### CS540 Introduction to Artificial Intelligence Lecture 17

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## Lecture Indices

- Lecture 13 will be recorded and posted after the final exam.
- Lecture 14 is a midterm review so it will be skipped.
- Lecture 15 is this lecture.
- Pre-recorded lectures are official lectures for Summer 2020 and covers all relevant materials for homework and exams.
- BBCU lectures are only summaries, reviews, additional examples, and discussions.
- Quizzes are alternatives for putting all the weights on the exams, they are not mandatory, just a way to encourage interactions and discussions.

## Learning vs Search

- In reinforcement learning, the reward and state transition need to be learned by taking actions.
- In search problems, the reward and state transitions are given.
- The problem is to find a sequence of actions that lead to the goal with minimum cost.

## Search Problem Applications Motivation

- Puzzles and games.
- Navigation: route finding.
- Motion planning.
- Scheduling.

# Wolf, Sheep, Cabbage Example Motivation

### Search Problem

#### Motivation

- State space S is the set of all valid configurations.
- Initial states I and goal states G are subsets of S.
- Successor function s'(s) given the current state s is the set of states reachable in one step from s.
- There is a cost (or negative reward) associated with moving from s to s'(s).
- The search problem is the problem of finding a solution path from a state in I to a state in G, usually with minimum total cost.

## State Space

- The states need to represent all necessary information about the game.
- The actions are discrete and deterministic and are determined by the successor function.
- Each possible action at state s is associated with a state in the set s'(s).

### 8 Puzzle Example

Motivation

# Sizes of State Space

• Tic Tac Toe:  $10^3$ 

• Checkers:  $10^{20}$ 

• Chess: 10<sup>50</sup>

• Go: 10<sup>170</sup>

## State Space Graph

- A state space can be represented by a weighted directed graph (V, E, c).
- V is the set of vertices (also called nodes).
- *E* is the set of edges (also called arcs). Each edge is directed from one vertex to another vertex and represents an action.
- *c* is the cost (also called weights) associated with each edge. The costs are positive.

# Search Problem on Graph Definition

- Search starts at an initial state and finishes if one of the goal states is reached.
- The solution is a path in the graph from an initial state to a goal state.
- The cost of a solution is the sum of edge costs on the solution path.
- The optimal solution is the solution with the lowest cost.

# Expansion Definition

- Vertices that are explored so far are stored in a tree called the state space search tree.
- Expanding a vertex means to generate all successor vertices and add them (and the associated edges) to the state space search tree.
- The leaves of the search tree are unexpanded and are called the frontier (sometimes called the fringe).
- The search strategies differ in the order in which the vertices are expanded.

# Water Jugs Example Definition

### Performance

Definition

- A search strategy is complete if it finds at least one solution.
- A search strategy is optimal if it finds the optimal solution.
- For uninformed search, the costs are assumed to be 1 for all edges c=1.

### Complexity

#### Definition

- The time complexity of a search strategy is the worst case maximum number of vertices expanded.
- The space complexity of a search strategy is the worst case maximum number of states stored in the frontier at a single time.
- Notation: the goals are *d* edges away from the initial state. This means assuming a constant cost of 1, the optimal solution has cost *d*. The maximum depth of the graph is *D*.
- Notation: the branching factor is *b*, the maximum number of actions associated with a state.

$$b = \max_{s \in V} \left| s'(s) \right|$$

# Search Tree Diagram Definition

# Breadth First Search Description

- Use Queue (FIFO) for the frontier.
- Remove from the front, add to the back.

### Maze BFS Example

Definition

### Breadth First Search

#### Algorithm

- Input: a weighted digraph (V, E, c), initial states I and goal states G.
- Output: a path from I to G.
- EnQueue initial states.

$$Q = I$$

 While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_0$$
$$Q = Q + s'(s)$$

### Breadth First Search Performance

Discussion

- BFS is complete.
- BFS is optimal with c = 1.

### Breadth First Search Complexity

#### Discussion

• Time complexity: the worst case occurs when the goal is the last vertex at depth *d*.

$$T = b + b^2 + ... + b^d$$

 Space complexity: the worst case is storing all vertices at depth d is in the frontier.

$$S = b^d$$

### BiDirectional Search

#### Discussion

- BFS from the initial states and goal states at the same time.
- The search stops when the two frontiers meet (have non-empty intersection) in the middle.
- The time and space complexity is the same as BFS with depth  $\frac{d}{2}$ .

# Depth First Search Description

- Use Stack (LIFO) for the frontier.
- Remove from the front, add to the front.

### Maze DFS Example

Definition

### Depth First Search

#### Algorithm

- Input: a weighted digraph (V, E, c), initial states I and goal states G.
- Output: a path from I to G.
- Push initial states.

$$S = I$$

• While S is not empty and goal is not popped, pop S and push its successors.

$$s = S_0$$
$$S = s'(s) + S$$

## Depth First Search Performance

- DFS is incomplete if  $D = \infty$ .
- DFS is not optimal.

### Depth First Search Complexity

Discussion

• Time complexity: the worst case occurs when the goal is the root of the last subtree expanded in the whole graph.

$$T = b^{D-d+1}... + b^{D-1} + b^{D}$$

• Space complexity: the worst case is storing all vertices sharing the parents with vertices in the current path.

$$S = (b-1)D + 1$$

# Iterative Deepening Search Description

- ullet DFS but stop if path length > 1
- repeat DFS but stop if path length > 2
- ...
- repeat DFS but stop if path length > d

# Maze IDS Example Definition

# Iterative Deepening Search

- Input: a weighted digraph (V, E, c), initial states I and goal states G.
- Output: a path from I to G.
- Perform DFS on the digraph restricted to vertices with depth
   \$ 1 from the initial state.
- Perform DFS on the digraph restricted to vertices with depth
   \$ 2 from the initial state.
- Repeat until the goal is deQueued.

## Iterative Deepening Search Performance

- IDS is complete.
- IDS is optimal with c = 1.

### Iterative Deepening Search Complexity

Discussion

• Time complexity: the worst case occurs when the goal is the last vertex at depth *d*.

$$T = db + (d-1)b^{2} + ... + 3b^{d-2} + 2b^{d-1} + 1b^{d}$$

• Space complexity: it has the same space complexity as DFS.

$$S=(b-1)\,d$$

### Non-Tree Search

#### Discussion

- If the state space is not a tree, search strategies need to remember the states that are already expanded.
- A vertex should be removed from the frontier if it is already expanded.

### Configuration Space

Discussion

## Uniformed vs Informed Search

- Uninformed search means only the goal G and the successor function s' are given.
- Informed search means which non-goal states are better is also known.
- Usually, iterative deepening is used for uninformed search.