

CS540 Introduction to Artificial Intelligence

Lecture 22

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Based on lecture slides by Jerry Zhu and Yingyu Liang

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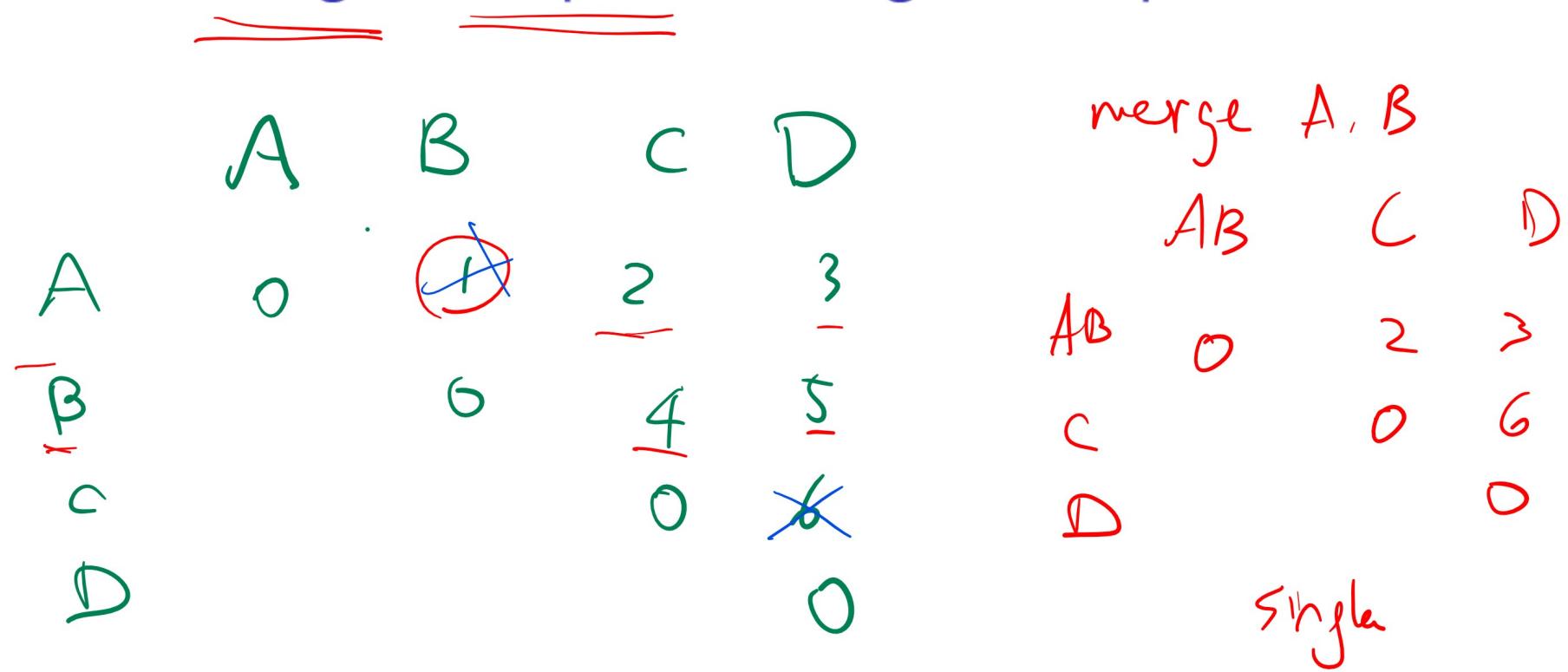
Coverage

- Unsupervised Learning
- Search

Unsupervised Learning

- Hierarchical Clustering
- K Means Clustering (no Gaussian mixture)

Single Complete Linkage Example



single
linkage
dist

$$\begin{aligned}
 & \text{single linkage dist} \\
 & (\underline{A, B}) \text{ and } (C, D) = \min(2, 3, 4, 5) \\
 & = 2
 \end{aligned}$$

complete

$$\begin{aligned}
 & \text{complete} \\
 & \text{---} \quad \text{---} : \max\{2, 3, 4, 5\} \\
 & = 5
 \end{aligned}$$

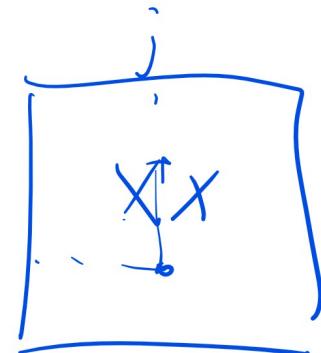
Principal Component Analysis

- Basic Linear (Matrix) Algebra
- Compute Projection and Variance (use formula on the formula sheet)
- Feature Reconstruction

Matrix Multiplication Example

i,j entry of $\frac{X^T X}{A}$

$$e_i = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{pmatrix} \quad i \text{ position}$$



$$e_i^T X^T X e_j$$

$$e = \begin{pmatrix} 1 \\ 1 \\ \vdots \\ 1 \end{pmatrix} \quad \text{all } 1's$$

$$\left(\begin{matrix} 0 & 0 & \dots & 0 & 0 \\ \vdots & \vdots & & \vdots & \vdots \\ 0 & 0 & \dots & 0 & 0 \end{matrix} \right) \left(\begin{matrix} | & | & | & | & | \\ e_i & e_2 & e_3 & \dots & e_j & \dots & e_n \end{matrix} \right) \left(\begin{matrix} 0 \\ 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{matrix} \right)$$

i-th row

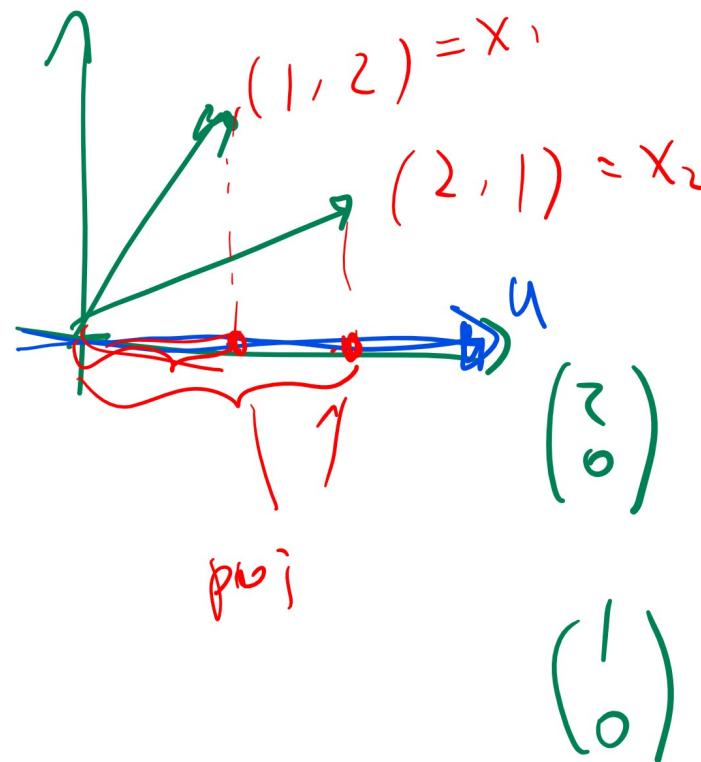
j-th entry of i-th row

$$\text{sum of row } i - e_i^T X^T X e$$

i,j entry

Sum of every + by
 $e^T X^T X e$

Projected Variance Example



variance of magnitude of projection

$$u = \frac{x_1^T x_1}{x_1^T x_1 + x_2^T x_2} \cdot \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \frac{2 \cdot 1 + 0 \cdot 2}{2 \cdot 2 + 0 \cdot 0} \cdot \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$u^T x_1 u = 1 \cdot \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

with $x = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$ $u^T x u = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$

mag of proj 2, 1

Variance: MLE, $\frac{1}{n} \sum (x_i - \hat{M})^2$ $\frac{1}{2} (0.5^2 + 0.5^2)$
 unbiased, $\frac{1}{n-1} \sum (x_i - \hat{M})^2$ $\frac{1}{1} (0.5^2 + 0.5^2)$

Search

- Uniformed: no heuristic
- Informed: heuristic
- Local: optimization
- Adverserial: sequential move game
- Equilibrium: simultaneous move game

Uninformed Search

- BFS
- DFS
- IDS
- UCS

Informed Search

- Greedy
- A (or A Star)

Uniformed Search Counting Example

Stages are 1 to 1024 = 2^{10}

of vertices expanded | history

1
1024

BFS : 1024 change to 1025

DFS: 11 / BFS: 1025

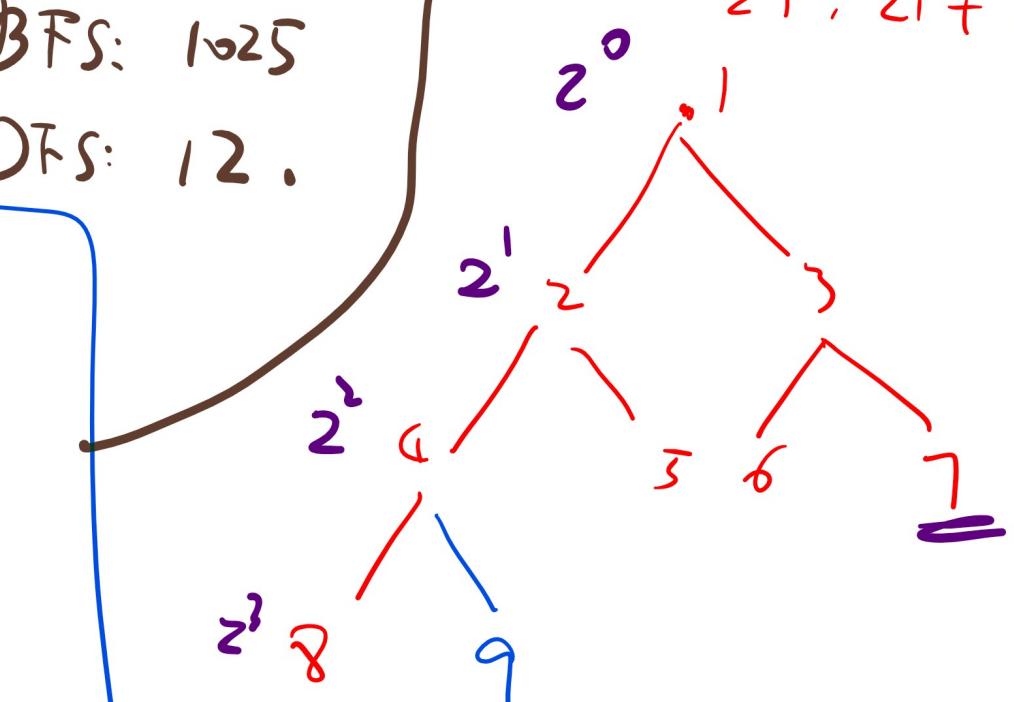
Successor of i , j

z_i^j, z_{i+1}^j

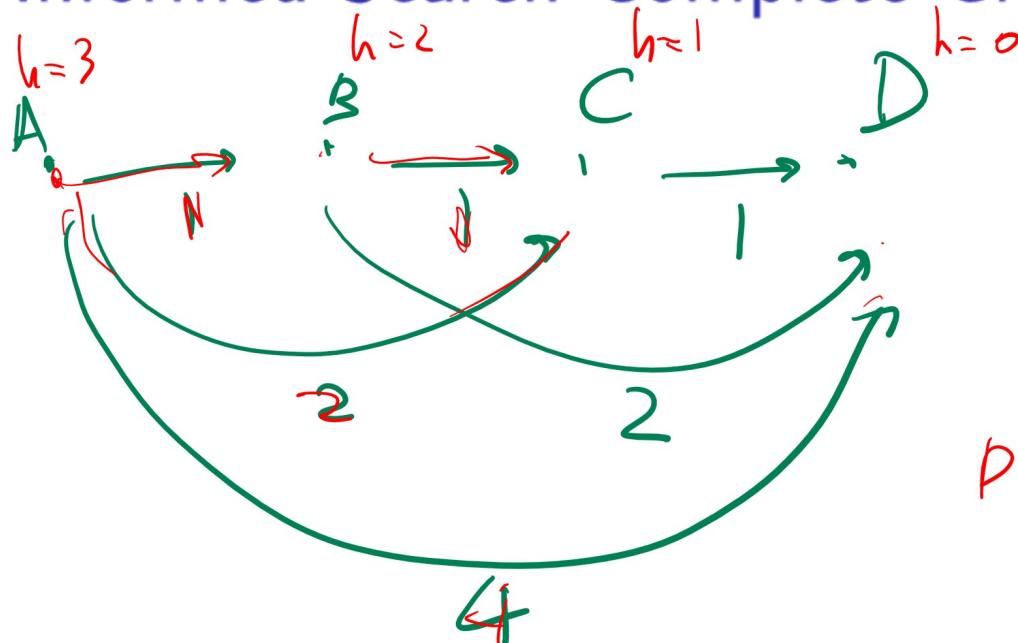
change 1024 to 1023

BFS; 1023

DFS: 1023,



Informed Search Complete Graph Example



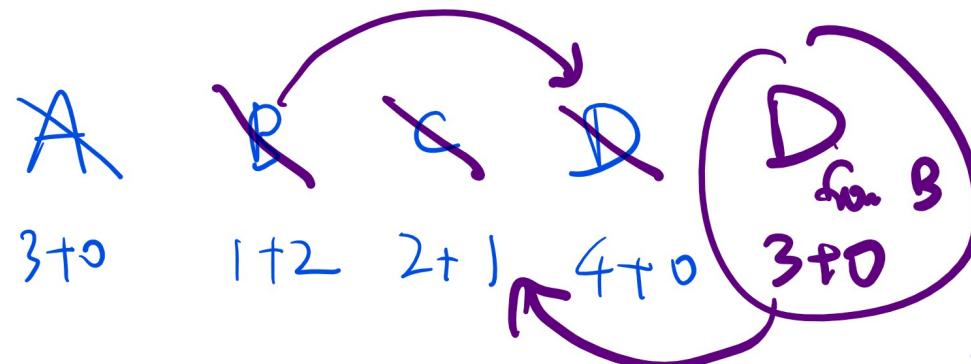
Expansion path
A, B, C, D

UCS A, B, C, D

Greedy
A* A D

A, B, C, D

PQ:



cost = cost of parent + cost between parent and this state,

Local Search

- Hill Climbing
- Simulated Annealing
- Genetic Algorithm

SAT Local Search Example

fitness score = # of clauses satisfied

$$\begin{array}{c}
 (x \vee y) \wedge (x \vee z) \wedge (y \vee z) \\
 \hline
 \text{clause 1} \qquad \qquad \qquad \text{2} \qquad \qquad \qquad \text{3}
 \end{array}$$

initial $s_0 = x, y, z = \text{False}$. $f(s_0) = 0$

→ find a random neighbor. $s_1 = (x = T, y, z = F)$

$$f(s_1) = 2 > 0$$

$$\text{prob}_{SA}(\text{move to } s_1) = 1$$

→ find another random neighbor. $s_2 = (x, y, z = F)$

$$f(s_2) = 0 < 2$$

$$\text{prob}_{SA}(\text{move to } s_2) \stackrel{\text{assume}}{=}$$

$$= e^{-\frac{|f(s_2) - f(s_1)|}{\text{temp}}} = e^{-\frac{2}{\text{temp}}}$$

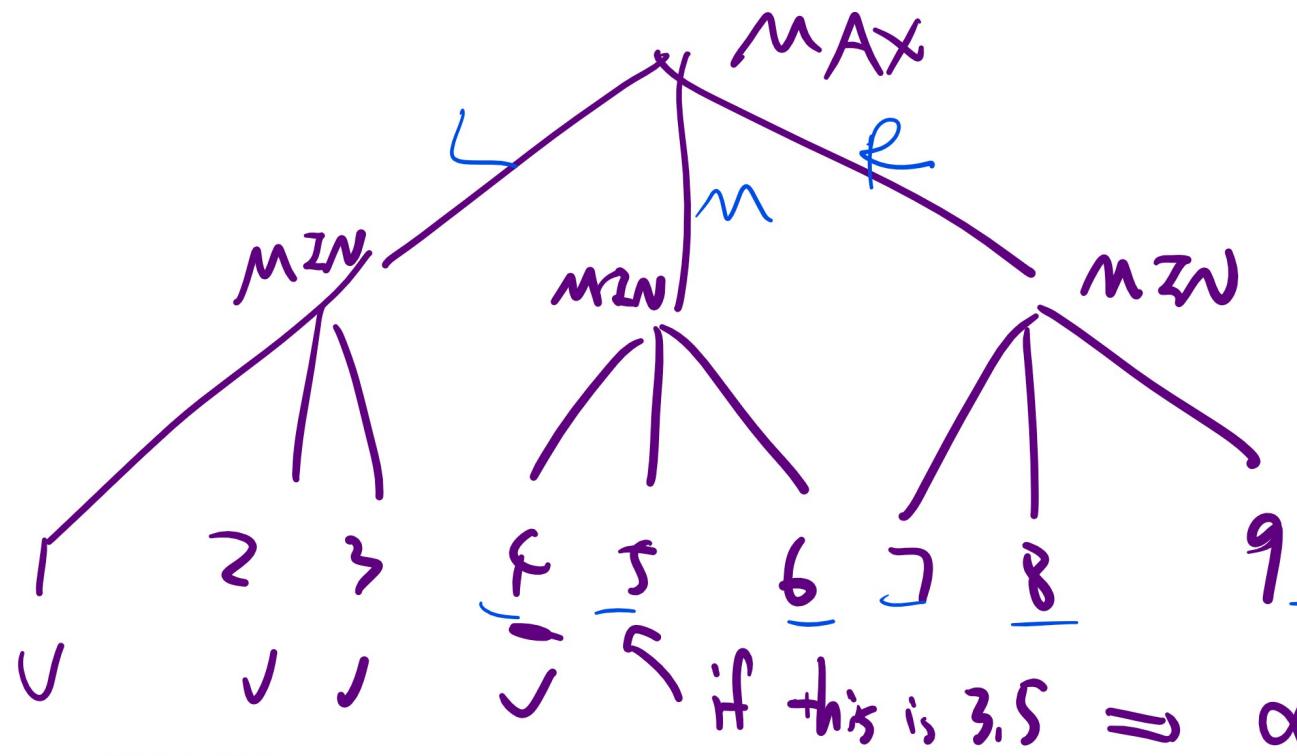
$$\text{Temp}_0 = 1$$

$$\text{Temp}_{t+1} = 0.9 \text{ Temp}_t$$

~~Adversarial Search~~

- Minimax
- Alpha Beta

Alpha Beta Max Pruning

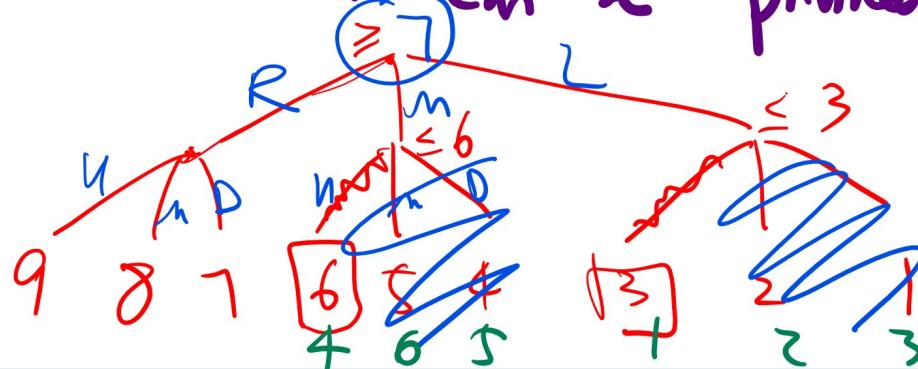


Flat pruned states = 0

DFS, leftmost first

Cannot prune.

Max # of states can be pruned after re order branches,



Max # = 4,

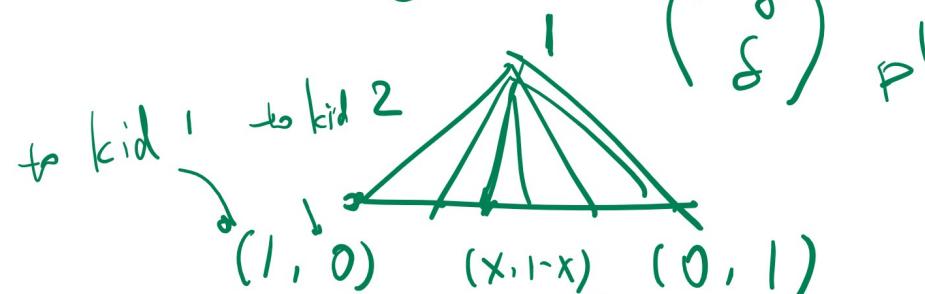
Equilibrium Search

- IESDS
- Best Responses (Nash)
- Fixed Point (Nash)

Bargaining Example (see Lee 21)

2 kids try to divide cake

each time they disagree dad eats δ portion of cake



player 1

propose $1 - \delta$

player 2

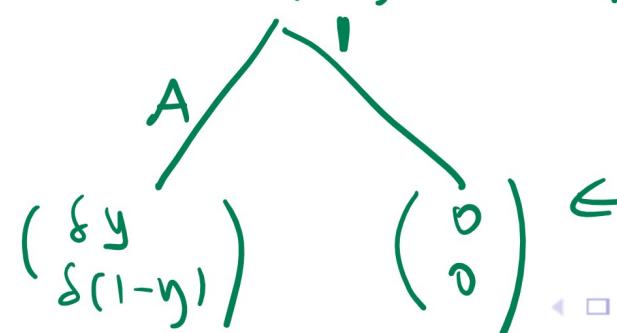
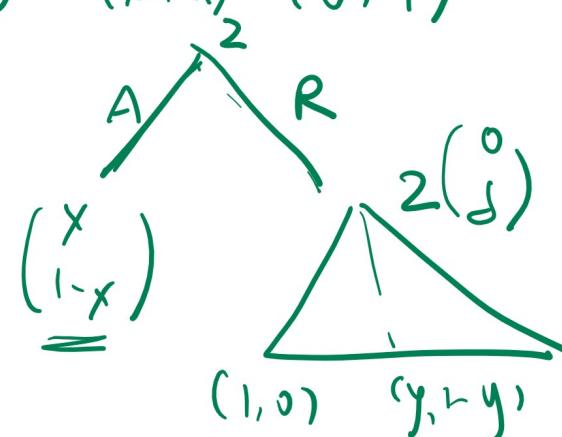
accept iff

$$1 - x \geq \delta$$

propose $(0, 1)$
 $(0, \delta)$

player 1

accept iff $\delta y \geq 0$
 $y \geq 0$



assumption

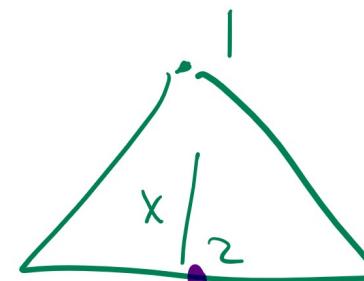
$N=2$

$$\begin{pmatrix} 1-\delta \\ \delta \end{pmatrix}$$

$N=\infty$

P1 propose

$$x = 1 - \delta + \delta^2 z$$

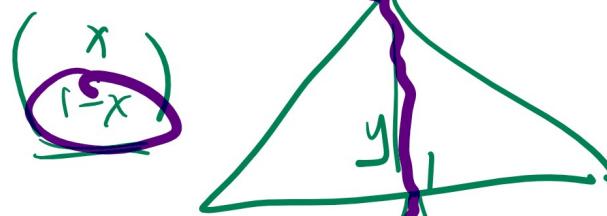


← size = 1

P2 accept iff $1-x \geq \delta(1-\delta z)$

P2 propose $y = \delta z$

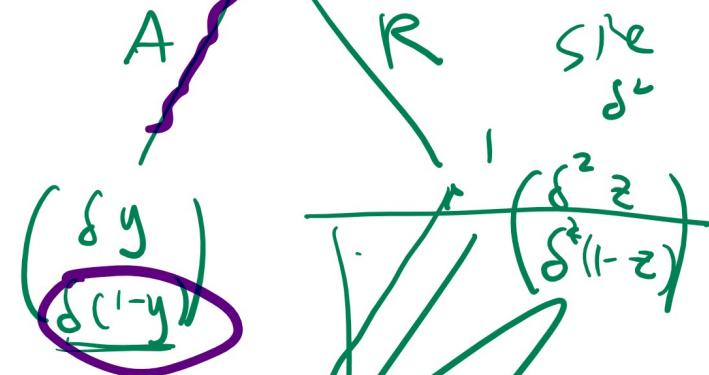
$$(1-x)$$



← size = δ

P1 accept iff $\delta y \geq \delta^2 z$

$$x < 1 - \delta + \delta^2 x$$



size
 δ^2

$$x = \frac{1}{1-\delta}$$

$N-2$
period
game.

$$(\delta y)$$

$$(\delta^2 z)$$



$$N=2, \quad \begin{pmatrix} 1-\delta \\ \delta \end{pmatrix}$$

$$N=4, \quad \begin{pmatrix} 1 - \delta + \delta^2 \begin{pmatrix} 1-\delta \\ 1-\delta \end{pmatrix} \\ \delta - \delta^2(1-\delta) \end{pmatrix}$$

$$\begin{pmatrix} 1 - \delta + \delta^2 - \delta^3 \\ \delta - \delta^2 + \delta^3 \end{pmatrix}$$

$$N=6, \quad \begin{pmatrix} 1 - \delta + \delta^2 - \delta^3 + \delta^4 - \delta^5 \\ \delta - \delta^2 + \delta^3 - \delta^4 + \delta^5 \end{pmatrix}$$

:

$$N=\infty \quad \left((1-\delta)(1+\delta^2+\delta^4+\delta^6+\dots) \right)$$

$$= \left((1-\delta) \cdot \frac{1}{1-\delta^2} \right) = \left(\begin{array}{c} 1 \\ \hline 1+\delta \\ \hline \delta \end{array} \right)$$

$$1+x+x^2+\dots = \frac{1}{1-x}, \quad x = \delta^2$$