

# ARP and ICMP Protocols

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# References

- ARP

<https://info.support.huawei.com/info-finder/encyclopedia/en/ARP.html>

- ICMP

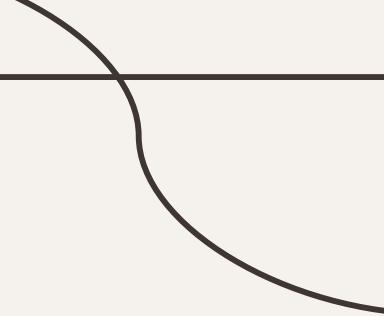
<https://info.support.huawei.com/info-finder/encyclopedia/en/ICMP.html>

# Introduction

- ❑ **Address Resolution Protocol (ARP)** is used to discover the **Layer 2 (MAC)** address of a device on your **local network** when you only know its **Layer 3 (IP) address**.
- ❑ **Internet Control Message Protocol (ICMP)** is used for **sending error** and **control messages between network devices**. It's a fundamental support protocol that **helps routers and hosts communicate about** the **health** and **status** of the network, making it possible to **diagnose problems and ensure efficient operation**.

# Comparison: ARP vs ICMP

Feature	ARP	ICMP
Layer	2 / 3	3
Purpose	Map IP → MAC	Control, errors, diagnostics
Works in LAN only?	Yes	No (Internet-wide)
Key message	Request / Reply	Echo, Time Exceeded, Unreachable
Security	Weak	Often blocked by firewalls



# ARP

# Why Do We Need ARP?

- ❑ ARP is required to map IP addresses to MAC addresses.
- ❑ For hosts or other Layer 3 network devices to communicate on a LAN, the sender must know the destination IP address to which it will send IP packets.
- ❑ Sender must first be encapsulated with MAC addresses before they can be transmitted over the physical network.
- ❑ It is therefore necessary for hosts or Layer 3 network devices to maintain an ARP table for storing the mapping information of IP and MAC addresses.

# What Are the Types of ARP?

## 1. Dynamic ARP

- It **entries** are **automatically** generated and maintained when **ARP packets** are **sent** and **received**.
- **They can be aged, updated, or overwritten by Static ARP entries.**
- **Dynamic ARP applies to complex networks that transmit delay-sensitive services.**

# What Are the Types of ARP?

## 2. Static ARP

- It allows a network administrator to manually create the fixed mappings between **IP** and **MAC** addresses.
- **Static ARP entries cannot be aged or overwritten by dynamic ARP entries, ensuring system security.**
- **In most cases**, devices on a network can use **ARP** to **dynamically** learn **ARP entries** and **age** or **update** the **generated dynamic ARP entries**.
- **However**, when a network **encounters** an **ARP attack**, the **dynamic ARP entries** may be **incorrectly updated** or **aged**.
- **As a result**, the communication between authorized users becomes **abnormal**.

# What Are the Types of ARP?

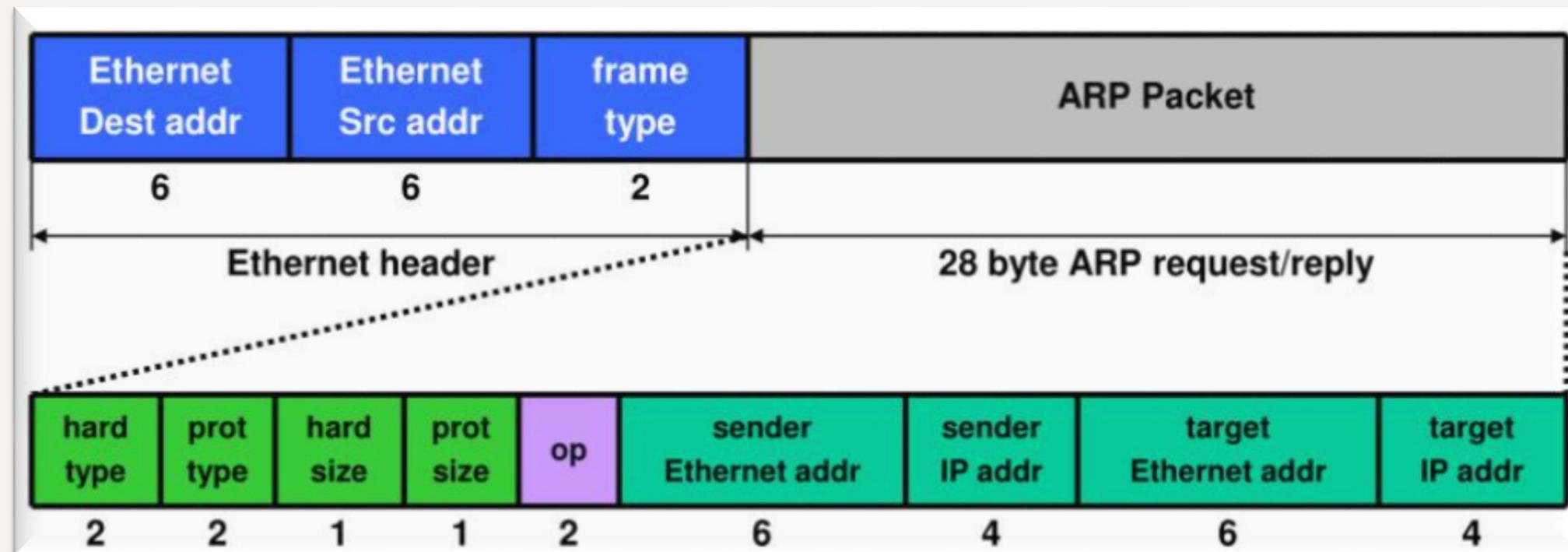
- If a **static ARP** entry is **configured** on a **device**, the **device** can communicate with the **peer device** using **only** the **specified MAC address**.
- **Network attackers** cannot modify the **mapping** between the **IP** and **MAC addresses** using **ARP packets**, **ensuring** communication between the two devices.
- **Static ARP entries are generally configured on gateways.**

# What Are the Types of ARP?

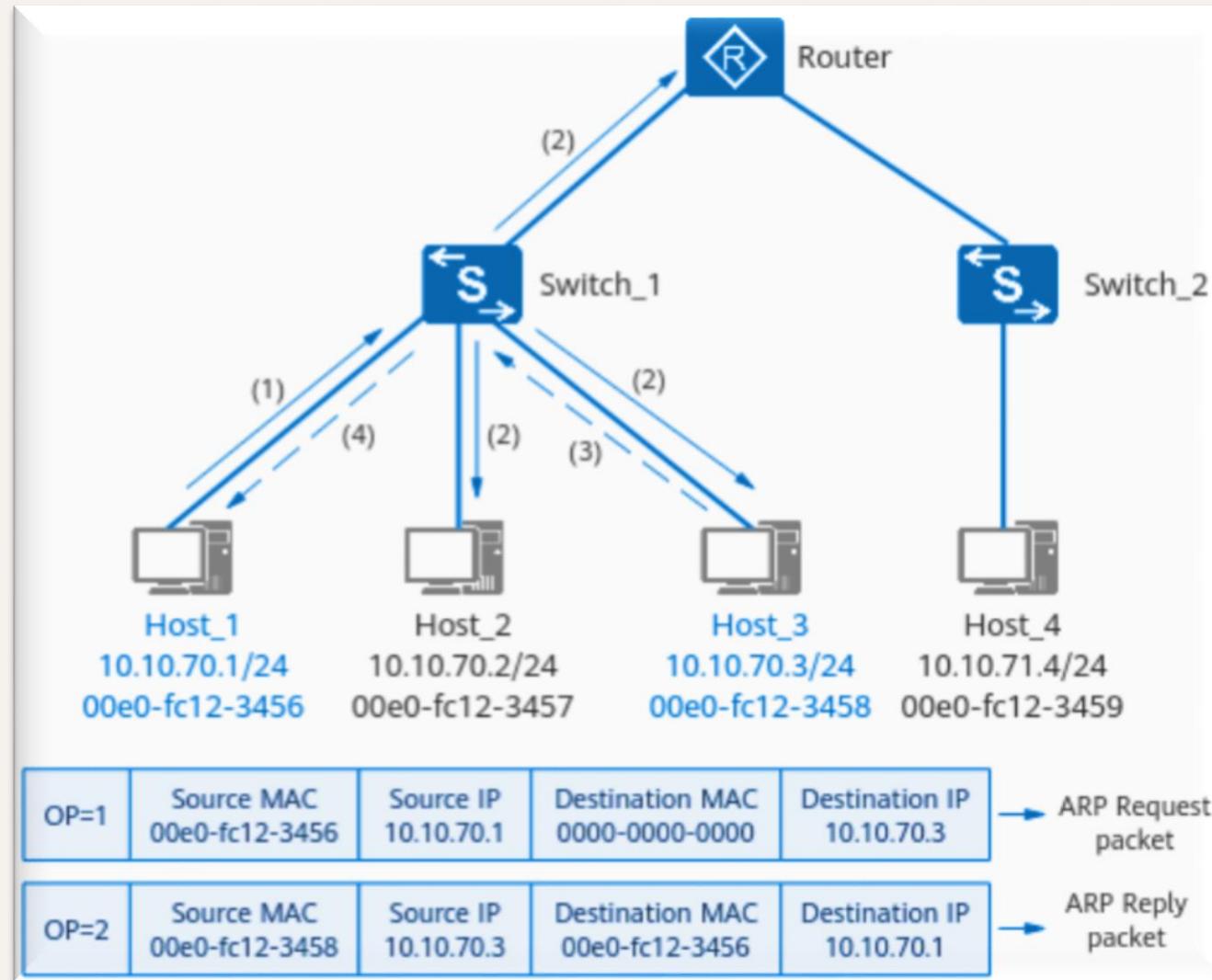
## 3. Gratuitous ARP

- It allows a device to send an **ARP Request** packet that **carries its own IP address** as the **destination IP address**.
- **Gratuitous ARP has the following functions:**
  - **Checks for IP address conflicts.**
  - **Advertises a new MAC address.**

# ARP Format



# How Does ARP Work?



# ARP Operation (Example)

## □ Assume:

- **Host A: 192.168.1.10**
- **Host B: 192.168.1.20**
- **Host A** wants to send packet to **Host B**.

## □ Step 1: ARP Request (Broadcast)

- **Host A** doesn't know **B's MAC**, so it sends **broadcast frame**:
- **Destination MAC: FF:FF:FF:FF:FF:FF**
- **Contents:** "Who has **192.168.1.20**?"

## □ Step 2: ARP Reply (Unicast)

- **Host B** responds directly to **A: 192.168.1.20** is at **MAC: AA:BB:CC:DD:EE:FF**.

# ARP Operation (Cont.)

## □ Step 3: ARP Cache Update

- Host A saves **mapping**:

IP	MAC	Timer
192.168.1.20	AA:BB:CC:DD:EE:FF	~2–20 mins

- This Timer Handling MAC Address Changes (The device using a specific IP address can change, and thus its MAC address will be different)

## □ Step 4: Communication Begins

- Now Host A sends actual data to Host B using Layer 2 MAC + Layer 3 IP.

# Problems with ARP

## 1. Security Problems

- **ARP Spoofing/Poisoning:** Attackers send **fake ARP replies to associate their MAC address with a legitimate IP**, enabling:
  - **Man-in-the-Middle (MitM) attacks for data interception.**
  - **Denial-of-Service (DoS) by redirecting traffic to a non-existent MAC.**
- **MAC Flooding:** Attackers flood a switch with **fake MAC addresses**, causing the switch to **broadcast traffic to all ports (like a hub)**.
- **No Authentication:** ARP has no mechanism to verify that an **ARP reply** is from the **legitimate owner** of an **IP address**. **It is a trusting, stateless protocol.**

# Problems with ARP

## 2. Management & Troubleshooting Problems

- **Lack of Scalability:** In **large Layer 2 networks**, the **volume of ARP broadcast traffic can become a significant overhead.**
- **No Built-in Loop Prevention:** ARP frames are **bridged** like any other traffic, so **network loops can cause ARP frames to circulate endlessly, exacerbating storms.**

## 3. Reliability & Network Problems

- **IP Address Conflicts:** If two devices have the **same IP**, **ARP caches** on other hosts become **unstable, flipping between two MAC addresses and causing broken connectivity.**

# **ICMP**

# What Is ICMP?

- ❑ ICMP is a **network layer protocol** used to **transmit control messages** between **hosts** and **routers** to **report whether hosts are reachable and routes are available**.
- ❑ It plays an important role in the TCP/IP protocol suite and is typically **used by** the **IP** or **higher layer protocols** (TCP or UDP).
- ❑ Although these **control messages** **do not** transmit user data.

# Why ICMP is Needed

- ❑ ICMP is essential because IP (Internet Protocol) itself is "unreliable" and lacks built-in mechanisms for diagnostics and error reporting. ICMP fills this critical gap.

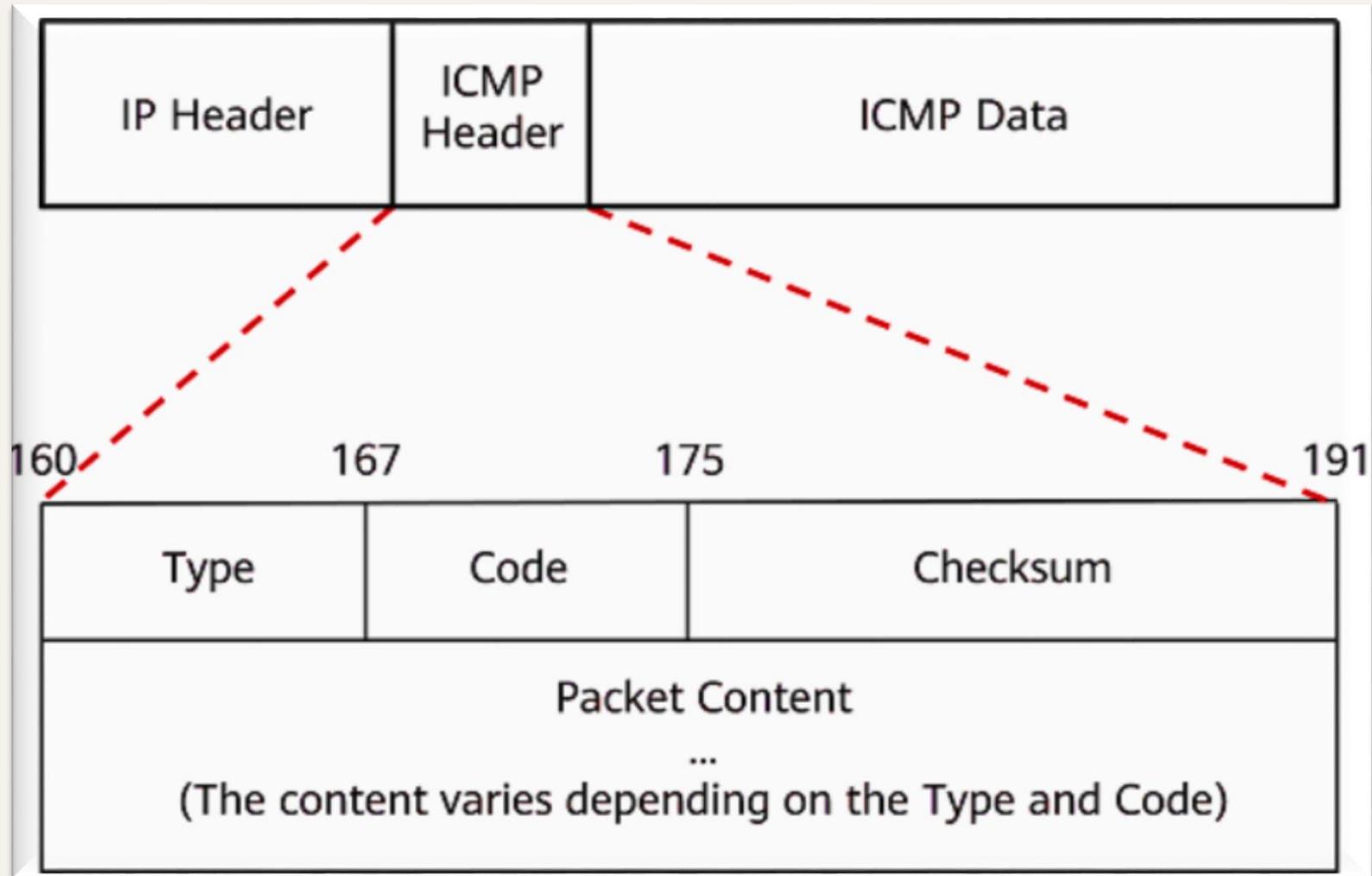
## 1. For Error Reporting and Signaling.

- **Informs Senders of Delivery Failures.**
- **Prevents Packet Loops:** Sends "**Time Exceeded**" messages when a packet's TTL (Time to Live) reaches zero.

## 2. For Network Diagnostics and Troubleshooting

- **Connectivity Testing:** The **ping** command uses **ICMP Echo Request** and **Reply** to **verify** if a **remote host** is **online** and **reachable**, and to measure **round-trip latency**.
- **Path Discovery:** The **traceroute** (or **tracert**) **command** uses **ICMP Time Exceeded** and **Destination Unreachable** messages to **map** the **entire network path between source and destination**.

# ICMP Message Format



# ICMP Message Format (Cont.)

Field	Size (Bytes)	Description
Type	1	<p>Identifies the general category of the ICMP message.</p> <ul style="list-style-type: none"><li>➤ 0 = Echo Reply</li><li>➤ 8 = Echo Request</li><li>➤ 3 = Destination Unreachable</li><li>➤ 11 = Time Exceeded</li></ul>
Code	1	<p>Provides more specific detail within the given Type. e.g., For Type 3 (Destination Unreachable):</p> <ul style="list-style-type: none"><li>➤ 0 = Network unreachable</li><li>➤ 1 = Host unreachable</li><li>➤ 3 = Port unreachable</li></ul>
Checksum	2	<p>Error-checking value for the entire ICMP message. Used to detect corruption during transmission.</p>

# Common ICMP Messages

## 1. Echo Request / Echo Reply (ping)

- Used to test: **Host reachability**, **Round-trip time**, and **Packet loss**.
- **Ping = ICMP Type 8 (request) and Type 0 (reply)**.
- **Example:**

```
C:\Users\Islam Zakaria>ping www.youtube.com

Pinging youtube-ui.l.google.com [142.250.180.174] with 32 bytes of data:
Reply from 142.250.180.174: bytes=32 time=66ms TTL=111
Reply from 142.250.180.174: bytes=32 time=98ms TTL=111
Reply from 142.250.180.174: bytes=32 time=66ms TTL=111
Reply from 142.250.180.174: bytes=32 time=67ms TTL=111

Ping statistics for 142.250.180.174:
    Packets: Sent = 4, Received = 4, Lost = 0 (0% loss),
    Approximate round trip times in milli-seconds:
        Minimum = 66ms, Maximum = 98ms, Average = 74ms
```

**Fields explained:**

**bytes** → **packet payload**

**time** → **RTT**

**TTL** → **remaining hops**

# Common ICMP Messages (Cont.)

## 2. Destination Unreachable (Type 3)

- Generated when a router or host cannot deliver a packet.
- Common codes:

Code	Meaning
0	Network unreachable
1	Host unreachable
3	Port unreachable
4	Fragmentation needed (MTU issue)

# Common ICMP Messages (Cont.)

## 3. Time Exceeded (Type 11)

- Used by traceroute.
- Occurs when:
  - TTL becomes **zero**, so **packet is dropped**.
  - Router sends **ICMP Time Exceeded** back to sender
- **Routers decrement TTL by 1 each hop.**

```
C:\Users\Islam Zakaria>tracert www.youtube.com

Tracing route to youtube-ui.l.google.com [142.251.209.46]
over a maximum of 30 hops:

 1  <1 ms    <1 ms    <1 ms  DESKTOP-1LD9KH9 [192.168.1.1]
 2  29 ms    29 ms    30 ms  197.161.132.1
 3  31 ms    28 ms    28 ms  172.24.144.73
 4  28 ms    28 ms    28 ms  172.24.144.53
 5  28 ms    30 ms    28 ms  172.19.1.241
 6  29 ms    29 ms    29 ms  81.10.87.222
 7  70 ms    67 ms    71 ms  93.186.129.182
 8  66 ms    66 ms    65 ms  93.186.129.45
 9  66 ms    66 ms    67 ms  72.14.221.64
10  66 ms    67 ms    66 ms  192.178.96.161
11  99 ms    66 ms    66 ms  142.251.235.177
12  66 ms    66 ms    66 ms  mil04s51-in-f14.1e100.net [142.251.209.46]

Trace complete.
```

# ICMP in Real-World Tools

## 1. ping (Connectivity Test)

- Sends **ICMP Echo Request (Type 8)** to target.
- Target replies with **ICMP Echo Reply (Type 0)**.
- **Measures: Reachability, Round-trip Time (Latency), Packet Loss.**

## 2. traceroute/tracert (Path Discovery)

- Uses **ICMP Time Exceeded messages (Type 11)**
- Sends packets with increasing **TTL** values
- Each router along path sends **ICMP Time Exceeded** when **TTL=0**
- **Maps entire network path to destination**

## 3. Path MTU Discovery

- Uses **ICMP Fragmentation Needed messages (Type 3, Code 4)**
- **Discovers maximum packet size without fragmentation**
- **Automatically adjusts packet size for optimal performance**

# ICMP in Real-World Tools (Cont.)

## 4. Destination Unreachable Diagnostics

- Helps identify why connections fail

## 5. Network Monitoring Systems

- Use ICMP Echo for continuous availability monitoring
- Alert when hosts become unreachable
- Track latency and jitter over time

## 6. Router Redirect Function

- Routers inform hosts of better network paths
- Optimizes local network routing

## 7. Error Reporting in Applications

- Operating systems use ICMP messages to report connection failures
- e.g., "Connection refused" from Port Unreachable messages
- "No route to host" from Host/Network Unreachable messages

# ICMP Attacks

## 1. Reconnaissance Attacks

- **Host Discovery:** Attackers use ping to identify live hosts on a network
- **Network Mapping:** traceroute reveals network topology and firewall locations
- **OS Fingerprinting:** Analyzing ICMP response behavior to identify operating systems.

## 2. Denial of Service (DoS) Attacks

- **ICMP Flood:** Overwhelm target with massive amounts of ICMP Echo Requests
- **Ping of Death:** Oversized ICMP packets cause buffer overflow and system crashes

## 3. Redirect Attacks

- **ICMP Redirect:** Attackers send fake redirect messages to poison routing tables

# ICMP Security Measures

## 1. Filtering & Blocking

- Block unnecessary ICMP types at network perimeter
- Use stateful inspection to only allow replies to outbound requests
- Implement rate limiting to prevent flood attacks

## 2. ICMP Type Management

- **Allow:** Echo Reply (0), Destination Unreachable (3), Time Exceeded (11)
- **Block:** Redirect (5), external Echo Requests (8), Address Mask Requests (17)

## 3. Monitoring & Detection

- Detect unusual ICMP patterns and volumes
- Inspect payloads for hidden data
- Log suspicious ICMP messages

# ICMP Security Measures (Cont.)

## 4. Host Hardening

- Disable ICMP on critical servers when possible
- Harden OS ICMP stack settings
- Use host-based firewalls for granular control

Any  
Questions