## Final Formula Sheet

## CS540

## August 19, 2019

## 1 Formulas

- Euclidean Distance:  $\rho_2(x_i, x_{i'}) = \sqrt{\sum_{j=1}^{m} (x_{ij} x_{i'j})^2}$
- Manhattan Distance:  $\rho_1(x_i, x_{i'}) = \sum_{j=1}^{m} |x_{ij} x_{i'j}|$
- Chebyshev Distance:  $\rho_{\infty}(x_i, x_{i'}) = \max_{j=1,2,\dots,m} \{|x_{ij} x_{i'j}|\}$
- Single Linkage Distance:  $\rho\left(C_{k},C_{k'}\right)=\min\left\{\rho\left(x_{i},x_{i'}\right):x_{i}\in C_{k},x_{i'}\in C_{k'}\right\}$
- Length of a vector (in Euclidean space): length of  $x=(x_1,x_2,...,x_m)$  is  $\|x\|_2=\sqrt{x^Tx}=\sqrt{x_1^2+x_2^2+...+x_m^2}$
- Projection onto any vector  $v_k$ : proj $_{v_k} x_i = \left(\frac{v_k^T x_i}{v_k^T v_k}\right) v_k$ , and its length is  $\|\operatorname{proj}_{v_k} x_i\|_2 = \frac{v_k^T x_i}{\sqrt{v_k^T v_k}}$
- Projection onto a unit vector: first  $u_k = \frac{v_k}{\sqrt{v_k^T v_k}}$  to convert  $v_k$  into a unit vector  $u_k$  and proj  $u_k x_i = u_k^T x_i u_k$ , and its length is  $\| \operatorname{proj} u_k x_i \|_2 = u_k^T x_i$
- Sample Variance (MLE):  $\hat{\Sigma} = \frac{1}{n} \sum_{i=1}^{n} (x_i \hat{\mu}) (x_i \hat{\mu})^T$ , with  $\hat{\mu} = \frac{1}{n} \sum_{i=1}^{n} x_i$
- PCA Max Variance Formulation:  $\max_{u_k} u_k^T \hat{\Sigma} u_k$  such that  $u_k^T u_k = 1$
- PCA-reconstruction:  $x_i = \sum_{k=1}^{m} (u_k^T x_i) u_k \approx \sum_{k=1}^{K} (u_k^T x_i) u_k$
- Discounted Expected Reward (Cost):  $\mathbb{E}\left[r_t + \gamma r_{t+1} + \gamma^2 r_{t+2} + \ldots\right], \gamma \in [0, 1], \mathbb{E}\left[r_t\right] = \sum_{r_t \in R} r_t \mathbb{P}\left\{r_t | s_{t-1}, a_{t-1}\right\}$
- $\bullet$  Logic Notations:  $x_i$  or  $x_j: x_i \vee x_j, x_i$  and  $x_j: x_i \wedge x_j$  , not  $x_i: \neg x_i$
- Simulated Annealing Update Probability: if f(s') is worse than  $f(s_i)$  then  $p_i = \exp\left(-\frac{|f(s') f(s_i)|}{\text{Temp}}\right)$ .

e.g. if the 
$$f$$
 is the cost (minimization problem), then  $s_{t+1} = \begin{cases} s' & \text{if } f\left(s'\right) < f\left(s_{t}\right), \text{ else} \\ s' & \text{with probability } \exp\left(-\frac{|f\left(s'\right) - f\left(s_{t}\right)|}{\text{Temp }(t)}\right) \end{cases}$ 

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e.g. if the  $f$  is the reward (score) (maximization problem), then  $s_{t+1} = \begin{cases} s' & \text{if } f\left(s'\right) > f\left(s_{t}\right), \text{ else} \\ s' & \text{with probability } \exp\left(-\frac{|f\left(s'\right) - f\left(s_{t}\right)|}{\text{Temp }(t)}\right) \end{cases}$ 

$$s_{t} & \text{otherwise}$$

- Genetic Algorithm Reproduction Probability: if F is fitness then  $p_i = \frac{F\left(s_i\right)}{\sum\limits_{i=1}^{N} F\left(s_j\right)}$
- Alpha Beta Pruning: prune if  $\alpha \geqslant \beta$  where  $\alpha\left(s\right) = \max_{s' \in s'\left(s\right)} \beta\left(s'\right), \beta\left(s\right) = \min_{s' \in s'\left(s\right)} \alpha\left(s'\right)$
- Geometric Sum:  $1 + x + x^2 + ... + x^n = \frac{1 x^n}{1 x}, x \in (0, 1)$  and  $1 + x + x^2 + ... = \frac{1}{1 x}$
- Arithmetic Sum:  $1 + 2 + 3 + ... + n = \frac{(1+n) \cdot n}{2}$