

Chapter 15

Normalization for Relational Database

Data Normalization

- Primarily a tool to validate and improve a logical design so that it satisfies certain constraints that **avoid unnecessary duplication of data**
- The process of decomposing relations with anomalies to produce **smaller, well-structured relations**

Well-Structured Relations

- A relation that contains minimal data redundancy and allows users to insert, delete, and update rows without causing data inconsistencies
- Goal is to avoid anomalies
 - **Insertion Anomaly**—adding new rows forces user to create duplicate data
 - **Deletion Anomaly**—deleting rows may cause a loss of data that would be needed for other future rows
 - **Modification Anomaly**—changing data in a row forces changes to other rows because of duplication

General rule of thumb: A table should not pertain to more than one entity type

Example

<u>EmpID</u>	Name	Salary	<u>Course#</u>	CourseTitle	Date
100	Alaa	32000	459	SPSS	9/9/2016
100	Alaa	32000	876	Surveys	7/8/2016
140	Atheer	40000	333	Visual Basic	1/1/2016
150	Aisha	23000	459	SPSS	9/9/2016
150	Aisha	23000	901	C++	12/8/2016
140	Atheer	40000	901	C++	12/8/2016

Is this a relation?

Yes: Unique rows and no multivalued attributes

What's the primary key?

Composite: Emp_ID, Course#

Anomalies in this Table

Insertion—can't enter a new employee without having the employee take a class. In addition, adding course for existing employee duplicate employees' data and course data

Deletion—if we remove employee 140, we lose information about the existence of a Visual Basic class

Modification—giving a salary increase to employee 100 forces us to update multiple records

Why do these anomalies exist?

Because there are two themes (entity types) in this one relation. This results in data duplication and an unnecessary dependency between the entities

Functional Dependency

- Functional dependency, denoted by $X \longrightarrow Y$ between two sets of attributes X, Y means that value of Y is determined by the value of X .
- The value of the X of a tuple uniquely (or **functionally**) *determine the value of Y*
- Y is functionally dependent on X

Functional Dependency (2)

- $Ssn \rightarrow Ename$
- The value of an employee's Social security number uniquely determines the employee name
- $Pnumber \rightarrow \{Pname, Plocation\}$
- The value of a project's number (Pnumber) uniquely determines the project name and location
- $\{Ssn, Pnumber\} \rightarrow Hours$
- A combination of Ssn and Pnumber values uniquely determines the number of hours the employee currently works on specific project

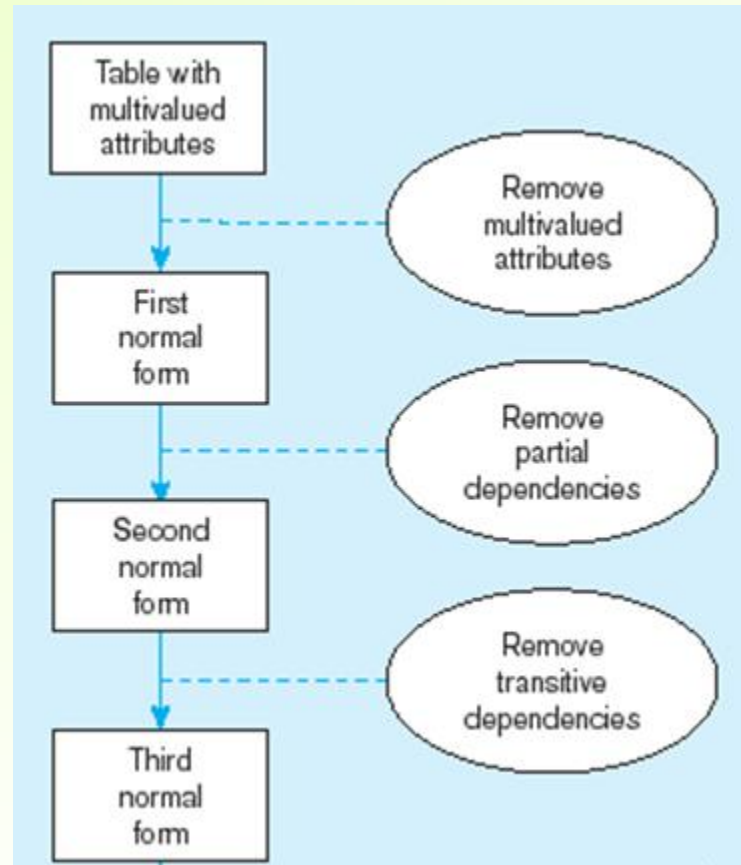
Normal Forms Based on Primary Key

- Most practical relational design projects take one of the following two approaches:
 - Perform a conceptual schema design using a conceptual model such as ER or EER and map the conceptual design into a set of relations
 - Design the relations based on external knowledge derived from an existing implementation of files or forms or reports

Normalization of Relations

- The normalization process takes a relation schema through a series of tests to certify whether it satisfies a certain **normal form**.
- There are three normal forms, which are *first, second, and third normal form*
- It is a purifying process that makes the design have better quality and minimizes redundancy

Normalization Steps




First Normal Form

- For a relation to be in 1st normal form, redundant groups or multivalued attributes should be removed
- To change to 1NF:
 - ▶ Remove nested relation attributes into a new relation
 - ▶ Propagate the primary key into it

Removing multivalued attributes

(a)

DEPARTMENT

Dname	<u>Dnumber</u>	Dmgr_ssn	Dlocations
			

(b)

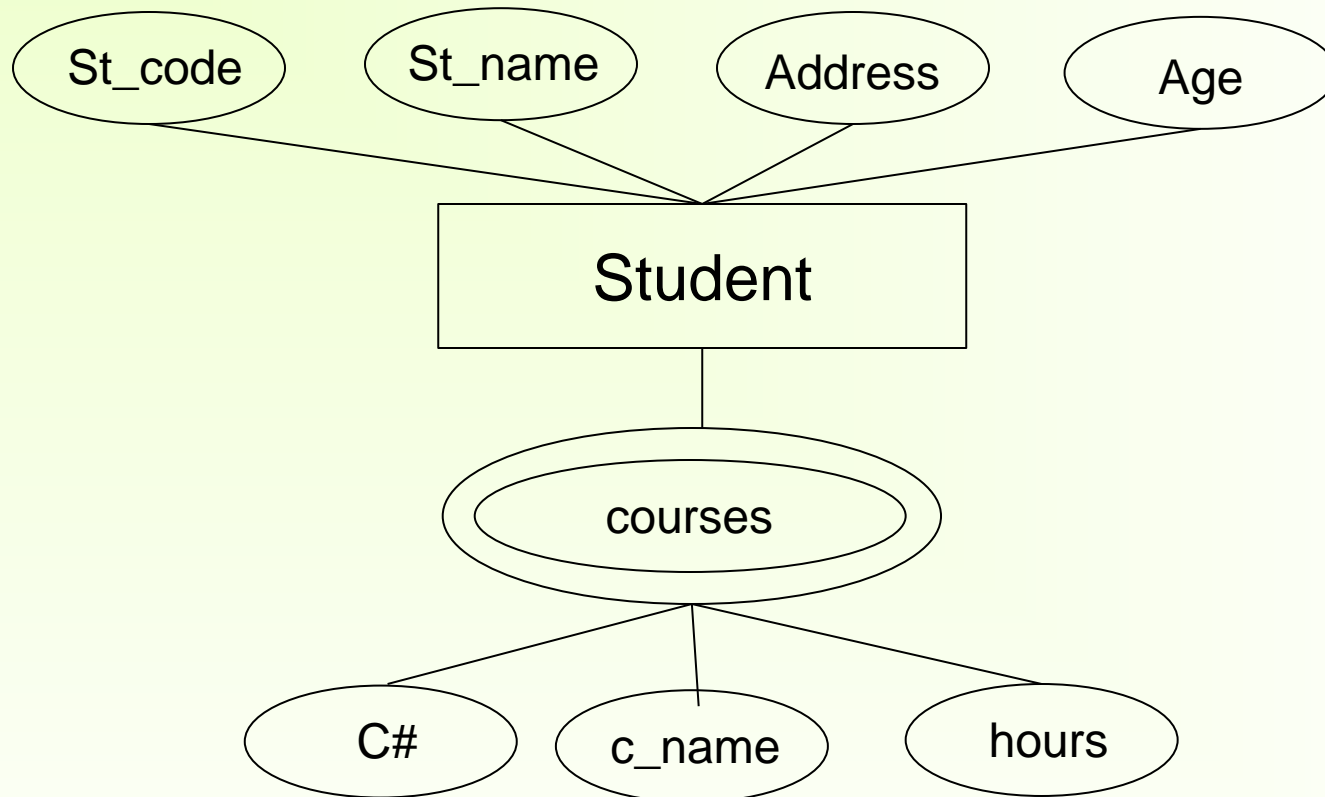
DEPARTMENT

Dname	<u>Dnumber</u>	Dmgr_ssn	Dlocations
Research	5	333445555	{Bellaire, Sugarland, Houston}
Administration	4	987654321	{Stafford}
Headquarters	1	888665555	{Houston}

New relation to be in 1st normal form

<u>Dnumber</u>	<u>Dlocations</u>
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Removing Relation within Relation or Repeating Groups



Example Repeating Groups

<u>St_code</u>	St_name	Address	Age	<u>C#</u>	C_name	Hours
20111	Karim	AlRawda	18			
				DS34	Data Str.	3
				OR23	Ope. Rese	3
				DB12	Database	3
20112	Mona	Alfayhaa	17			
				N22	Network	2
				OR23	Ope. Rese	3
				Db12	Database	3
				DS34	Data Str.	3

Example Repeating Groups

<u>Order_ID</u>	Order_ Date	Customer_ ID	Customer_ Name	Customer_ Address	<u>Product_ID</u>	Product_ Description	Product_ Finish	Unit_ Price	Ordered_ Quantity
1006	10/24/2006	2	Value Furniture	Plano, TX	7	Dining Table	Natural Ash	800.00	2
1006	10/24/2006	2	Value Furniture	Plano, TX	5	Writer's Desk	Cherry	325.00	2
1006	10/24/2006	2	Value Furniture	Plano, TX	4	Entertainment Center	Natural Maple	650.00	1
1007	10/25/2006	6	Furniture Gallery	Boulder, CO	11	4-Dr Dresser	Oak	500.00	4
1007	10/25/2006	6	Furniture Gallery	Boulder, CO	4	Entertainment Center	Natural Maple	650.00	3

To Be in First Normal Form

The previous data to be in 1st normal form, split the repeating group in another table with the primary key of the relation to be as foreign key in the new relation

<u>St_code</u>	St_name	Address	Age
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<u>St_code</u>	<u>C#</u>	C_name	Hours
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Second Normal Form

- Based on concept of full functional dependency
 - ▶ Versus **partial dependency**- Remove partial dependency

Definition. A relation schema R is in 2NF if every nonprime attribute A in R is *fully functionally dependent* on the primary key of R .

- Second normalize into a number of 2NF relations
 - ▶ Nonprime attributes are associated only with part of primary key on which they are fully functionally dependent

Second Normal Form (2)

<u>SSN</u>	<u>P#</u>	Hours	Ename	Pname	Plocation
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- Ename is functionally dependent on SSN only, no need for P#
- Pname and Plocation are functionally dependent on P#
- So this relation **is not in the second normal form**
- **Solution:**

<u>SSN</u>	<u>P#</u>	Hours
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<u>SSN</u>	Ename
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<u>P#</u>	Pname	Plocation
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Third Normal Form

- To be in 3rd normal form you should remove **transitive dependency**
- Non key attribute is dependent on non key attribute

Third Normal Form (2)

<u>SSN</u>	Ename	Address	Dept_code	Dept_name	Mgr_SSN
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- Dept_name and mgr_SSN are functionally dependent on dept_code which is non key attribute
- **Solution:**

<u>SSN</u>	Ename	Address	<i>Dept_code</i>
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<u>Dept_code</u>	Dept_name	Mgr_SSN
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Exercises

Consider the following report, suppose sales order number with item ordered is the primary key, normalize this report to reach better design

Sales Order

Fiction Company
202 N. Main
Mahattan, KS 66502

CustomerNumber: 1001
Customer Name: ABC Company
Customer Address: 100 Points
Manhattan, KS 66502

Sales Order Number: 405
Sales Order Date: 2/1/2000
Clerk Number: 210
Clerk Name: Martin Lawrence

Item Ordered	Description	Quantity	Unit Price	Total
800	widgit small	40	60.00	2,400.00
801	tingimajigger	20	20.00	400.00
805	thingibob	10	100.00	1,000.00
Order Total				3,800.00

Relation

R = salesOrderNo, salesOrderDate, custNo, custName, address clerkNo, clerkName, {itemsOrdered, description, quantity, unitPrice}

To be in 1st normal form, remove repeating group:

R1 = salesOrderNo, salesOrderDate, custNo, custName, address clerkNo, clerkName

R2 = salesOrderNo, itemsOrdered, description, quantity, unitPrice

2nd Normal Form

R1= salesOrderNo, salesOrderDate, custNo,
custName, address clerkNo, clerkName

R2.1= salesOrderNo, itemsOrdered, quantity

R2.2= itemsOrdered, description, unitPrice

3rd Normal Form

R1.1= salesOrderNo, salesOrderDate, custNo, clerkNo

R1.2= custNo, custName, address

R1.3= clerkNo, clerkName

R2.1= salesOrderNo, itemsOrdered, quantity

R2.2= itemsOrdered, description, unitPrice

Exercises

- Normalize the following schemas into 3rd normal form:
- **BRANCH** (Branch#, Branch_Addr, {ISBN, Title, Author, Publisher, Num_copies})

1st Normal Form:

R1: Branch#, Branch_addr

R2: Branch#, ISBN, title, author, publisher, num_copies

2nd Normal Form:

R1: Branch#, Branch_addr

R2.1: Branch#, ISBN, num_copies

R2.2: ISBN, title, author, publisher

3rd Normal Form:

No change in the previous relations.

Exercises

- Project code, project title, project manager, project budget {employeeNo, employeename, completed_hour, departmentNo, department_name, rate_per_hour}
- **Note:** rate per hour for each employee is fixed regardless of the project. Completed hour means the number of hours employee accomplished in this project

■ 1st Normal Form:

R1: Pcode, ptitle, pmgr, pbudget

R2: Pcode, empNo, empName, c_hours,
dNo, dName, rate

■ 2nd Normal Form:

R1: Pcode, ptitle, pmgr, pbudget

R2.1: Pcode, empNo, c_hours

R2.2: empNo, empName, dNo, dName, rate

■ 3rd Normal Form:

R1: Pcode, ptitle, pmgr, pbudget

R2.1: Pcode, empNo, c_hours

R2.2.1: empNo, empName, dNo, rate

R2.2.2: dNo, dName

Exercises

Al Salam Hospital - Doctor's report Date: 10/5/2004

Doctor Id. : A121

Doctor Name : Dr. Ahmed

Department Id : A

Department Name : Internal Diseases

P#	Pat-name	Address	Given Treatments			
			Item#	Description	Quantity	Unit Price
10	Saleh	Maadi	A01	Aspirin	10	1.5
			A03	Panadol	6	3.5
			B01	Vitamin C	12	4.0

Relation

$R = \underline{\text{docId}}, \text{docname}, \text{deptId}, \text{deptName}, \{ \text{patient}(\underline{\text{p\#}}, \text{pname}, \text{address}) \{ \text{given treatment}(\underline{\text{item\#}}, \text{description}, \text{quantity}, \text{unit_price}) \} \}$

First Normal Form

- $R1 = \underline{\text{docId}}, \text{docname}, \text{deptId}, \text{deptName}$
- $R2 = \underline{\text{docId}}, \underline{\text{p\#}}, \text{pname}, \text{address}$
- $R3 = \underline{\text{docId}}, \underline{\text{p\#}}, \underline{\text{item\#}}, \text{description}, \text{quantity}, \text{unit_price}$

Second Normal Form

- R1= docId, docname, deptId, deptName
- R2.1= docId, p#
- R2.2= p#, pname, address
- R3= docId, p#, item#, qantity
- R3.2= item#, description, unit_price

Third Normal Form

- R1.1= docId, docname, deptId
- R1.2=deptId, deptName
- R2.2= p#, pname, address
- R3=docId, p#, item#, qantity
- R3.2= item#, description, unit_price