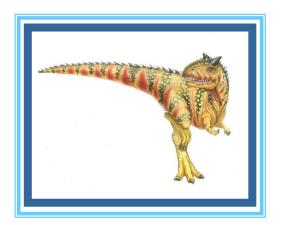
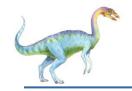
# Chapter 13: File-System Interface

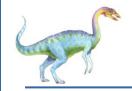




#### **Outline**

- File Concept
- Access Methods
- Disk and Directory Structure
- File-System Mounting
- File Sharing
- Protection

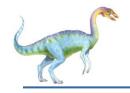




## **Objectives**

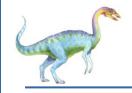
- To explain the functions of file systems
- To describe the interfaces to file systems
- To discuss file-system design tradeoffs, including access methods, file sharing, file locking, and directory structures
- To explore file-system protection





## File Concept

- Contiguous logical address space
- Types:
  - Data
    - numeric
    - character
    - binary
  - Program
- Contents defined by file's creator
  - Many types
    - Consider text file, source file, executable file



#### File Attributes

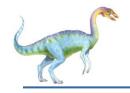
- Name only information kept in human-readable form
- Identifier unique tag (number) identifies file within file system
- Type needed for systems that support different types
- Location pointer to file location on device
- Size current file size
- Protection controls who can do reading, writing, executing
- Time, date, and user identification data for protection, security, and usage monitoring
- Information about files are kept in the directory structure, which is maintained on the disk
- Many variations, including extended file attributes such as file checksum



#### File info Window on Mac OS X

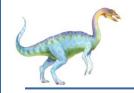






## **File Operations**

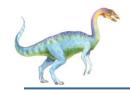
- File is an abstract data type
- 6 basic operations
  - Create
  - Write at write pointer location
  - Read at read pointer location
  - Reposition within file seek
  - Delete
  - Truncate
- To avoid constant searching of the file entry in the directory,
  - $Open(F_i)$  search the directory structure on disk for entry  $F_i$ , and move the content of entry to memory
  - Close  $(F_i)$  move the content of entry  $F_i$  in memory to directory structure on disk



## **Open Files**

- Several pieces of data are needed to manage open files:
  - Open-file table: tracks open files
  - File pointer: pointer to last read/write location, per process that has the file open
  - File-open count: counter of number of times a file is open – to allow removal of data from openfile table when the last process closes it
  - Disk location of the file: cache of data access information
  - Access rights: per-process access mode information

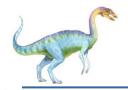




## **Open File Locking**

- Provided by some OS and file systems
  - Similar to reader-writer locks
  - Shared lock similar to reader lock several processes can acquire concurrently
  - Exclusive lock similar to writer lock
- Mediates access to a file
- Mandatory or advisory:
  - Mandatory access is denied depending on locks held and requested
  - Advisory processes can find status of locks and decide what to do





## File Locking Example – Java API

```
import java.io.*;
import java.nio.channels.*;
public class LockingExample {
    public static final boolean EXCLUSIVE = false;
    public static final boolean SHARED = true;
    public static void main(String arsg[]) throws IOException {
          FileLock sharedLock = null:
          FileLock exclusiveLock = null;
          try {
                     RandomAccessFile raf = new RandomAccessFile("file.txt",
    "rw");
                    // get the channel for the file
                     FileChannel ch = raf.getChannel();
                    // this locks the first half of the file - exclusive
                     exclusiveLock = ch.lock(0, raf.length()/2, EXCLUSIVE);
                    /** Now modify the data . . . */
                    // release the lock
                     exclusiveLock.release();
```



## File Locking Example – Java API (Cont.)

```
// this locks the second half of the file - shared
          sharedLock = ch.lock(raf.length()/2+1, raf.length(),
                                SHARED):
          /** Now read the data . . . */
          // release the lock
          sharedLock.release();
} catch (java.io.IOException ioe) {
          System.err.println(ioe);
}finally {
          if (exclusiveLock != null)
          exclusiveLock.release();
          if (sharedLock != null)
          sharedLock.release();
```

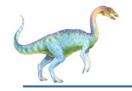




# File Types – Filename Extension

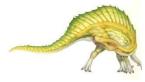
file type	usual extension	function	
executable	exe, com, bin or none	ready-to-run machine- language program	
object	obj, o	compiled, machine language, not linked	
source code	c, cc, java, pas, asm, a	source code in various languages	
batch	bat, sh	commands to the command interpreter	
text	txt, doc	textual data, documents	
word processor	wp, tex, rtf, doc	various word-processor formats	
library	lib, a, so, dll	libraries of routines for programmers	
print or view	ps, pdf, jpg	ASCII or binary file in a format for printing or viewing	
archive	arc, zip, tar	related files grouped into one file, sometimes com- pressed, for archiving or storage	
multimedia	mpeg, mov, rm, mp3, avi	binary file containing audio or A/V information	

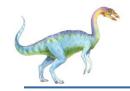




#### File Structure

- None sequence of words, bytes
- Simple record structure
  - Lines
  - Fixed length
  - Variable length
- Complex Structures
  - Formatted document
  - Relocatable load file
- Who decides:
  - OS
  - Program





#### **Access Methods**

Sequential Access

```
read next
write next
reset
no read after last write
(rewrite)
```

Direct Access – file is fixed length logical records

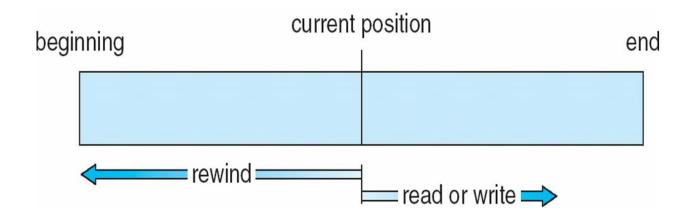
```
read n
write n
position to n
          read next
          write next
rewrite n
```

n = relative block number

- Relative block numbers allow OS to decide where file should be placed
  - See allocation problem in Ch.14



## Sequential-access File





## Simulation of Sequential Access on Direct-access File

sequential access	implementation for direct access	
reset	cp = 0;	
read next	read cp; cp = cp + 1;	
write next	write $cp$ ; $cp = cp + 1$ ;	





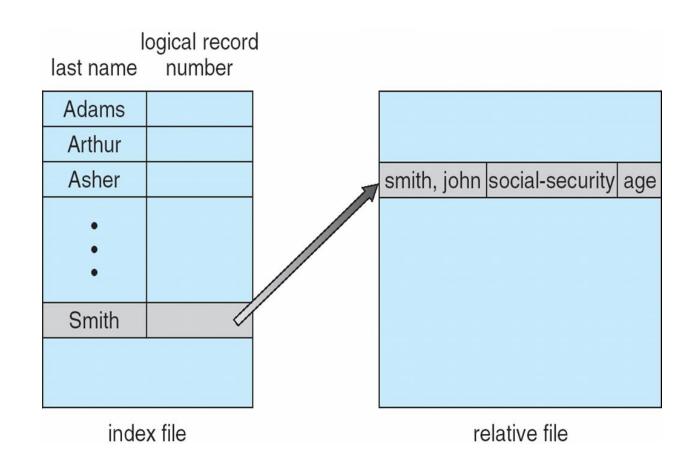
#### **Other Access Methods**

- Can be built on top of base methods
- Generally involve creation of an index for the file
  - Keep index in memory for fast determination of location of data to be operated on (consider UPC code plus record of data about that item)
- If too large, index (in memory) of the index (on disk)
  - E.g. IBM indexed sequential-access method (ISAM)
    - Small master index, points to disk blocks of secondary index
    - File kept sorted on a defined key
    - All done by the OS
- VMS operating system provides index and relative files as another example (see next slide)





## **Example of Index and Relative Files**

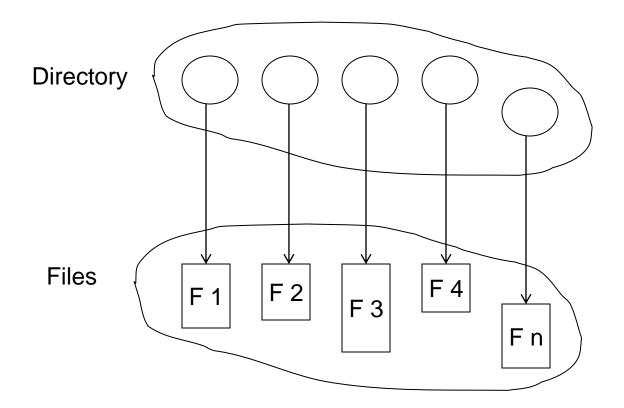




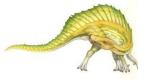


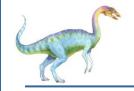
## **Directory Structure**

A collection of nodes containing information about all files



Both the directory structure and the files reside on disk

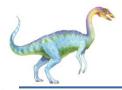




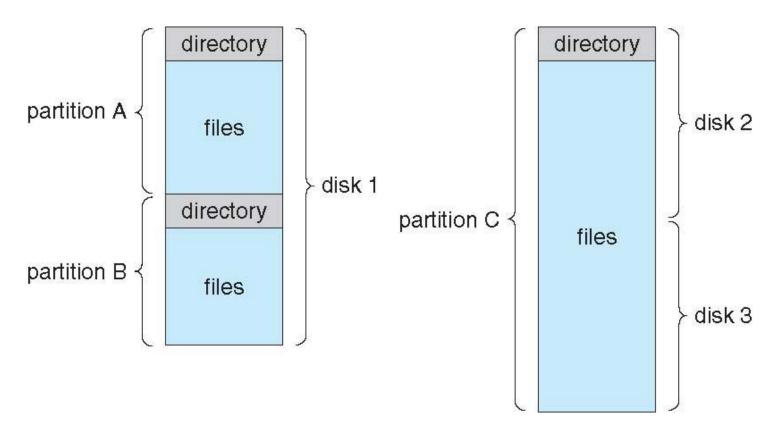
#### **Disk Structure**

- Disk can be subdivided into partitions
  - Disks or partitions can be RAID protected against failure
  - Disk or partition can be used raw without a file system, or formatted with a file system
  - Partitions also known as minidisks, slices
- Entity containing file system known as a volume
  - Each volume containing file system also tracks that file system's info in device directory or volume table of contents
- As well as general-purpose file systems there are many special-purpose file systems, frequently all within the same OS or computer

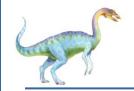




# A Typical File-system Organization



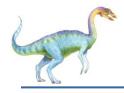




## **Types of File Systems**

- We mostly talk of general-purpose file systems
  - But systems frequently have may file systems, some general- and some special- purpose
- Consider Solaris has
  - tmpfs memory-based volatile FS for fast, temporary I/O
  - objfs interface into kernel memory to get kernel symbols for debugging
  - ctfs contract file system for managing daemons
  - lofs loopback file system allows one FS to be accessed in place of another
  - procfs kernel interface to process structures
  - ufs, zfs general purpose file systems

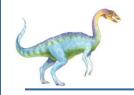




# **Operations Performed on Directory**

- Search for a file
- Create a file
- Delete a file
- List a directory
- Rename a file
- Traverse the file system





## **Directory Organization**

#### The directory is organized logically to obtain

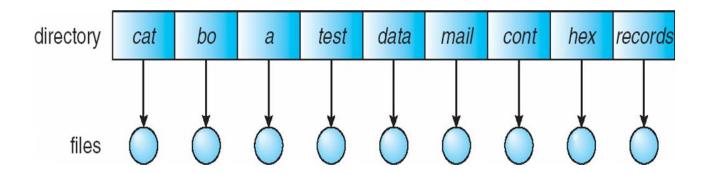
- Efficiency locating a file quickly
- Naming convenient to users
  - Two users can have the same name for different files
  - The same file can have several different names
- Grouping logical grouping of files by properties, (e.g., all Java programs, all games, ...)





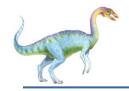
## **Single-Level Directory**

A single directory for all users



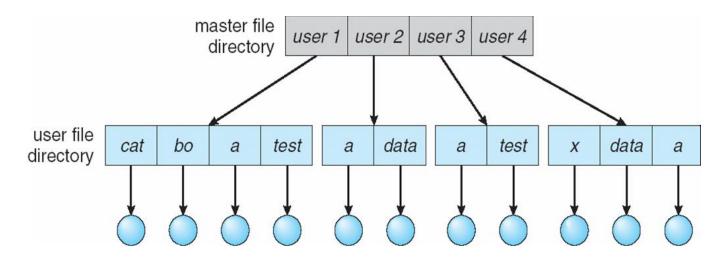
- Naming problem
- Grouping problem



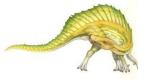


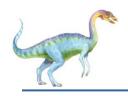
## **Two-Level Directory**

Separate directory for each user

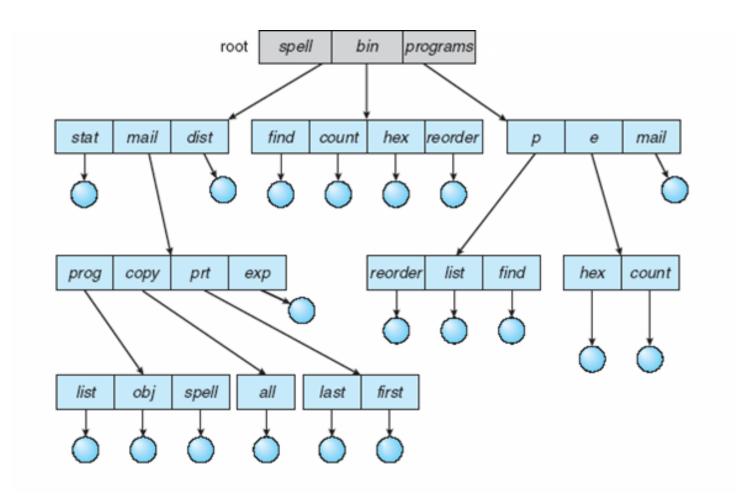


- Path name
- Can have the same file name for different user
- Efficient searching
- No grouping capability

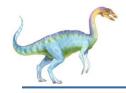




## **Tree-Structured Directories**







## **Tree-Structured Directories (Cont.)**

- Efficient searching
- Grouping Capability
- Current directory (working directory)
  - cd /spell/mail/prog
  - type list





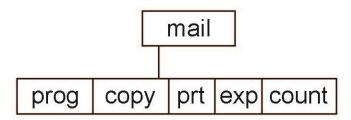
## **Tree-Structured Directories (Cont.)**

- Absolute or relative path name
- Creating a new file is done in current directory
- Delete a file

Creating a new subdirectory is done in current directory

Example: if in current directory /mail

mkdir count

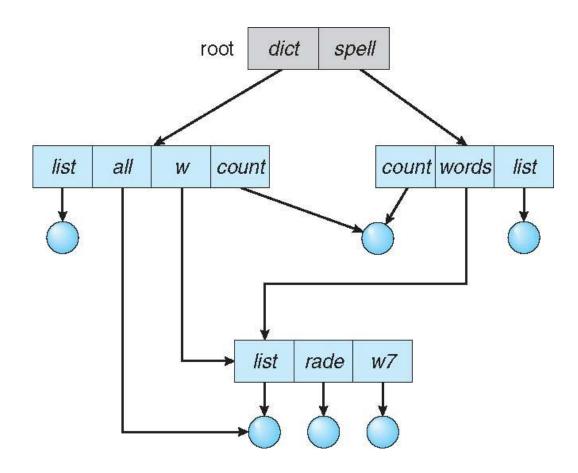


Deleting "mail" ⇒ deleting the entire subtree rooted by "mail"

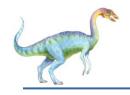


## **Acyclic-Graph Directories**

Have shared subdirectories and files





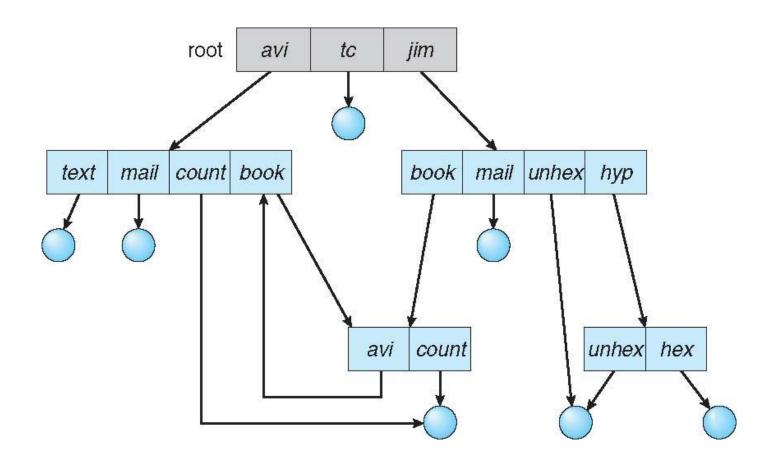


## **Acyclic-Graph Directories (Cont.)**

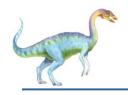
- Two different names (aliasing)
- If dict deletes list ⇒ dangling pointer Solutions:
  - Backpointers, so we can delete all pointers
    - Variable size records problem
  - Backpointers using a daisy chain organization
    - Entry-hold-count solution
- New directory entry type
  - Link another name (pointer) to an existing file
  - Resolve the link follow pointer to locate the file



## **General Graph Directory**



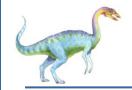




## **General Graph Directory (Cont.)**

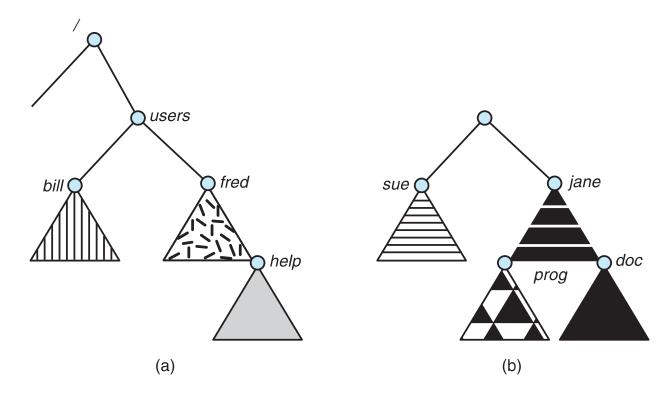
- How do we guarantee no cycles?
  - Allow only links to file not subdirectories
  - Garbage collection
  - Every time a new link is added use a cycle detection algorithm to determine whether it is OK





## **File System Mounting**

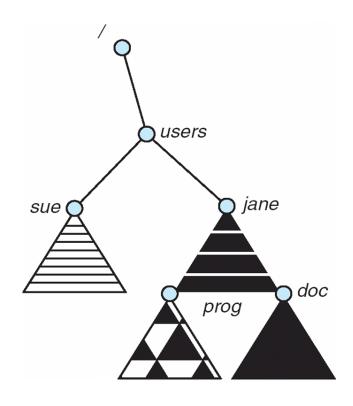
- A file system must be mounted before it can be accessed
- A unmounted file system (i.e., Fig. 11-11(b)) is mounted at a mount point



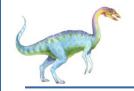




## **Mount Point**

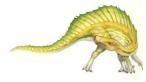






## File Sharing

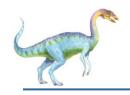
- Sharing of files on multi-user systems is desirable
- Sharing may be done through a protection scheme
- On distributed systems, files may be shared across a network
- Network File System (NFS) is a common distributed filesharing method
- If multi-user system
  - User IDs identify users, allowing permissions and protections to be per-user
     Group IDs allow users to be in groups, permitting group access rights
  - Owner of a file / directory
  - Group of a file / directory





## File Sharing – Remote File Systems

- Uses networking to allow file system access between systems
  - Manually via programs like FTP
  - Automatically, seamlessly using distributed file systems
  - Semi automatically via the world wide web
- Client-server model allows clients to mount remote file systems from servers
  - Server can serve multiple clients
  - Client and user-on-client identification is insecure or complicated
  - NFS is standard UNIX client-server file sharing protocol
  - CIFS is standard Windows protocol
  - Standard OS file system calls are translated into remote calls
- Distributed Information Systems (distributed naming services) such as LDAP, DNS, NIS, Active Directory implement unified access to information needed for remote computing



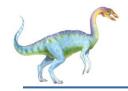
## File Sharing – Failure Modes

- All file systems have failure modes
  - For example corruption of directory structures or other nonuser data, called metadata
- Remote file systems add new failure modes, due to network failure, server failure
- Recovery from failure can involve state information about status of each remote request
- Stateless protocols such as NFS v3 include all information in each request, allowing easy recovery but less security





- Specify how multiple users are to access a shared file simultaneously
  - Similar to Ch.6 process synchronization algorithms
    - Tend to be less complex due to disk I/O and network latency (for remote file systems)
  - Unix file system (UFS) implements UNIX semantics:
    - Writes to an open file visible immediately to other users of the same open file
    - Sharing file pointer to allow multiple users to read and write concurrently
  - Andrew File System (AFS) implemented complex remote file sharing semantics - session semantics
    - Writes only visible to sessions starting after the file is closed



#### **Protection**

- File owner/creator should be able to control:
  - what can be done
  - by whom
- Types of access
  - Read
  - Write
  - Execute
  - Append
  - Delete
  - List



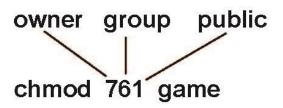


## **Access Lists and Groups**

- Mode of access: read, write, execute
- Three classes of users on Unix / Linux

		RWX
7	$\Rightarrow$	111
		RWX
6	$\Rightarrow$	110
		RWX
1	$\Rightarrow$	0 0 1
	7 6 1	

- Ask manager to create a group (unique name), say G, and add some users to the group
- For a particular file (say game) or subdirectory, define an appropriate access



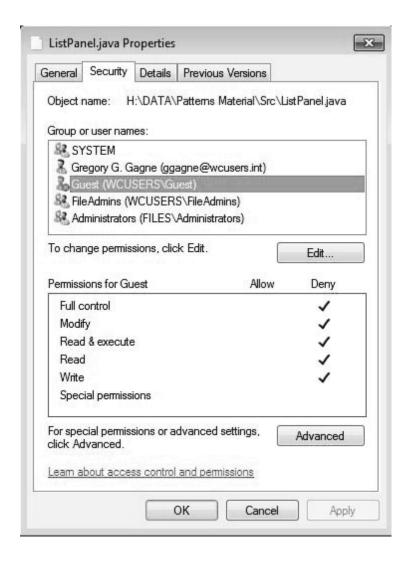
Attach a group to a file

chgrp G game

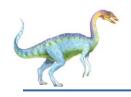




## **Windows Access-Control List Management**







# **A Sample UNIX Directory Listing**

-rw-rw-r	1 pbg	staff	31200	Sep 3 08:30	intro.ps
drwx	5 pbg	staff	512	Jul 8 09.33	private/
drwxrwxr-x	2 pbg	staff	512	Jul 8 09:35	doc/
drwxrwx	2 pbg	student	512	Aug 3 14:13	student-proj/
-rw-rr	1 pbg	staff	9423	Feb 24 2003	program.c
-rwxr-xr-x	1 pbg	staff	20471	Feb 24 2003	program
drwxxx	4 pbg	faculty	512	Jul 31 10:31	lib/
drwx	3 pbg	staff	1024	Aug 29 06:52	mail/
drwxrwxrwx	3 pbg	staff	512	Jul 8 09:35	test/



# **End of Chapter 13**

