Note of Succinct: Enabling Queries on Compressed Data

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1 Motivation

Evaluation of popular open-source data stores (MongoDB, Cassandra) using real-world datasets shows that indexes can be much as 8x larger than the input data size. Existing data stores either **resort** to using complex memory management techniques for identifying and caching "hot" data or **simply executing queries off-disk or off-SSD**. In either case, latency and throughput advantages of indexes drop compared to in-memory query execution.

2 Succinct Overview

- a distributed data store
- operates at a new point in the design space
- memory efficiency close to data scans and latency close to indexes
- is able to store more data in memory, avoiding latency and throughput degradation due to off-disk or off-SSD query execution

Contributions:

- Enables efficient queries directly on a compressed representation of the input data
 - a new data structure
 - a new query algorithm
- Efficiently supports data appends by chaining multiple stores
 - a small log-structure store optimized for fine-grained appends
 - an intermediate store optimized for query efficiency while supporting bulk appends
 - an immutable store that stores most of the data
- Exposes a minimal, yet powerful, API that operates on flat unstructured files

2.1 Succinct Interface

For string abbcdeabczabgz,

- search(f, "ab") returns [0, 6, 10].
- wildcardsearch(f, prefix="ab", suffix="z", dist=2) returns [6, 9] for abcz, [10, 13] for abgz.

For semi-structured data (KV), like figure 1, transforming the input data into flat files.



Figure 1: Succinct supports queries on semi-structured data by transforming the input data into flat files

3 Querying on Compressed Data

3.1 Compressed Suffix array and Succinct data representation

Suffix array Wikipedia

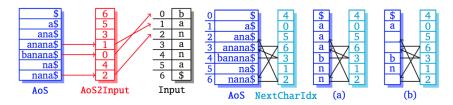


Figure 2

Figure 2 (a) shows that AoS stores the first char of each suffix only. Since suffixes are sorted, only the first AoS index at which each character occurs (e.g., $\{(\$,0),(a,1),(b,4),(n,5)\}$) need be stored. Succinct uses *Skewed Wavelet Tree* to compress each row independently.

Succinct uses sampling by "value" strategy. For sampling rate α , Succinct stores all AoS2Input values that are a multiple of α . This allows storing each sampled value val as val/ α , leading to a more space-efficient representation. Using $\alpha = 2$ for example of Figure 3a, for instance, the sampled

AoS2Input values are $\{6,0,4,2\}$, which can be stored as $\{3,0,2,1\}$. Sampled Input2AoS then becomes $\{1,3,2,0\}$ with i-th value being the index into sampled AoS2Input where i is stored. Succinct stores a small amount of additional information to locate sampled AoS2Input indexes.

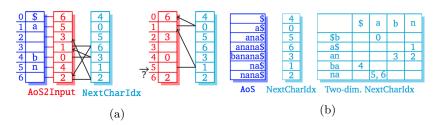


Figure 3: 3a for reducing the space usage of AoS2Input. 3b for two-dimensional NextCharIdx representation.

For instance, consider the query <code>search(anan)</code>; all occurrences of string <code>"nan"</code> are contained in the cell <code><n,an></code>. To find all occurrences of string <code>anan</code>, our algorithm performs a binary search only in the cell <code><a,na></code> in the next step. After this step, the algorithm has the indexes for which suffixes start with <code>"a"</code> and are followed by <code>"nan"</code>, the desired string. For a string of length m, the above algorithm performs 2(mt1) binary searches, two per NextCharIdx cell.

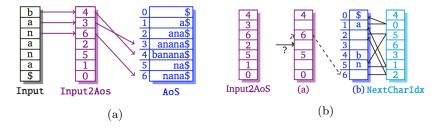


Figure 4: 4a for The Input2AoS provides the inverse mapping of AoS2Input. 4b for reducing the space usage of Input2AoS.

extract functionality is shown in Figure 4. For instance, to execute extract(3, 3), we find the next smaller sampled index (Input2AoS[2]) and corresponding suffix (AoS[2]="nana\$"). We then remove the first character since the difference between the desired index and the closest sampled index was 1; hence the result "ana\$".