Key idea: perform output query access $(q')\sigma_2\eta$

If outputs for $\sigma_2\eta$ from r' and q' are different in \mathcal{T} then $r' \# q'$

Else access $(q')\sigma_2$ leads to a conflict!