

CS540 Introduction to Artificial Intelligence

Lecture 16

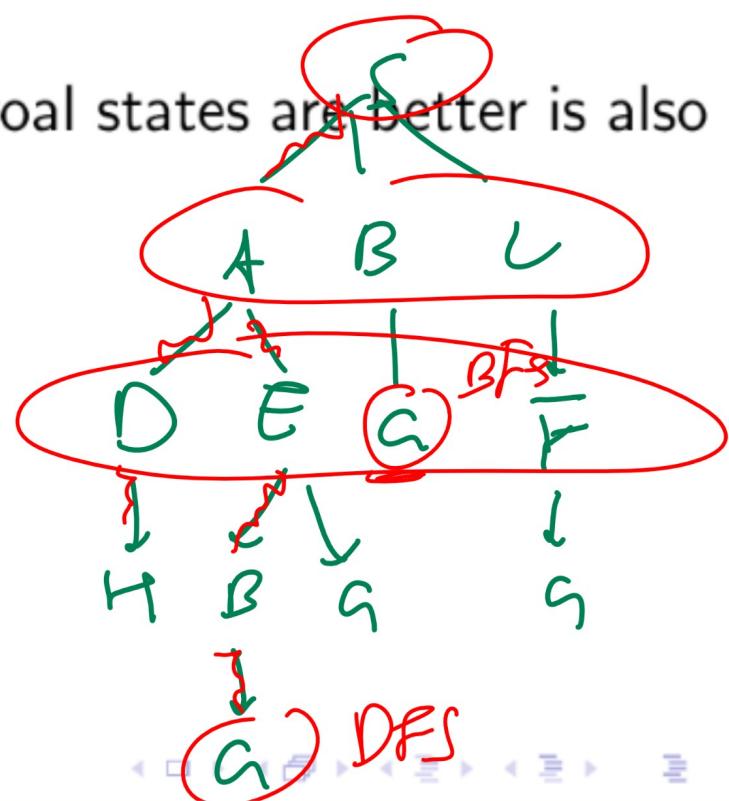
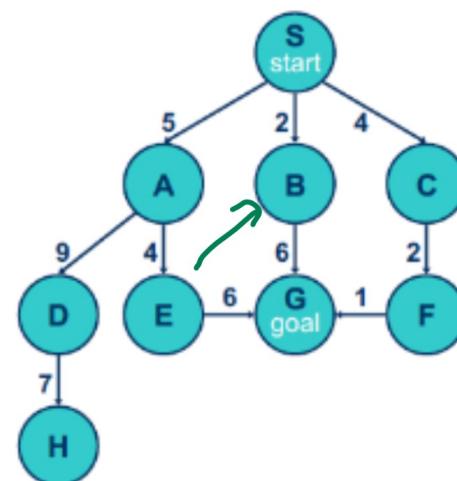
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Based on lecture slides by Jerry Zhu, Yingyu Liang, and Charles Dyer

July 20, 2020

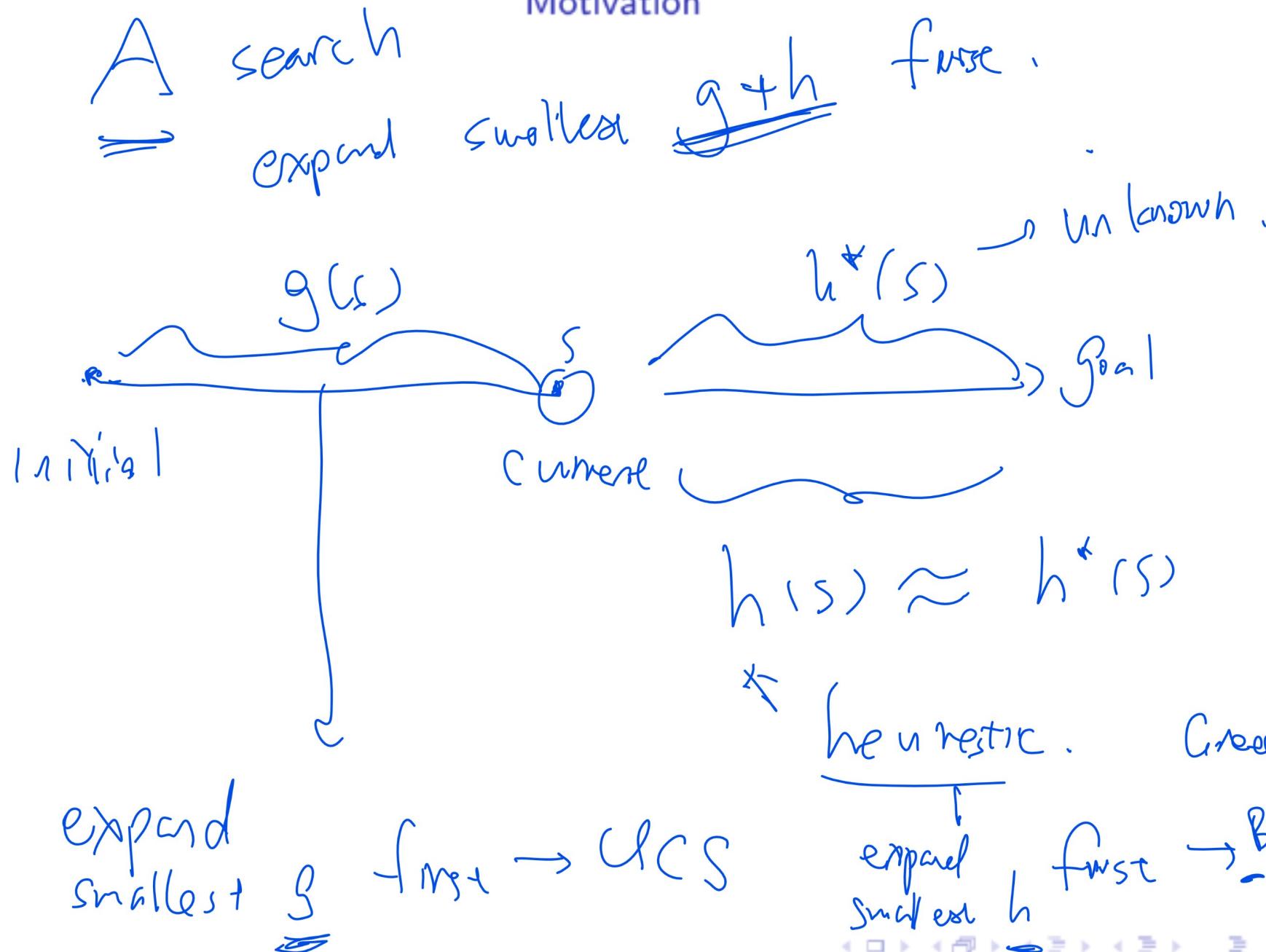
Uninformed vs. Informed Search

- Uninformed search means only the goal G and the successor functions s' are given.
 - Informed search means which non-goal states are better is also known.



Heuristic Diagram

Motivation



Uniform Cost Search

Description

- Expand the vertices with the lowest current path cost $g(s)$ first.
- It is BFS with a priority queue based on $g(s)$.
- It is equivalent to BFS if $c = 1$ is constant on all edges.
- It is also called Dijkstra's Algorithm.

UCS Example 1

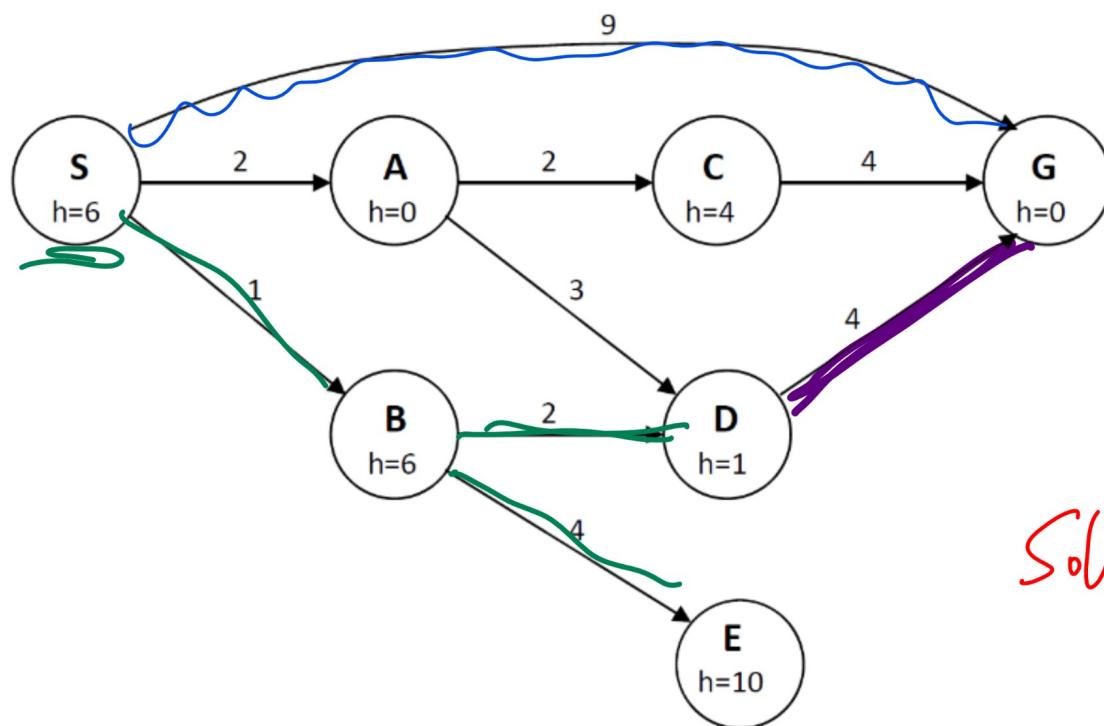
Quiz

- Spring 2017 Midterm Q1
- Given the following adjacency matrix. Find UCS expansion path.

-	S	A	B	C	D	E	G
S	$h = 6$	2	1	-	-	-	9
A	-	$h = 0$	-	2	3	-	-
B	-	-	$h = 6$	-	2	4	-
C	-	-	-	$h = 4$	-	-	4
D	-	-	-	-	$h = 1$	-	4
E	-	-	-	-	-	$h = 10$	-
G	-	-	-	-	-	-	$h = 0$

UCS Example 1 Diagram

Quiz



Expansion path:

S, B, A, D

C, D, E, G

Solution: S, B, D, G

Priority Queue: ~~B_s~~ ~~A_s~~

g	0	1	2
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~~D_B~~ ~~C_A~~ ~~D_E~~ ~~G_C~~ ~~G_D~~ ~~G_S~~

$1+2$	$2+3$	$1+4$	$3+4$	$4+4$	9
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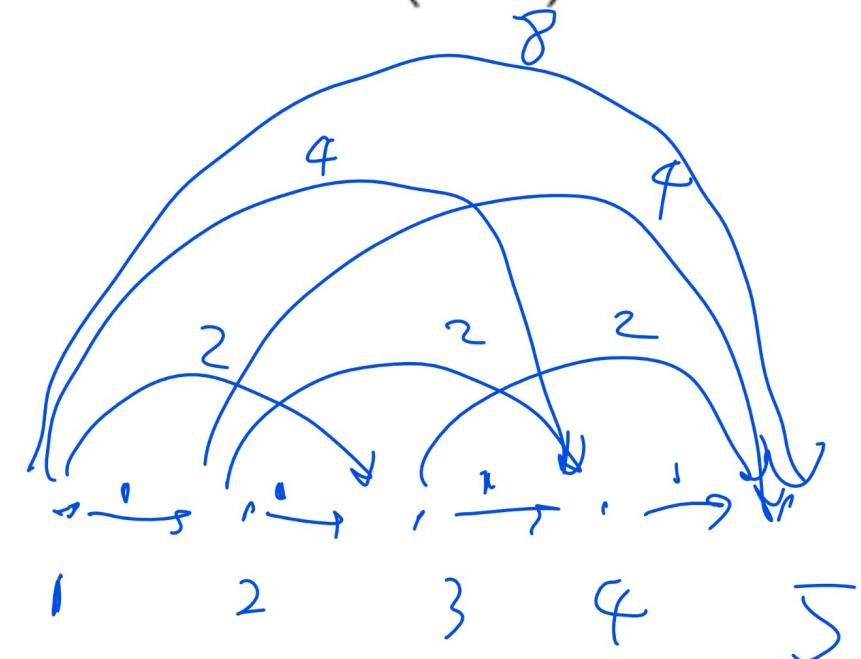
UCS Example 2

Quiz

Q 11

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. The initial state is 1 and goal state is 5. What is a vertex expansion sequence if Uniform Cost Search (UCS) is used?
 - A: 1, 5
 - B: 1, 2, 3, 4, 5
 - C: 1, 2, 3, 4, 4, 5
 - D: 1, 2, 3, 3, 4, 4, 5
 - E: 1, 2, 3, 3, 4, 4, 4, 5
- ?

$\frac{1}{4}$



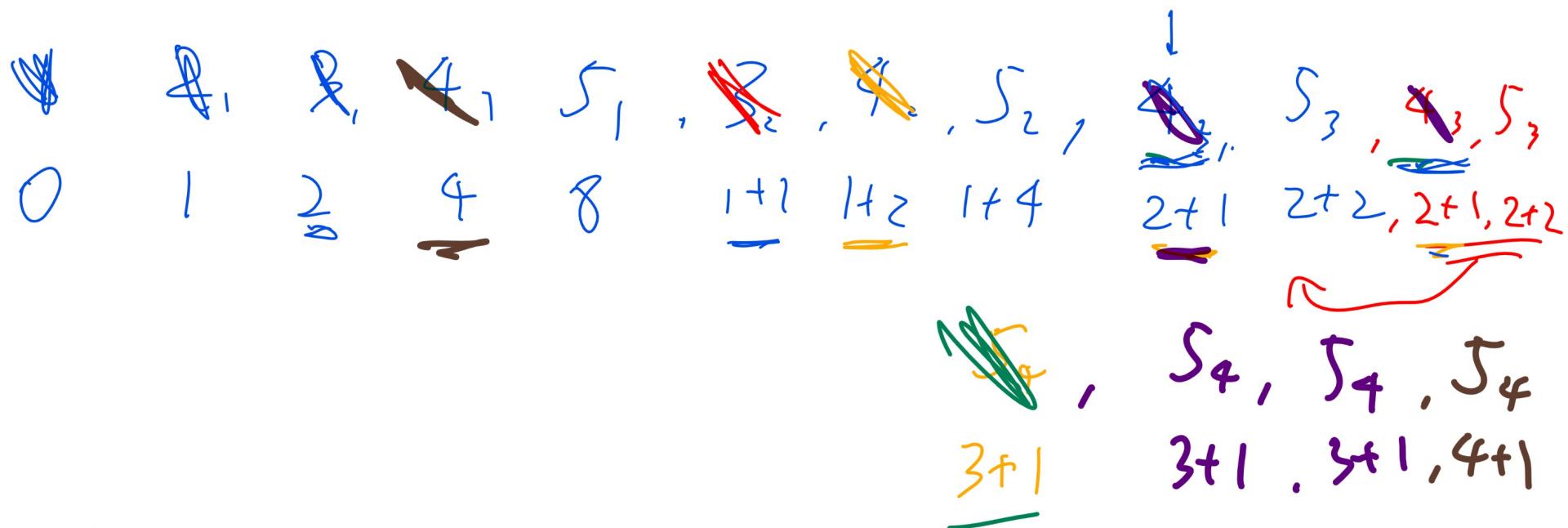
if $c = \text{constant}$, BFS = UCS.

UCS Example 2 Diagram

Queue:

Quiz

selecting smallest to deQueue;

 $1, 2, 3, 4, 4, 4, 4, 5$

Post video later

Uniform Cost Search

Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G .
- Output: a path from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $g(s)$ for $s \in Q$.

$$Q = I$$

- While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s)$$

$$Q = Q + s' (s)$$

Uniform Cost Search Performance

Discussion

- UCS is complete.
- UCS is optimal with any c .

Best First Greedy Search

Description

- Expand the vertices with the lowest heuristic cost $\underline{h(s)}$ first.
- Use a priority queue based on $h(s)$.

Greedy Example 1

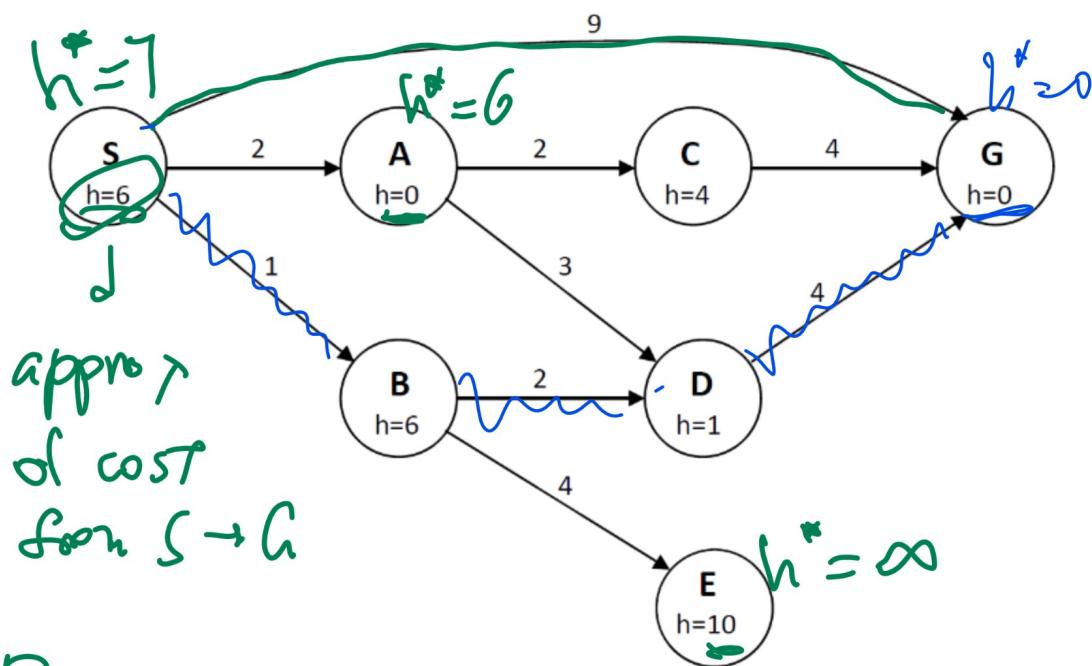
Quiz

- Given the following adjacency matrix. Find Greedy Search expansion path.

-	S	A	B	C	D	E	G
S	$h = 6$	2	1	-	-	-	9
A	-	$h = 0$	-	2	3	-	-
B	-	-	$h = 6$	-	2	4	-
C	-	-	-	$h = 4$	-	-	4
D	-	-	-	-	$h = 1$	-	4
E	-	-	-	-	-	$h = 10$	-
G	-	-	-	-	-	-	$h = 0$

Greedy Example 1 Diagram

Quiz



Priority Queue:

~~A_s~~

G_s D C A J

B_s

h

0

0 1 4 6

6

expansion:

S, A, G

Solver

$S \rightarrow G$

BF GS

Greedy
is wrong
in general.

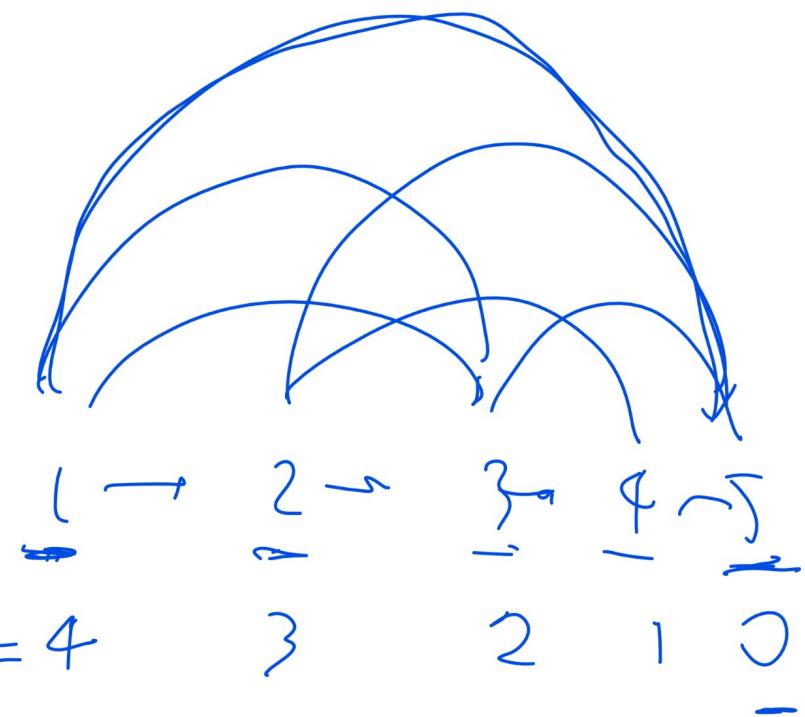
Greedy Example 2

Quiz

Q12

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. The heuristic is $h(i) = \underline{5 - i}$. The initial state is 1 and goal state is 5. What is a vertex expansion sequence if Best First Greedy Search is used?

- A: 1, 5
- B: 1, 2, 3, 4, 5
- C: 1, 2, 3, 4, 4, 5
- D: 1, 2, 3, 3, 4, 4, 5
- E: 1, 2, 3, 3, 4, 4, 4, 5



$$h^* = h = 4$$

Greedy Example 2 Diagram

Quiz

Best First Greedy Search

Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G , and the heuristic function $h(s), s \in V$.
- Output: a path from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $h(s)$ for $s \in Q$.

$$Q = I$$

- While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} h(s)$$

$$Q = Q + s' (s)$$

Best First Greedy Search Performance

Discussion

- Greedy is incomplete.
- Greedy is not optimal.

A Search

Description

- Expand the vertices with the lowest total cost $g(s) + h(s)$ first.
- Use a priority queue based on $g(s) + h(s)$.
- A stands for Always be optimistic?

A Search Example 1

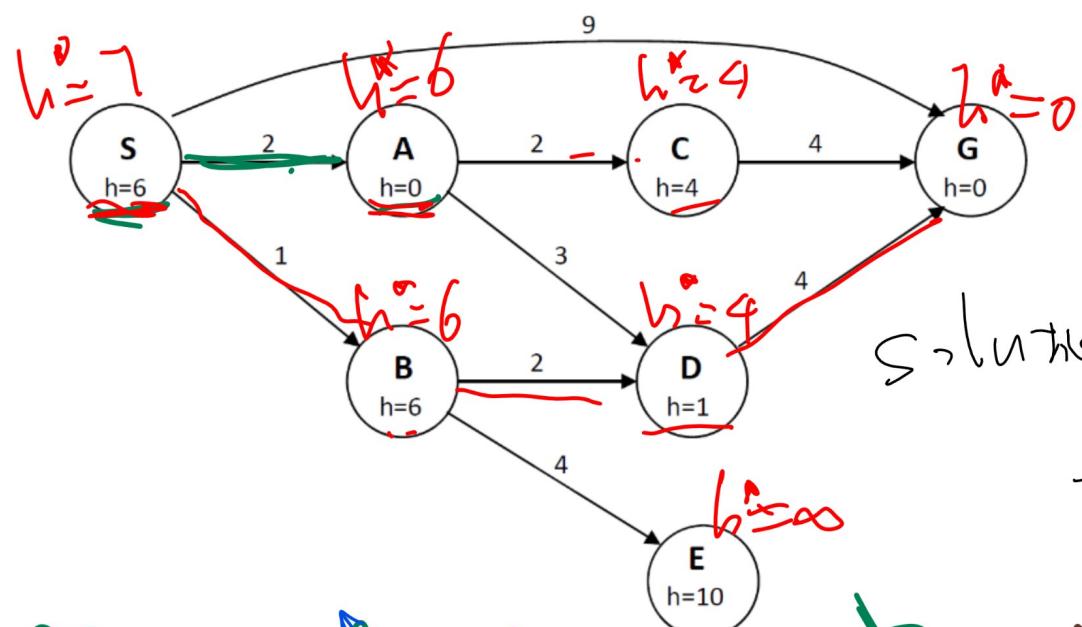
Quiz

- Given the following adjacency matrix. Find A Search expansion path.

-	S	A	B	C	D	E	G
S	$h = 6$	2	1	-	-	-	9
A	-	$h = 0$	-	2	3	-	-
B	-	-	$h = 6$	-	2	4	-
C	-	-	-	$h = 4$	-	-	4
D	-	-	-	-	$h = 1$	-	4
E	-	-	-	-	-	$h = 10$	-
G	-	-	-	-	-	-	$h = 0$

A Search Example 1 Diagram

Quiz



Expansion

S, A, D, B, D, G

Solution:

S → B → D → G

pa.	A	D	B	C	G	E
g	2	3	5	0	1	7
h	0	1	1	6	6	0
g+h	2	4	6	6	7	15

A Search

Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G , and the heuristic function $h(s), s \in V$.
- Output: a path from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $g(s) + h(s)$ for $s \in Q$.

$$Q = I$$

- While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

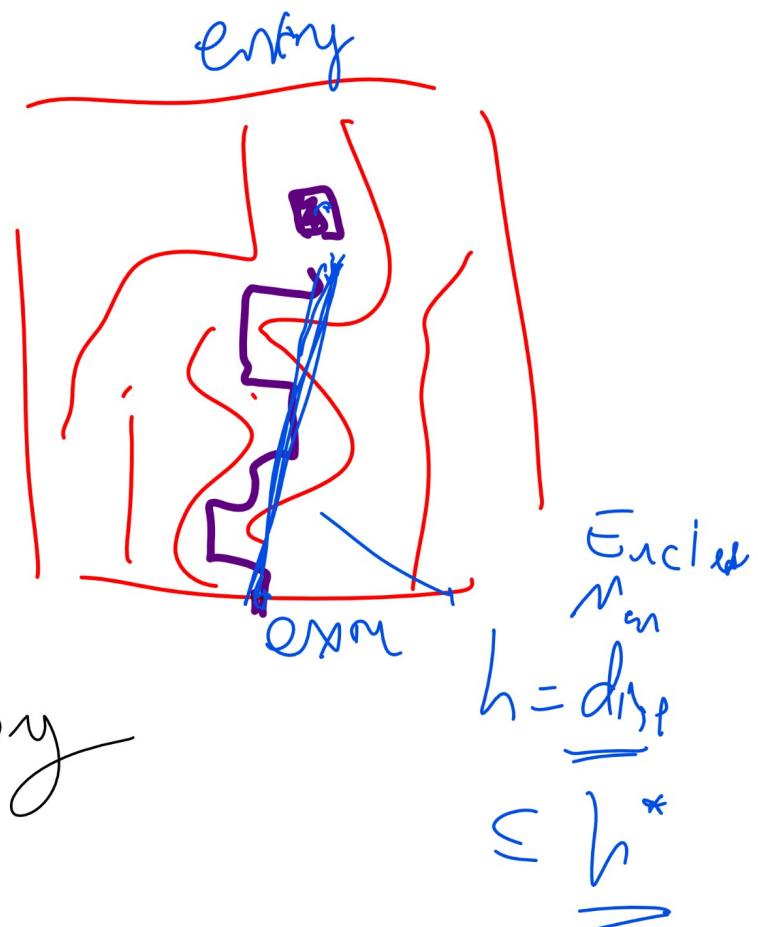
$$Q = Q + s'(s)$$

A Search Performance

Discussion

- A is complete.
- A is not optimal

h can be wrong

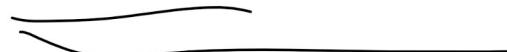


A Star Search

Description

$$0 \leq h \leq h^*$$

- A^* search is A search with an **admissible heuristic**.



Admissible Heuristic

Definition

- A heuristic is admissible if it never over estimates the true cost.

$$\boxed{0 \leq h(s) \leq h^*(s)}$$

Dominated Heuristic

Definition

- One heuristic, h_1 , is dominated by another, h_2 , if:

$$h_1(s) \leq h_2(s) \leq h^*(s), \forall s \in S$$

- If h_2 dominates h_1 , then h_2 is better than h_1 since A^* using h_1 expands at least as many states (or more) than A^* using h_2 .
- If h_2 dominated h_1 , A^* with h_2 is better informed than A^* with h_1 .

Admissible Heuristic 8 Puzzle Example

Quiz

$$0 \leq h \leq h^*$$

- Which ones (select multiple) of the following are admissible heuristic function for the 8 Puzzle?

A: $h(s)$ = number of tiles in the wrong position.



B: $h(s) = 0$.



C: $h(s) = 1$.

$$h(\text{goal}) = 1 > h^*(\text{goal}) = 0$$

D: $h(s)$ = sum of Manhattan distance between each tile and its goal location.

$$\leq h^*$$



E: $h(s)$ = sum of Euclidean distance between each tile and its goal location.

$$\leq h_D \leq h^*$$



Admissible Heuristic General Example 1

Quiz

$$0 \leq h \leq h^*$$

- Which ones (select multiple) of the following are admissible heuristic function?

- A: $h(s) = h^*(s)$.
- B: $h(s) = \max\{2, h^*(s)\}$.
- C: $h(s) = \min\{2, h^*(s)\}$.
- D: $h(s) = h^*(s) - 2$.
- E: $h(s) = \sqrt{h^*(s)}$.

$$h^*(\text{goal}) = 0$$

$$\begin{aligned}h(\text{goal}) &= \max(2, 0) \\&= 2 > 0\end{aligned}$$

$$= h^*(s)$$

$$h^*(\text{goal}) = 0$$

$$h(\text{goal}) = 0 - 2 < 0$$

$$h^*(s) = 0.25, \quad h(s) = \sqrt{0.25} = 0.5 > 0.25 = h^*$$

A Star Search Example 2

Quiz

Q13(hse)

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. How many of the following heuristic functions are admissible? For $i = 1, 2, 3, 4, 5$:

$$0 \leq h \leq h^*$$



$$h(i) = 5 - i$$

$$\textcircled{2} h(i) = \sqrt{5 - i}$$

$$\textcircled{3} h(i) = \log_2(6 - i)$$

$$\textcircled{4} h(i) = 1 - \frac{1}{i=5}$$

$$h(i) = 0$$

$$\min(1, h^*)$$

- A: 1, B: 2, C: 3, D: 4, E: 5

	h_2	h_4	i	$b^*(i)$	h^*
2	1	1			$\log_2 5$
$\sqrt{3}$	2	2			2
$\sqrt{2}$	3	3			$\log_2 3$
1	4	4			1
0	5	5			0

$$\left\{ \begin{array}{l} 1 \\ 0 \end{array} \right. \begin{array}{l} i=j \\ i \neq 5 \end{array}$$

A Star Search Example 2

Quiz

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. Which one of the following heuristic functions is not dominated (among the admissible ones)? For $i = 1, 2, 3, 4, 5$:
 - A: $h(i) = 5 - i$
 - B: $h(i) = \sqrt{5 - i}$
 - C: $h(i) = \log_2 (6 - i)$
 - D: $h(i) = 1 - \mathbb{1}_{i=5}$
 - E: $h(i) = 0$

A Star Search with Revisit, Part I

Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G , and the heuristic function $h(s), s \in V$.
- Output: a path with minimum cost from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $g(s) + h(s)$ for $s \in Q$.

$$Q = I$$

$$g(I) = 0$$

$$g(s) = \infty, \text{ for } s \notin I$$

- Initialize the list of visited vertices, P .

$$P = \emptyset$$

A Star Search with Revisit, Part II

Algorithm

- While Q is not empty and goal is not deQueued, deQueue Q , put it on P and enQueue its successors to Q , and update the cost functions.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$P = P + s$$

$$Q = Q + s'(s), \text{ update } g(s') = \min \{g(s'), g(s) + c(s, s')\}$$

A Search Performance Discussion

- A^* with $h \leq h^\star$
- A^* is complete.
 - A^* is optimal.

Iterative Deepening A Star Search

Discussion

- A* can use a lot of memory.
- Do path checking without expanding any vertex with $g(s) + h(s) > 1$.
- Do path checking without expanding any vertex with $g(s) + h(s) > 2$.
- ...
- Do path checking without expanding any vertex with $g(s) + h(s) > d$.

IDA*

Iterative Deepening A Star Search Performance

Discussion

- IDA* is complete.
- IDA* is optimal.
- IDA* is more costly than A^* .

Beam Search

Discussion

- Version 1: Keep a priority queue with fixed size k . Only keep the top k vertices and discard the rest.
- Version 2: Only keep the vertices that are at most ε worse than the best vertex in the queue. ε is called the beam width.

$$\underbrace{j \text{th}}_{\downarrow} \leq \varepsilon$$

Beam Search Performance

Discussion

- Beam is incomplete.
- Beam is not optimal.