Operating System Homework Report 2

1. How to implemented the program in detail?

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- use comment to explain

```
#define _GNU_SOURCE
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <pthread.h>
#include <unistd.h>
void *thread func(void *arg) {
      pthread_barrier_wait(&barrier);
```

```
pthread mutex lock(&mutex);
                  if (my_clock() - sttime > 1* time_wait)
      pthread_mutex_unlock(&mutex);
  pthread exit(NULL);
int main(int argc, char **argv) {
  int opt;
  while ((opt = getopt(argc, argv, "n:t:s:p:")) != -1) {
      switch (opt) {
              n = atoi(optarg);
              t = atof(optarg);
              s = optarg;
              p = optarg;
```

```
char *policy input = strdup(s);
char *priority_input = strdup(p);
char policy array[n][10];
char priority_array[n][10];
int policy[n];
int priority[n];
char *policy tokens = strtok(policy input, ",");
strcpy(policy_array[id++], policy_tokens);
    strcpy(policy_array[id++], policy_tokens);
char *priority_tokens = strtok(priority_input, ",");
strcpy(priority array[id++], priority tokens);
    strcpy(priority array[id++], priority tokens);
free(policy input);
free(priority input);
    if (strcmp(policy_array[i], "NORMAL") == 0) {
        policy[i] = SCHED OTHER;
    } else if (strcmp(policy_array[i], "FIFO") == 0) {
        fprintf(stderr, "Error: Unknown scheduling policy.\n");
    exit(EXIT FAILURE);
    priority[i] = atoi(priority array[i]);
```

```
pthread t *tid = (pthread t *)malloc(sizeof(pthread t) * n);
      threadInfo[i].policy = policy[i];
      threadInfo[i].priority = priority[i];
  cpu set t set; CPU ZERO(&set); CPU SET(0, &set);
  sched setaffinity(getpid(), sizeof(set), &set);
  pthread barrier init(&barrier, NULL, n + 1);
      pthread_attr_init(&attr);
      pthread attr setinheritsched(&attr, PTHREAD EXPLICIT SCHED);
      pthread attr setschedpolicy(&attr, threadInfo[i].policy);
      param.sched priority = threadInfo[i].priority;
      pthread_attr_setschedparam(&attr, &param);
      int result = pthread_create(&tid[i], &attr, thread_func, (void
           fprintf(stderr, "Error creating thread %ld: %s\n", i,
strerror(result));
```

```
}
}
//wait until all threads are ready
pthread_barrier_wait(&barrier);

/* 5. Start all threads at once */
for (int i = 0; i < n; i++) {
    //join thread
    int result = pthread_join(tid[i], NULL);
    if (result != 0) {
        fprintf(stderr, "Error joining thread %d: %s\n", i,
strerror(result));
        exit(EXIT_FAILURE);
    }
}
/* 6. Wait for all threads to finish */
//destroy barrier
pthread_barrier_destroy(&barrier);
free(tid);
free(threadInfo);

return 0;
}</pre>
```

This is the result of the test cases.

```
zhuyan1228@zhuyan1228:~$ sudo ./sched_test.sh ./sched_demo ./a.out
Running testcase 1: ./sched_demo -n 1 -t 0.5 -s NORMAL -p -1 .....
Result: Success!
Running testcase 2: ./sched_demo -n 2 -t 0.5 -s FIFO,FIFO -p 10,20 .....
Result: Success!
Running testcase 3: ./sched_demo -n 3 -t 1.0 -s NORMAL,FIFO,FIFO -p -1,10,30 ...
...
Result: Success!
```

- 2. Describe the results of ./sched_demo -n 3 -t 1.0 -s NORMAL,FIFO,FIFO -p -1,10,30 and what causes that:
 - Threads 1 and 2 use FIFO policy, so they have real-time policy priority and the bigger priority value has higher priority; Thread 0 uses Normal policy so it doesn't have real-time policy priority(-1), assigned lower priorities as it is a standard scheduling policy.
 - Use setaffinity let them run on the same cpu.
 - When a SCHED_FIFO thread becomes runnable, it will always immediately preempt any currently running thread.
 - In that case, most likely, thread 2 will be run first (priority(30) > thread 1(10) > thread 0(-)), and thread 0 will be run at the end.
 - Thread 2 -> Thread 1 -> Thread 0

```
zhuyan1228@zhuyan1228:~$ sudo ./a.out -n 3 -t 1.0 -s NORMAL,FIFO,FIFO -p -1,10,30
Thread 2 is running
Thread 2 is running
Thread 2 is running
Thread 1 is running
Thread 1 is running
Thread 1 is running
Thread 0 is running
Thread 0 is running
Thread 0 is running
zhuyan1228@zhuyan1228:~$ sudo ./a.out -n 3 -t 1.0 -s NORMAL,FIFO,FIFO -p -1,10,30
Thread 2 is running
Thread 2 is running
Thread 2 is running
Thread 1 is running
Thread 1 is running
Thread 1 is running
Thread 0 is running
Thread 0 is running
Thread 0 is running
```

- 3. Describe the results of ./sched_demo -n 4 -t 0.5 -s NORMAL,FIFO,NORMAL,FIFO -p -1,10,-1,30, and what causes that:
 - Thread 3 and thread 1 use FIFO policy, and the others use Normal policy, so they may run in front of them.
 - Thread 3 have a higher priority than thread 1, it will run before thread 1 because it FIFO policy.
 - Use setaffinity let them run on the same cpu.
 - Normal (SCHED_OTHER) policy: assign a time slice to the thread during which it can execute. It's possible that thread 0 and 2 will fight for using CPU. May cause preempt.
 - Thread 3 -> Thread 1 -> Thread 0 or 2

```
zhuyan1228@zhuyan1228:~$ sudo ./a.out -n 4 -t 0.5 -s NORMAL,FIFO,NORMAL,FIFO -p -1,10,-1,30
Thread 3 is running
Thread 3 is running
Thread 1 is running
Thread 1 is running
Thread 1 is running
Thread 2 is running
Thread 0 is running
Thread 2 is running
Thread 0 is running
Thread 0 is running
Thread 0 is running
Thread 2 is running
```

```
zhuyan1228@zhuyan1228:~$ sudo ./a.out -n 4 -t 0.5 -s NORMAL,FIFO,NORMAL,FIFO -p -1,10,-1,30
Thread 3 is running
Thread 3 is running
Thread 1 is running
Thread 1 is running
Thread 1 is running
Thread 2 is running
Thread 2 is running
Thread 0 is running
```

- 4. Describe how did you implement n-second-busy-waiting?
 - my_clock This function gets the current thread's CPU time in seconds, used for measuring time.

```
//get the CPU time of the current thread in seconds function
static double my_clock(void) {
    struct timespec t;
    // ensure clock_gettime can excute successfully
    assert(clock_gettime(CLOCK_THREAD_CPUTIME_ID, &t) == 0);
    //calculate total seconds
    double seconds = t.tv_sec + 1e-9 * t.tv_nsec;
    return seconds;
}
```

- The busy-wait loop simulates some computation or delay in each thread

- In thread_func, the busy-wait loop in the section doing the task, enters a busy-wait loop for time_wait(t) seconds using the my_clock function, ensure n-second-busy-waiting.

- Thread_func