

CS540 Introduction to Artificial Intelligence

Lecture 16

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Based on lecture slides by Jerry Zhu, Yingyu Liang, and Charles Dyer

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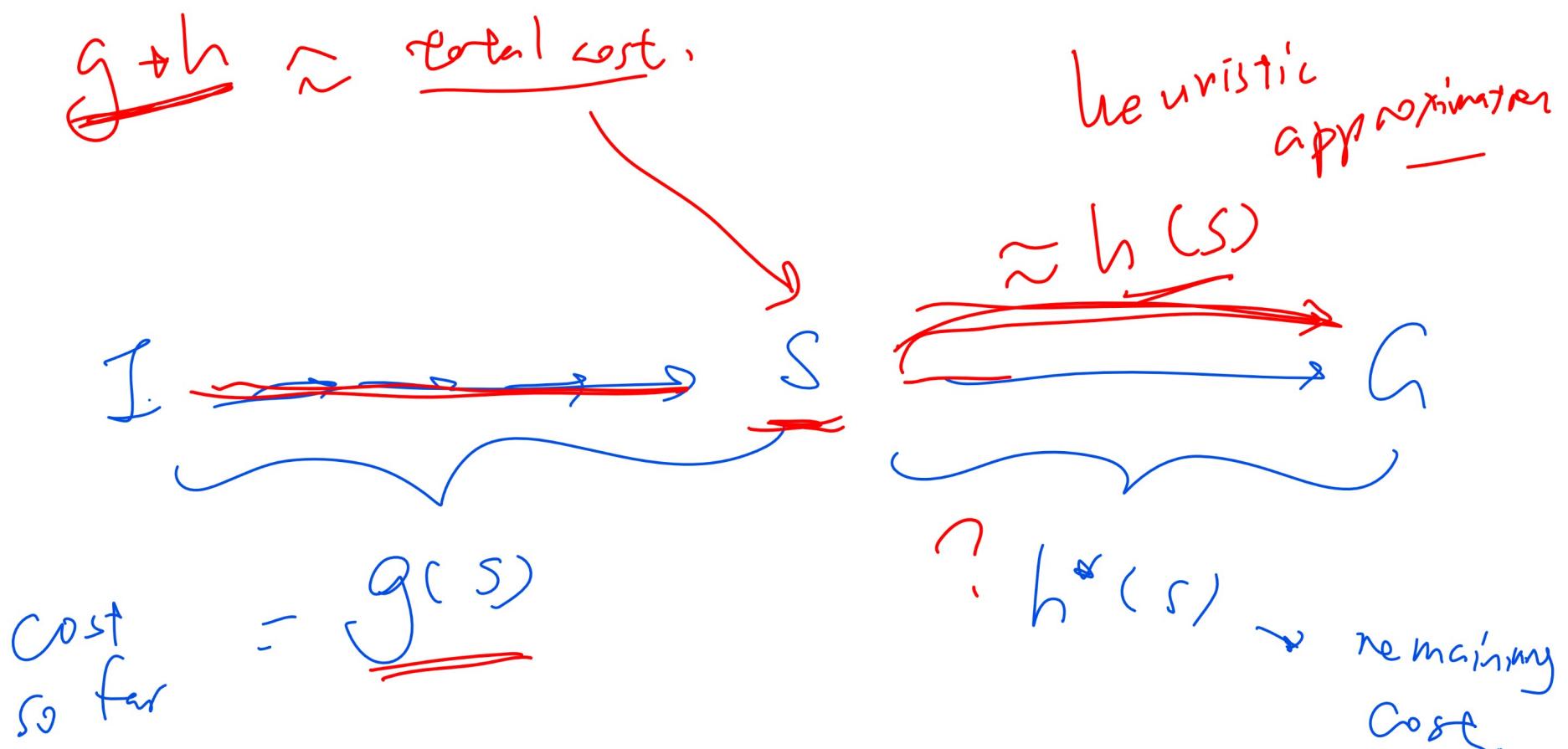
Uniformed vs. Informed Search

Motivation

- Uninformed search means only the goal G and the successor functions s' are given.
 - Informed search means which non-goal states are better is also known.

Heuristic Diagram

Motivation



Uniform Cost Search

Description

- Expand the vertices with the lowest current path cost $g(s)$ first.
- It is BFS with a priority queue based on $g(s)$.
- It is equivalent to BFS if $c = 1$ is constant on all edges.
- It is also called Dijkstra's Algorithm.

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UCS Example 1

Quiz

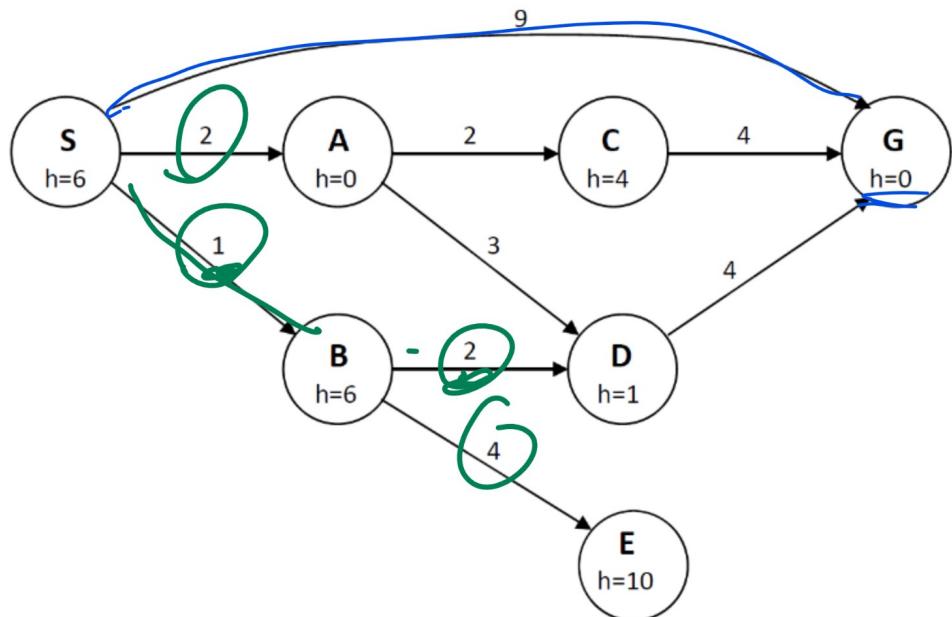
- Spring 2017 Midterm Q1
- Given the following adjacency matrix. Find UCS expansion path.



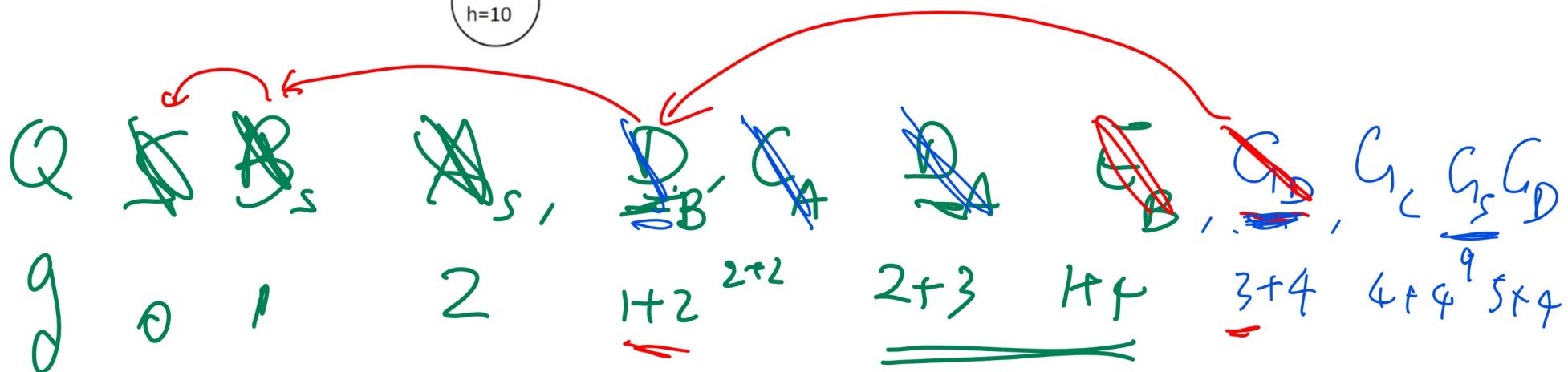
-	S	A	B	C	D	E	G
S	$h = 6$	2	1	-	-	-	9
A	-	$h = 0$	-	2	3	-	-
B	-	-	$h = 6$	-	2	4	-
C	-	-	-	$h = 4$	-	-	4
D	-	-	-	-	$h = 1$	-	4
E	-	-	-	-	-	$h = 10$	-
G	-	-	-	-	-	-	$h = 0$

UCS Example 1 Diagram

Quiz



SB DG Schriften
apt.



Expansion path S, B, A, D, C, P, E, G.

UCS Example 2

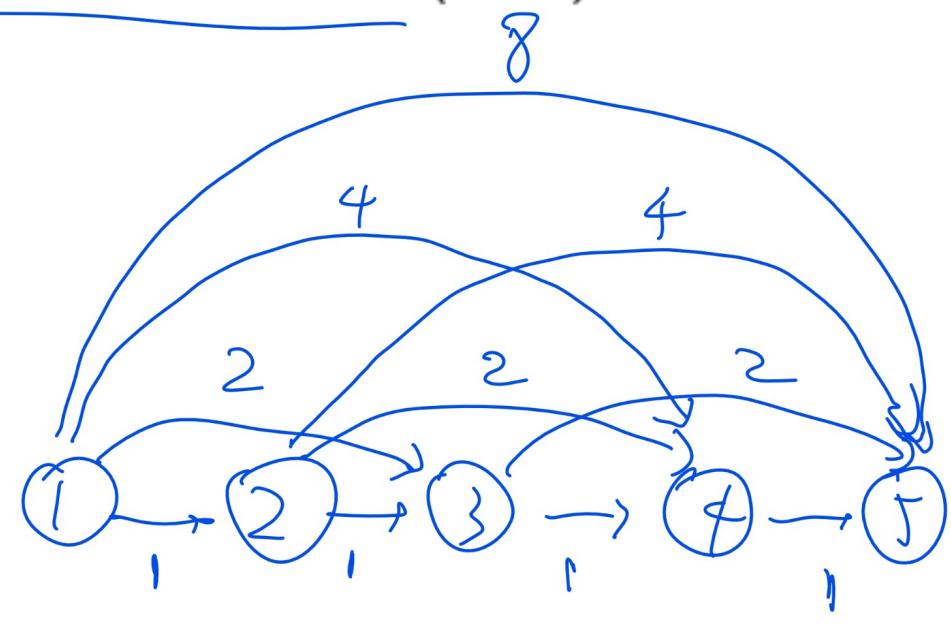
Q5 (last)

Quiz

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. The initial state is 1 and goal state is 5. What is a vertex expansion sequence if Uniform Cost Search (UCS) is used?

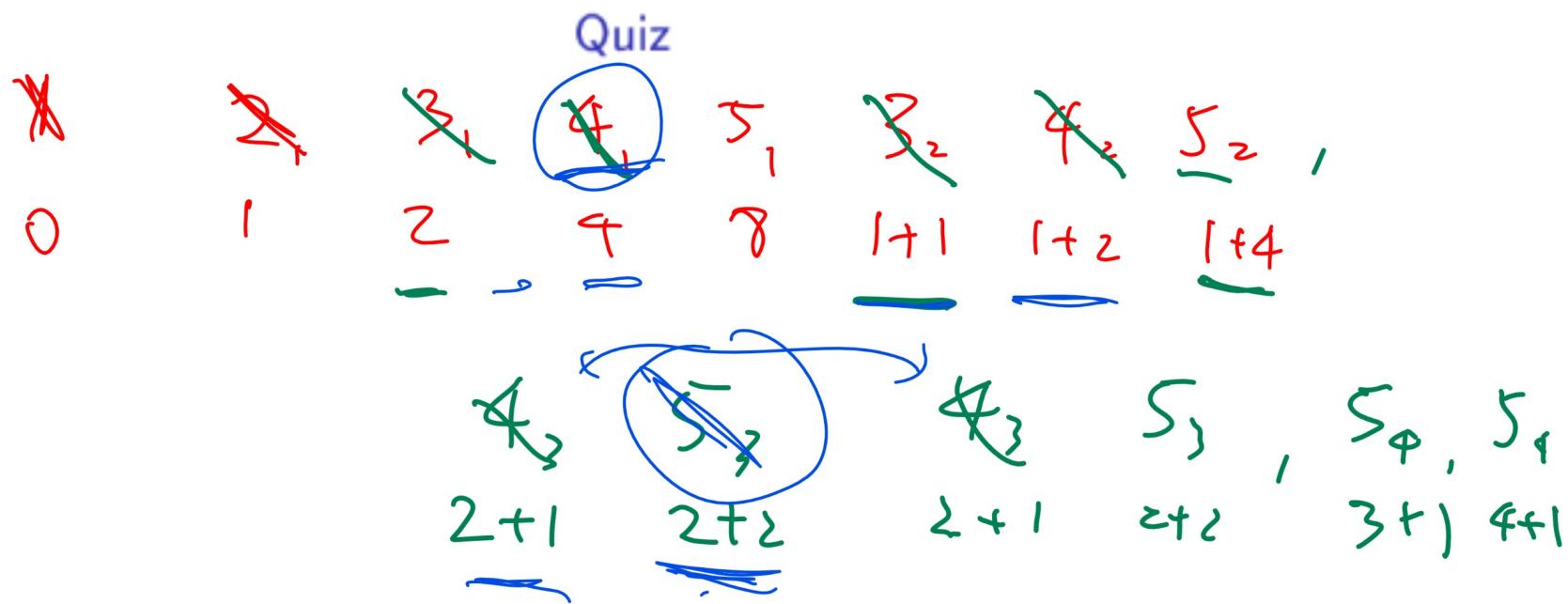
- A: 1, 5
- B: 1, 2, 3, 4, 5
- C: 1, 2, 3, 4, 4, 5
- D: 1, 2, 3, 3, 4, 4, 5
- E: 1, 2, 3, 3, 4, 4, 4, 5

1, 2



UCS Example 2 Diagram

Q
g



Record a video.

Uniform Cost Search

Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G .
- Output: a path from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $g(s)$ for $s \in Q$.

$$Q = I$$

- While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s)$$

$$Q = Q + s' (s)$$

Uniform Cost Search Performance

Discussion

- UCS is complete.
- UCS is optimal with any c .

M9, M10

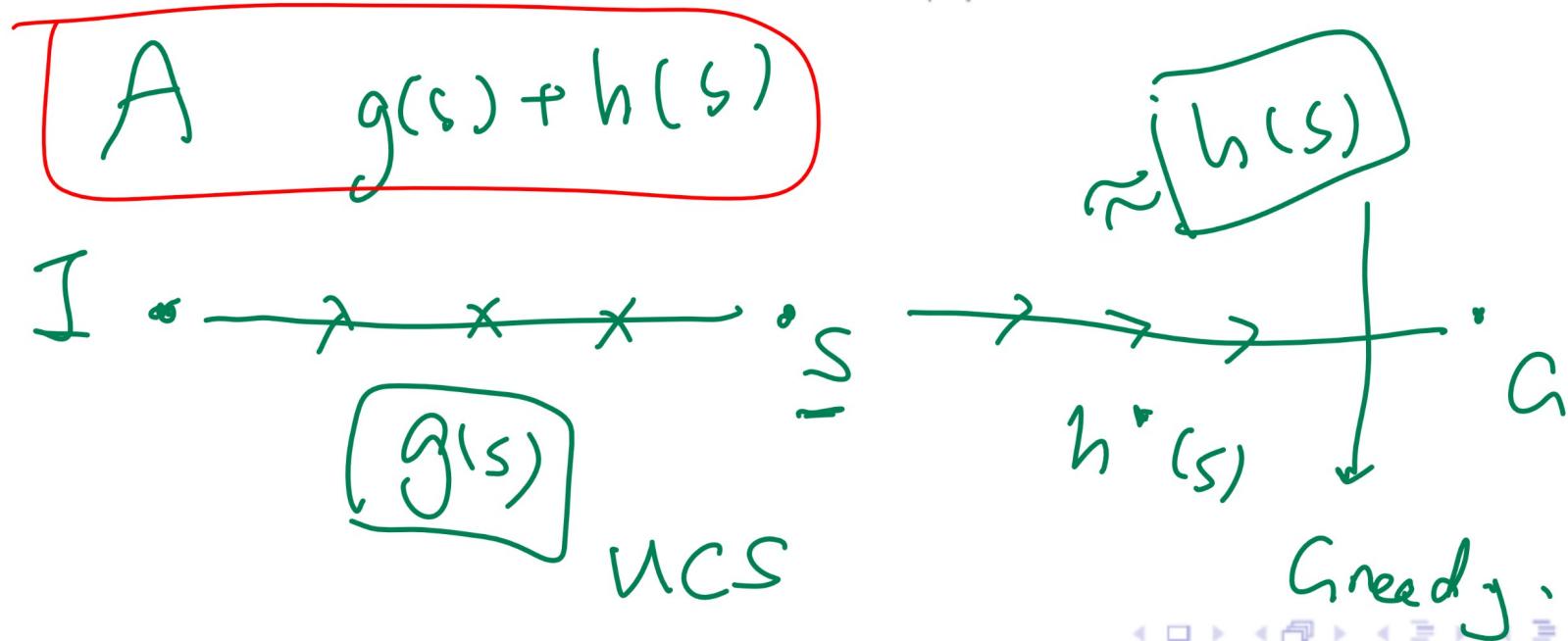
PS \rightarrow b/w.

wants for SGT.

Best First Greedy Search

Description

- Expand the vertices with the lowest heuristic cost $h(s)$ first.
- Use a priority queue based on $h(s)$.



Greedy Example 1

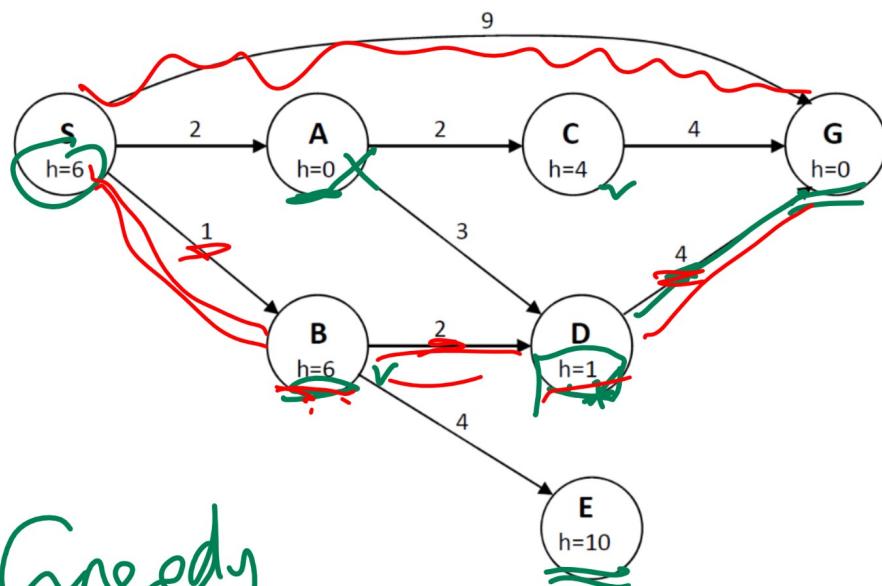
Quiz

- Given the following adjacency matrix. Find Greedy Search expansion path.

-	S	A	B	C	D	E	G
S	$h = 6$	2	1	-	-	-	9
A	-	$h = 0$	-	2	3	-	-
B	-	-	$h = 6$	-	2	4	-
C	-	-	-	$h = 4$	-	-	4
D	-	-	-	-	$h = 1$	-	4
E	-	-	-	-	-	$h = 10$	-
G	-	-	-	-	-	-	$h = 0$

Greedy Example 1 Diagram

Quiz



UCS

SBDGC lowest cost
7

Q: ~~S~~ ~~A~~ ~~S~~ D ~~C~~ ~~G~~ ~~A~~ ~~S~~ BS
h: 0 0 1 4 6 6

expansion. S A G
→ not a solution

solution

~~SG~~

Greedy Example 2

Quiz

Q)

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. The heuristic is $h(i) = 5 - i$, The initial state is 1 and goal state is 5. What is a vertex expansion sequence if Best First Greedy Search is used?

- A: 1, 5
- B: 1, 2, 3, 4, 5
- C: 1, 2, 3, 4, 4, 5
- D: 1, 2, 3, 3, 4, 4, 5
- E: 1, 2, 3, 3, 4, 4, 4, 5

UCS

Q :

X

X

X

X

X

4

0

1

2

3

h :

4

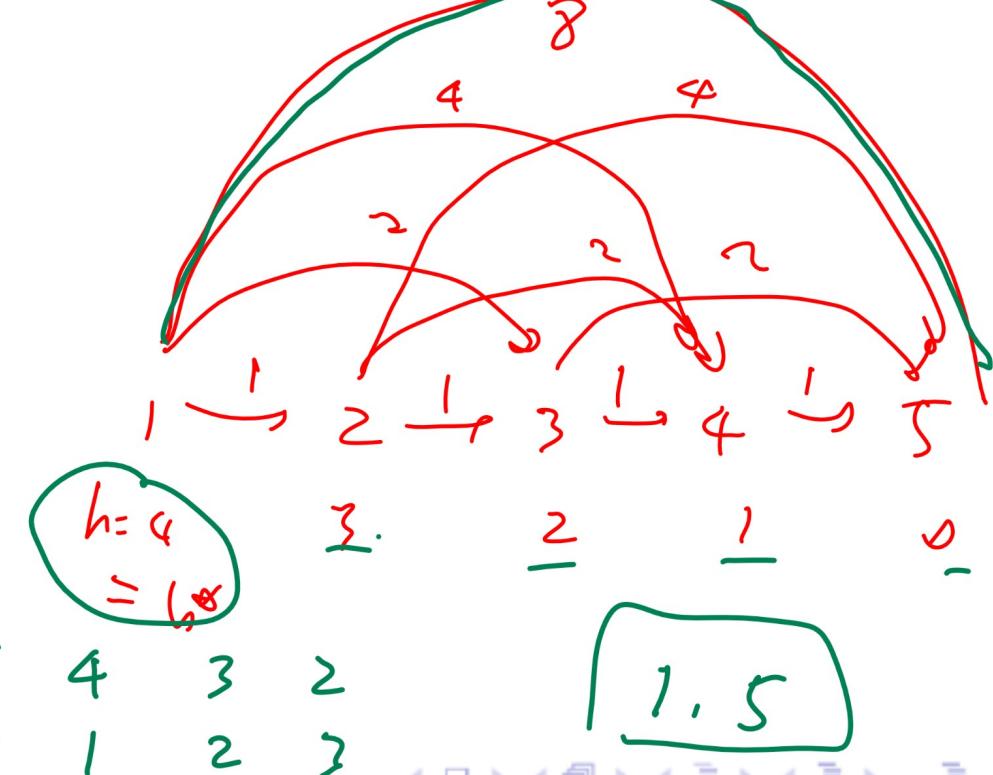
0

1

2

3

Quiz



Greedy Example 2 Diagram

Quiz

Best First Greedy Search

Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G , and the heuristic function $h(s), s \in V$.
- Output: a path from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $h(s)$ for $s \in Q$.

$$Q = I$$

- While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} h(s)$$

$$Q = Q + s' (s)$$

Best First Greedy Search Performance

Discussion

- Greedy is incomplete.
- Greedy is not optimal.

A Search

Description

expansion path = list of deleted popped \rightarrow always ends with Goal

- Expand the vertices with the lowest total cost $g(s) + h(s)$ first.
- Use a priority queue based on $g(s) + h(s)$.
- A stands for Always be optimistic?

A Search Example 1

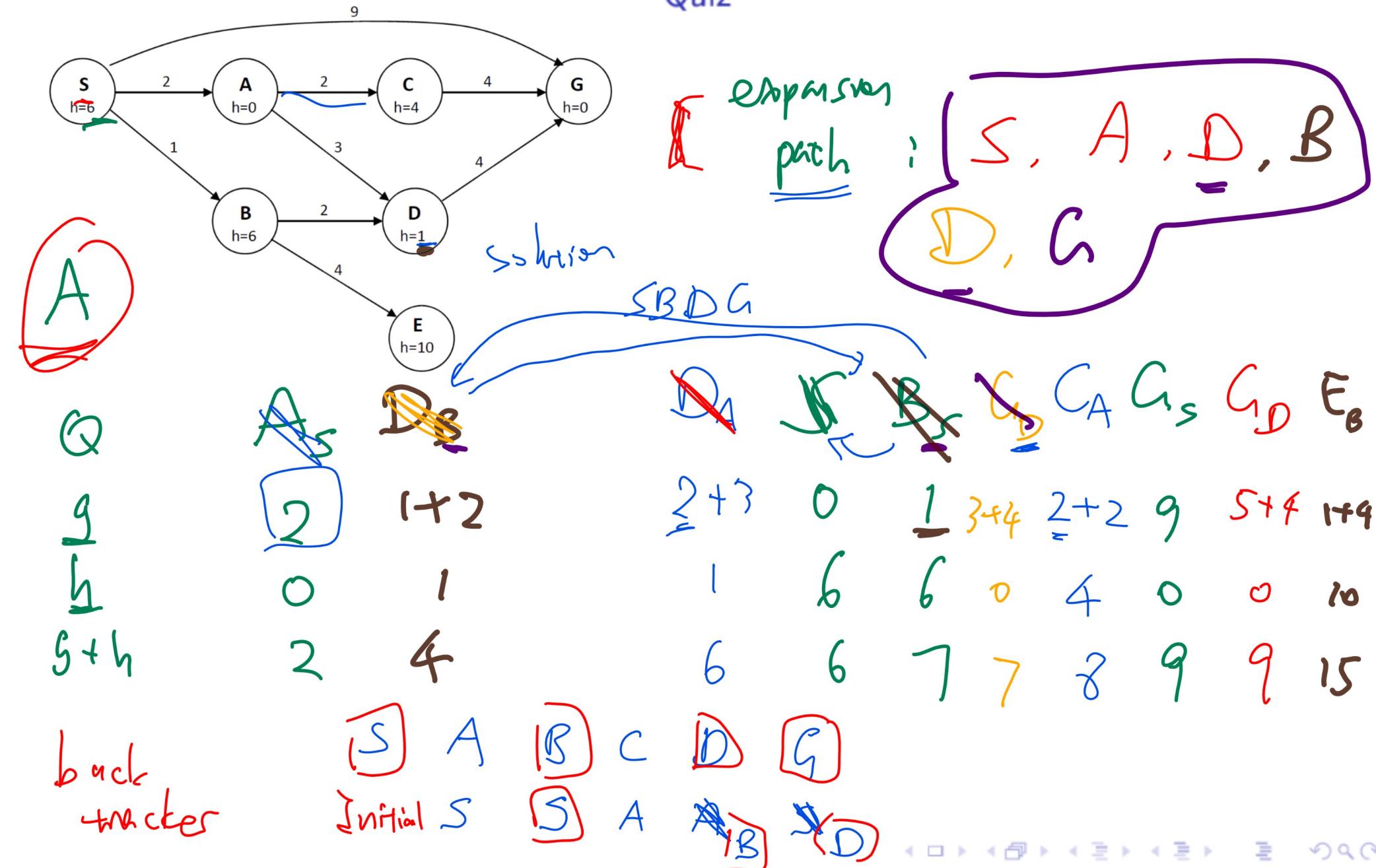
Quiz

- Given the following adjacency matrix. Find A Search expansion path.

-	S	A	B	C	D	E	G
S	$h = 6$	2	1	-	-	-	9
A	-	$h = 0$	-	2	3	-	-
B	-	-	$h = 6$	-	2	4	-
C	-	-	-	$h = 4$	-	-	4
D	-	-	-	-	$h = 1$	-	4
E	-	-	-	-	-	$h = 10$	-
G	-	-	-	-	-	-	$h = 0$

A Search Example 1 Diagram

Quiz



A Search Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G , and the heuristic function $h(s), s \in V$.
- Output: a path from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $g(s) + h(s)$ for $s \in Q$.

$$Q = I$$

- While Q is not empty and goal is not deQueued, deQueue Q and enQueue its successors.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$Q = Q + s'(s)$$

A Search Performance Discussion

- A is complete. ✓
- A is not optimal. ✗

A Star Search

Description

h

- A^* search is A search with an admissible heuristic.



Admissible Heuristic

Definition

- A heuristic is admissible if it never over estimates the true cost.

$$0 \leq h(s) \leq h^*(s)$$

A search $\underline{h} \Rightarrow A^*$ Optimal

Dominated Heuristic

Definition

- One heuristic, h_1 , is dominated by another, h_2 , if:

$$h_1(s) \leq h_2(s) \leq h^*(s), \forall s \in S$$

more better h_1

- If h_2 dominates h_1 , then h_2 is better than h_1 since A^* using h_1 expands at least as many states (or more) than A^* using h_2 .
- If h_2 dominated h_1 , A^* with h_2 is better informed than A^* with h_1 .

Admissible Heuristic 8 Puzzle Example

Quiz

$$C = 1$$

- Which ones (select multiple) of the following are admissible heuristic function for the 8 Puzzle? $h(\text{goal}) = 0$
- A: $h(s) = \text{number}$ of tiles in the wrong position.
- B: $h(s) = 0.$ ✓ UCS admissible
- C: $h(s) = 1.$ ~~$h(\text{goal}) = 0 < 1$~~ worst possible admissible $h.$
- D: $h(s) = \text{sum of Manhattan distance}$ between each tile and its goal location. ✓ $\checkmark h^M$ best
- E: $h(s) = \text{sum of Euclidean distance}$ between each tile and its goal location. ✓

Admissible Heuristic General Example 1

Quiz

$s \rightarrow$ Any config of the puzzle.

- Which ones (select multiple) of the following are admissible heuristic function?

- 0 < h(s) $\leq h^*$
- A: $h(s) = h^*(s)$. $\leq h^*$
 - B: $h(s) = \max\{2, h^*(s)\}$. ≥ 0
 - C: $h(s) = \min\{2, h^*(s)\} \leq h^*$
 - D: $h(s) = h^*(s) - 2 \leq -2$
 - E: $h(s) = \sqrt{h^*(s)}$.

$$h^* > 1$$

$$h^* \approx 0.25$$

$$\sqrt{h^*} < h^*$$

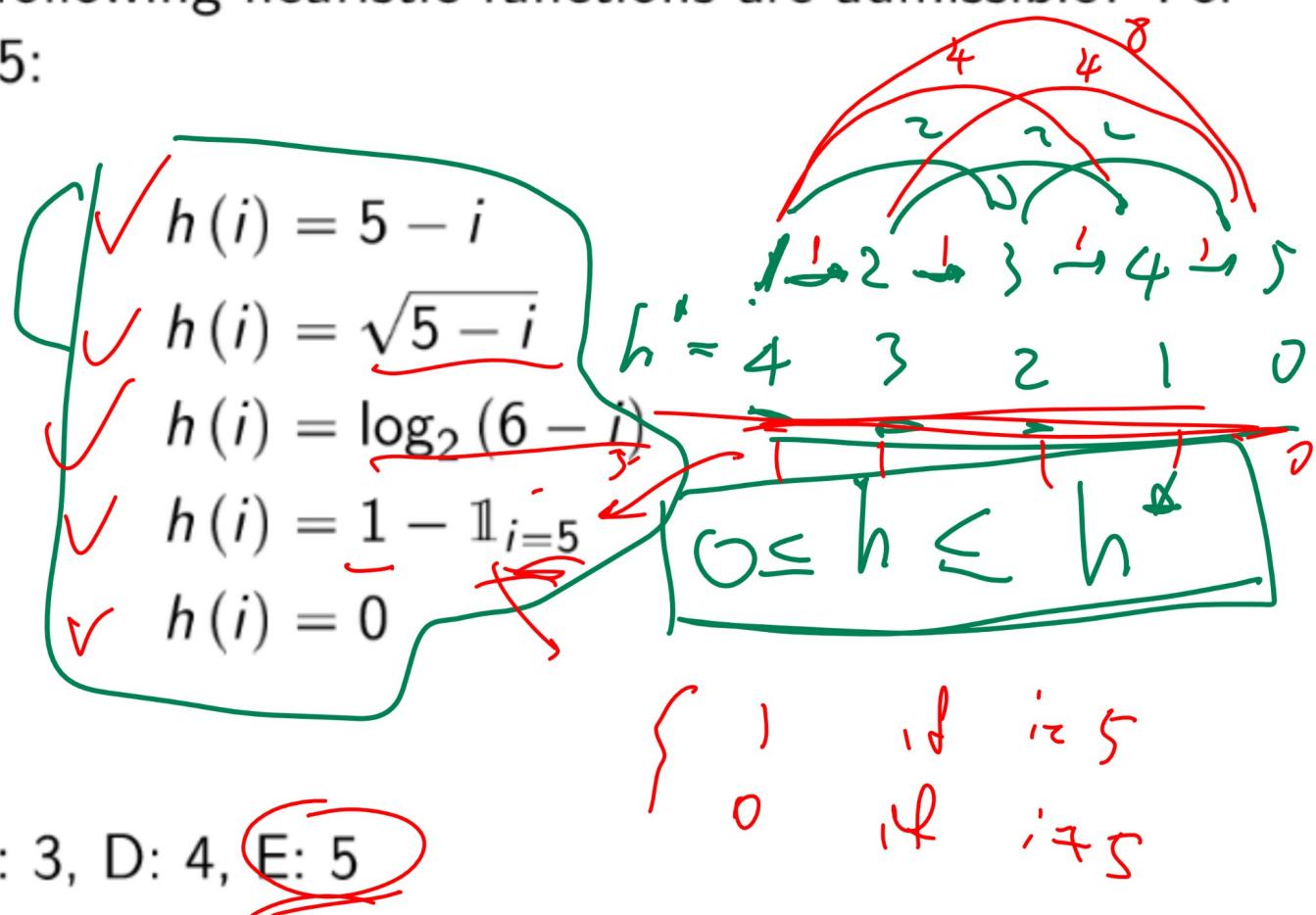
$$\sqrt{h^*} = h = 0, J > h^*$$

A Star Search Example 2

Quiz

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. How many of the following heuristic functions are admissible? For $i = 1, 2, 3, 4, 5$:

Q2



A Star Search Example 2

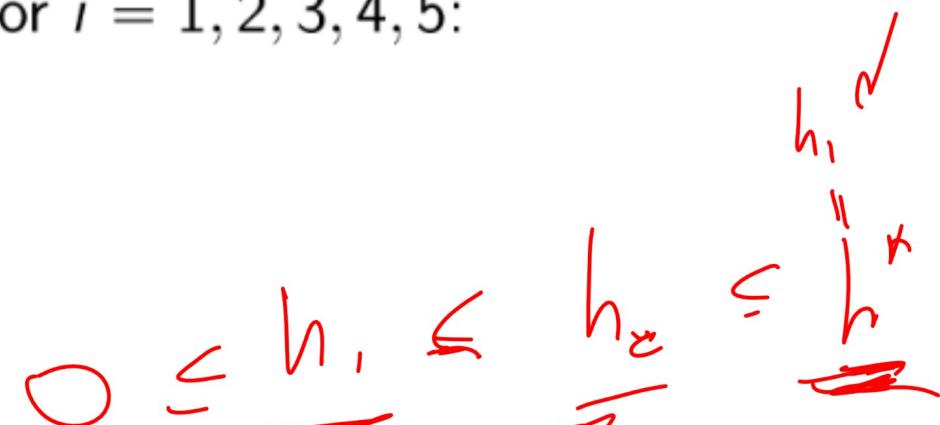
Quiz

(last)
Q3, A

- Given that cost from state i to j is 2^{j-i-1} for $j > i$. Which one of the following heuristic functions is not dominated (among the admissible ones)? For $i = 1, 2, 3, 4, 5$:

- A: $h_1(i) = 5 - i$
- B: $h_2(i) = \sqrt{5 - i}$
- C: $h_3(i) = \log_2(6 - i)$
- D: $h_4(i) = 1 - \mathbb{1}_{i=5}$
- E: $h_5(i) = 0$

\subseteq is dom by D,



h_1 dominates by h_2

h_2 dominates h_3

A Star Search with Revisit, Part I

Algorithm

- Input: a weighted digraph (V, E, c) , initial states I and goal states G , and the heuristic function $h(s), s \in V$.
- Output: a path with minimum cost from I to G .
- EnQueue initial states into a priority queue Q . Here, Q is ordered by $g(s) + h(s)$ for $s \in Q$.

$$Q = I$$

$$g(I) = 0$$

$$g(s) = \infty, \text{ for } s \notin I$$

- Initialize the list of visited vertices, P .

$$P = \emptyset$$

A Star Search with Revisit, Part II

Algorithm

- While Q is not empty and goal is not deQueued, deQueue Q , put it on P and enQueue its successors to Q , and update the cost functions.

$$s = Q_{(0)} = \arg \min_{s \in Q} g(s) + h(s)$$

$$P = P + s$$

$$Q = Q + s'(s), \text{ update } g(s') = \min \{g(s'), g(s) + c(s, s')\}$$

A Search Performance Discussion

- A^* is complete.
- A^* is optimal.

Iterative Deepening A Star Search

Discussion

- A^* can use a lot of memory.
- Do path checking without expanding any vertex with $g(s) + h(s) > 1$.
- Do path checking without expanding any vertex with $g(s) + h(s) > 2$.
- ...
- Do path checking without expanding any vertex with $g(s) + h(s) > d$.

Iterative Deepening A Star Search Performance

Discussion

- IDA* is complete.
- IDA* is optimal.
- IDA* is more costly than A^* .

Beam Search

Discussion

- Version 1: Keep a priority queue with fixed size k . Only keep the top k vertices and discard the rest.
- Version 2: Only keep the vertices that are at most ε worse than the best vertex in the queue. ε is called the beam width.

Beam Search Performance

Discussion

- Beam is incomplete.
- Beam is not optimal.