

Assignment 2

Question 1

1)

Given

$F = \{A \rightarrow DE, B \rightarrow GI, E \rightarrow CD, CE \rightarrow ADH, H \rightarrow G, AH \rightarrow I\}$,

derive $A \rightarrow I$:

1. $CE \rightarrow ADH$ (given)
 2. $CE \rightarrow AH$ (by F5(Projectivity) from 1)
 3. $AH \rightarrow I$ (given)
 4. $CE \rightarrow I$ (by F3(Transitivity) from 2 and 3)
 5. $A \rightarrow DE$ (given)
 6. $A \rightarrow E$ (by F5(Projectivity) from 5)
 7. $E \rightarrow CD$ (given)
 8. $E \rightarrow C$ (by F5(Projectivity) from 7)
 9. $A \rightarrow C$ (by F3(Transitivity) from 6 and 8)
 10. $AA \rightarrow CE$ (by F2((Augmentation) from 6 and 9); that is, $A \rightarrow CE$)
 11. $A \rightarrow I$ (by F3(Transitivity) from 10 and 4)
- Therefore, $A \rightarrow I \in F^+$.

2)

Step 1:

Let $X := \{A, B, E, C, H, J\}$

Step 2:

Try to remove A, H, C

$\{B, E, C, H, J\}^+ = \{A, B, C, D, E, G, H, I, J\}$

Thus $X := \{B, E, C, H, J\}$

$\{B, E, C, J\}^+ = \{A, B, C, D, E, G, H, I, J\}$

Thus $X := \{B, E, C, J\}$

$\{B, E, J\}^+ = \{A, B, C, D, E, G, H, I, J\}$

Thus $X := \{B, E, J\}$

Step 3:

Try to remove B, E, J

$\{E, J\}^+ = \{A, C, D, E, G, H, I, J\}$

$\{B, J\}^+ = \{B, G, I, J\}$

$\{B, E\}^+ = \{A, B, C, D, E, G, H, I\}$

Thus cannot be removed

So $\{B, E, J\}$ is a candidate key for R

3)

1. According to 1):

We can prove:

1. $CE \rightarrow ADH$ (given)
2. $CE \rightarrow AH$ (by F5(Projectivity) from 1)
3. $AH \rightarrow I$ (given)
4. $CE \rightarrow I$ (by F3(Transitivity) from 2 and 3)
5. $E \rightarrow CD$ (given)
6. $E \rightarrow C$ (by F5(Projectivity) from 5)
7. $EE \rightarrow CE$ (by F2((Augmentation) from 6); that is, $E \rightarrow CE$)
8. $E \rightarrow I$ (by F3 (Transitivity) from 7 and 4)

Thus:

The attributes of I are transitively dependent on E .

Therefore:

R with respect to F does not satisfies Third Normal Form (3NF)

2. Because:

$E \rightarrow CD \in F^+$ (According to 1))

$\{B, E, J\}$ is a candidate key for R (According to 2)

$E \subset BEJ$

Thus,

CD is partially dependent on BEJ

Therefore,

R with respect to F does not satisfies Second Normal Form (2NF)

3. All of attribute values are atomic

Therefore,

R with respect to F satisfies First Normal Form (1NF)

Hence, the highest normal form of R with respect to F is First Normal Form (1NF)

4)

$R = \{A, B, C, D, E, G, H, I, J\}$

$F = \{A \rightarrow DE, B \rightarrow GI, E \rightarrow CD, CE \rightarrow ADH, H \rightarrow G, AH \rightarrow I\}$

Step 1:

$F' = \{A \rightarrow D, A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D,$

$CE \rightarrow A, CE \rightarrow D, CE \rightarrow H, H \rightarrow G, AH \rightarrow I\}$

Step 2:

$CE \rightarrow A$

$\{C\}^+ = \{C\}$; thus $C \rightarrow A$ is not inferred by F' .

Hence, $CE \rightarrow A$ cannot be replaced by $C \rightarrow A$.

$\{E\}^+ = \{A, C, D, E, G, H, I\}$; thus $E \rightarrow A$ is inferred by F' .

Hence, $CE \rightarrow A$ can be replaced by $E \rightarrow A$.

$F'' = \{A \rightarrow D, A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D,$

$E \rightarrow A, CE \rightarrow D, CE \rightarrow H, H \rightarrow G, AH \rightarrow I\}$

$CE \rightarrow D$

$\{C\}^+ = \{C\}$; thus $C \rightarrow D$ is not inferred by F' .
 Hence, $CE \rightarrow D$ cannot be replaced by $C \rightarrow D$.
 $\{E\}^+ = \{A, C, D, E, G, H, I\}$; thus $E \rightarrow D$ is inferred by F' .
 Hence, $CE \rightarrow D$ can be replaced by $E \rightarrow D$.
 $F'' = \{A \rightarrow D, A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D,$
 $E \rightarrow A, E \rightarrow D, CE \rightarrow H, H \rightarrow G, AH \rightarrow I\}$

$CE \rightarrow H$

$\{C\}^+ = \{C\}$; thus $C \rightarrow H$ is not inferred by F' .
 Hence, $CE \rightarrow A$ cannot be replaced by $C \rightarrow H$.
 $\{E\}^+ = \{A, C, D, E, G, H, I\}$; thus $E \rightarrow H$ is inferred by F' .
 Hence, $CE \rightarrow H$ can be replaced by $E \rightarrow H$.
 $F'' = \{A \rightarrow D, A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D,$
 $E \rightarrow A, E \rightarrow D, E \rightarrow H, H \rightarrow G, AH \rightarrow I\}$

$AH \rightarrow I$

$\{H\}^+ = \{H, G\}$; thus $H \rightarrow I$ is not inferred by F' .
 Hence, $AH \rightarrow I$ cannot be replaced by $H \rightarrow I$.
 $\{A\}^+ = \{A, C, D, E, G, H, I\}$; thus $A \rightarrow I$ is inferred by F' .
 Hence, $AH \rightarrow I$ can be replaced by $A \rightarrow I$.
 $F'' = \{A \rightarrow D, A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D,$
 $E \rightarrow A, E \rightarrow H, H \rightarrow G, A \rightarrow I\}$

Step 3:

$A^+|_{F'' - \{A \rightarrow D\}} = \{A, C, D, E, G, H, I\}$; thus $A \rightarrow D$ is inferred by

$$F'' - \{A \rightarrow D\}$$

That is, $A \rightarrow D$ is redundant.

Thus, we can remove $A \rightarrow D$ from F'' to obtain F''' .

$A^+|_{F'' - \{A \rightarrow E\}} = \{A, D\}$; thus $A \rightarrow D$ is not inferred by

$$F'' - \{A \rightarrow E\}$$

That is, $A \rightarrow E$ is not redundant.

Iteratively,

$$F''' = \{A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D,$$

$$E \rightarrow A, E \rightarrow H, H \rightarrow G, A \rightarrow I\}$$

$$\text{Thus, } F_{min} = \{A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D,$$

$$E \rightarrow A, E \rightarrow H, H \rightarrow G, A \rightarrow I\}$$

5)

$R = (A, B, C, D, E, G, H, I, J)$

$F_{min} = \{A \rightarrow E, B \rightarrow G, B \rightarrow I, E \rightarrow C, E \rightarrow D, \\ E \rightarrow A, E \rightarrow H, H \rightarrow G, A \rightarrow I\}$

Candidate key: (B, E, J)

$R_1 = (B, G, I)$

$R_2 = (A, E, I)$

$R_3 = (E, C, D, A, H)$

$R_4 = (H, G)$

$R_5 = (B, E, J)$

	A	B	C	D	E	G	H	I	J
(B, G, I)	b	a	b	b	b	a	b	a	b
(A, E, I)	a	b	b	b	a	b	b	a	b
(E, C, D, A, H)	a	b	a	a	a	b	a	b	b
(H, G)	b	b	b	b	b	a	a	b	b
(B, E, J)	b	a	b	b	a	a	b	b	b

According to the rule of testing for the lossless join property:

Base on key (B, E, J) :

	A	B	C	D	E	G	H	I	J
(B, G, I)	b	a	b	b	b	a	b	a	b
(A, E, I)	a	b	b	b	a	b	b	a	b
(E, C, D, A, H)	a	b	a	a	a	b	a	b	b
(H, G)	b	b	b	b	b	a	a	b	b
(B, E, J)	a	a	a	a	a	a	a	a	a

Therefore,

The decomposition of R_1, R_2, R_3, R_4, R_5 is lossless-join.

R_1 can infer $F_1 = \{B \rightarrow G, B \rightarrow I\}$

R_2 can infer $F_2 = \{A \rightarrow E, A \rightarrow I\}$

R_3 can infer $F_3 = \{E \rightarrow C, E \rightarrow D, E \rightarrow A, E \rightarrow H\}$

R_4 can infer $F_4 = \{H \rightarrow G\}$

R_5 can infer $F_5 = \emptyset$

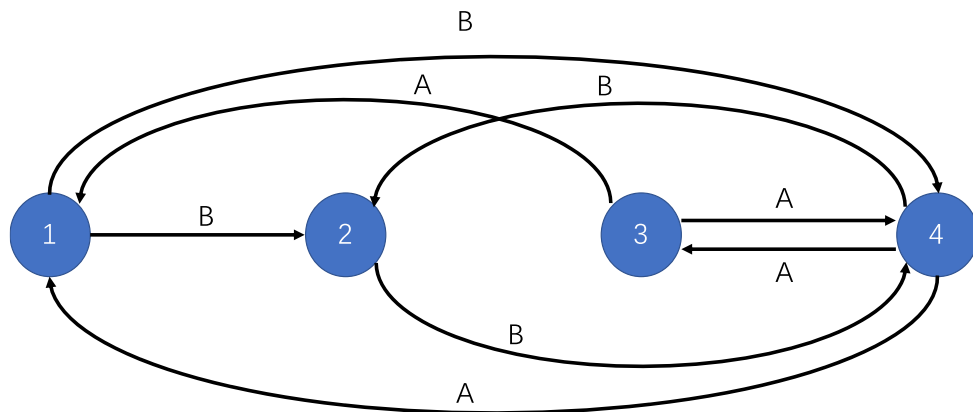
$F^+ = (F_1 \cup F_2 \cup F_3 \cup F_4 \cup F_5)^+$

Therefore,

The decomposition of R_1, R_2, R_3, R_4, R_5 is dependency-preserving.

Question 2

1)



According to the schedule graph, the graph is a cyclic.

Therefore, the transaction schedule is not conflict serializable.

2)

T ₁	T ₂	T ₃	T ₄
R(B)			
W(B)			
R(A)			
W(A)			
	R(B)		
	W(B)		
		R(A)	
		W(A)	
			R(A)
			W(A)
			R(B)
			W(B)

3)

T ₁	T ₂
Write_lock(B)	
Write_lock(A)	
R(B)	
R(A)	
W(B)	
Unlock(B)	
W(A)	Write_lock(B)
Unlock(A)	R(B)
	W(B)
	Unlock(B)