

15.053 Tutorial 04

LP Transformation Tricks

Table of Contents

- Introduction & definitions slide 3
- Example 1: Absolute value constraints slides 4-6
- Example 2: Maximin and minimax problems slides 7-8
- Example 3: Maximin objective function slides 9-10
- Example 4: Ratio constraints slide 11
- Summary slide 12

Introduction and Definitions

- The tutorial will show three different types of non-linear constraints that can be transformed into linear constraints:
 - **Absolute value constraints**, in which the absolute value of a linear combination of decision variables is constrained
 - **Maximizing minimum constraints** (and vice versa), in which we find the largest value of a “min” function (taking the minimum value from a set), or find the smallest value of a “max” function (taking the maximum value from a set)
 - **Ratio constraints**, in which the constraint is non-linear due to fractional coefficients



This is important since linear programs are so much easier to solve than non-linear programs

Example 1

Marketing Problem (1 of 2)

From Tutorial 02: Algebraic Formulations

Problem Statement

- Begin with the same marketing problem as seen in Tutorial 02, plus an absolute value constraint:

$$\begin{aligned} \bullet \quad & \min \quad 500 x_1 + 200 x_2 + 250 x_3 + 125 x_4 \\ & \text{s.t.} \quad 50,000 x_1 + 25,000 x_2 + 20,000 x_3 + 15,000 x_4 \geq 1,500,000 \\ & \quad 0 \leq x_1 \leq 20 \\ & \quad 0 \leq x_2 \leq 15 \\ & \quad 0 \leq x_3 \leq 10 \\ & \quad 0 \leq x_4 \leq 15 \\ & \quad |x_1 + x_2 - x_3 - x_4| \leq 5 \end{aligned}$$

	TV	Radio	Mail	Newspaper
Audience Size	50,000	25,000	20,000	15,000
Cost/Impression	\$500	\$200	\$250	\$125
Maximum # Ads	20	15	10	15

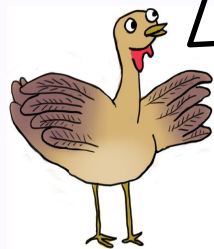
Task

- Solve with the added absolute value constraint (total number of ads in electronic media is within 5 of the number in paper-based media), and solve

Marketing Problem (2 of 2)

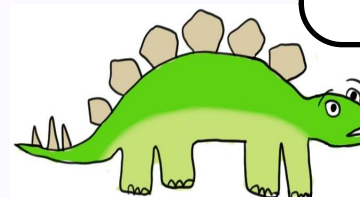
- There is a “trick” we can use to change the constraint $|x_1 + x_2 - x_3 - x_4| \leq 5$ into a linear constraint
- $|x_1 + x_2 - x_3 - x_4| \leq 5$ is equivalent to a combination of the following two constraints:
 - $x_1 + x_2 - x_3 - x_4 \leq 5$ and $-x_1 - x_2 + x_3 + x_4 \leq 5$
- The new problem is linear, and still equivalent to the original non-linear program:

$$\begin{aligned}
 & \bullet \text{min} && 500 x_1 + 200 x_2 + 250 x_3 + 125 x_4 \\
 & \text{s.t.} && 50,000 x_1 + 25,000 x_2 + 20,000 x_3 + 15,000 x_4 \geq 1,500,000 \\
 & && 0 \leq x_1 \leq 20, \quad 0 \leq x_2 \leq 15, \quad 0 \leq x_3 \leq 10, \quad 0 \leq x_4 \leq 15 \\
 & && x_1 + x_2 - x_3 - x_4 \leq 5 \\
 & && -x_1 - x_2 + x_3 + x_4 \leq 5
 \end{aligned}$$



What do you mean by equivalent? The two problems look different to me...

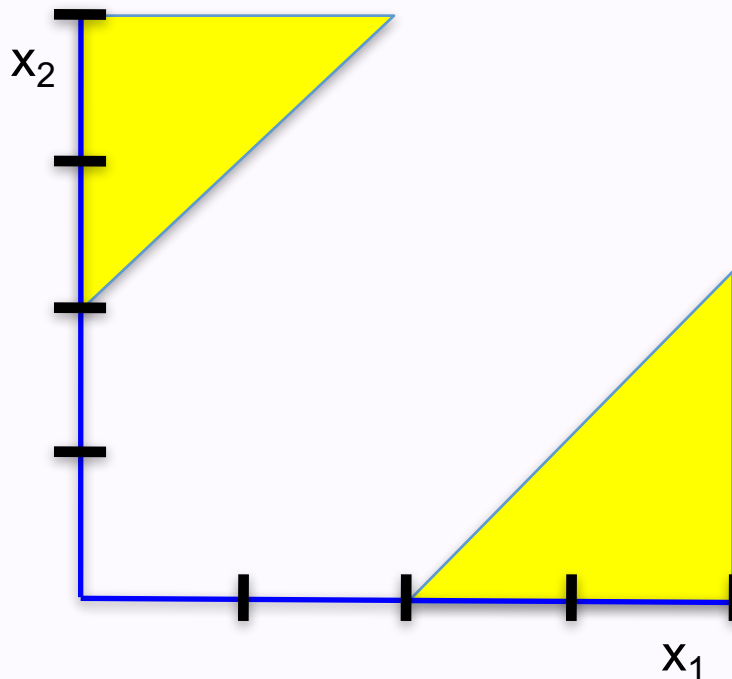
The problems are equivalent in the sense that any **feasible solution** for the non-linear program is feasible for the linear program, and vice versa. That is, **the feasible regions are exactly the same**



I'm extinct... The set of living dinosaurs is equivalent to a linear program with no feasible solutions!

Absolute Value Constraints

- We now know how to transform constraints of the form $|ax_1 + bx_2| \leq c$
- However, our trick will not work on absolute value constraints of the form $|ax_1 + bx_2| \geq c$
 - We can still transform this into two equivalent constraints, but the feasible region will no longer be convex, meaning we can't solve the LP by traditional methods
 - For example, try out $|x_1 - x_2| \geq 2$:
 - This is equivalent to $x_1 - x_2 \geq 2$ OR $-x_1 + x_2 \geq 2$, which can't be made linear:



- The feasible region is in yellow, and it's in two separate pieces
- Recall from class that a linear programming feasible region is always connected, and convex. In fact, it's always **convex**: if two points are feasible, then so is the line segment joining the two points.

Maximin Problem (1 of 2)

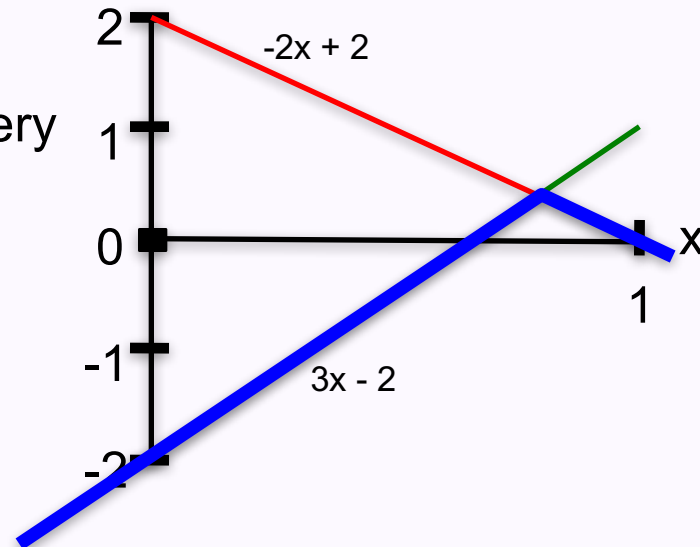
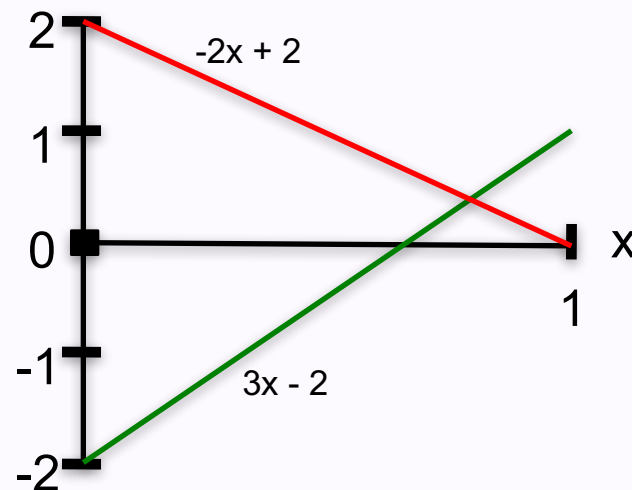
- The next constraint we will make linear appears when we maximize the value of a **min function**, which looks like $\min \{3x - 2, -2x + 2\}$

Problem Statement

- | | |
|----------|----------------------------|
| maximize | $\min \{3x - 2, -2x + 2\}$ |
| s.t. | $0 \leq x \leq 4$ |

Task

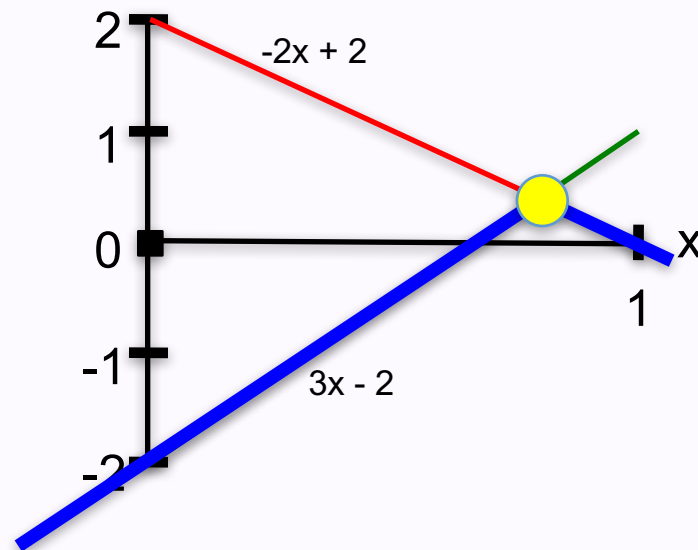
- Turn this into an LP
- To start, we can look at the graph of the problem shown above, and identify the minimum value of the two functions for every value of x (shown here in blue):



Example 2

Maximin Problem (2 of 2)

- Looking at these minimum values, we then find the maximum value of the blue line (shown as a yellow dot) to solve the problem
- The max of the min has a value of $2/5$ at $x = 4/5$:



Marketing Problem (1 of 2)

Problem Statement

- Begin with the same marketing problem as seen in Tutorial 02, but with a new objective function:
 - $$\begin{aligned} \bullet \quad & \max \quad \min \{50x_1, 25x_2, 20x_3, 15x_4\} \\ & \text{s.t.} \quad 50,000 x_1 + 25,000 x_2 + 20,000 x_3 + 15,000 x_4 \geq 1,500,000 \\ & \quad 0 \leq x_1 \leq 20 \\ & \quad 0 \leq x_2 \leq 15 \\ & \quad 0 \leq x_3 \leq 10 \\ & \quad 0 \leq x_4 \leq 15 \end{aligned}$$
- The minimum of $\{50x_1, 25x_2, 50x_3, 15x_4\}$ is the smallest number of persons reached by one of the four media (in 1000's)

Task

- Formulate this as an LP whose solution will find the maximum value of z such that each medium reaches at least $1000z$ people

Try to figure this out yourself before looking at the answer!



Marketing Problem (2 of 2)

- Now, we want to maximize z subject to the constraint that z is at most the number of ads seen for each media
- z is upper bounded by each expression in the min function, which we can represent with the constraints:
 - $z \leq 50x_1$
 - $z \leq 25x_2$
 - $z \leq 20x_3$
 - $z \leq 15x_4$
- Then, we can maximize z to get our final LP:
 - max** z
 - s.t. $50x_1 + 25x_2 + 20x_3 + 15x_4 \geq 1,500$
 - $0 \leq x_1 \leq 20, \quad 0 \leq x_2 \leq 15, \quad 0 \leq x_3 \leq 10, \quad 0 \leq x_4 \leq 15$
 - $z \leq 50x_1, \quad z \leq 25x_2, \quad z \leq 20x_3, \quad z \leq 15x_4$
- Note:* This technique works whenever you need to maximize the minimum (“**maximin**”) of linear functions. A similar trick works whenever you want to minimize the maximum (“**minimax**”) of linear functions

Ratio Constraints

Problem Statement

- Consider the Marketing Problem again
- Suppose that we wanted to add the **ratio constraint** that at least 20% of all ads had to be by mail:
 - $x_1 / (x_1 + x_2 + x_3 + x_4) \geq .2$

Task

- Formulate this as an LP
-
- The “trick” here is to multiply by the denominator so that the constraint becomes linear:
 - $x_1 \geq .2 (x_1 + x_2 + x_3 + x_4)$
 - $.8x_1 - .2 x_2 - .2 x_3 - .2 x_4 \geq 0$
 - *Note:* You can only do this if you know that the denominator’s value is positive for all possible x
 - If you multiply both sides of an inequality by a negative number, the direction of the inequality reverses
 - The new constraint is valid if $x = 0$, so we don’t need to consider that case

Summary on Non-linear Constraints

- Most of the time, if there is a constraint or objective that isn't linear, it *cannot* be transformed into a linear constraint or objective
- In the 3 cases presented above, transformations into LPs *can* be done:
 - Certain absolute value constraints
 - Remember to check if the transformation works by graphing the feasible region!
 - Maximizing the minimum or minimizing the maximum
 - Ratio constraints



These techniques are really useful and easy once you use them a couple of times!