

Geometry Lab - Computer Graphics & Robotics - Sept 13 - Blog

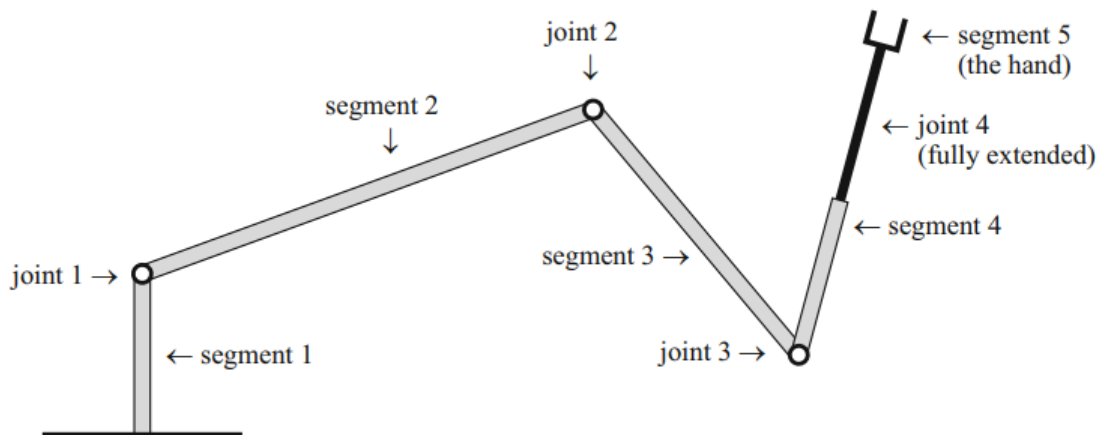
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Forward Kinematic Problem #1

1 Problem statement:

Given a robot with a list of segments, joints, and prismatic joint, can we give an explicit description of the areas that the robot can reach:



Example 1

2 Reading summary:

We can parametrize the circular area of reach of the i th joint to be S^i or by the interval $[0, 2\pi]$ with the endpoints identified.

Similarly, we can use the I^i to describe the reach of the i th prismatic joint.

With that, possible settings of the whole collection of joints in a robot with r revolute joints and p prismatic joints can be parametrized by the Cartesian product:

$$J = S^1 \times \cdots \times S^r \times I_1 \times \cdots \times I_p$$

We will call J the joint space of the robot.

We also represent possible positions of the “hand” by the points (a, b) of a region U .

We can also call $C = U \times V$ the configuration space or operational space of the robot’s hand.

3 Exercise Problem:

Give descriptions of the joint space J and the configuration space C for the planar robot picture in Example 1 in the text. For your description of C , determine a bounded subset of $U \in \mathbb{R}^2$ containing all possible hand positions. Hint: The description of U will depend on the lengths of the segments.

4 Discussion and discovery:

4.1 Describing U with S and I :

By the description of J , we can use S and I to describe the coordinate (x, y) :

$$\begin{aligned} x &= S^1 + S^2 + S^3 + I_1 \\ y &= S^1 + S^2 + S^3 + I_1 + L \text{ (Segment 1 offset)} \end{aligned}$$

4.2 Describing U with r and θ :

Let θ_i be the i th angle formed between joints

Let l_i be the length of i th prismatic joints

Let S_i be the length of the i th segment length

In this problem we can define our set of variables $J = (\theta_1, \theta_2, \theta_3, L_1)$ where $\theta_1, \theta_2, \theta_3$ ranges are $[0, 2\pi]$, and L_1 has range of $[0, l_1]$

Using trigonometry, we can obtain that the difference of (x, y) coordinate of the i th joints be $\Delta x = S_i \cos(\theta_i)$ and $\Delta y = S_i \sin(\theta_i)$

Therefore, we can represent the final position of the robot arm in the example with:

$$\begin{aligned} x &= S_2 \cos(\theta_1) + S_3 \cos(\theta_2) + (S_4 + L_1) \cos(\theta_3) \\ y &= S_1 + S_2 \sin(\theta_1) + S_3 \sin(\theta_2) + (S_4 + L_1) \sin(\theta_3) \end{aligned}$$

4.3 Describing U by inequalities

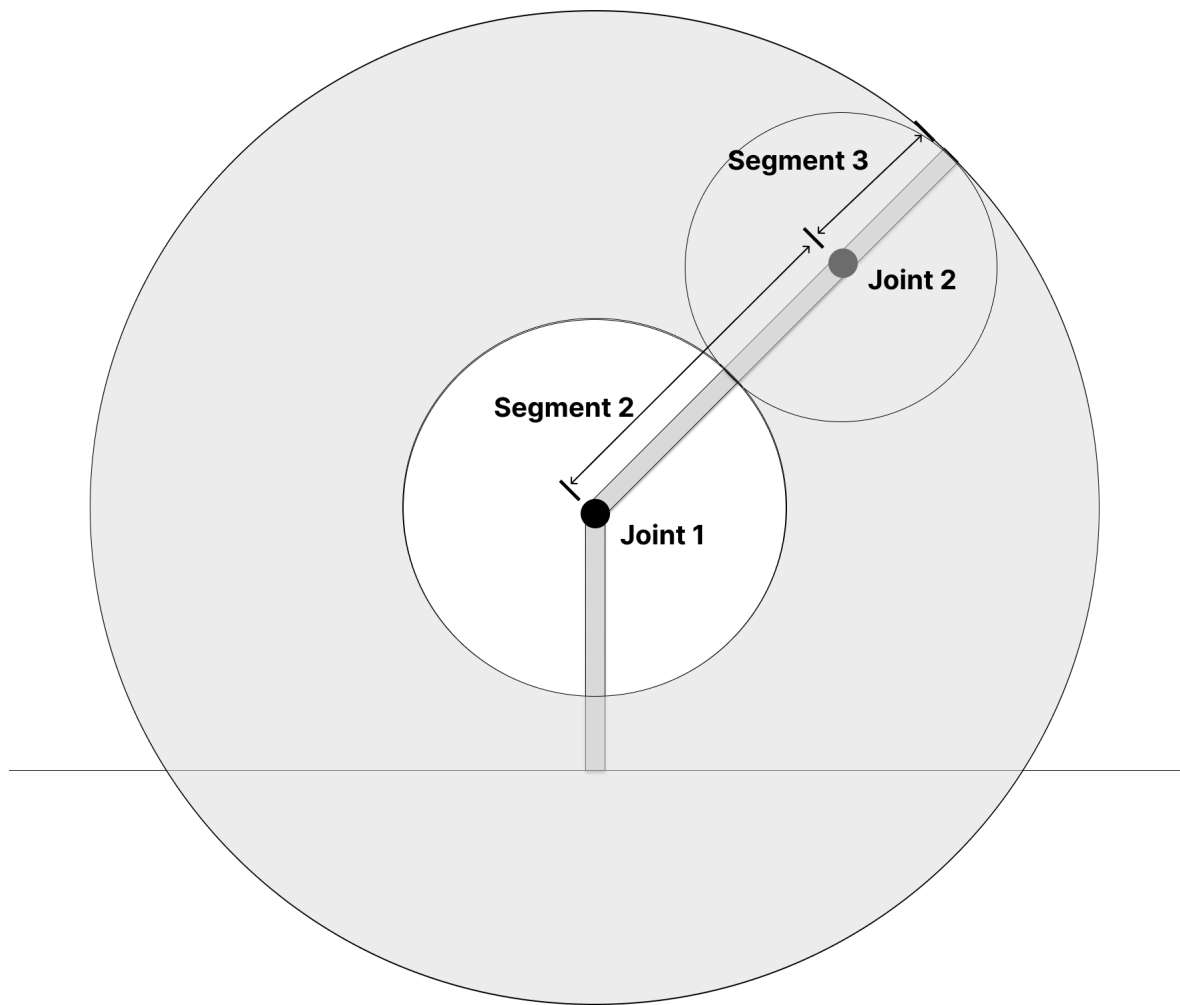
Assume that all joints are able to rotate from 0 to 2π Using inequalities, we can obtain U as:

$$U = \{(x, y) \in \mathbb{R}^2 | \text{lower bound} \leq (x - x_0)^2 + (y - y_0)^2 \leq \text{upper bound}\}$$

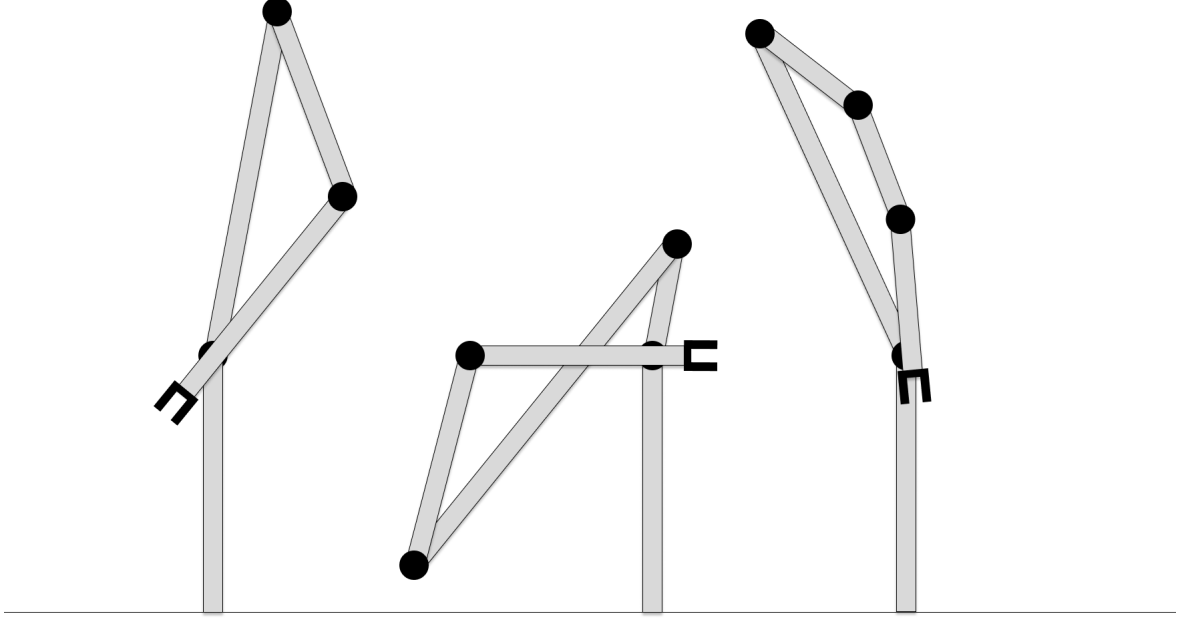
For the upper bound, we can observe that the farthest distance the robot can reach is when all of the segment of the robot points to the same direction, where their length and upper bound becomes:

$$\begin{aligned} length &= \sum_{i=1}^n l_i \\ upperbound &= (length)^2 = \left(\sum_{i=1}^n l_i\right)^2 \end{aligned}$$

For the lower bound, we are considering the case where due to the length of one particular segment, we are not able to reach the center space. In the picture below, segment 3 is shorter than segment 2, causing only the shaded area to be the reachable area, leaving a small circular area unreachable in the middle:



We notice that the case of unreachable area only appears when the sum of all smaller segments are less than the longest segments. Because if the largest side is smaller than the sum of the other sides then we can construct a polygon or a crossed polygon to reach back to the origin:



When constructing the lower bound, we also want to consider that the lower bound cannot be negative, the least it can be is 0, or no lower bound. Therefore, we can define our lower bound to be:

$$lowerbound = (\max(2 \cdot \max(l) - \sum^i l_i, 0))^2$$

Then we define U to be:

$$U = \{(x, y) \in \mathbb{R}^2 | (\max(2 \cdot \max(l) - \sum^i l_i, 0))^2 \leq (x - x_0)^2 + (y - y_0)^2 \leq (\sum^i l_i)^2\}$$