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(a) min(R₁, R₂, R₃) = R₁ = 500 kbps (b) min(R₁, R₂, R₃) = R₃ = 200 kbps

(a)
$$\frac{2M}{100k} = \frac{2000}{100} = 20 \text{ users}$$

(b) $\frac{20\%}{00} = \frac{2000}{100} = 20 \text{ users}$
(c) $\frac{40}{00} = \frac{2000}{100} = 20 \text{ users}$
(d) $\frac{20}{000} = \frac{2000}{100} = 20 \text{ users}$
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$$\frac{3}{\sqrt{N}} \left(\frac{N + \frac{1}{N} + 2\frac{1}{N} + \dots + (N - \lambda)\frac{1}{N} + (N - 1)\frac{1}{N}}{\sqrt{N}} \right) \cdot \frac{1}{N}$$

$$= \frac{N - 1}{\sqrt{N}} \hat{\lambda} \frac{1}{NR}$$

$$= \frac{(N - 1) \cdot N}{2} \cdot \frac{1}{NR}$$

$$= \frac{(N - 1) \cdot L}{2R}$$

4

end-to-end delay (ms) =
$$d_{proc} + d_{queue} + d_{trans} + d_{prop}$$
 $d_{1} = 4000 \text{ km}$ $S_{1} = S_{2} = 2 \cdot 10^{8} \text{ m/s}$

$$= 1 + \frac{L}{R_{1}} + \frac{L}{R_{2}} + \frac{d_{1}}{S_{1}} + \frac{d_{2}}{S_{2}} \text{ (msec)}$$

$$= 1 + \log v \left(\frac{8000}{2M} + \frac{8000}{2 \cdot 10^{8}} + \frac{1000 \text{ k}}{2 \cdot 10^{8}} \right)$$

$$= 1 + \log v \left(\frac{8000}{2M} + \frac{8000}{2M} + \frac{4000 \text{ k}}{2 \cdot 10^{8}} + \frac{1000 \text{ k}}{2 \cdot 10^{8}} \right)$$

$$= 1 + \log v \left(0.004 + 0.004 + 0.005 + 0.005 \right)$$

$$= 1 + 33$$

$$= 34 \text{ msec}$$

5

d=10000km, R=1Mbps, S=2.5×108 m/sec

$$d_{prop} = \frac{d}{S} = \frac{10000 \, \text{k}}{2.5 \times 10^8} = \frac{1}{2.5 \times 10} = \frac{1}{25} = 0.04 \, \text{s} \qquad R \cdot d_{prop} = 1 \, \text{M} \cdot 0.04 = 40 \, \text{kbits}$$

(b) 400 kbits > bandwidth-delay product, maximum bits = 40 kbits

(c)
Bandwidth-delay product is the maximum number of bits that can be filled in a network link.

(d) width of a bit = $\frac{d}{R \cdot dprop} = \frac{10000k}{40k} = \frac{1000}{4} = 250 \text{ m}$

(e) width of a bit = $\frac{d}{R \cdot d_{\text{prop}}} = \frac{d}{R \cdot \frac{d}{S}} = \frac{S}{R}$