

1 Number Systems

Decimal

$$\begin{array}{ccc} 1 & 5 & 7 \\ \hline 100 & 10 & 1 \end{array} \quad \text{weights}$$

$$1 \times 100 + 5 \times 10 + 7 \times 1 = 157$$

base-10 number system

Binary

Binary is a *base-2* number system.

$$\begin{array}{ccccc} 0 & 1 & 0 & 1 & 1 \\ \hline 16 & 8 & 4 & 2 & 1 \end{array} \text{ weights}$$

$$0 \times 16 + 1 \times 8 + 0 \times 4 + 1 \times 2 + 1 \times 1 = 11 \text{ dec}$$

$$8 + 2 + 1 = 11 \text{ decimal.}$$

We call a sequence of eight bits a *byte*.

To *complement* a bit means to flip it.

Hexadecimal

Hexadecimal (or hex for short) is the *base-16* number system.

16 symbols: 0 to 9 and A, B, C, D, E, and F

$$\begin{array}{ccc} 2 & C & 5 \\ \hline 256 & 16 & 1 \end{array} \text{ weights (in decimal)}$$

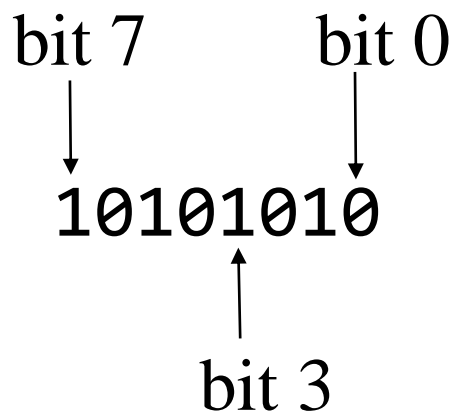
$$2 \times 256 + C \times 16 + 5 \times 1$$

$$2 \times 256 + 12 \times 16 + 5 \times 1 = 512 + 192 + 5 = 709$$

Decimal	Binary	Hex
0	0000	0
1	0001	1
2	0010	2
3	0011	3
4	0100	4
5	0101	5
6	0110	6
7	0111	7
8	1000	8
9	1001	9
10	1010	A (or a)
11	1011	B (or b)
12	1100	C (or c)
13	1101	D (or d)
14	1110	E (or e)
15	1111	F (or f)

Rule: Adding a 0 on the right side of a positional whole number multiplies its value by its base.

Numbering Bits



most significant bit

least significant bit

Adding Positional Numbers

decimal addition:

$$\begin{array}{r} 1 \text{ carries} \\ 157 \\ + 238 \\ \hline 395 \end{array}$$

binary addition:

$$\begin{array}{r} 11 \text{ carries} \\ 0011 \\ + 0011 \\ \hline 0110 \end{array}$$

hex addition:

$$\begin{array}{r} 1 \text{ carries} \\ 1B \\ + 37 \\ \hline 52 \end{array}$$

Representing Negative Numbers

$$\begin{array}{rcl}
 0011 & = & +3 \\
 + 1011 & = & -3 \\
 \hline
 1110 & = & -6
 \end{array}
 \quad \text{Sign-magnitude}$$

Carry out of leftmost column

$$\begin{array}{rcl}
 1 & 111 & \text{carries} \\
 & 1111 & \\
 + & 0001 & \\
 \hline
 & 0000 &
 \end{array}$$

Rule: Adding 1 to a binary number with a fixed number of bits all of which are 1 results in all zeros.

$$\begin{array}{rcl}
 0011 & = & +3 \\
 + 1100 & = & +3 \text{ with each bit flipped} \\
 \hline
 1111 & &
 \end{array}$$

$$\begin{array}{rcl}
 1100 & = & +3 \text{ with each bit flipped} \\
 + 0001 & \text{add 1} & \\
 \hline
 1101 & \text{Is this -3?} &
 \end{array}$$

Is 1101 the representation of -3 that we want? Let's add it to +3 to see if it gives zero:

$$\begin{array}{rcl}
 1 & 111 & \text{carries} \\
 & 0011 & = +3 \\
 + & 1101 & \text{Is this -3?} \\
 \hline
 & 0000 &
 \end{array}$$

Rule: To negate a binary number in the two's complement system, flip its bits and add 1.

Rule: In the two's complement system, all 1's represents -1.

Signed and Unsigned Numbers

Unsigned	Value	Signed	Value
0000	0	1000	-8
0001	1	1001	-7
0010	2	1010	-6
0011	3	1011	-5
0100	4	1100	-4
0101	5	1101	-3
0110	6	1110	-2
0111	7	1111	-1
1000	8	0000	0
1001	9	0001	1
1010	10	0010	2
1011	11	0011	3
1100	12	0100	4
1101	13	0101	5
1110	14	0110	6
1111	15	0111	7

↑
sign bit

Range of Numbers

n	Unsigned	Signed
3	0 to 7	-4 to 3
4	0 to 15	-8 to 7
5	0 to 31	-16 to 15
6	0 to 63	-32 to 31
7	0 to 127	-64 to 63
8	0 to 255	-128 to 127
9	0 to 511	-256 to 255
10	0 to 1023	-512 to 511
12	0 to 4095	-2048 to 2047
16	0 to 65535	-32768 to 32767
k	0 to $2^k - 1$	-2^{k-1} to $2^{k-1} - 1$

Powers of 2

n	2^n
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
10	1,024 (aka 1K)
11	2,048 (aka 2K)
12	4,096 (aka 4K)
15	32,768 (aka 32K)
16	65,536 (aka 64K)
20	1,048,576 (aka 1M)
30	1,073,741,824 (aka 1G)

Converting Between Binary and Hex

1	1010	1110	0000	1100
↓	↓	↓	↓	↓
1	A	E	0	C

Converting Decimal to Binary

	remainders
	↓
	0 1
10) <u>1</u>	5
10) <u>15</u>	7
10) <u>157</u>	

↑

Start here and work up

	0 1
2) <u>1</u>	0
2) <u>2</u>	0
2) <u>4</u>	1
2) <u>9</u>	1
2) <u>19</u>	1
2) <u>39</u>	0
2) <u>78</u>	1
2) <u>157</u>	

↑

Start here and work up

Converting Fractions

Diagram illustrating the conversion of the decimal $.153$ to the whole number 3 through successive multiplication by 10:

$$\begin{array}{r} .153 \\ \times 10 \\ \hline 1.530 \\ \times 10 \\ \hline 15.300 \\ \times 10 \\ \hline 153.000 \end{array}$$

Annotations:

- whole parts are the digits \rightarrow 1, 5, 3
- multiply fractional part of each product \leftarrow .530, .300, .000

Converting Binary Fraction

$$\begin{array}{r} .375 \\ \times \quad 2 \\ \hline 0 \quad .750 \\ \times \quad 2 \\ \hline 1 \quad .500 \\ \times \quad 2 \\ \hline 1 \quad .000 \end{array}$$

.375 decimal = .011 binary

Infinite Length Equivalents

$$\begin{array}{r} .1 \\ \times \quad 2 \\ \hline 0.2 \\ \times \quad 2 \\ \hline 0.4 \end{array} \quad \begin{array}{l} \swarrow \\ .2 \text{ repeats} \\ \searrow \end{array}$$
$$\begin{array}{r} \times \quad 2 \\ \hline 0.8 \\ \times \quad 2 \\ \hline 1.6 \\ \times \quad 2 \\ \hline 1.2 \end{array}$$

The diagram illustrates the multiplication of 0.1 by 2, showing the result 0.2, 0.4, 0.8, 1.6, and 1.2. The text ".2 repeats" is placed to the right of the results, with arrows pointing to the decimal parts of the results (0.2, 0.8, and 1.2) to indicate that the digit 2 repeats in the decimal part of the result.

Zero and Sign Extension

Extend 1111111 to 16 bits

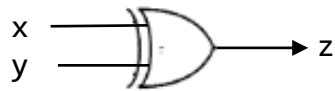
000000001111111

1111111111111111

Rule: Always sign extend signed numbers.

Rule: Always zero extend unsigned numbers.

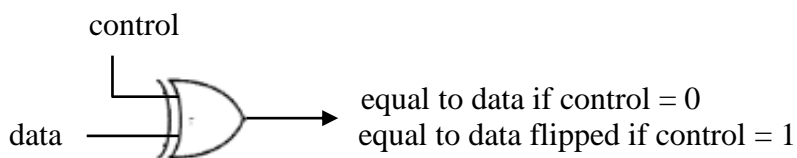
Exclusive OR Gate



Exclusive OR Gate

x	y	z
0	0	0
0	1	1
1	0	1
1	1	0

difference detecting gate



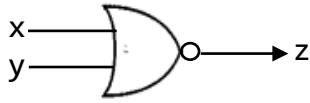
Control	Data	Output of exclusive OR
0	0	0
0	1	1
1	0	1
1	1	0

Data unaffected when control = 0

Data flipped when control = 1

Rule: If the control line of an exclusive OR is 0, the data passes through gate unchanged. If the control line is 1, the data is complemented.

NOR Gate

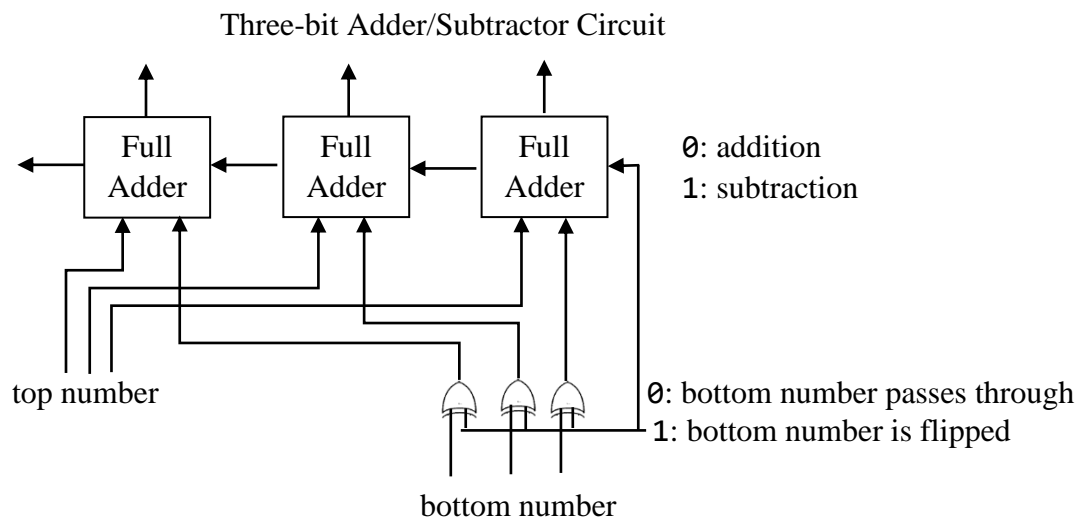
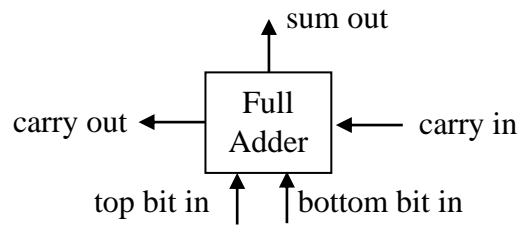


x	y	z
0	0	1
0	1	0
1	0	0
1	1	0

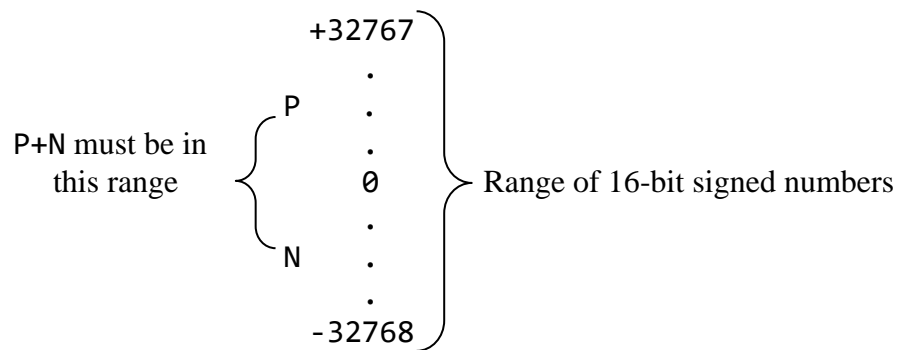
1 when all
inputs = 0

zero detecting gate

Addition and Subtraction



Signed Overflow



Rule: If two signed numbers *with different signs* are added or two signed numbers *with the same sign* are subtracted, overflow cannot occur.

Overflow scenarios:

Adding two non-negative numbers

(a)

$$\begin{array}{r} 0 \text{-----} \\ + 0 \text{-----} \\ \hline 1 \text{-----} \end{array}$$

Sign bit should
be 0

Adding two negative numbers

(b)

$$\begin{array}{r} 1 \text{-----} \\ + 1 \text{-----} \\ \hline 0 \text{-----} \end{array}$$

Sign bit should
be 1

Here are all the possible scenarios for no overflow:

(c)

$$\begin{array}{r} 0 \text{-----} \\ + 0 \text{-----} \\ \hline 0 \text{-----} \end{array}$$

(d)

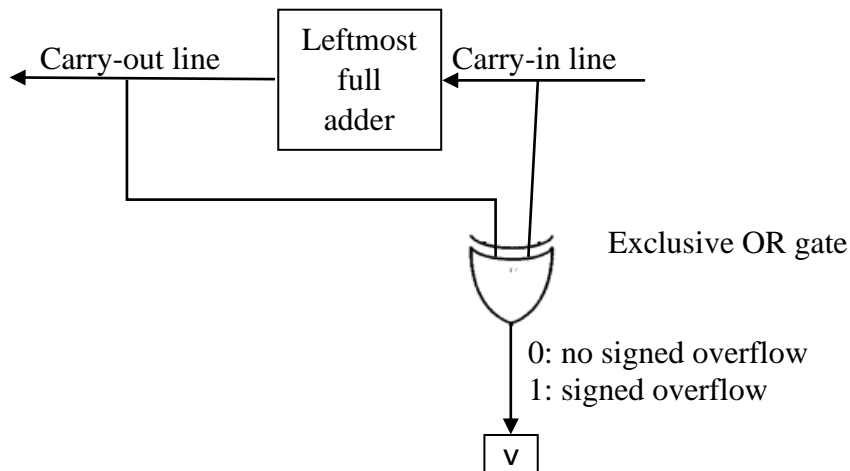
$$\begin{array}{r} 1 \text{-----} \\ + 1 \text{-----} \\ \hline 1 \text{-----} \end{array}$$

(e)

$$\begin{array}{r} 0 \text{-----} \\ + 1 \text{-----} \\ \hline 0 \text{-----} \end{array}$$

(f)

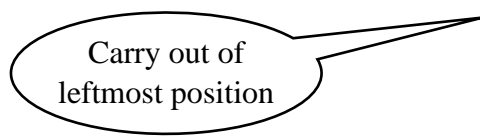
$$\begin{array}{r} 0 \text{-----} \\ + 1 \text{-----} \\ \hline 1 \text{-----} \end{array}$$



Rule: In an addition or subtraction of signed numbers, overflow has occurred if the carry into the leftmost position does not match the carry out.

Rule: If overflow occurs during the addition of two signed numbers, the sign bit of the computed result is wrong.

Unsigned Overflow

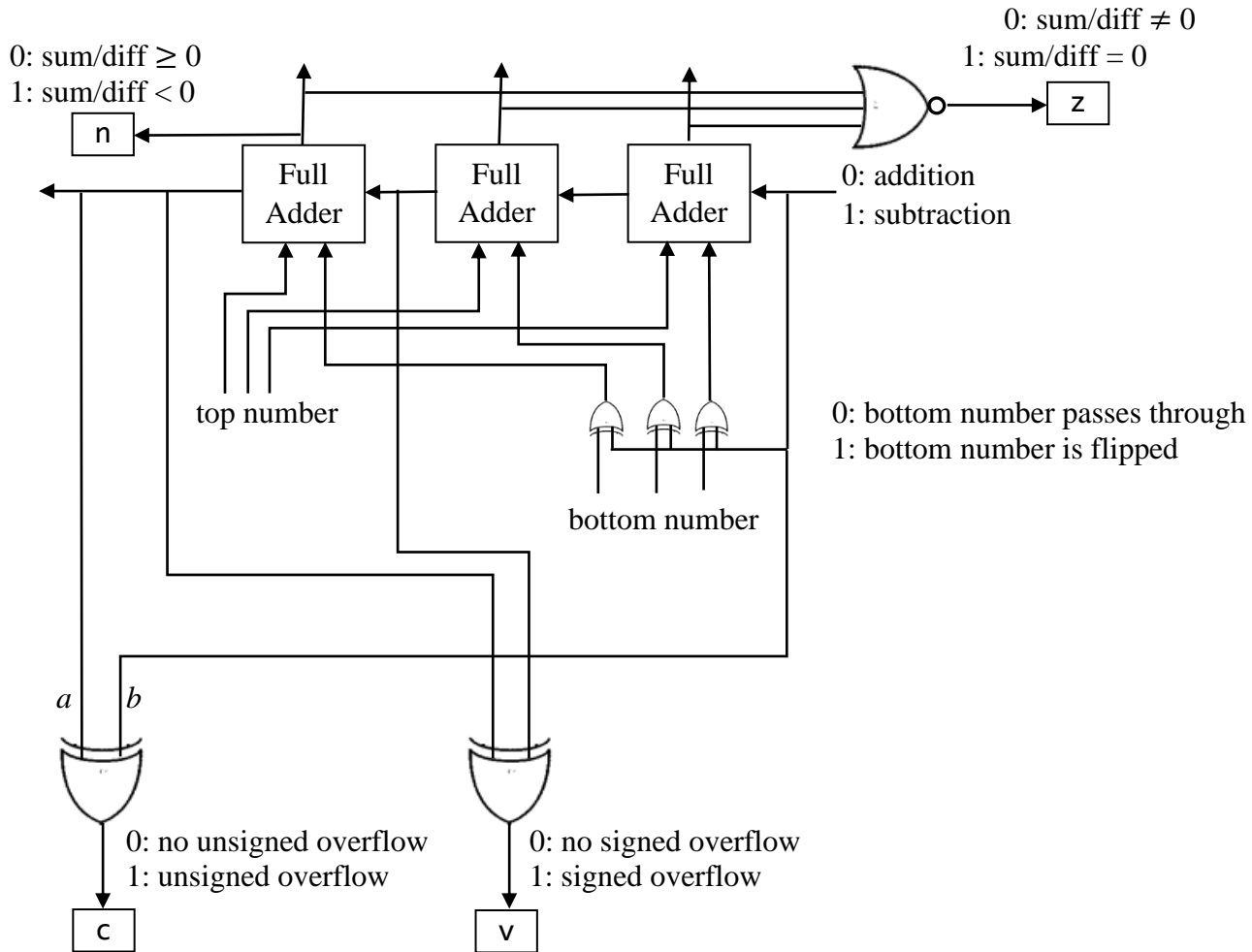


Carry out of leftmost position

$$\begin{array}{r} 1 \text{ carries} \\ 1111111111111111 \\ 1111111111111111 = 65535 \\ + 0000000000000001 = 1 \\ \hline 0000000000000000 = 0 \end{array}$$

Rule: In an addition of unsigned numbers, a carry out of the leftmost position indicates overflow. In a subtraction, no carry out of the leftmost position (or a borrow in if the borrow technique is used) indicates overflow.

Three-bit Adder/Subtractor Circuit with n, z, c, and v Flags



Floating Point

hidden bit exponent
 ↓ ↙
 1.10101 x 2³ (which equals 13.25 decimal)
 fractional
 part

- One bit: the sign bit (0 for zero or positive, 1 for negative)
- Eight bits: the exponent *to which 127 is added*
- 23 bits: the fractional part, zero extended on the right

101010000000000000000000.

Thus, the 32-bit floating-point representation of 13.25 decimal is

sign
0 1000010 101010000000000000000000
 adjusted fractional part
 exponent

Special Representations

Exponent field	Fractional part	Value
all 0's	all 0's	0.0
all 0's	nonzero	denormal number (exponent = -126, hidden bit 0)
all 1's	all 0's	infinity
all 1's	nonzero	NaN (Not a Number)

Denormal Numbers

Smallest positive magnitude that can be represented by a regular floating-point number is

0 00000001 00000000000000000000000000000000

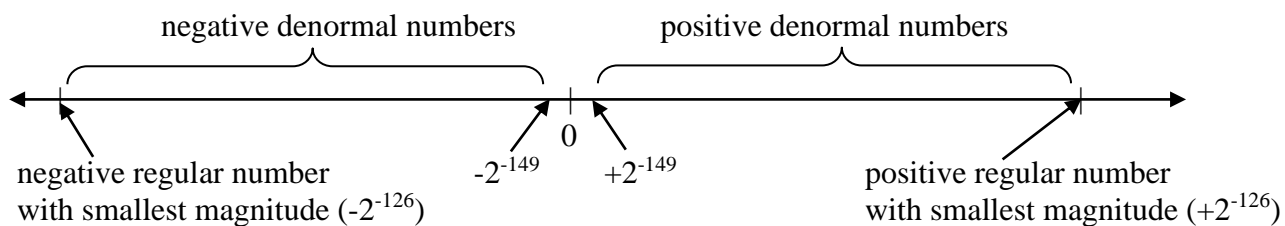
which equals

↗ 1.00000000000000000000000000000000 $\times 2^{-126}$
hidden bit

$$0\ 00000000\ 11111111111111111111111111111111 = 0.11111111111111111111111111111111 \times 2^{-126}$$

which is just under 2^{-126} , down to

$$0\ 00000000\ 00000000000000000000000000000001 = 0.00000000000000000000000000000001 \times 2^{-126}$$



Floating-point numbers (`float` and `double` in C) are *only an approximation of the set of real numbers in mathematics*

```
1 // ex0101.c  Finite precision of floating point numbers
2 #include <stdio.h>
3 int main()
4 {
5     double sum = 0.0;
6     int i;
7     for (i = 1; i <= 10; i++)
8         sum = sum + 0.1;
9     printf("%.17f\n", sum);    // sum is not 1.0
10    return 0;
11 }
```

```
1 // ex0102.c Floating-point errors
2 #include <stdio.h>
3 int main()
4 {
5     double x = 1E70;    //  $10^{70}$ 
6     double y = 1E5;     //  $10^5$ 
7     if (x == x + y)
8         printf("y must be zero\n"); // displayed but y is not zero
9     else
10        printf("y must be nonzero\n");
11    return 0;
12 }
```

Observation: Numbers can act like zero when added to numbers with significantly larger magnitudes.

```

1 // ex0103.c Lack of associativity in floating-point computations
2 #include <stdio.h>
3 int main(void)
4 {
5     double sum = 0.0;
6     int i;
7     for (i = 1; i <= 100000000; i++)
8         sum = sum + 1.0/i;
9     printf("%.17f\n", sum);
10
11     sum = 0.0;
12     for (i = 100000000; i >= 1; i--)           // better for loop
13         sum = sum + 1.0/i;
14     printf("%.17f\n", sum);
15     return 0;
16 }

```

Observation: Addition of floating-point numbers is not associative (i.e., order of evaluation matters).