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## 1 Basic

### 1.1 Default code

```
#include<bits/stdc++.h>
// #pragma GCC optimize("O3,unroll-loops")
// #pragma target optimize("avx2,bmi,bmi2,lzcnt,popcnt")
#define IO ios_base::sync_with_stdio(0);cin.tie(0);cout.tie(0);
#define pii pair<int,int>
#define ft first
#define sd second
#define int long long
#define ld long double
#define PI acos(-1)
#define SZ(x) (int)x.size()
#define all(v) (v).begin(), (v).end()
#define _for(i,a,b) for(int i=(a);i<(b);++i)
using namespace std;
template<typename T>
ostream& operator<<(ostream& os,const vector<T>& vn){
    for(int i=0;i<vn.size();i++)os<<vn[i]<<" ";
    return os;
}
template<typename T>
ostream& operator<<(ostream& os,const set<T>& vn){
    for(typename set<T>::iterator it=vn.begin();it!=vn.end();it++)os<<*it<<" ";
    return os;
}
mt19937 mt(hash<string>{}("Mashu")); //mt();
// mt19937 mt(chrono::steady_clock::now().time_since_epoch().count());
// g++ a.cpp -Wall -Wshadow -fsanitize=undefined -o a.exe
// ./a.exe
const int INF=INT_MAX;
void sol() {}
signed main() {
    // auto start=chrono::high_resolution_clock::now();
    // #ifdef LOCAL
    // freopen("input.txt","r",stdin);
    // freopen("output.txt","w",stdout);
    // #endif
    IO
    int t=1;
    // cin>>t;
    while(t--){sol();}
    // auto stop = chrono::high_resolution_clock::now();
    // auto duration = chrono::duration_cast<chrono::milliseconds>(stop - start);
    // cerr<<"Time:"<<duration.count()<<" ms\n";
}
```

### 1.2 Misc

```
iota(vec.begin(),vec.end(),1);// 產生1~size的整數列
stoi(s.begin(),s.end(),k);// 法1,字串轉成k進位int
string s;cin>>s;
int x=stoi(s,0,2); // 法2,2可以改其他進位
int bbb = bitset<10>(bb).to_ulong();//二進位轉十進位
__builtin_popcountll // 二進位有幾個1
__builtin_clzll // 左起第一個1前0的個數
__builtin_parityll // 1的個數的奇偶性
__builtin_mul_overflow(a,b,&res) // a*b是否溢位

// double 轉整數 請加 int b=round(a)
// 或是 int b =floor(a+0.5) (floor向下取整)
```

### 1.3 Fast read & write

```
inline int read() {
    char c = getchar(); int x = 0, f = 1;
    while(c < '0' || c > '9') {if(c == '-') f = -1; c = getchar();}
```

```

while(c >= '0' && c <= '9') x = x * 10 + c - '0', c = getchar();
return x * f;
}
inline void write(int x){
    if(x<0) putchar('-'),x=-x;
    if(x>9) write(x/10);
    putchar(x%10+'0');
}

```

## 1.4 Sort cmp

```

struct cmp{inline bool operator()(const int a,const int b){return a<b;}}; //common use
auto cmp= [](vector<int> a, vector<int> b) {return a[1]<b[1];}; //for set use
set<vector<int>, decltype(cmp)> prepare, done;

```

## 1.5 Discretization

```

vector<int> vec;
sort(vec.begin(),vec.end());
vec.resize(unique(vec.begin(),vec.end())-vec.begin());
for(int i=0;i<n;++i){ //+1是讓 index是1到N 可以不要
    arr[i]=lower_bound(vec.begin(),vec.end(),ll[i])-vec.begin()+1;
}

```

## 1.6 Custom unordered\_map

```

struct Type{
    int x;
    string y;
    bool operator==(const Type &other) const {
        return (x == other.x && y == other.y);
    }
};
struct hashes{
    size_t operator()(const Type &o) const {
        return ((hash<int>()(o.x)^(hash<string>()(o.y)<<1))>>1);
    }
};
//unordered_map<Type,int,hashes> map;

```

## 1.7 \_\_int128 read

```

// __int128_t p;
// lll n=qr(p);
#define lll __int128
template<class type_name> inline type_name qr(type_name sample)
{
    type_name ret=0,sgn=1;
    char cur=getchar();
    while(!isdigit(cur))
        sgn=(cur=='-'?-1:1),cur=getchar();
    while(isdigit(cur))
        ret=(ret<<1)+(ret<<3)+cur-'0',cur=getchar();
    return sgn==-1?-ret:ret;
}
inline void print(__int128 x){
    if(x < 0){
        putchar('-');
        x = -x;
    }
    if(x > 9)
        print(x / 10);
    putchar(x % 10 + '0');
}

```

## 1.8 字典序 a 嚴格小於 b

```

template<class T> //字典序a嚴格小於b
bool lexicographicallySmaller(const vector<T> &a,const vector<T> &b){
    int n=a.size();
    int m=b.size();
    int i;
    for(int i=0;i<n && i<m;++i){
        if(a[i]<b[i])return true;
        else if(b[i]<a[i])return false;
    }
    return (i==n && i<m);
}

```

## 1.9 生成 n 位數的二進制組合

```

//產生1~n位數的二進位組合
vector<int> comb;
void genBinComb(int n){
    for(int i=1;i<=(1<n);++i){
        string s = bitset<5>(i).to_string();//bitset大
        //小要記得設
        comb.push_back(stoi(s));
    }
}

```

## 1.10 Radom

```

mt19937 gen(0x5EED);
int randint(int lb, int ub)
{ return uniform_int_distribution<int>(lb, ub)(gen); }

```

## 2 對拍

### 2.1 run.bat

```

@echo off
g++ ac.cpp -o ac.exe
g++ wa.cpp -o wa.exe
g++ gen1.cpp -o gen.exe

:loop
    echo %%x
    gen.exe > input
    ac.exe < input > ac
    wa.exe < input > wa
    fc ac wa
    if not errorlevel 1 goto loop

```

### 2.2 run.sh

```

for ((i=0;;i++))
do
    echo "$i"
    python3 gen.py > input
    ./ac < input > ac.out
    ./wa < input > wa.out
    diff ac.out wa.out || break
done

```

## 3 Flow & Matching

### 3.1 Dinic

```

// flow.init(n,s,t):有n個點(0~n-1)，起點s終點t
// flow.add_edge(u,v,f):建一條邊，從u點到v點流量為f
// flow.solve():回傳網路最大流答案
//時間複雜度: O(V^2 * E)
struct Dinic{
    struct Edge{ int v,f,re; };
    int n,s,t,level[MXN];
    vector<Edge> E[MXN];
}

```

```

void init(int _n, int _s, int _t){
    n = _n; s = _s; t = _t;
    for (int i=0; i<n; i++) E[i].clear();
}
void add_edge(int u, int v, int f){
    E[u].push_back({v,f,(int)(E[v].size())});
    E[v].push_back({u,0,(int)(E[u].size()-1)});
}
bool BFS(){
    for (int i=0; i<n; i++) level[i] = -1;
    queue<int> que;
    que.push(s);
    level[s] = 0;
    while (!que.empty()){
        int u = que.front(); que.pop();
        for (auto it : E[u]){
            if (it.f > 0 && level[it.v] == -1){
                level[it.v] = level[u]+1;
                que.push(it.v);
            }
        }
    }
    return level[t] != -1;
}
int DFS(int u, int nf){
    if (u == t) return nf;
    int res = 0;
    for (auto &it : E[u]){
        if (it.f > 0 && level[it.v] == level[u]+1){
            int tf = DFS(it.v, min(nf,it.f));
            res += tf; nf -= tf; it.f -= tf;
            E[it.v][it.re].f += tf;
            if (nf == 0) return res;
        }
    }
    if (!res) level[u] = -1;
    return res;
}
int solve(int res=0){
    while (BFS())
        res += DFS(s,2147483647);
    return res;
} }flow;

```

```

        if(!vis[it.v]) {vis[it.v] = true; q
            .push(it.v);}
    }
}
if(dis[t]==LLONG_MAX) return false;
// 不管流量是多少，花費不能是正數時加上這行（最
// 小花費可行流）
// if(dis[t] >= 0) return false;
return true;
}
int DFS(int u, int nf) {
    if(u==t) return nf;
    int res = 0; vis[u] = true;
    for(int &i=ptr[u] ; i<(int)E[u].size() ; i++) {
        auto &it = E[u][i];
        if(it.f>0 && dis[it.v]==dis[u]+it.w && !vis
            [it.v]) {
            int tf = DFS(it.v, min(nf, it.f));
            res += tf;
            nf-=tf;
            it.f-=tf;
            E[it.v][it.re].f += tf;
            if(nf==0) { vis[u]=false; break; }
        }
    }
    return res;
}
pair<int,ll> solve(){
    int flow = 0; ll cost = 0;
    while (SPFA()){
        fill_n(ptr, n, 0);
        int f = DFS(s, INT_MAX);
        flow += f;
        cost += dis[t]*f;
    }
    return {flow, cost};
} // reset: do nothing
} flow;

```

### 3.2 最大流最小花費

```

//最大流量上的最小花費
//最大流量優先，相同才是找最小花費，複雜度 $O(V^2 \cdot E^2)$ 
// flow.init(n,s,t):有n個點(0~n-1)，起點s終點t
// flow.add_edge(u,v,f,c):建一條邊，從u點到v點流量為f，
// 每一單位流量的花費為c
// flow.solve():回傳一個pair(maxFlow,minCost)
// 限制：圖不能有負環
// 網路最大流的add_edge(u,v,f)可以無痛轉成最大流量上的
// 最小花費add_edge(u,v,1,f)即建立一條從u到v的邊流量為
// 1，單位流量花費為f
// $O(V^2 \cdot E^2)$ 
#define ll long long
struct zkwflow{
    static const int maxN=20000;
    struct Edge{ int v,f,re; ll w;};
    int n,s,t,ptr[maxN]; bool vis[maxN]; ll dis[maxN];
    vector<Edge> E[maxN];
    void init(int _n,int _s,int _t){
        n=_n,s=_s,t=_t;
        for(int i=0;i<n;i++) E[i].clear();
    }
    void add_edge(int u,int v,int f,ll w){
        E[u].push_back({v,f,(int)E[v].size(),w});
        E[v].push_back({u,0,(int)E[u].size()-1,-w});
    }
    bool SPFA() {
        fill_n(dis, n, LLONG_MAX);
        fill_n(vis, n, false);
        queue<int> q;
        q.push(s); dis[s]=0;
        while(!q.empty()) {
            int u = q.front(); q.pop();
            vis[u] = false;
            for(auto &it : E[u]){
                if(it.f>0 && dis[it.v]>dis[u]+it.w){
                    dis[it.v] = dis[u]+it.w;

```

### 3.3 匈牙利匹配

```

//匈牙利演算法-二分圖最大匹配
//記得每次使用需清空vis數組
// $O(nm)$ 
//其中Map為鄰接表(Map[u][v]為u和v是否有連接) S為紀錄這
//個點與誰匹配(S[i]為答案i和誰匹配)
const int M=505, N=505;
bool Map[M][N] = {0};
int S[N];
bool vis[N];
bool dfs(int u){
    for(int i=0;i<N;i++){
        if(Map[u][i]&&!vis[i]){ //有連通且未拜訪
            vis[i]=1; //紀錄是否走過
            if(S[i]==-1||dfs(S[i])){ //紀錄匹配
                S[i]=u;
                return true; //反轉匹配邊以及未匹配邊
                的狀態
            }
        }
    }
    return false;
}
//此二分圖為左邊M個點右邊N個點，跑匈牙利只要跑1~M就可以
//了，(S[右邊的點] -> 左邊的點)
memset(S,-1,sizeof(S));
int ans = 0;
for(int i=0;i<M;i++){
    memset(vis,0,sizeof(vis));
    if(dfs(i)) ans++;
    //跑匈牙利
}
cout<<ans<<"\n";
for(int i=0 ; i<N ; i++) {
    if(S[i]!=-1) cout<<"pair: "<<S[i]<<" "<<i<<"\n";
}

```

### 3.4 KM

```
//二分圖最大權完美匹配
//二分圖左邊的點都要匹配到右邊的點，且每條邊都有權重，
//求權重最大值，複雜度 $O(V^3)$ 
// graph.init(n):二分圖左右各n個點
// graph.add_edge(u,v,w):建一條邊，從u點到v點權重為w
// graph.solve():回傳最大權重
struct KM{ // max weight, for min negate the weights
    int n, mx[MXN], my[MXN], pa[MXN];
    ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
    bool vx[MXN], vy[MXN];
    void init(int _n) { // 1-based, N個節點
        n = _n;
        for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0);
    }
    void add_edge(int x, int y, ll w) {g[x][y] = w;} //
    // 左邊的集合節點x連邊右邊集合節點y權重為w
    void augment(int y) {
        for(int x, z; y; y = z)
            x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
    }
    void bfs(int st) {
        for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]=0;
        queue<int> q; q.push(st);
        for(;;) {
            while(q.size()) {
                int x=q.front(); q.pop(); vx[x]=1;
                for(int y=1; y<=n; ++y) if(!vy[y]){
                    ll t = lx[x]+ly[y]-g[x][y];
                    if(t==0){
                        pa[y]=x;
                        if(!my[y]){augment(y);return;}
                        vy[y]=1, q.push(my[y]);
                    }else if(sy[y]>t) pa[y]=x,sy[y]=t;
                }
            }
            ll cut = INF;
            for(int y=1; y<=n; ++y)
                if(!vy[y]&&cut>sy[y]) cut=sy[y];
            for(int j=1; j<=n; ++j){
                if(vx[j]) lx[j] -= cut;
                if(vy[j]) ly[j] += cut;
                else sy[j] -= cut;
            }
            for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]==0){
                if(!my[y]){augment(y);return;}
                vy[y]=1, q.push(my[y]);
            }
        }
    }
    ll solve(){ // 回傳值為完美匹配下的最大總權重
        fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
        fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
        for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)
            // 1-base
            lx[x] = max(lx[x], g[x][y]);
        for(int x=1; x<=n; ++x) bfs(x);
        ll ans = 0;
        for(int y=1; y<=n; ++y) ans += g[my[y]][y];
        return ans;
    }
} graph;
```

## 4 Graph

### 4.1 Dijkstra

```
int dis[N],vis[N];
void dijkstra(int s){// $O(V^2+E)$ 
    memset(dis,0x3f,sizeof(dis));
    memset(vis,0,sizeof(vis));
    dis[s] = 0;
    priority_queue<pii,vector<pii>,greater<pii>> pq;
    pq.push({0,s});
    while(!pq.empty()){
        auto [nowd,now] = pq.top();
```

```
        pq.pop();
        if(vis[now])
            continue;
        vis[now] = 1;
        for(auto [nxt,w]: graph[now]){
            if(dis[nxt] > dis[now] + w){
                dis[nxt] = dis[now] + w;
                pq.push({dis[nxt],nxt});
            }
        }
    }
}
```

### 4.2 Bellman-Ford

```
//總共m條邊，鬆弛n-1次-> $O(nm)$ 
//在第n次做一次鬆弛，如果有點被鬆弛到，代表這張圖存在負環
for(int j = 0; j < n-1; j++){
    for(int i = 0; i < m; i++){ // 對於所有邊都嘗試鬆弛
        if(dis[ edge[i].to ] > dis[ edge[i].from ] +
            edge[i].weight){
            dis[ edge[i].to ] = dis[ edge[i].from ] +
                edge[i].weight;
        }
    }
}
```

### 4.3 SPFA

```
#define MXN 200005
struct SPFA{// $O(kE)$  k:小常數
    int n;
    ll inq[MXN], len[MXN];
    vector<ll> dis;
    vector<pair<int, ll>> edge[MXN];
    void init(int _n){
        n = _n;
        dis.clear(); dis.resize(n, 1e18);
        for(int i = 0; i < n; i++){
            edge[i].clear();
            inq[i] = len[i] = 0;
        }
    }
    void addEdge(int u, int v, ll w){
        edge[u].push_back({v, w});
    }
    vector<ll> solve(int st = 0){
        deque<int> dq; //return {-1} if has negative cycle
        dq.push_back(st); //otherwise return dis from st
        inq[st] = 1; dis[st] = 0;
        while(!dq.empty()){
            int u = dq.front(); dq.pop_front();
            inq[u] = 0;
            for(auto [to, d] : edge[u]){
                if(dis[to] > d+dis[u]){
                    dis[to] = d+dis[u];
                    len[to] = len[u]+1;
                    if(len[to] > n) return {-1};
                    if(inq[to]) continue;
                    (!dq.empty()&&dis[dq.front()] > dis[to]?
                        dq.push_front(to) : dq.push_back(to));
                    inq[to] = 1;
                }
            }
        }
        return dis;
    }
} spfa;
```

### 4.4 Floyd-Warshall

```
for(int k=0;k<n;++k){// $O(N^3)$ 
    for(int i=0;i<n;++i){
        for(int j=0;j<n;++j){
            if(graph[i][k]!=MAX&&graph[k][j]!=MAX)//避免不連通圖
                graph[i][j]=min(graph[i][j],graph[i][k]
                    +graph[k][j]);
        }
    }
}
```

```

    }
}
for(int k=0;k<n;++k){//判斷負環
    for(int i=0;i<n;++i){
        for(int j=0;j<n;++j){
            if(graph[i][k]!=MAX&&graph[k][j]!=MAX&&
                graph[k][k]<0){//避免不連通圖&&負環
                graph[i][j]=-MAX;
            }
        }
    }
}
// if(graph[a][b]==-MAX)
//     cout<<"-Infinity\n";
// else if(graph[a][b]==MAX)
//     cout<<"Impossible\n";
// else
//     cout<<graph[a][b]<<"\n";

```

## 4.5 歐拉路徑

```

const int MXN = 2e5+7;
struct ola{
    vector<pair<int, int>> edge[MXN];
    int ind[MXN], ru[MXN], use[MXN*3], es=0, n, go;
    void init(int _n) {
        n = _n;
        for(int i=0 ; i<n ; i++) edge[i].clear(), ind[i]=0;
    }
    void add_edge(int a, int b) {
        use[es] = 0;
        edge[a].push_back({b, es});
        edge[b].push_back({a, es});
        ru[a]++; ru[b]++;
    }
    bool check() {
        int cnt = 0;
        for(int i=0 ; i<n ; i++) {
            if(ru[i]&1) go = i, cnt++;
        }
        if(cnt==0 || cnt==2) return true;
        return false;
    }
    vector<int> ans;
    void dfs(int x) {
        for(int i=ind[x] ; i<edge[x].size() ; i=ind[x]) {
            int u = edge[x][i].first, pos = edge[x][i].second;
            ind[x]++;
            if(!use[pos]) {
                use[pos]=1;
                dfs(u);
            }
        }
        ans.push_back(x);
    }
    vector<int> solve() {
        if(!check()) return {-1};
        ans.clear(); dfs(go);
        if(ans.size()!=es+1) return {-1};
        return ans;
    }
} euler;

```

## 4.6 BCC

```

//無向圖上，不會產生割點的連通分量稱為點雙連通分量，
//base
#define PB push_back
#define REP(i, n) for(int i = 0; i < n; i++)
struct BccVertex {
    int n, nScc, step, dfn[MXN], low[MXN];
    vector<int> E[MXN], sccv[MXN];
    int top, stk[MXN];
    void init(int _n) {
        n = _n;

```

```

        nScc = step = 0;
        for (int i = 0; i < n; i++)
            E[i].clear();
    }
    void addEdge(int u, int v) {
        E[u].PB(v); E[v].PB(u);
    }
    void DFS(int u, int f) {
        dfn[u] = low[u] = step++;
        stk[top++] = u;
        for (auto v : E[u]) {
            if (v == f) continue;
            if (dfn[v] == -1) {
                DFS(v, u);
                low[u] = min(low[u], low[v]);
                if (low[v] >= dfn[u]) {
                    int z;
                    sccv[nScc].clear();
                    do {
                        z = stk[--top];
                        sccv[nScc].PB(z);
                    } while (z != v);
                    sccv[nScc++].PB(u);
                }
            }
            else low[u] = min(low[u], dfn[v]);
        }
    }
    vector<vector<int>> solve() { //回傳每個點雙聯通分量
        vector<vector<int>> res;
        for (int i = 0; i < n; i++)
            dfn[i] = low[i] = -1;
        for (int i = 0; i < n; i++)
            if (dfn[i] == -1) {
                top = 0;
                DFS(i, i);
            }
        REP(i, nScc) res.PB(sccv[i]);
        return res;
    }
} graph;

```

## 4.7 SCC

```

//在有向圖裡的任兩點u、v，皆存在至少一條u到v的路徑
//以及v到u的路徑
//fill zero 注意多筆測資要改fill
//注意要0base
#define PB push_back
#define FZ(x) memset(x, 0, sizeof(x))
const int MXN = 1e5;
struct Scc {
    int n, nScc, vst[MXN], bln[MXN]; //nScc 有幾個強連通分量
    vector<int> E[MXN], rE[MXN], vec;
    void init(int _n) {
        n = _n;
        for (int i = 0; i < MXN; i++)
            E[i].clear(), rE[i].clear();
    }
    void addEdge(int u, int v) {
        E[u].PB(v); rE[v].PB(u);
    }
    void DFS(int u) {
        vst[u] = 1;
        for (auto v : E[u])
            if (!vst[v]) DFS(v);
        vec.PB(u);
    }
    void rDFS(int u) {
        vst[u] = 1;
        bln[u] = nScc;
        for (auto v : rE[u])
            if (!vst[v]) rDFS(v);
    }
    void solve() {
        nScc = 0;
        vec.clear();
        FZ(vst);
        for (int i = 0; i < n; i++)

```



```

        if (!vst[i]) DFS(i);
        reverse(vec.begin(), vec.end());
        FZ(vst);
        for (auto v : vec)
            if (!vst[v]) {rDFS(v); nScc++;}
    }
} scc;

```

## 4.8 2SAT

[illegible]

## 4.9 MaximalClique

```
//極大團
//對於一張圖選任意的點子集，如果不能在多選一個點使得選
//的點子集為更大的團
#define N 80
struct MaxClique{ // 0-base
    typedef bitset<N> Int;
    Int lnk[N], v[N];
    int n;
    void init(int _n){
        n = _n;
        for(int i = 0 ; i < n ; i ++){
            lnk[i].reset(); v[i].reset();
        }
    }
    void addEdge(int a , int b)
    { v[a][b] = v[b][a] = 1; }
    int ans , stk[N], id[N] , di[N] , deg[N];
    Int cans;
    void dfs(int elem_num, Int candi, Int ex){
        if(candi.none() && ex.none()){
            cans.reset();
            for(int i = 0 ; i < elem_num ; i ++){
                cans[id[stk[i]]] = 1;
            }
            ans = elem_num; //cans=1 is in maximal clique
            return;
        }
        int pivot = (candi|ex)._Find_first();
        Int smaller_candi = candi & (~lnk[pivot]);
        while(smaller_candi.count()){
            int nxt = smaller_candi._Find_first();
            candi[nxt] = smaller_candi[nxt] = 0;
            ex[nxt] = 1;
            stk[elem_num] = nxt;
            dfs(elem_num+1, candi&lnk[nxt], ex&lnk[nxt]);
        }
    }
    int solve(){
        for(int i = 0 ; i < n ; i ++){
```

```

        id[i] = i; deg[i] = v[i].count();
    }
    sort(id , id + n , [&](int id1, int id2){
        return deg[id1] > deg[id2]; });
    for(int i = 0 ; i < n ; i ++ ) di[id[i]] = i;
    for(int i = 0 ; i < n ; i ++ )
        for(int j = 0 ; j < n ; j ++ )
            if(v[i][j]) lnk[di[i]][di[j]] = 1;
    ans = 1; cans.reset(); cans[0] = 1;
    dfs(0, Int(string(n, '1')), 0);
    return ans;
} }solver;

```

#### 4.10 MaximumClique

```
//最大團：圖上最多可以選幾個點，使選的彼此之間都有連邊
//最大獨立集：圖上最多可以選幾個點，使選的彼此之間都沒有連邊
//最大獨立集通常會轉換為用補圖做最大團
//O(1.1888^n)
#define N 111
struct MaxClique{ // 0-base
    typedef bitset<N> Int;
    Int linkto[N] , v[N];
    int n;
    void init(int _n){
        n = _n;
        for(int i = 0 ; i < n ; i++){
            linkto[i].reset(); v[i].reset();
        }
    }
    void addEdge(int a , int b)
    { v[a][b] = v[b][a] = 1; }
    int popcount(const Int& val)
    { return val.count(); }
    int lowbit(const Int& val)
    { return val._Find_first(); }
    int ans , stk[N];
    int id[N] , di[N] , deg[N];
    Int cans;
    void maxclique(int elem_num, Int candi){
        if(elem_num > ans){
            ans = elem_num; cans.reset();
            for(int i = 0 ; i < elem_num ; i++){
                cans[id[stk[i]]] = 1;
            }
        }
        int potential = elem_num + popcount(candi);
        if(potential <= ans) return;
        int pivot = lowbit(candi);
        Int smaller_candi = candi & (~linkto[pivot]);
        while(smaller_candi.count() && potential > ans){
            int next = lowbit(smaller_candi);
            candi[next] = !candi[next];
            smaller_candi[next] = !smaller_candi[next];
            potential --;
            if(next == pivot || (smaller_candi & linkto[next]
                ).count()){
                stk[elem_num] = next;
                maxclique(elem_num + 1, candi & linkto[next]);
            }
        }
    }
    int solve(){//回傳值為最大團的點數量
        for(int i = 0 ; i < n ; i++){
            id[i] = i; deg[i] = v[i].count();
        }
        sort(id , id + n , [&](int id1, int id2){
            return deg[id1] > deg[id2]; });
        for(int i = 0 ; i < n ; i++) di[id[i]] = i;
        for(int i = 0 ; i < n ; i++){
            for(int j = 0 ; j < n ; j++){
                if(v[i][j]) linkto[di[i]][di[j]] = 1;
            }
        }
        Int cand; cand.reset();
        for(int i = 0 ; i < n ; i++) cand[i] = 1;
        ans = 1;
        cans.reset(); cans[0] = 1;
        maxclique(0, cand);
        return ans;
    }
} solver;
```

#### 4.11 Minimum Mean Cycle

```
//給定一張有向圖，邊上有權重，要找到一個環其平均權重最小
/* minimum mean cycle O(VE) */
struct MMC{
#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
    struct Edge { int v,u; double c; };
    int n, m, prv[V][V], prve[V][V], vst[V];
    Edge e[E];
    vector<int> edgeID, cycle, rho;
    double d[V][V];
    void init( int _n )
    { n = _n; m = 0; }
    // WARNING: TYPE matters
    //建一條單向邊 (u, v) 權重為 w
    void addEdge( int vi , int ui , double ci )
    { e[ m ++ ] = { vi , ui , ci }; }
    void bellman_ford() {
        for(int i=0; i<n; i++) d[0][i]=0;
        for(int i=0; i<n; i++) {
            fill(d[i+1], d[i+1]+n, inf);
            for(int j=0; j<m; j++) {
                int v = e[j].v, u = e[j].u;
                if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                    d[i+1][u] = d[i][v]+e[j].c;
                    prv[i+1][u] = v;
                    prve[i+1][u] = j;
                }
            }
        }
    }
    double solve(){//回傳值為最小平均權重 (小數)
        // returns inf if no cycle, mmc otherwise
        double mmc=inf;
        int st = -1;
        bellman_ford();
        for(int i=0; i<n; i++) {
            double avg=-inf;
            for(int k=0; k<n; k++) {
                if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])/(n-k));
                else avg=max(avg,inf);
            }
            if (avg < mmc) tie(mmc, st) = tie(avg, i);
        }
        fill(vst,0); edgeID.clear(); cycle.clear(); rho.clear();
        for (int i=n; !vst[st]; st=prv[i--][st]) {
            vst[st]++;
            edgeID.PB(prve[i][st]);
            rho.PB(st);
        }
        while (vst[st] != 2) {
            if(rho.empty()) return inf;
            int v = rho.back(); rho.pop_back();
            cycle.PB(v);
            vst[v]++;
        }
        reverse(ALL(edgeID));
        edgeID.resize(SZ(cycle));
        return mmc;
    }
} }mmc;
```

## 4.12 Dominator Tree

```
// 給一張有向圖，圖上有一個起點 s 可以走到所有點。
// 定義 "支配" 為從起點 s 出發，所有能走到節點 x 的路徑
// 的最後一個必經點
// 最後 idom[i] 為點 i 的支配點
struct DominatorTree{ // O(n+m)
#define REP(i,s,e) for(int i=(s);i<=(e);i++)
#define REPD(i,s,e) for(int i=(s);i>=(e);i--)
    int n, s;
    vector< int > g[ MAXN ], pred[ MAXN ];
    vector< int > cov[ MAXN ];
    int dfn[ MAXN ], nfd[ MAXN ], ts;
    int par[ MAXN ]; //idom[u] s到u的最後一個必經點
    int sdom[ MAXN ], idom[ MAXN ];
    int mom[ MAXN ], mn[ MAXN ];
    inline bool cmp( int u , int v )
    { return dfn[u] < dfn[ v ]; }
};
```

```
int eval( int u ){
    if( mom[ u ] == u ) return u;
    int res = eval( mom[ u ] );
    if(cmp( sdom[ mn[ mom[ u ] ] ], sdom[ mn[ u ] ] ))
        mn[ u ] = mn[ mom[ u ] ];
    return mom[ u ] = res;
}
//節點數量，起點編號 1-base
void init( int _n , int _s ){
    ts = 0; n = _n; s = _s;
    REP( i, 1, n ) g[ i ].clear(), pred[ i ].clear();
}
void addEdge( int u , int v ){
    g[ u ].push_back( v );
    pred[ v ].push_back( u );
}
void dfs( int u ){
    ts++;
    dfn[ u ] = ts;
    nfd[ ts ] = u;
    for( int v : g[ u ] ) if( dfn[ v ] == 0 ){
        par[ v ] = u;
        dfs( v );
    }
}
void build(){// 建立支配樹
    REP( i, 1, n ){
        dfn[ i ] = nfd[ i ] = 0;
        cov[ i ].clear();
        mom[ i ] = mn[ i ] = sdom[ i ] = i;
    }
    dfs( s );
    REPD( i, n, 2 ){
        int u = nfd[ i ];
        if( u == 0 ) continue;
        for( int v : pred[ u ] ) if( dfn[ v ] ){
            eval( v );
            if( cmp( sdom[ mn[ v ] ], sdom[ u ] ) )
                sdom[ u ] = sdom[ mn[ v ] ];
        }
        cov[ sdom[ u ] ].push_back( u );
        mom[ u ] = par[ u ];
        for( int w : cov[ par[ u ] ] ){
            eval( w );
            if( cmp( sdom[ mn[ w ] ], par[ u ] ) )
                idom[ w ] = mn[ w ];
            else idom[ w ] = par[ u ];
        }
        cov[ par[ u ] ].clear();
    }
    REP( i, 2, n ){
        int u = nfd[ i ];
        if( u == 0 ) continue;
        if( idom[ u ] != sdom[ u ] )
            idom[ u ] = idom[ idom[ u ] ];
    }
} domT;
```

## 4.13 ManhattanMST

```
// 出來的是保證能做出MST的邊 邊是亂的 ~~魔術師~~
// 需要swap(u.first,first,u.first.second) + sort unique
#define ld long long
struct Pt{
    ld x,y;
    Pt(ld x=0,ld y=0):x(x),y(y){}
    Pt operator+(const Pt &a) const {return Pt(x+a.x, y+a.y); }
    Pt operator-(const Pt &a) const {return Pt(x-a.x, y-a.y); }
};
vector<pair<pair<int,int>,int>> ManhattanMST(vector<Pt> P) {
    vector<int> id(P.size());
    iota(id.begin(),id.end(), 0);
    vector<pair<pair<int,int>, int>> edg;
    for (int k = 0; k < 4; k++) {
        sort(id.begin(),id.end(), [&](int i, int j) {
            return (P[i] - P[j]).x < (P[j] - P[i]).y;
        });
        map<int, int> sweep;
        for (int i : id) {
```

```

    auto it = sweep.lower_bound(-P[i].y);
    while (it != sweep.end()) {
        int j = it->second;
        Pt d = P[i] - P[j];
        if (d.y > d.x) {
            break;
        }
        edg.push_back({{i, j}, d.x + d.y});
        it = sweep.erase(it);
    }
    sweep[-P[i].y] = i;
}
for (Pt &p : P) {
    if (k % 2) {
        p.x = -p.x;
    } else {
        swap(p.x, p.y);
    }
}
return edg;
}
}

```

## 5 DP

### 5.1 數位 DP

```

// dp[位數][狀態]
// dp[pos][state]: 定義為目前位數在前導狀態為state的時
// 候的計數
// ex: 求數字沒有出現66的數量 L~r
// -> dp[pos][1] 可表示計算pos個位數在前導出現一個6的計
// 數 -> dp[3][1] 則計算 6XXX
// 模板的pos是反過來的，但不影響(只是用來dp記憶用)

// pos: 目前位數
// state: 前導狀態
// lead: 是否有前導0 (大部分題目不用但有些數字EX:00146
// 如果有影響時要考慮)
// Limit: 是否窮舉有被num限制
vector<int> num;
int dp[20][state];
int dfs(int pos, int state, bool lead, bool limit) {
    if(pos==num.size()) {
        //有時要根據不同state回傳情況
        return 1;
    }
    if(limit==false && lead==false && dp[pos][state]
        !=-1) return dp[pos][state];
    int up = limit?num[pos]:9;
    int ans = 0;
    for(int i=0; i<=up; i++) {
        //有時要考慮那些狀況要continue
        ans += dfs(pos+1, state|(check[i]==2), lead&&i
            ==0, limit&&i==num[pos]);
    }
    if(limit==false && lead==false) dp[pos][state] =
        ans;
    return ans;
}

```

### 5.2 SOS DP

```

for (int mask = 0; mask < (1 << N); mask++)
    F[mask] = A[mask];

for (int i = 0; i < N; i++)
    for (int mask = 0; mask < (1 << N); mask++)
        if (mask & (1 << i))
            F[mask] += F[mask ^ (1 << i)];

```

## 6 Math

### 6.1 Formulas

```

//五次方幂次和 1, 33, 276, 1300, 4425, 12201, 29008,
// 61776
a(n) = n^2*(n+1)^2*(2*n^2+2*n-1)/12
//四次方幂次和 1, 17, 98, 354, 979, 2275, 4676, 8772,
// 15333
a(n) = n*(1+n)*(1+2*n)*(-1+3*n+3*n^2)/30
//錯位排列 0, 1, 2, 9, 44, 265, 1854, 14833, 133496,
// 1334961
dp[1]=0;dp[2]=1;
for(int i=3;i<=n;++i){dp[i]=(i-1)*(dp[i-2]+dp[i-1])%MOD;
}

```

### 6.2 llladdmul

```

#define ll long long
#define lll __int128
ll mypow(lll n, lll k, ll p){
    ll res=1;
    for(;k>=1,n=n*n%p;if(k&1)res=res*n%p;
    return res;
}
ll mul(ll a, ll b, ll mod){
    lll c=a%mod;
    return c*b%mod;
}
ll add(ll x, ll y, ll mod){
    lll c=x; c+=y;
    return c%mod;
}

```

### 6.3 Primes

```

mashu lucky prime : 91145149
1097774749, 1076767633, 100102021, 999997771
1001010013, 1000512343, 987654361, 999991231
999888733, 98789101, 987777733, 999991921, 1010101333

```

### 6.4 Coprime (互質 Pair)

```

// dp[i] 為認挑2個數字，最大公因數為i的組合數
const int mxn = 1e6+7;
int cnt[mxn], dp[mxn];
int sol(vector<int> vec) {
    for(int u: vec) cnt[u]++;
    for(int i=mxn-1; i>=1; i--) {
        int a=0, b=0;
        for(int j=i; j<mxn; j+=i) {
            a+=cnt[j]; b+=dp[j];
        }
        dp[i] = (a*(a-1)/2) - b;
    }
    return dp[1];
}

```

### 6.5 Quick Pow

```

// a^b
const int MOD = 1e9+7;
int qpow(int n, int k, int p) {
    int ret = 1;
    for(;k>=1,n=n*n%p;if(k&1)ret=ret*n%p;
    return ret;
}
// a^(b^c) = a^(q*(p-1)+r) = a^r so let b^c mod p-1
bc=qpow(b,c,p-1);
ans=qpow(a,bc,p);

```

### 6.6 Mat quick Pow



```
// mat t(r,c);
struct mat{
    long long a[200][200],r,c; // resize
    mat(int _r,int _c){r=_r;c=_c;memset(a,0,sizeof(a))
    ;}
    void build(){for(int i=0;i<r;++i)a[i][i]=1;} // 單位矩陣
};
mat operator * (mat x,mat y){
    mat z(x.r,y.c);
    for(int i=0;i<x.r;++i)for(int j=0;j<y.c;++j)for(int k=0;k<y.c;++k)
        z.a[i][j]=(z.a[i][j]+x.a[i][k]*y.a[k][j]%MOD)%MOD;
    return z;
}
mat qpow(mat a,int k){
    mat r(a.r,a.r);r.build();while(k){if(k&1)r=r*a;a=a*a;k>>=1;}return r;
}
```

## 6.7 Primes Table

```
int np[MXN];
vector<int> vec;
void sol(){
    np[0]=np[1]=1;
    for(int i=2;i<MXN;++i){
        if(!np[i]){
            for(int j=i;j<MXN;j+=i){
                np[j]=1;
            }
            vec.push_back(i);
        }
    }
}
```

## 6.8 Phi 函數

```
// 計算小於n的數中與n互質的有幾個
// O(sqrtN)
// a^b mod c = a^(b%phi(c) + phi(c)) mod c : if b>= phi(c)
//          = a^b mod c : if b< phi(c)
int phi(int n){
    int res = n, a=n;
    for(int i=2;i*i<=a;i++){
        if(a%i==0){
            res = res/i*(i-1);
            while(a%i==0) a/=i;
        }
    }
    if(a>1) res = res/a*(a-1);
    return res;
}
// 建表 最大1e7
int phi[MXN];
void phitable(int n){
    phi[1]=1;
    for(int i=2;i<=n;++i){
        if(phi[i])continue;
        for(int j=i;j<=n;j+=i){
            int &pj=phi[j];
            if(pj==0)pj=j;
            pj=pj/i*(i-1);
        }
    }
}
```

## 6.9 Factor Table

```
const int MXN = 1e7+7; //if>1e7 TLE
int np[MXN],fac[MXN],num[MXN];
// isprime, 最大質因數, 質因數數量
void table(){
    np[1]=1;
```

```
for(int i=2;i<MXN;++i){
    if(np[i])continue;
    for(int j=i;j<MXN;j+=i){
        if(i!=j)np[j]=1;
        fac[j]=i;
        num[j]++;
    }
}
//質因數分解
vector<int> res;
void div(int x){
    for(;x>1;x/=fac[x])res.push_back(fac[x]);
}
```

## 6.10 卡塔蘭數

```
//1 1 2 5 14 42 132 429 1430 4862 16796
// O(N), 要記得開Long Long 跟設定 MOD
cat[0]=1; cat[1]=1;
for(ll i=1 ; i<N ; i++)
    cat[i] = cat[i-1]*(i*4-2)%MOD*qpow(i+1, MOD-2,MOD)%MOD;
```

## 6.11 Miller Rabin

```
// n < 4,759,123,141      3 : 2, 7, 61
// n < 1,122,004,669,633  4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383  6 : pimes <= 13
// n < 2^64               7 :
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
LL magic[]={}
bool witness(LL a,LL n,LL u,int t){
    if(!a) return 0;
    LL x=mypow(a,u,n);
    for(int i=0;i<t;i++){
        LL nx=mul(x,x,n);
        if(nx==1&&x!=1&&x!=n-1) return 1;
        x=nx;
    }
    return x!=1;
}
bool miller_rabin(LL n) {
    int s=(magic number size)
    // iterate s times of witness on n
    if(n<2) return 0;
    if(!(n&1)) return n == 2;
    ll u=n-1; int t=0;
    // n-1 = u*2^t
    while(!(u&1)) u>>=1, t++;
    while(s--){
        LL a=magic[s]%n;
        if(witness(a,n,u,t)) return 0;
    }
    return 1;
}
```

## 6.12 PollarRho

```
// does not work when n is prime O(n^(1/4))
ll f(ll x, ll mod){ return add(mul(x,x,mod),1,mod); }
ll pollard_rho(ll n) {
    if(!(n&1)) return 2;
    while(true){
        ll y=2, x=rand()%(n-1)+1, res=1;
        for(int sz=2; res==1; sz*=2) {
            for(int i=0; i<sz && res<=1; i++) {
                x = f(x, n);
                res = __gcd(abs(x-y), n);
            }
            y = x;
        }
        if (res!=0 && res!=n) return res;
    }
}
```

```
// 如果被卡隨機 用下面的
ll f(ll x, ll c, ll mod){return add(mul(x,x,mod),c,mod);}
ll pollard_rho(ll n){
    ll c=1,x=0,y=0,p=2,q,t=0;
    while(t++%128 or __gcd(p,n)==1){
        if(x==y)c++,y=f(x=2,c,n);
        if(q=mul(p,abs(x-y),n))p=q;
        x=f(x,c,n);y=f(f(y,c,n),c,n);
    }
    return __gcd(p,n);
}
```

```
int nthHarmonicSum9(int n){
    int inv2=qpow(2,MOD-2,MOD),ans=0;
    for(int i=1;i<=n;){
        int v = n/i; int j = n/v;
        int area=((j-i+1)%MOD)*((j+i)%MOD)%MOD*
            inv2%MOD; //梯形
        ans=(ans+v*area%MOD)%MOD;
        i=j+1;
    }
    return ans;
};
```

### 6.13 PrimeFactorO(logn)

```
vector<ll> ret;
void fac(ll x){
    if(x<2)return;
    if (miller_rabin(x)) {
        ret.push_back(x);
        return;
    }
    ll f = pollard_rho(x);
    fac(f); fac(x/f);
}
```

### 6.14 O(1)mul

```
// < Long Long
ll mul(ll x, ll y, ll mod){
    ll ret=x*y-(ll)((long double)x/mod*y)*mod;
    // LL ret=x*y-(ll)((long double)x*y/mod+0.5)*mod;
    return ret<0?ret+mod:ret;
}
```

### 6.15 Josephus Problem

```
//base1 n people count k find Lastone O(n)
int jo(int n, int k){return n>1?(jo(n-1,k)+k-1)%n+1:1;}
//base0 when k<n O(klogn)
int jo(int n, int k) {
    if (n == 1) return 0;
    if (k == 1) return n - 1;
    if (k > n) return (jo(n - 1, k) + k) % n;
    int f = jo(n - n / k, k) - n % k;
    return f + (f < 0 ? n : (f / (k - 1)));
}
//base1 when k=2 fast find mth
int jo2(int n, int m, int f=0){
    if(n == 1) return 1;
    int kill = (n + f) / 2;
    if(m <= kill) return 2 * m - f;
    return 2 * jo2(n - kill, m - kill, (n ^ f) & 1) -
        (1 ^ f);
}
```

### 6.16 Harmonic Sum

```
struct Harmonic{
    const double gamma = 0.5772156649;
    //求第N個調和級數
    double nthHarmonic(int n){
        double result = log(n)+gamma;
        return result;
    }
    //求項數n的Sn>k
    int findNearstN(int k){
        int n = exp(k-gamma)+0.5;
        return n;
    }
    // 16n
    // n/1 + n/2 + n/3 + ... + n/n
    //就是這東西
    [20,10,6,5,4,3,2,2,2,2,1,1,1,1,1,1,1,1]
    //這是N以下的全因數和
}
```

### 6.17 Polya

```
// n個點的環 m種染色 求不重複的方案數
int polya(int n,int m){
    int ans=0;
    for(int i=1;i*i<=n;++i){
        if(n%i)continue;
        ans=(ans+qpow(m,i)*phi(n/i)%MOD)%MOD;
        if(i*i!=n)ans=(ans+qpow(m,n/i)*phi(i)%MOD)%MOD;
    }
    return (ans*qpow(n,MOD-2))%MOD; //除掉 可以用旋轉得
    到的方案
}
```

### 6.18 FFT

```
// MAXN一定要2的幕次 先跑pre_fft()
// (must be 2^k)
// before any usage, run pre_fft() first
const int MAXN = 1024*1024;
typedef long double ld;
typedef complex<ld> cplx; //real() , imag()
const ld PI = acos(-1);
const cplx I(0, 1);
cplx omega[MAXN+1];
void pre_fft(){
    for(int i=0; i<=MAXN; i++)
        omega[i] = exp(i * 2 * PI / MAXN * I);
}
// n must be 2^k
void fft(int n, cplx a[], bool inv=false){
    int basic = MAXN / n;
    int theta = basic;
    for (int m = n; m >= 2; m >>= 1) {
        int mh = m >> 1;
        for (int i = 0; i < mh; i++) {
            cplx w = omega[inv ? MAXN-(i*theta%MAXN) :
                i*theta%MAXN];
            for (int j = i; j < n; j += m) {
                int k = j + mh;
                cplx x = a[j] - a[k];
                a[j] += a[k];
                a[k] = w * x;
            }
            theta = (theta * 2) % MAXN;
        }
    }
    int i = 0;
    for (int j = 1; j < n - 1; j++) {
        for (int k = n >> 1; k > (i ^= k); k >>= 1);
        if (j < i) swap(a[i], a[j]);
    }
    if(inv) for (i = 0; i < n; i++) a[i] /= n;
}
cplx arr[MAXN+1];
inline void mul(int _n, ll a[], int _m, ll b[], ll ans[]){
    int n=1,sum=_n+_m-1;
    while(n<sum)
        n<<=1;
    for(int i=0;i<n;i++){
        double x=(i<_n?a[i]:0),y=(i<_m?b[i]:0);
        arr[i]=complex<double>(x+y,x-y);
    }
    fft(n,arr);
    for(int i=0;i<n;i++)
        arr[i]=arr[i]*arr[i];
}
```

```
fft(n,arr,true);
for(int i=0;i<sum;i++)
    ans[i]=(long long)(arr[i].real()/4+0.5);
}
```

## 7 Data Structure

### 7.1 BIT

```
//注意值域
const int N = 1e5+5;
int bit[N];
struct BIT {
    int n;
    void init(int _n){ n = _n;}
    int low(int x) {return x&-x;}
    void update(int x, int val) {
        while(x<n) bit[x]+=val, x+=low(x);
    }
    int query(int x) {
        int res = 0;
        while(x) res += bit[x], x-=low(x);
        return res;
    }
    int query(int l, int r) {return query(r) - query(l - 1); }
};
```

### 7.2 BIT 二維

```
struct BIT {
    static const int mxn = 2005;
    int bit[mxn][mxn] = {0};
    int low(int x) {return x&-x;}
    void add(int x, int y, int val) {
        for(int i=x ; i<mxn ; i+=low(i)) for(int j=y ; j<mxn ; j+=low(j)) bit[i][j]+=val;
    }
    int query(int x, int y) {
        int ans = 0;
        for(int i=x ; i ; i-=low(i)) for(int j=y ; j ; j-=low(j)) ans+=bit[i][j];
        return ans;
    }
    int range_query(int a, int b, int x, int y) {
        return query(x, y) - query(x, b-1) - query(a-1, y) + query(a-1, b-1);
    }
} bit;
```

### 7.3 並查集

```
int bin[mxn];
vector<tuple<int, int, int, int>> timing;
// int ans = n; 連通快數量
int fa(int x) {return bin[x]<0?x:fa(bin[x]);}
void union(int x, int y) {
    x = fa(x); y = fa(y);
    timing.push_back({x, bin[x], y, bin[y]});
    if(x==y) return;
    //ans--;
    if(-bin[x] > -bin[y]) {bin[x]+=bin[y]; bin[y]=x;}
    else {bin[y]+=bin[x]; bin[x]=y;}
}
void undo() {
    auto [a,b,c,d] = timing.back();
    timing.pop_back();
    //if(a!=c) ans++;
    bin[a] = b; bin[c] = d;
}
```

### 7.4 稀疏表 O(1) 區間最大最小值

```
//st[i][j]表示[i,i+2^j-1]的最值,區間最大長度為log2(n)
//i為base
const int N = 5e4+5;
int stMax[N][20],stMin[N][20],a[N];
struct ST{
    int k;
    void build(int n,int a[]){
        k=log2(n);
        for(int i = 1; i <= n; i++) stMin[i][0] =
            stMax[i][0] = a[i];
        for(int j = 1; j <= k; j++){
            for(int i = 1; i + (1 << j) - 1 <= n; i++){
                stMax[i][j] = max(stMax[i][j - 1],
                    stMax[i + (1 << (j - 1))][j - 1]);
                stMin[i][j] = min(stMin[i][j - 1],
                    stMin[i + (1 << (j - 1))][j - 1]);
            }
        }
    }
    int queryMax(int l,int r){
        int j = log2(r-l+1);
        return max(stMax[l][j],stMax[r-(1<<j)+1][j]);
    }
    int queryMin(int l,int r){
        int j = log2(r-l+1);
        return min(stMin[l][j],stMin[r-(1<<j)+1][j]);
    }
}st;
```

### 7.5 Segment Tree

```
struct seg {
    #define left (index<<1)
    #define right (index<<1|1)
    static const int MXN = 200005;
    int val[MXN*4], tag[MXN*4];
    int a[MXN];
    void push(int index, int l, int r) {
        if(tag[index]!=0) {
            val[index]+=tag[index]*(r-l+1);
            if(l==r) {
                tag[left] += tag[index];
                tag[right] += tag[index];
            }
            tag[index]=0;
        }
    }
    void pull(int index, int l, int r) {
        int mid = l+r>>1;
        push(left, l, mid);
        push(right, mid+1, r);
        val[index] = val[left]+val[right];
    }
    void build(int index, int l, int r) {
        if(l==r) {
            val[index] = a[l];
            return;
        }
        int mid = (l+r)>>1;
        build(left, l, mid);
        build(right, mid+1, r);
        pull(index, l, r);
    }
    void add(int index, int s, int e, int l, int r, int v) {
        if(e<l || r<s) return;
        if(l<=s && e<=r) {
            tag[index] += v;
            push(index, s, e);
            return;
        }
        int mid = (s+e)>>1;
        push(index, s, e);
        add(left, s, mid, l, r, v);
        add(right, mid+1, e, l, r, v);
        pull(index, s, e);
    }
    int query(int index, int s, int e, int l, int r) {
        if(e<l || r<s) return 0;
        if(l<=s && e<=r) {
            push(index, s, e);

```

```

        return val[index];
    }
    push(index, s, e);
    int mid = (s+e)>>1;
    return query(right, mid+1, e, l, r)
        +query(left, s, mid, l, r);
}
} tree;

```

```

    }
    int query(int ql, int qr) {
        return query(root, 0, n, ql, qr);
    }
    void add(int ql, int qr, int val) {
        add(root, 0, n, ql, qr, val);
    }
} tree;

```

## 7.6 動態開點線段數

```

// tree.init(區間大小 0~n)
// tree.add(ql, qr, val) 區間加值
// tree.query(ql, qr) 區間總和查詢
struct seg {
    struct Node {
        int val, tag;
        Node *l, *r;
        Node(int v=0) : val(v), tag(0), l(nullptr), r(
            nullptr) {}
    };
    Node* root;
    int n;
    void init(int _n) {
        n = _n;
        root = new Node();
    }
    void push(Node* cur, int l, int r) {
        if(cur->tag) {
            cur->val += (r-l+1)*cur->tag;
            if(l!=r) {
                if(!cur->l) cur->l = new Node();
                if(!cur->r) cur->r = new Node();
                cur->l->tag += cur->tag;
                cur->r->tag += cur->tag;
            }
        }
        cur->tag = 0;
    }
    void pull(Node* node, int l, int r) {
        int mid = l+r>>1;
        push(node->l, l, mid);
        push(node->r, mid+1, r);
        node->val = node->l->val + node->r->val;
    }
    void add(Node* cur, int l, int r, int ql, int qr,
        int val) {
        if (ql <= l && r <= qr) {
            cur->tag += val;
            push(cur, l, r);
            return;
        }
        if (!cur->l) cur->l = new Node();
        if (!cur->r) cur->r = new Node();

        int mid = (l + r) / 2;
        push(cur, l, r);
        if(ql<=mid) add(cur->l, l, mid, ql, qr, val);
        if(mid+1<=qr) add(cur->r, mid + 1, r, ql, qr,
            val);
        pull(cur, l, r);
    }
    int query(Node* cur, int l, int r, int ql, int qr)
    {
        if(ql<=l && r<=qr) {
            push(cur, l, r);
            return cur->val;
        }
        if (!cur->l) cur->l = new Node();
        if (!cur->r) cur->r = new Node();
        int mid = l+r>>1;
        int ans = 0;
        push(cur, l, r);
        if(ql<=mid) ans+=query(cur->l, l, mid, ql, qr);
        if(mid+1<=qr) ans+=query(cur->r, mid+1, r, ql,
            qr);
        pull(cur, l, r);
        return ans;
    }
}

```

## 7.7 動態開點線段數 2D

```

// tree.init(n,m) 二維大小
// tree.add(qx, qy, val) 座標 (qx, qy) 加值 val
// tree.query(qlx, qly, qrx, qry) 座標 左下角(qlx, qly)
// 到 座標 右上角(qrx, qry) 的矩陣總和
// 單點加值
// 區間查詢
// O(Lg(n)Lg(m))
struct segy {
    int n;
    struct Node {
        int val;
        Node *l, *r;
        Node(int v=0) : val(v), l(nullptr), r(nullptr)
            {}
    };
    segy(int _n=1e9) {n=_n; root=new Node();}
    Node* root;
    void init(int _n=1e9) {
        n = _n;
        root = new Node();
    }
    void pull(Node* node) {
        node->val = (node->l?node->l->val:0) + (node->r
            ?node->r->val:0);
    }
    void add(Node* cur, int l, int r, int pos, int val)
    {
        if (l==r) {
            cur->val += val;
            return;
        }
        int mid = l+r>>1;
        if(pos<=mid) {if(!cur->l) cur->l = new Node();
            add(cur->l, l, mid, pos, val);}
        else {if(!cur->r) cur->r = new Node();add(cur->
            r, mid + 1, r, pos, val);}
        pull(cur);
    }
    int query(Node* cur, int l, int r, int ql, int qr)
    {
        if(ql<=l && r<=qr) {
            return cur->val;
        }
        int mid = l+r>>1;
        int ans = 0;
        if(ql<=mid) {if(!cur->l) cur->l = new Node();
            ans+=query(cur->l, l, mid, ql, qr);}
        if(mid+1<=qr) {if(!cur->r) cur->r = new Node();
            ans+=query(cur->r, mid+1, r, ql, qr);}
        pull(cur);
        return ans;
    }
}
int query(int ql, int qr) {
    return query(root, 0, n, ql, qr);
}
void add(int pos, int val) {
    add(root, 0, n, pos, val);
}
}
struct segx {
    struct Node {
        segy * tree_y;
        Node *l, *r;
        Node(int m) : tree_y(new segy(m)), l(nullptr),
            r(nullptr) {}
    };
    Node* root;
    int n,m;
}

```

```

segx(int _n=1e9, int _m=1e9) {n=_n; m=_m; root=new
Node(m);}
void init(int _n=1e9, int _m=1e9) {n=_n; m=_m; root=
new Node(m);}
void add(Node* cur, int l, int r, int qx, int qy,
int val) {
if(l!=r) {
int mid = (l + r) / 2;
if(qx<=mid) {if(!cur->l) cur->l = new Node(
m); add(cur->l, l, mid, qx, qy, val);}
else {if(!cur->r) cur->r = new Node(m); add
(cur->r, mid + 1, r, qx, qy, val);}
}
cur->tree_y->add(qy, val);
}
int query(Node* cur, int lx, int rx, int qlx, int
qrx, int qly, int qry) {
if(qlx<=lx && rx<=qrx) {
return cur->tree_y->query(qly, qry);
}
int mid = lx+rx>>1;
int ans = 0;
if(qlx<=mid) {if(!cur->l) cur->l = new Node(m);
ans+=query(cur->l, lx, mid, qlx, qrx, qly,
qry);}
if(mid+1<=qrx) {if(!cur->r) cur->r = new Node(m)
; ans+=query(cur->r, mid+1, rx, qlx, qrx,
qly, qry);}
return ans;
}
int query(int qlx, int qly, int qrx, int qry) {
return query(root, 0, n, qlx, qrx, qly, qry);
}
void add(int qx, int qy, int val) {
add(root, 0, n, qx, qy, val);
}
} tree;

```

## 7.8 持久化線段樹

```

struct seg {
// 加值持久化線段樹
struct Node {
int val;
Node *l, *r;
};
vector<Node*> version;
void pull(Node* node) {
node->val = node->l->val+node->r->val;
}
Node* build(int l,int r) {
Node* node=new Node;
if(l==r) {
node->val = 0; //初始值
return node;
}
int mid = (l+r)/2;
node->l = build(l,mid);
node->r = build(mid+1,r);
pull(node);
return node;
}
Node* update(Node* cur,int l,int r,int pos,int v) {
Node* node=new Node;
if(l==r){
//改成加值換這行
//node->val=cur->val + v;
node->val=v;
return node;
}
int mid=(l+r)/2;
if(pos<=mid) {
node->l=update(cur->l,l,mid,pos,v);
node->r=cur->r;
} else {
node->l=cur->l;
node->r=update(cur->r,mid+1,r,pos,v);
}
pull(node);
return node;
}

```

```

}
int query(Node* cur,int s, int e, int ql, int qr){
if(ql<=s && e<=qr) return cur->val;
int ans = 0;
int mid = (s+e)/2;
if(ql<=mid) ans += query(cur->l, s, mid, ql, qr
);
if(mid+1<=qr) ans += query(cur->r, mid+1, e, ql
, qr);
return ans;
}
} tree;
// push 初始的樹
// tree.version.push_back(tree.build(1, n));
// update(舊版, 1, n, pos, v) return 新版
// 把pos值改成v

```

## 7.9 Time Segment Tree

```

#include <bits/stdc++.h>
#define int long long int
using namespace std;
int n, q;
struct node{
int val;
node *l, *r;
node(int v) {val=v; l=r=nullptr;}
node() {val=0; l=r=nullptr;}
};
vector<node*> timing;
node* build(int s, int e) {
node *ret = new node();
if(s==e) return ret;
int mid = (s+e)>>1;
ret->l = build(s, mid);
ret->r = build(mid+1, e);
ret->val = ret->l->val + ret->r->val;
return ret;
}
node* update(node* pre, int s, int e, int pos, int v) {
node *ret = new node();
if(s==e) {ret->val=pre->val+v; return ret;}
int mid = (s+e)>>1;
if(pos<=mid) {
ret->l = update(pre->l, s, mid, pos, v);
ret->r = pre->r;
} else {
ret->r = update(pre->r, mid+1, e, pos, v);
ret->l = pre->l;
}
ret->val = ret->l->val + ret->r->val;
return ret;
}
void add(int pos, int v) {
timing.push_back(update(timing.back(), 1, n, pos, v
));
}
int que(node* pre, node* now, int l, int r, int k) {
if(l==r) return r;
int mid = (l+r)>>1;
int diff = now->l->val - pre->l->val;
//printf("now %d~%d diff %d\n", l, r, diff);
if(diff>=k) return que(pre->l, now->l, l, mid, k);
else return que(pre->r, now->r, mid+1, r, k-diff);
return -1;
}
int query(int l, int r, int k) {
l--;
return que(timing[l], timing[r], 1, n, k);
}
int num[100005];
vector<int> sor;
map<int, int> mp;
signed main() {
cin>>n>>q;
timing.push_back(build(1, n));
for(int i=0,a ; i<n ; i++) {
cin>>a; num[i] = a; sor.push_back(a);
}
}

```



```

// add: 1 1 1 2 1
// num: 3 3 3 4 3
// sor: 3 4
sort(sor.begin(), sor.end());
sor.erase(unique(sor.begin(), sor.end()), sor.end());
}
for(int i=0 ; i<n ; i++) {
    int pos = lower_bound(sor.begin(), sor.end(),
        num[i]) - sor.begin() + 1;
    //printf("mp[%d] = %d\n", pos, num[i]);
    mp[pos] = num[i];
    num[i] = pos;
    add(num[i], 1);
}
while(q--) {
    int a, b, c; cin>>a>>b>>c;
    cout<<mp[query(a, b, c)]<<endl;
}
}

```

## 7.10 Treap

```

struct Treap {
    int sz, val, pri, tag;
    Treap *l, *r;
    Treap(int _val){
        val=_val; sz=1;
        pri=rand(); l=r=NULL; tag=0;
    }
};
int Size(Treap *a) {return a?a->sz:0;}
void pull(Treap *a) {
    a->sz = Size(a->l) + Size(a->r) + 1;
}
//val of a is always bigger than val of b
Treap* merge(Treap *a, Treap *b) {
    if(!a || !b) return a ? a : b;
    if(a->pri>b->pri) {
        a->r = merge(a->r, b);
        pull(a);
        return a;
    } else {
        b->l = merge(a, b->l);
        pull(b);
        return b;
    }
}
// a<k, b>=k
void split(Treap *t, int k, Treap*&a, Treap*&b){
    if(!t) {a=b=NULL; return;}
    if(k <= t->val) {
        b = t;
        split(t->l, k, a, b->l);
        pull(b);
    }
    else {
        a = t;
        split(t->r, k, a->r, b);
        pull(a);
    }
}
Treap* add(Treap *t, int v) {
    Treap *val = new Treap(v);
    Treap *l = NULL, *r = NULL;
    split(t, v, l, r);
    return merge(merge(l, val), r);
}
Treap* del(Treap *t, int v) {
    Treap *l, *mid, *r, *temp;
    split(t, v, l, temp);
    split(temp, v+1, mid, r);
    return merge(l, r);
}
// base 1
int position(Treap *t, int p) {
    if(Size(t->l)+1==p) return t->val;
    if(Size(t->l)<p) return position(t->r, p-Size(t->l)-1);
    else return position(t->l, p);
}

```

```

//num of >= k
int query(Treap *t, int k) {
    if(!t) return 0;
    if(t->val==k) return Size(t->l)+1;
    if(t->val>k) return query(t->l, k);
    return Size(t->l)+1+query(t->r, k);
}

```

## 7.11 PBDS

```

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#define ordered_set tree<int, null_type, less<int>,
    rb_tree_tag, tree_order_statistics_node_update>
using namespace __gnu_pbds;
// ordered_set s;
// s.insert(1); s.erase(s.find(1));
// order_of_key(k) : Number of items strictly smaller
    than k .
// find_by_order(k) : K-th element in a set (counting
    from zero). (return iterator)

```

## 8 String

### 8.1 SA

```

#pragma GCC optimize("O3,unroll-loops")
#pragma target optimize("avx2,bmi,bmi2,lzcnt,popcnt")
#include<bits/stdc++.h>
#include<chrono>
#define mid (l + r) / 2
using namespace std;
const int N = 100010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )
#define REP1(i,a,b) for ( int i=(a); i<=int(b); i++ )
    bool _t[N*2];
    int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
        hei[N], r[N];
    int operator [] (int i){ return _sa[i]; }
    void build(int *s, int n, int m){
        memcpy(_s, s, sizeof(int) * n);
        sais(_s, _sa, _p, _q, _t, _c, n, m);
        mkhei(n);
    }
    void mkhei(int n){
        REP(i,n) r[_sa[i]] = i;
        hei[0] = 0;
        REP(i,n) if(r[i]) {
            int ans = i>0 ? max(hei[r[i]-1] - 1, 0) : 0;
            while(_s[i+ans] == _s[_sa[r[i]-1]+ans]) ans++;
            hei[r[i]] = ans;
        }
    }
    void sais(int *s, int *sa, int *p, int *q, bool *t,
        int *c, int n, int z){
        bool uniq = t[n-1] = true, neq;
        int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
            lst = -1;
#define MS0(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
        memcpy(x, c, sizeof(int) * z); \
        XD; \
        memcpy(x + 1, c, sizeof(int) * (z - 1)); \
        REP(i,n) if(sa[i] && !t[sa[i]-1]) sa[x[s[sa[i]
            ]-1]]++ = sa[i]-1; \
        memcpy(x, c, sizeof(int) * z); \
        for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i]
            ]-1]) sa[--x[s[sa[i]-1]]] = sa[i]-1;
        MS0(c, z);
        REP(i,n) uniq &= ++c[s[i]] < 2;
        REP(i,z-1) c[i+1] += c[i];
        if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
        for(int i = n - 2; i >= 0; i--) t[i] = (s[i]==s[i
            +1] ? t[i+1] : s[i]<s[i+1]);
        MAGIC(REP1(i,1,n-1) if(t[i] && !t[i-1]) sa[--x[s[i]
            ]]=p[q[i]=nn++]=i);
        REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {

```

```

    neq=lst<0||memcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa
    [i])*sizeof(int));
    ns[q[lst=sa[i]]]=nmxz+=neq;
}
sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz
+ 1);
MAGIC(for(int i = nn - 1; i >= 0; i--) sa[--x[s[p[
    nsa[i]]]]) = p[nsa[i]]];
}
}sa;
int H[ N ], SA[ N ];
void suffix_array(int* ip, int len) {
    // should padding a zero in the back
    // ip is int array, len is array length
    // ip[0..n-1] != 0, and ip[len] = 0
    ip[len++] = 0;
    sa.build(ip, len, 128);
    for (int i=0; i<len; i++) {
        H[i] = sa.hei[i + 1];
        SA[i] = sa._sa[i + 1];
    }
    // resulting height, sa array \in [0, len)
}
bool check(string &s, string &t, int p) {
    for(int i=0; i<t.size() && i+p<s.size(); ++i) {
        if(t[i]<s[i+p]) return 1;
        else if(t[i]>s[i+p]) return 0;
    }
    if(t.size()>s.size()-p) return 0;
    return 1;
}
//example for finding patterns in a string
string s,t;
int ip[N],len;
int main(){
    int n;
    cin>>s>>n;
    len = s.length();
    for(int i=0; i<len; ++i) ip[i]=(int)s[i];
    ip[len] = 0;
    suffix_array(ip, len);
    int l,r;
    for(int i=0; i<n; ++i){
        cin>>t;
        l = 0, r = s.size()-1;
        while(l!=r){
            if(check(s,t,SA[mid])) r=mid;
            else l = mid+1;
        }
        bool f=1;
        if(t.size()>s.size()-SA[l]){
            cout<<"NO\n", f=0;
            continue;
        }
        for(int j=0; j<t.size(); ++j){
            if(t[j]!=s[j+SA[l]]){
                cout<<"NO\n", f=0;
                break;;
            }
        }
        if(f) cout<<"YES\n";
    }
}

```

## 8.2 KMP

```

// 回傳所有匹配成功的起始位置，s為文本，t為匹配字串
// nxt表示為匹配失敗時要退回的位置，也是t字串的相等前綴
// 後綴的最大長度
// *注意前綴後綴為長度最多為n-1的字子串
// nxt[j] = -1 if j=0
// 0 if 沒有相等的前綴後綴
// k k 為相等前綴後綴的最大長度
// 以下為例子
// j: 0 1 2 3 4 5 6
// t: a b a a b e
// nxt[j]: -1 0 0 1 1 2 0
// O(n+m)，n為s長，m為t長
const int MXN = 1e6+5;
int nxt[MXN];

```

```

vector<int> KMP(string s, string t){
    int slen = s.length(), tlen = t.length(), i=0, j=0, k
    =-1;
    nxt[0]=-1;
    while(j<tlen){//build nxt
        if(k==-1 || t[j]==t[k]) nxt[++j] = ++k;
        else k=nxt[k];
    }
    i=0, j=0;
    vector<int> ret;
    while(i<slen){// matching
        if(j==-1 || s[i]==t[j]) i++, j++;
        else j=nxt[j];
        if(j==tlen){
            ret.push_back(i-tlen+1); //1-base
            j=nxt[j];
        }
    }
    return ret;
}

//另一版
//if t is the substring of s:
//if t in s:
bool cmp(string s, string t) {
    vector<int> front(t.size(), 0);
    for(int i=1, j=0; i<t.size(); i++) {
        while(j>0 && t[i]!=t[j]) j = front[j-1];
        if(t[i]==t[j]) j++;
        front[i] = j;
    }
    int j=0, i=0;
    while(i<s.size()) {
        if(s[i]==t[j]) j++, i++;
        else {i += (j==0); j = (j<1?0:front[j-1]);}
        if(j==t.size()) return true;
    }
    return false;
}

```

## 8.3 Single Hash

```

//字串雜湊前的idx是0-base，雜湊後為1-base
//H[R] - H[L-1] * p^(R-L+1)
//cmp的+modl是為了防止負數
//記得build完之後要buildPow
//小心遇到hash出負數要記得+modl
#define int long long
const int p = 75577, modl = 1e9 + 7, MXN = 1e6+5;
int Hash[MXN], qpow[MXN];
void build(const string& s) {
    Hash[0]=0;
    for(int i=1; i<=s.size(); i++)
        Hash[i] = (Hash[i-1] * p + s[i-1]) % modl;
}
void buildPow(){
    qpow[0]=1;
    for(int i=1; i<MXN; ++i) qpow[i]=qpow[i-1]*p%modl;
}
bool cmp(int i, int j, int len) {
    return (Hash[i+len-1] - Hash[i-1] * qpow[len] %
        modl + modl) % modl ==
        (Hash[j+len-1] - Hash[j-1] * qpow[len] % modl +
        modl) % modl;
}
int get(int i, int j) {
    return (Hash[j]-Hash[i-1]*qpow[j-i+1]%modl+modl)%
        modl;
}

```

## 8.4 Double Hash

```

//字串雜湊前的idx是0-base，雜湊後為1-base
//即區間為 [0, n-1] -> [1, n]
//若要取得區間[L, R]的值則
//H[R] - H[L-1] * p^(R-L+1)
//cmp為比較從i開始長度為len的字串和從j開始長度為len的字
//串是否相同

```

```
//(h[i+len-1] - h[i-1] * qpow(p, Len) % modL + modL)
#define int long long
#define x first
#define y second
const int P1 = 75577, P2 = 17, MOD = 1e9 + 7, MXN = 1e6
+5;
pair<int,int> Hash[MXN];
int qpow[2][MXN];
void build(const string& s){
    pair<int,int> val = make_pair(0,0);
    Hash[0]=val;
    for(int i=1; i<s.size(); i++){
        val.x = (val.x * P1 + s[i-1]) % MOD;
        val.y = (val.y * P2 + s[i-1]) % MOD;
        Hash[i] = val;
    }
}
void buildPow(){
    qpow[0][0]=qpow[1][0]=1;
    for(int i=1; i<MXN; ++i){
        qpow[0][i]=qpow[0][i-1]*P1%MOD;
        qpow[1][i]=qpow[1][i-1]*P2%MOD;
    }
}
bool cmp( int i, int j, int len ) {
    return ((Hash[i+len-1].x-Hash[i-1].x*qpow[0][len]%
    MOD+MOD)%MOD == (Hash[j+len-1].x-Hash[j-1].x*
    qpow[0][len]%MOD+MOD)%MOD)
    && ((Hash[i+len-1].y-Hash[i-1].y*qpow[1][len]%MOD+
    MOD)%MOD == (Hash[j+len-1].y-Hash[j-1].y*qpow
    [1][len]%MOD+MOD)%MOD);
}
pair<int, int> get(int i, int j) {
    return {(Hash[j].x-Hash[i-1].x*qpow[0][j-i+1]%MOD+
    MOD)%MOD, (Hash[j].y-Hash[i-1].y*qpow[1][j-i
    +1]%MOD+MOD)%MOD};
}
```

## 8.5 Trie

```
//cnt為記錄有多少個一樣的單詞且end的時候才有數字
const int MXN=1e6+5; //MXN取文本長
int trie[MXN][26], cnt[MXN], tot=0; //0 base
void update(string s){
    int p=0; //0 base
    for(int i=0; i<s.size(); ++i){
        int ch = s[i] - 'a';
        if(!trie[p][ch]) trie[p][ch]=++tot;
        p = trie[p][ch];
    }
    cnt[p]++;
}
int query(string s){
    int p=0;
    for(int i=0; i<s.size(); ++i){
        int ch=s[i] - 'a';
        p = trie[p][ch];
        if(!p) return 0;
    }
    return cnt[p];
}
void visualizeTrie(int node = 0, int depth = 0) { //for
    debug
    for (int i = 0; i < 26; ++i) {
        if (trie[node][i]) {
            for (int j = 0; j < depth; ++j) cout << "
            ";
            cout << (char)('a' + i) << " (" << cnt[trie
            [node][i]] << ")\n";
            visualizeTrie(trie[node][i], depth + 1);
        }
    }
}
```

## 8.6 Z value

```
// O(n)
//z[i] = lcp(s[1...], s[i...])
//1base
```

```
int z[MXN];
void Z_value(const string& s) {
    int i, j, left, right, len = s.size();
    left=right=0; z[0]=len;
    for(i=1; i<len; i++){
        j=max(min(z[i-left], right-i), 0);
        for(; i+j<len&&s[i+j]==s[j]; j++);
        z[i]=j;
        if(i+z[i]>right) {
            right=i+z[i];
            left=i;
        }
    }
}
```

## 8.7 MinRotation

```
//rotate(begin(s), begin(s)+minRotation(s), end(s))
//For example, rotations of acab are acab, caba, abac,
//and baca.
//find lexicographically minimal rotation of a string
int minRotation(string s) {
    int a = 0, N = s.size(); s += s;
    for(int b=0; b<N; b++) for(int k=0; k<N; k++) {
        if(a+k == b || s[a+k] < s[b+k])
            {b += max(0, k-1); break;}
        if(s[a+k] > s[b+k]) {a = b; break;}
    } return a;
}
```

## 8.8 Manacher 馬拉車回文

```
// O(N)求以每個字元為中心的最長回文半徑
// 頭尾以及每個字元間都加入一個
// 沒出現過的字元，這邊以'@'為例
// s為傳入的字串，len為字串長度
// z為儲存以每個字元為中心的回文半徑+1(有包含'@'要小心)
// ex: s = "abaac" -> "@a@b@a@a@c@"
// z = [12141232121]
const int MXN = 1e6+5;
int z[2*MXN];
char s[2*MXN];
void z_value_pal(char *s, int len, int *z){
    len=(len<<1)+1;
    for(int i=len-1; i>=0; i--){
        s[i]=i&1?s[i>>1]:'@';
        z[0]=1;
        for(int i=1, l=0, r=0; i<len; i++){
            z[i]=i<r?min(z[l+1-i], r-i):1;
            while(i-z[i]>=0&&i+z[i]<len&&s[i-z[i]]==s[i+z[i]])
                ++z[i];
            if(i+z[i]>r) l=i, r=i+z[i];
        }
    }
    // cin>>s;
    // z_value_pal(s, strlen(s), z);
    // int mx=-1, mxi=0;
    // for(int i=0; i<strlen(s); ++i)
    //     if(mx<z[i]) mx = z[i], mxi = i;
    // mx--;
    // for(int i=mxi-mx; i<=mxi+mx; ++i)
    //     if(s[i]!='@') cout<<s[i];
}
```

## 8.9 PalTree 回文樹

```
// Len[s]是對應的回文長度
// num[s]是有幾個回文後綴
// cnt[s]是這個回文字字串在整個字串中的出現次數
// fail[s]是他長度次長的回文後綴，aba的fail是a
const int MXN = 1000010;
struct PalT{
    int nxt[MXN][26], fail[MXN], len[MXN];
    int tot, lst, n, state[MXN], cnt[MXN], num[MXN];
    int diff[MXN], sfail[MXN], fac[MXN], dp[MXN];
    char s[MXN]={-1};
    int newNode(int l, int f){
        len[tot]=l, fail[tot]=f, cnt[tot]=num[tot]=0;
        memset(nxt[tot], 0, sizeof(nxt[tot]));
        diff[tot]=(l>0?l-len[f]:0);
    }
}
```

```

    sfail[tot]=(1>0&&diff[tot]==diff[f]?sfail[f]:f)
    ;
    return tot++;
}
int getfail(int x){
    while(s[n-len[x]-1]!=s[n]) x=fail[x];
    return x;
}
int getmin(int v){
    dp[v]=fac[n-len[sfail[v]]-diff[v]];
    if(diff[v]==diff[fail[v]])
        dp[v]=min(dp[v],dp[fail[v]]);
    return dp[v]+1;
}
int push(){
    int c=s[n]-'a',np=getfail(1st);
    if(!(1st=nxt[np][c])){
        1st=newNode(len[np]+2,nxt[getfail(fail[np])][c]);
        nxt[np][c]=1st; num[1st]=num[fail[1st]]+1;
    }
    fac[n]=n;
    for(int v=1st;len[v]>0;v=sfail[v])
        fac[n]=min(fac[n],getmin(v));
    return ++cnt[1st],1st;
}
void init(const char *_s){
    tot=1st=n=0;
    newNode(0,1),newNode(-1,1);
    for(;_s[n];) s[n+1]=_s[n],++n,state[n-1]=push();
    for(int i=tot-1;i>1;i--) cnt[fail[i]]+=cnt[i];
}
} palt;
// state 數組
// state[i] 代表第 i 個字元為結尾的最長回文編號(編號
// 是甚麼不重要)
//
// S = "abacaaba"
//
// 以第 2(0-base) 個字元為結尾的最長回文是 aba
// 以第 7(0-base) 個字元為結尾的最長回文是 aba
// 兩個最長回文都相同，因此 state[2] 會等於 state[7]
//
// Len 數組
// 求出某個 state 的長度
//
// S = "aababa"
//
// (0-base)
// len[state[1]] = 2 ("aa")
// len[state[3]] = 3 ("aba")
// len[state[5]] = 5 ("ababa")
//
// num 數組
// 某個state的回文有幾個回文後綴
//
// 假設某個 state 代表的回文為 = "ababa" 為例
// state 代表的回文的 num = 3
// -> ababa -> aba -> a
//
// cnt 數組
// 某個 state 的回文在整個字串中出現次數
//
// S = "aababaa"
// state[3] 代表的回文為 "aba" 在整個字串中出現 2
// 次
// 因此 cnt[state[3]] = 2
//
// fail數組
// 每個 state 的次長回文後綴的 state 編號
//
// S = "ababa"
// len[fail[4]] = 3 (fail[4] = "aba")
// len[fail[2]] = 1 (fail[2] = "a")
// len[fail[0]] = 0 (fail[0] = "" 空字串)
// 0 所代表的 state 是空字串

```

## 8.10 DistinctSubsequence

```

//預設為小寫字母
//return the number of distinct non-empty subsequences
//of sting
#define int long long
int mod = 1e9 + 7;
vector<int> cnt(26);
int distinct_subsequences(string s) {
    for (char c : s)
        cnt[c - 'a'] = accumulate(begin(cnt), end(cnt), 1LL
                                   ) % mod;
    return accumulate(begin(cnt), end(cnt), 0LL) % mod;
}

```

## 9 Tree

### 9.1 LCA

```

//先建edge[MXN]
//跑dfs，再跑makeanc
//之後才可以呼叫lca
// 0-base
const int MXN=1e5;
const int logN=__lg(MXN);
int tin[MXN],tout[MXN],anc[MXN][logN+1];
vector<int> edge[MXN];
int ti=0;
void dfs(int x,int f){
    anc[x][0]=f;
    tin[x]=ti++;
    for(int u:edge[x]){
        if(u==f)continue;
        dfs(u,x);
    }
    tout[x]=ti++;
}
// x is y's anc
inline bool isanc(int x,int y){
    return tin[x]<=tin[y] && tout[x]>=tout[y];
}
int lca(int x,int y){
    if(isanc(x,y))return x;
    if(isanc(y,x))return y;
    for(int i=logN;i>0;--i){
        if(!isanc(anc[y][i],x)){
            y=anc[y][i];
        }
    }
    return anc[y][0];
}
void makeanc(int n){
    for(int i=1;i<=logN;++i){
        for(int j=0;j<n;++j){
            anc[j][i] = anc[anc[j][i-1]][i-1];
        }
    }
}

```

### 9.2 TreeHash

```

// 1. dfs 先做子樹
// 2. 葉節點的hash值為1
// 3. 對於節點x，其hash值為紀錄x的所有子樹的hash值(紀錄
// 到temp)，然後由小排到大(排除子樹的隨機問題)
// 4. n表示節點x有幾個子樹，p和MOD通常為一個很大的質
// 數，由此算出x的hash值
// 5. 樹根的hash值即為整顆樹的hash值，若兩顆樹的hash值
// 相同，則兩棵樹就是同構
const int MXN = 200005;
int subtree_sz[MXN];
int hash_[MXN];
int base = 44560482149;
int MOD = 274876858367;
int dfs(int x, int fa, vector<int>* edge){
    vector<int> temp;
    subtree_sz[x] = 1;
    for(int child : edge[x]){

```

```

    if(child==fa) continue;
    temp.push_back(dfs(child, x, edge));
    subtree_sz[x] += subtree_sz[child];
}
sort(temp.begin(), temp.end());
int ret = subtree_sz[x];
for(int v : temp){
    ret = (((ret * base + v + ret) % MOD + ret) %
    MOD + v) % MOD ;
}
hash_[x] = ret;
return ret;
}

```

### 9.3 輕重鍵剖分

```

const int MXN = 2e5+7;
int top[MXN], son[MXN], dfn[MXN], rnk[MXN], dep[MXN],
    father[MXN];
vector<int> edge[MXN];
int dfs1(int v, int fa, int d) {
    int maxsz = -1, maxu, total = 1;
    dep[v] = d;
    father[v] = fa;
    for(int u: edge[v]) {
        if(fa == u) continue;
        int temp = dfs1(u, v, d+1);
        total += temp;
        if(temp>maxsz) {
            maxsz = temp;
            maxu = u;
        }
    }
    if(maxsz==-1) son[v] = -1;
    else son[v] = maxu;
    return total;
}

int times = 1;
void dfs2(int v, int fa) {
    rnk[times] = v;
    dfn[v] = times++;
    top[v] = (fa==-1 || son[fa] != v ? v : top[fa]);
    if(son[v]!=-1) dfs2(son[v], v);
    for(int u: edge[v]) {
        if(fa == u || u == son[v]) continue;
        dfs2(u, v);
    }
}

//rnk: 剖分後的編號 (rnk[時間] = 原點)
//dfn: 剖分後的編號 (dfn[原點] = 時間)
//top: 剖分的頭頭
//son: 剖分的重兒子

```

## 10 Geometry

### 10.1 2D Definition

```

#define ld long double
const ld eps=1e-10;
int dcmp(ld x){if(fabs(x)<eps) return 0;else return x
<0?-1:1;}
struct Pt{
    ld x,y;
    Pt(ld x=0,ld y=0):x(x),y(y){}
    Pt operator+(const Pt &a) const {
        return Pt(x+a.x, y+a.y); }
    Pt operator-(const Pt &a) const {
        return Pt(x-a.x, y-a.y); }
    Pt operator*(const ld &a) const {
        return Pt(x*a, y*a); }
    Pt operator/(const ld &a) const {
        return Pt(x/a, y/a); }
    ld operator*(const Pt &a) const { //dot
        return x*a.x + y*a.y; }
    ld operator^(const Pt &a) const { //cross
        return x*a.y - y*a.x; }
    bool operator<(const Pt &a) const {

```

```

        return x < a.x || (x == a.x && y < a.y); }
    //return dcmp(x-a.x) < 0 || (dcmp(x-a.x) == 0
    && dcmp(y-a.y) < 0); }
    bool operator>(const Pt &a) const {
        return x > a.x || (x == a.x && y > a.y); }
    //return dcmp(x-a.x) > 0 || (dcmp(x-a.x) == 0
    && dcmp(y-a.y) > 0); }
    bool operator==(const Pt &a) const {
        return dcmp(x-a.x) == 0 && dcmp(y-a.y) == 0; }
    // return x == other.x && y == other.y;
    bool operator!=(const Pt &a) const {
        return !(*this == a); }
};

typedef Pt Vec;
ld Dot(Vec a,Vec b){return a.x*b.x+a.y*b.y;}
ld Cross(Vec a,Vec b){return a.x*b.y-a.y*b.x;}
ld Length(Vec a){return sqrt(Dot(a,a));}
int Sgn(double x){ return (x > -eps) - (x < eps); }//
return 0: x==0, 1: x>0, -1: x<0
ld Angle(Vec a,Vec b){return acos(Dot(a,b)/Length(a)/
    Length(b));} //弧度
ld Degree(Vec a,Vec b){return Angle(a,b)*180/acos(-1);}
//角度
ld Ori(Pt a,Pt b,Pt c){return Cross(b-a,c-a);} //1.(a,b)
x(a,c)的面積 2. a在bc左側>0 3. a在bc右側<0 4. a在
bc上==0
Vec Rotate(Vec a,ld rad){return Vec(a.x*cos(rad)-a.y*
    sin(rad),a.x*sin(rad)+a.y*cos(rad));} //逆時針旋轉,
rad為弧度
Vec Normal(Vec a){ld L=Length(a);return Vec(-a.y/L,a.x/
    L);} //單位法向量,確保a不是零向量
Vec Unit(Vec x) { return x / Length(x); } //單位向量
Vec Perp( const Vec v ){ return { v.y , -v.x };} //垂直
向量
bool argcmp(const Pt &a, const Pt &b) { // 極角cmp: arg
(a) < arg(b)
    int f = (Pt{a.y, -a.x} > Pt{b.y, -b.x}) ? 1 : -1 * (a != Pt
    {b.y, -b.x});
    int g = (Pt{b.y, -b.x} > Pt{c.y, -c.x}) ? 1 : -1 * (b != Pt
    {c.y, -c.x});
    return f == g ? (a ^ b) > 0 : f < g;
}

struct Circle {
    Pt o; ld r;
    Circle(Pt _o=Pt(0, 0), ld _r=0):o(_o), r(_r) {}
};

```

### 10.2 Line Definition

```

struct Line {
    Pt a, b, v; // start, end, end-start
    ld ang;
    Line(Pt _a=Pt(0, 0), Pt _b=Pt(0, 0)):a(_a),b(_b) { v
    = b-a; ang = atan2(v.y, v.x); }
    bool operator<(const Line &L) const {
        return ang < L.ang;
    } };
int PtSide(Pt p, Line L) { //return 1:左側 0:線上 -1:右
    側
    return Sgn(Ori(L.a, L.b, p));
}
bool PtOnSeg(Pt p, Line L) { //點是否在線段上
    return Sgn(Ori(L.a, L.b, p)) == 0 && Sgn((p - L.a)
    * (p - L.b)) <= 0;
}
Pt Proj(Pt p, Line l) { //點到線段的投影點
    Pt dir = Unit(l.b - l.a);
    return l.a + dir * (dir * (p - l.a));
}
bool isInter(Line l, Line m) { //判斷兩線段是否相交
    if (PtOnSeg(m.a, l) || PtOnSeg(m.b, l) ||
        PtOnSeg(l.a, m) || PtOnSeg(l.b, m))
        return true;
    return PtSide(m.a, l) * PtSide(m.b, l) < 0 &&
        PtSide(l.a, m) * PtSide(l.b, m) < 0;
}
Pt LineInter(Line l, Line m) { //兩線段交點
    double s = Ori(m.a, m.b, l.a), t = Ori(m.a, m.b, l.
    b);

```



```

    return (l.b * s - l.a * t) / (s - t);
}

```

### 10.3 Basic

```

//確保兩直線 $P+tv$ 和 $Q+tw$ 有唯一交點且 $Cross(v,w)$ 非零
Pt getLineIntersect(Line a, Line b) {
    Pt p1 = a.a, p2 = a.b, q1 = b.a, q2 = b.b;
    ld f1 = (p2-p1)^(q1-p1), f2 = (p2-p1)^(p1-q2), f;
    if(dcmp(f=f1+f2) == 0)
        return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
    return q1*(f2/f) + q2*(f1/f);
}

```

```

//點到直線距離
double distanceToLine(Pt p, Pt a, Pt b) {
    Vec v1=b-a, v2=p-a;
    return fabs(Cross(v1,v2)/Length(v1));
}

```

```

//點到線段距離
double distanceToSegment(Pt p, Pt a, Pt b) {
    if(a==b) return Length(p-a);
    Vec v1=b-a, v2=p-a, v3=p-b;
    if(dcmp(Dot(v1,v2))<0) return Length(v2);
    else if(dcmp(Dot(v1,v3))>0) return Length(v3);
    else return fabs(Cross(v1,v2)/Length(v1));
}

```

```

//點到直線投影
Pt GetLineProjection(Pt p, Pt a, Pt b) {
    Vec v=b-a;
    return a+v*(Dot(v,p-a)/Dot(v,v));
}

```

```

//點 $p$ 於直線 $ab$ 的對稱點
Pt getSymmetryPoint(Pt p, Pt a, Pt b) {
    Pt q=getLineProjection(p,a,b);
    return q*2-p;
}

```

```

//判斷線段相交(剛好交一點), 若兩線段共線-> $c1=c2=0$ 
bool isSegmentProperIntersection(Pt a1, Pt a2, Pt b1, Pt b2) {
    double c1=Cross(a2-a1,b1-a1), c2=Cross(a2-a1,b2-a1),
           c3=Cross(b2-b1,a1-b1), c4=Cross(b2-b1,a2-b1);
    return dcmp(c1)*dcmp(c2)<0&&dcmp(c3)*dcmp(c4)<0;
}

```

```

//判斷線段相交(只要有交點即可)
bool isSegmentNotProperIntersection(Pt a1, Pt a2, Pt b1, Pt b2) {
    return max(a1.x,a2.x)>=min(b1.x,b2.x)&&max(b1.x,b2.x)>=min(a1.x,a2.x)&&max(a1.y,a2.y)>=min(b1.y,b2.y)&&max(b1.y,b2.y)>=min(a1.y,a2.y)
    &&dcmp(Cross(a1-b1,a2-b1))*dcmp(Cross(a1-b2,a2-b2))<=0&&dcmp(Cross(b1-a1,b2-a1))*dcmp(Cross(b1-a2,b2-a2))<=0;
}

```

```

//點是否在線段上
bool isOnSegment(Pt p, Pt a1, Pt a2) {
    return dcmp(Cross(a1-p,a2-p))==0&&dcmp(Dot(a1-p,a2-p))<=0;
}

```

### 10.4 PolygonArea

```

//須注意Long Long 及 加上絕對值
double polygonArea(Pt* p, int n) {
    double area=0;
    for(int i=1; i<n-1; ++i) {
        area+=Cross(p[i]-p[0], p[i+1]-p[0]);
    }
    return area/2;
}

```

### 10.5 IsPointInPolygon

```

//判斷點是否在多邊形內部
int isPointInPolygon(Pt p, Pt* poly, int n) {
    int wn=0;
    for(int i=0; i<n; ++i) {
        if(isOnSegment(p, poly[i], poly[(i+1)%n])) return -1; //在邊界上
        int k=dcmp(Cross(poly[(i+1)%n]-poly[i], p-poly[i]));
        int d1=dcmp(poly[i].y-p.y);
        int d2=dcmp(poly[(i+1)%n].y-p.y);
        if(k>0&&d1<=0&&d2>0) wn++;
        if(k<0&&d2<=0&&d1>0) wn--;
    }
    if(wn!=0) return 1; //內部
    return 0; //外部
}

```

### 10.6 ConvexHull

```

//若要求高精度用dcmp比較
//若是搞int點要記得定義改int
//輸入不能有重複點, 注意h的點未排序!
//若需保留共線點, 把while裡的Ori判斷式改成<=0
void hull(vector<Pt> &dots) { // n=1 => ans = {}
    sort(dots.begin(), dots.end());
    vector<Pt> ans(1, dots[0]);
    for (int ct = 0; ct < 2; ++ct, reverse(all(dots)))
        for (int i = 1, t = SZ(ans); i < SZ(dots); ans.push_back(dots[i++]))
            while (SZ(ans) > t && Ori(ans[SZ(ans)-2], ans.back(), dots[i]) <= 0)
                ans.pop_back();
    ans.pop_back(), ans.swap(dots);
}

```

### 10.7 MinkowskiSum

```

//定義: 給兩點集合 $A, B$ ,  $Minkowski\ Sum$ 是 $A+B=\{a+b | a \in A, b \in B\}$ 的凸包
//視覺化為一個凸包  $A$  繞著凸包  $B$  轉一圈,  $O(N)$ 
//兩個凸多邊形的  $Minkowski\ sum$ , 也會是凸多邊形
// $P$  和  $Q$  組成的  $Minkowski\ sum$  最多有  $|P|+|Q|$  個點
//在凸包  $A$  和  $B$  上的邊也會在  $Minkowski\ sum$  上出現
//傳入的點集合不用逆時針排序
vector<Pt> Minkowski(vector<Pt> A, vector<Pt> B) { // |A|, |B| >= 3
    hull(A), hull(B);
    vector<Pt> C(1, A[0] + B[0]), s1, s2;
    for (int i = 0; i < SZ(A); ++i)
        s1.push_back(A[(i+1) % SZ(A)] - A[i]);
    for (int i = 0; i < SZ(B); ++i)
        s2.push_back(B[(i+1) % SZ(B)] - B[i]);
    for (int i = 0, j = 0; i < SZ(A) || j < SZ(B); )
        if (j >= SZ(B) || (i < SZ(A) && Cross(s1[i], s2[j]) >= 0))
            C.push_back(B[j % SZ(B)] + A[i++]);
        else
            C.push_back(A[i % SZ(A)] + B[j++]);
    return hull(C), C;
}

```

### 10.8 Polygon Shortest Distance

```

//給兩多邊形, 求最短距離
double PolyDist(vector<Pt> A, vector<Pt> B) {
    for (auto &p : B) p = {-p.x, -p.y};
    auto C = Minkowski(A, B); // assert SZ(C) > 0
    if (isPointInPolygon(Pt{}, C.data(), C.size()))
        return 0;
    double ans = distanceToSegment(Pt{}, C.back(), C[0]);
    for (int i = 0; i + 1 < SZ(C); ++i) {

```

```

        ans = min(ans, distanceToSegment(Pt{ }, C[i], C[i
            + 1]));
    }
    return ans;
}

```

## 10.9 ConvexHullTrick

```

struct Convex {
    int n;
    vector<Pt> A, V, L, U;
    //init, pass convex hull points
    Convex(const vector<Pt> &_A) : A(_A), n(_A.size())
    { // n >= 3
        auto it = max_element(all(A));
        L.assign(A.begin(), it + 1);
        U.assign(it, A.end());
        U.push_back(A[0]);
        for (int i = 0; i < n; i++) {
            V.push_back(A[(i + 1) % n] - A[i]);
        }
    }
    int PtSide(Pt p, Line L) {
        return dcmp((L.b - L.a)^(p - L.a));
    }
    int inside(Pt p, const vector<Pt> &h, auto f) {
        auto it = lower_bound(all(h), p, f);
        if (it == h.end()) return 0;
        if (it == h.begin()) return p == *it;
        return 1 - dcmp((p - *prev(it))^(it - *prev(
            it)));
    }
    // 1. whether a given point is inside the Convex
    Hull
    // ret 0: out, 1: on, 2: in
    int inside(Pt p) {
        return min(inside(p, L, less{}), inside(p, U,
            greater{}));
    }
    static bool cmp(Pt a, Pt b) { return dcmp(a ^ b) >
        0; }
    // 2. Find tangent points of a given vector
    // ret the idx of far/closer tangent point
    int tangent(Pt v, bool close = true) {
        assert(v != Pt{ });
        auto l = V.begin(), r = V.begin() + L.size() -
            1;
        if (v < Pt{ }) l = r, r = V.end();
        if (close) return (lower_bound(l, r, v, cmp) -
            V.begin()) % n;
        return (upper_bound(l, r, v, cmp) - V.begin())
            % n;
    }
    // 3. Find 2 tang pts on CH of a given outside
    point
    // return index of tangent points
    // return {-1, -1} if inside CH
    array<int, 2> tangent2(Pt p) {
        array<int, 2> t{-1, -1};
        if (inside(p) == 2) return t;
        if (auto it = lower_bound(all(L), p); it != L.
            end() and p == *it) {
            int s = it - L.begin();
            return {(s + 1) % n, (s - 1 + n) % n};
        }
        if (auto it = lower_bound(all(U), p, greater{});
            it != U.end() and p == *it) {
            int s = it - U.begin() + L.size() - 1;
            return {(s + 1) % n, (s - 1 + n) % n};
        }
        for (int i = 0; i != t[0]; i = tangent((A[t[0]
            = i] - p), 0));
        for (int i = 0; i != t[1]; i = tangent((p - A[t
            [1] = i]), 1));
        return t;
    }
    int find(int l, int r, Line L) {
        if (r < l) r += n;
        int s = PtSide(A[l % n], L);
        return *ranges::partition_point(views::iota(l,
            r),
            [&](int m) {

```

```

                return PtSide(A[m % n], L) == s;
            }) - 1;
    };
    // 4. Find intersection point of a given Line
    // intersection is on edge (i, next(i))
    vector<int> intersect(Line L) {
        int l = tangent(L.a - L.b), r = tangent(L.b - L
            .a);
        if (PtSide(A[l], L) == 0) return {l};
        if (PtSide(A[r], L) == 0) return {r};
        if (PtSide(A[l], L) * PtSide(A[r], L) > 0)
            return {};
        return {find(l, r, L) % n, find(r, l, L) % n};
    }
};

```

## 10.10 Polar Sort

```

//極角排序，從270度開始逆時針排序
bool cmp(const Pt& lhs, const Pt& rhs) {
    if (Cross((lhs < Pt{ }), (rhs < Pt{ })))
        return (lhs < Pt{ }) < (rhs < Pt{ });
    return Cross(lhs, rhs) > 0;
}
/* 若要以p[i]為原點排序->計算v=p[j]-p[i]
for(int j=0;j<n;j++){
    if(i!=j){
        Vector v = p[j]-p[i];
        node[nodeSz++] = {v,j};
    }
}
sort(node,node+nodeSz,cmp);
*/

```

## 10.11 PickTheorm

```

int area, in, on; //area:多邊形面積 in:內部格點數 on:邊界
格點數
void PickTheorm(Pt* p, int n) {
    area = polygonArea(p, n);
    for (int i = 0; i < n; ++i) {
        on += __gcd(abs((int)p[i].x - (int)p[(i + 1) % n].x),
            abs((int)p[i].y - (int)p[(i + 1) % n].y));
    }
    in = abs(area) + 1 - on / 2;
}

```

## 10.12 ShortestPair

```

//最近點對距離注意若整數要define double Long Long
double closestEuclideanDistance(Pt* p, int n) {
    sort(p, p + n);
    set<Pt> s = {{p[0].y, p[0].x}};
    int j = 0;
    Pt t;
    double dd = LLONG_MAX, d;
    for (int i = 1; i < n; ++i) {
        d = sqrt(dd);
        while (j < i && p[j].x < p[i].x - d) {
            s.erase({p[j].y, p[j].x});
        }
        auto l = s.lower_bound({p[i].y - d, p[i].x - d});
        auto u = s.upper_bound({p[i].y + d, p[i].x + d});
        for (auto it = l; it != u; ++it) {
            t = {it->y, it->x};
            dd = min(dd, Dot(p[i] - t, p[i] - t));
        }
        s.emplace(p[i].y, p[i].x);
    }
    return dd;
}

```

## 10.13 FarthestPair

```

ld FarthestPair(vector<Pt>& poly){//旋轉卡尺
    ld ret=0;
    for(int i = 0, j = i+1; i<poly.size()&& j<poly.size()
        (); i++){
        while( Length(poly[i]-poly[j]) <= Length(poly[i]
            ]-poly[(j+1)%poly.size()]) && j!=i)
            j = (j+1) % poly.size();
        ret = max(ret, Length(poly[i]-poly[j]));
    }
    return ret;
}

```

## 10.14 幾何中位數

```

//回傳為到每個頂點距離和最小的點
Pt weiszfeld(const Pt *p,int n){
    double nn=n;
    Pt cur = p[0], next;
    for(int i=1;i<n;++i)
        cur.x+=p[i].x, cur.y+=p[i].y;
    cur.x/=nn, cur.y/=nn;
    double w,numerX,numerY,denomin;
    while(1){
        numerX=numerY=denomin=0;
        bool update=0;
        double d;
        for(int i=0;i<n;++i){
            d=Length(cur-p[i]);
            if(d>eps){
                w = 1.0/d;
                numerX+=w*p[i].x;
                numerY+=w*p[i].y;
                denomin+=w;
                update=1;
            }else{
                next = p[i];
                break;
            }
        }
        if(update){
            next.x = numerX/denomin;
            next.y = numerY/denomin;
        }
        if(Length(cur-next)<eps) break;
        cur = next;
    }
    return next;
}

```

## 10.15 矩陣掃描線

```

#include <bits/stdc++.h>
#define int long long int
using namespace std;
int n, st[1000005<<2], lazy[1000005<<2], old
    [1000005<<2];
vector<tuple<int, int, int, int>> v;
vector<int> sor;
void pull(int index, int l, int r) {
    if(lazy[index]) st[index] = old[index];
    else if(l==r) st[index] = 0;
    else st[index] = st[index<<1]+st[index<<1];
    // printf("pull %lld~%lld, %lld\n", l, r, st[index]
    );
    return;
}
void insert(int index, int s, int e, int l, int r, int
    k) {
    //printf("insert: range %lld~%lld, query %lld~%lld\
    n", s, e, l, r);
    if(l<=s && e<=r) {
        lazy[index] +=k;
        pull(index, s, e);
        return;
    }
    int mid = (s+e)/2;
    if(l<=mid) insert(index<<1, s, mid, l, r, k);

```

```

    if(mid<r) insert(index<<1|1, mid+1, e, l, r, k);
    pull(index, s, e);
}
void input(int index, int l, int r) {
    if(l==r) {
        old[index] = sor[l]-sor[l-1];
        return;
    }
    int mid = (l+r)/2;
    input(index<<1, l, mid);
    input(index<<1|1, mid+1, r);
    old[index] = old[index<<1] + old[index<<1|1];
    //cout<<l<<" to "<<r<<" is "<<old[index]<<endl;
    return;
}
// int diff=1000005;
signed main(){
    cin >> n;
    int l, r, d, u;
    for (int i = 0; i < n; i++){
        cin >> l >> d >> r >> u;
        // l+=diff;
        // d+=diff;
        // r+=diff;
        // u+=diff;
        sor.push_back(d);
        sor.push_back(u);
        v.push_back({l, d, u, 1});
        v.push_back({r, d, u, -1});
    }
    set<int> temp(sor.begin(), sor.end());
    sor = vector<int>(temp.begin(), temp.end());
    sort(sor.begin(), sor.end());
    for(int i=0 ; i<v.size() ; i++) {
        auto [a, b, c, k] = v[i];
        v[i] = make_tuple(a, (int)(lower_bound(sor.
            begin(), sor.end(), b)-sor.begin()), (int)(
            lower_bound(sor.begin(), sor.end(), c)-sor.
            begin()), k);
    }
    input(1, 1, sor.size()-1);
    // cout<<"get: ";
    // for(int i: sor) cout<<i<<" "; cout<<endl;
    sort(v.begin(), v.end());
    int pre=0;
    int ans=0;
    for(auto [pos, a, b, k]: v) {
        if(pre!=pos) {
            ans+=(pos-pre)*st[1];
            pre = pos;
        }
        insert(1, 1, sor.size()-1, a+1, b, k);
        // printf("now act: pos %lld, %lld~%lld, act:
        %lld\n", pos, a+1, b, k);
        // printf("now ans: %lld\n", st[1]);
    }
    cout<<ans<<endl;
}

```

## 10.16 Polygon Circle intersection area

```

//給兩圓，回傳兩圓的交點，若相交回傳的兩點一樣的話要用
dcmp判斷
vector<Pt> interCircle(Circle c1, Circle c2) {
    Pt o1 = c1.o, o2 = c2.o;
    ld r1 = c1.r, r2 = c2.r;
    if( Length( o1 - o2 ) > r1 + r2 ) return {};
    if( Length( o1 - o2 ) < max(r1, r2) - min(r1, r2) )
        return {};
    ld d2 = ( o1 - o2 ) * ( o1 - o2 );
    ld d = sqrt(d2);
    if( d > r1 + r2 ) return {};
    Pt u = (o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
    ld A = sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d))
        ;
    Pt v = Pt( o1.y-o2.y , -o1.x + o2.x ) * A / (2*d2);
    return {u+v, u-v};
}

```

## 10.17 Tangent line of two circles

```
//給兩圓，求兩圓的外切線或內切線
vector<Line> go( const Circle& c1 , const Circle& c2 ,
    int sign1 ){
    // sign1 = 1 for outer tang, -1 for inter tang
    vector<Line> ret;
    double d_sq = Dot( c1.o - c2.o , c1.o - c2.o );
    if( d_sq < eps ) return ret;
    double d = sqrt( d_sq );
    Pt v = ( c2.o - c1.o ) / d;
    double c = ( c1.r - sign1 * c2.r ) / d;
    if( c * c > 1 ) return ret;
    double h = sqrt( max( 0.0 , 1.0 - c * c ) );
    for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
        Pt n = { v.x * c - sign2 * h * v.y ,
            v.y * c + sign2 * h * v.x };
        Pt p1 = c1.o + n * c1.r;
        Pt p2 = c2.o + n * ( c2.r * sign1 );
        if( fabs( p1.x - p2.x ) < eps and
            fabs( p1.y - p2.y ) < eps )
            p2 = p1 + Perp( c2.o - c1.o );
        ret.push_back( { p1 , p2 } );
    }
    return ret;
}
```

## 10.18 Circle intersection Point

```
//傳入多邊形和圓形，回傳多邊形和圓形的交集面積
ld PCIntersect(vector<Pt> v, Circle cir) {
    for(int i = 0 ; i < (int)v.size() ; ++i) v[i] = v[i]
        - cir.o;
    ld ans = 0, r = cir.r;
    int n = v.size();
    for(int i = 0 ; i < n ; ++i) {
        Pt pa = v[i], pb = v[(i+1)%n];
        if(Length(pa) < Length(pb)) swap(pa, pb);
        if(dcmp(Length(pb)) == 0) continue;
        ld s, h, theta;
        ld a = Length(pb), b = Length(pa), c = Length(pb-pa);
        ld cosB = (pb*(pb-pa))/a/c, B = acos(cosB);
        if(cosB > 1) B = 0;
        else if(cosB < -1) B = PI;
        ld cosC = (pa*pb)/a/b, C = acos(cosC);
        if(cosC > 1) C = 0;
        else if(cosC < -1) C = PI;
        if(a > r) {
            s = (C/2)*r*r;
            h = a*b*sin(C)/c;
            if(h < r && B < PI/2) s -= (acos(h/r)*r*r - h*
                sqrt(r*r-h*h));
        }
        else if(b > r) {
            theta = PI - B - asin(sin(B)/r*a);
            s = 0.5*a*r*sin(theta) + (C-theta)/2*r*r;
        }
        else s = 0.5*sin(C)*a*b;
        ans += abs(s)*dcmp(v[i]^v[(i+1)%n]);
    }
    return abs(ans);
}
```

## 10.19 CircleCover

```
#define N 100
#define D long double
struct CircleCover{//O(N^2LogN)
    int C; Circle c[ N ]; //填入C(圓數量),c(圓陣列,0base)
    bool g[ N ][ N ], overlap[ N ][ N ];
    // Area[i] : area covered by "at least" i circles
    D Area[ N ];
    void init( int _C ){ C = _C; } //總共 _C 個圓
    bool CCinter( Circle& a , Circle& b , Pt& p1 , Pt& p2
        ){
        Pt o1 = a.o , o2 = b.o;
        D r1 = a.r , r2 = b.r;
```

```
if( Length( o1 - o2 ) > r1 + r2 ) return {};
if( Length( o1 - o2 ) < max(r1, r2) - min(r1, r2) )
    return {};
D d2 = ( o1 - o2 ) * ( o1 - o2 );
D d = sqrt(d2);
if( d > r1 + r2 ) return false;
Pt u=(o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
D A=sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d));
Pt v=Pt( o1.y-o2.y , -o1.x + o2.x ) * A / (2*d2);
Pt p1 = u + v; p2 = u - v;
return true;
}
struct Teve {
    Pt p; D ang; int add;
    Teve() {}
    Teve(Pt _a, D _b, int _c):p(_a), ang(_b), add(_c){}
    bool operator<(const Teve &a)const
    {return ang < a.ang;}
}eve[ N * 2 ];
// strict: x = 0, otherwise x = -1
bool disjuct( Circle& a, Circle &b, int x )
{return dcmp( Length( a.o - b.o ) - a.r - b.r ) > x;}
bool contain( Circle& a, Circle &b, int x )
{return dcmp( a.r - b.r - Length( a.o - b.o ) ) > x;}
bool contain(int i, int j){
    /* c[j] is non-strictly in c[i]. */
    return (dcmp(c[i].r - c[j].r) > 0 ||
        (dcmp(c[i].r - c[j].r) == 0 && i < j) ) &&
        contain(c[i], c[j], -1);
}
void solve(){
    for( int i = 0 ; i <= C + 1 ; i ++ )
        Area[ i ] = 0;
    for( int i = 0 ; i < C ; i ++ )
        for( int j = 0 ; j < C ; j ++ )
            overlap[i][j] = contain(i, j);
    for( int i = 0 ; i < C ; i ++ )
        for( int j = 0 ; j < C ; j ++ )
            g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                disjuct(c[i], c[j], -1));
    for( int i = 0 ; i < C ; i ++ ){
        int E = 0, cnt = 1;
        for( int j = 0 ; j < C ; j ++ )
            if( j != i && overlap[j][i] )
                cnt ++;
        for( int j = 0 ; j < C ; j ++ )
            if( i != j && g[i][j] ){
                Pt aa, bb;
                CCinter(c[i], c[j], aa, bb);
                D A=atan2(aa.y - c[i].o.y, aa.x - c[i].o.x);
                D B=atan2(bb.y - c[i].o.y, bb.x - c[i].o.x);
                eve[E ++] = Teve(bb, B, 1);
                eve[E ++] = Teve(aa, A, -1);
                if(B > A) cnt ++;
            }
        if( E == 0 ) Area[ cnt ] += PI * c[i].r * c[i].r;
        else{
            sort( eve , eve + E );
            eve[E] = eve[0];
            for( int j = 0 ; j < E ; j ++ ){
                cnt += eve[j].add;
                Area[cnt] += (eve[j].p ^ eve[j + 1].p) * 0.5;
                D theta = eve[j + 1].ang - eve[j].ang;
                if( theta < 0 ) theta += 2.0 * PI;
                Area[cnt] +=
                    (theta - sin(theta)) * c[i].r*c[i].r * 0.5;
            }
        }
    }
}
```

## 10.20 Mec

```
// O(N) Minimum enclosing circle
struct Mec{ // return pair of center and r
    int n;
    Pt p[ MXN ], cen;
    double r2;
    void init( int _n , Pt _p[] ){
        n = _n;
        memcpy( p , _p , sizeof(Pt) * n );
    }
    double sqr(double a){ return a*a; }
    Pt center(Pt p0, Pt p1, Pt p2) {
```

```

Pt a = p1-p0;
Pt b = p2-p0;
double c1=Dot(a,a) * 0.5;
double c2=Dot(b,b) * 0.5;
double d = a ^ b;
double x = p0.x + (c1 * b.y - c2 * a.y) / d;
double y = p0.y + (a.x * c2 - b.x * c1) / d;
return Pt(x,y);
}
pair<Pt,double> solve(){
    random_shuffle(p,p+n);
    r2=0;
    for (int i=0; i<n; i++){
        if (Dot(cen-p[i],cen-p[i]) <= r2) continue;
        cen = p[i];
        r2 = 0;
        for (int j=0; j<i; j++){
            if (Dot(cen-p[j],cen-p[j]) <= r2) continue;
            cen=Pt((p[i].x+p[j].x)/2,(p[i].y+p[j].y)/2);
            r2 = Dot(cen-p[j],cen-p[j]);
            for (int k=0; k<j; k++){
                if (Dot(cen-p[k],cen-p[k]) <= r2) continue;
                cen = center(p[i],p[j],p[k]);
                r2 = Dot(cen-p[k],cen-p[k]);
            }
        }
        return {cen,sqrt(r2)};
    }
}mec;

```

## 10.21 半平面交

```

//O(NLgN)
// for point or line solution, change > to >=
bool onleft(Line L, Pt p) {
    return dcmp(L.v^(p-L.a)) > 0;
} // segment should add Counterclockwise
// assume that Lines intersect
// 傳入每條方程式的兩點方程式
// 回傳形成的凸多邊形頂點
// (半平面為像量 ab 的逆時針方向)
//注意題目輸入的點要是逆時針排序
vector<Pt> HPI(vector<Line>& L) {
    sort(L.begin(), L.end()); // sort by angle
    for(auto l:L){
        cerr<<l.a<<" "<<l.a.y<<" "<<l.b.x<<" "<<l.b.y<<"
            "<<l.ang<<"\n";
    }
    int n = L.size(), fir, las;
    Pt *p = new Pt[n];
    Line *q = new Line[n];
    q[fir=las=0] = L[0];
    for(int i = 1; i < n; i++) {
        while(fir < las && !onleft(L[i], p[las-1])) las--;
        while(fir < las && !onleft(L[i], p[fir])) fir++;
        q[++las] = L[i];
        if(dcmp(q[las].v^q[las-1].v) == 0) {
            las--;
            if(onleft(q[las], L[i].a)) q[las] = L[i];
        }
        if(fir < las) p[las-1] = getLineIntersect(q[las-1], q[las]);
    }
    while(fir < las && !onleft(q[fir], p[las-1])) las--;
    if(las-fir <= 1) return {};
    p[las] = getLineIntersect(q[las], q[fir]);
    int m = 0;
    vector<Pt> ans(las-fir+1);
    for(int i = fir; i <= las; i++) ans[m++] = p[i];
    return ans;
}

```

## 10.22 PolygonUnion

```

//O(N^2LgN)
//傳入二維vector，每個vector代表一個多邊形，每個多邊形
//的點必須按照順時針或逆時針順序
//回傳聯集多邊形的面積
ld tri(Pt o, Pt a, Pt b){ return (a-o) ^ (b-o);}

```

```

double polyUnion(vector<vector<Pt>> py){ //py[0~n-1]
    must be filled
    int n = py.size();
    int i,j,ii,jj,ta,tb,r,d; double z,w,s,sum=0,tc,td,
        area;
    vector<pair<double,int>> c;
    for(i=0;i<n;i++){
        area=py[i][py[i].size()-1]^py[i][0];
        for(int j=0;j<py[i].size()-1;j++) area+=py[i][j]^py[i][j+1];
        if((area/2)<0) reverse(py[i].begin(),py[i].end());
        py[i].push_back(py[i][0]);
    }
    for(i=0;i<n;i++){
        for(ii=0;ii+1<py[i].size();ii++){
            c.clear();
            c.emplace_back(0.0,0); c.emplace_back(1.0,0);
            for(j=0;j<n;j++){
                if(i==j) continue;
                for(jj=0;jj+1<py[j].size();jj++){
                    ta=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj]))
                        ;
                    tb=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj+1]))
                        ;
                    if(ta==0 && tb==0){
                        if((py[j][jj+1]-py[j][jj])*(py[i][ii+1]-py[i][ii])>0&&j<i){
                            c.emplace_back(segP(py[j][jj],py[i][ii],
                                py[i][ii+1]),1);
                            c.emplace_back(segP(py[j][jj+1],py[i][ii],
                                py[i][ii+1]),-1);
                        }
                    }
                    else if(ta>=0 && tb<0){
                        tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
                        td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
                        c.emplace_back(tc/(tc-td),1);
                    }
                    else if(ta<0 && tb>=0){
                        tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
                        td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
                        c.emplace_back(tc/(tc-td),-1);
                    }
                }
            }
            sort(c.begin(),c.end());
            z=min(max(c[0].first,0.0),1.0); d=c[0].second; s
                =0;
            for(j=1;j<c.size();j++){
                w=min(max(c[j].first,0.0),1.0);
                if(!d) s+=w-z;
                d+=c[j].second; z=w;
            }
            sum+=(py[i][ii]^py[i][ii+1])*s;
        }
    }
    return sum/2;
}

```

## 10.23 PolygonCover

```

//傳入二維vector，每個vector代表一個多邊形，每個多邊形
//的點必須按照順時針或逆時針順序
// return Area[i] : area covered by "at Least" i
// polygon
//O(N^2LgN)
vector<double> PolyUnion(const vector<vector<Pt>> &P) {
    auto reorder = [&](vector<Pt> &v) { //排序成逆時針
        且最左下角的點在最前面
        rotate(v.begin(), min_element(all(v)), v.end())
            ;
        if (Ori(v[0], v[1], v[2]) < 0) {
            reverse(all(v));
        }
    };
    for(auto &i: py){ reorder(i);}
    const int n = P.size();
    vector<double> Area(n + 1);
    vector<Line> Ls;
    for (int i = 0; i < n; i++)
        for (int j = 0; j < P[i].size(); j++)
            Ls.push_back({P[i][j], P[i][(j + 1) % P[i].size()]});
    auto cmp = [&](Line &l, Line &r) {
        Pt u = l.b - l.a, v = r.b - r.a;
    };
}

```



```

    if (argcmp(u, v)) return true;
    if (argcmp(v, u)) return false;
    return PtSide(l.a, r) < 0;
};
sort(all(Ls), cmp);
for (int l = 0, r = 0; l < Ls.size(); l = r) {
    while (r < Ls.size() and !cmp(Ls[l], Ls[r])) r
        ++;
    Line L = Ls[l];
    vector<pair<Pt, int>> event;
    for (auto &ls : Ls) {
        Pt c = ls.a, d = ls.b;
        if (Sgn((L.a - L.b) ^ (c - d)) != 0) {
            int s1 = PtSide(c, L) == 1;
            int s2 = PtSide(d, L) == 1;
            if (s1 ^ s2) event.emplace_back(
                LineInter(L, {c, d}), s1 ? 1 : -1);
        } else if (PtSide(c, L) == 0 and Sgn((L.a -
            L.b) * (c - d)) > 0) {
            event.emplace_back(c, 2);
            event.emplace_back(d, -2);
        }
    }
    sort(all(event), [&](auto i, auto j) {
        return (L.a - i.ft) * (L.a - L.b) < (L.a -
            j.ft) * (L.a - L.b);
    });
    int cov = 0, tag = 0;
    Pt lst{0, 0};
    for (auto [p, s] : event) {
        if (cov >= tag) {
            Area[cov] += lst ^ p;
            Area[cov - tag] -= lst ^ p;
        }
        if (abs(s) == 1) cov += s;
        else tag += s / 2;
        lst = p;
    }
    for (int i = n - 1; i >= 0; i--) Area[i] += Area[i
        + 1];
    for (int i = 1; i <= n; i++) Area[i] /= 2;
    return Area;
};

```

## 11 特殊題目

### 11.1 包含子字串計數

```

// * 給一個字串s
// * 求長度為len且有包含s的字串有幾種
// * 呼叫solve(s, len)
const int len = 1005;
int aut[len][26];
int dp[len][len];
const int mod = 1e9+7;
void prefix(string &s, vector<int> &pi) {
    for(int i=1, j=0 ; i<s.size() ; i++) {
        while(j>0 && s[i]!=s[j]) j = pi[j-1];
        if(s[i]==s[j]) j++;
        pi[i] = j;
    }
}
void automata(string &s, vector<int> &pi) {
    for(int i=0 ; i<s.size() ; i++) {
        for(int c=0 ; c<26 ; c++) {
            if(i>0 && c+'A' != s[i]) aut[i][c] = aut[pi
                [i-1]][c];
            else aut[i][c] = i + (c + 'A'==s[i]);
        }
    }
}
int quai(int x, int n) {
    if(n==0) return 1;
    int mid = quai(x,n/2);
    mid = mid*mid%mod;
    if(n&1) return mid*x%mod;
    return mid;
}
int solve(string s, int len) {
    vector<int> pi(s.size(), 0);

```

```

    prefix(s, pi);
    automata(s, pi);
    int n = s.size(), ans = quai(26, len);
    dp[0][0] = 1;
    for(int i=0 ; i<len ; i++) {
        for(int j=0 ; j<n ; j++) {
            for(int c=0 ; c<26 ; c++) {
                dp[i+1][aut[j][c]] += dp[i][j];
                dp[i+1][aut[j][c]] %= mod;
            }
        }
    }
    for(int i=0 ; i<n ; i++) ans = (ans - dp[len][i] +
        mod)%mod;
    return ans;
}

```

### 11.2 三維偏序

```

// vec
// {{a, b, c},
// {a, b, c},
// ...
// {a, b, c}}
// 貼上 BIT 模板
// 三維偏序
// a <= a, b <= b, c <= c
map<vector<int>, int> cnt;
int cdq(vector<vector<int>> &vec, int l, int r) {
    if(l==r) return 0;
    int mid = l+r>>1;
    int ans = cdq(vec, l, mid)+cdq(vec, mid+1, r);
    vector<vector<int>> temp;
    for(int i=l, j=mid+1 ; i<=mid || j<=r ; ) {
        while(i<=mid && (j>r || vec[i][1] <= vec[j][1]))
            {bit.add(vec[i][2], cnt[vec[i]]); temp.
                push_back(vec[i++]);}
        if(j<=r) {
            temp.push_back(vec[j]);
            ans += bit.query(vec[j][2]);
        }
    }
    for(int i=l ; i<=mid ; i++) bit.add(vec[i][2], -cnt[
        vec[i]]);
    for(int i=l ; i<=r ; i++) vec[i] = temp[i-1];
    return ans;
}
int solve(vector<vector<int>> &vec) {
    bit.init(2e5+5);
    for(vector<int> v: vec) cnt[v]++;
    sort(vec.begin(), vec.end());
    vec.erase(unique(vec.begin(), vec.end()), vec.end()
        );
    return cdq(vec, 0, vec.size()-1);
}

```

## 12 Python

### 12.1 時間日期 Datetime

```

from datetime import datetime, date, time, timedelta

# 閏年 2024
# 平年 2023

# 日期相減
start_time = datetime(2024, 8, 31)
end_time = datetime(2024, 9, 1)
delta = end_time - start_time
print(delta.days)
# 1

# 時間相減
start_time = datetime(1000, 1, 1, 14, 30)
end_time = datetime(1000, 1, 1, 16, 50)
delta = end_time - start_time
print(delta) # 2:20:00
print(delta.total_seconds()) # 8400.0

```

```
# 時間日期相加
specific_date = datetime(1000, 1, 1)
new_date = specific_date + timedelta(days=3, hours=1,
    minutes=1)
print(new_date)
# 1000-01-04 01:01:00
print(new_date.year, new_date.month, new_date.day,
    new_date.hour, new_date.minute, new_date.second)
# 1000 1 4 1 1 0
```

## 12.2 Decimal

```
from decimal import Decimal, getcontext, ROUND_FLOOR
getcontext().prec = 250 # set precision (MAX_PREC)
getcontext().Emax = 250 # set exponent limit (MAX_EMAX)
getcontext().rounding = ROUND_FLOOR # set round floor
itwo,two,N = Decimal(0.5),Decimal(2),200
pi = angle(Decimal(-1))
```

## 12.3 Fraction

```
from fractions import Fraction
import math
"""專門用來表示和操作有理數，可以進行算"""
frac1 = Fraction(1) # 1/1
frac2 = Fraction(1, 3) # 1/3
frac3 = Fraction(0.5) # 1/2
frac4 = Fraction('22/7') # 22/7
frac5 = Fraction(8, 16) # 自動約分為 1/2
frac9 = Fraction(22, 7)
frac9.numerator # 22
frac9.denominator # 7
x = Fraction(math.pi)
y2 = x.limit_denominator(100) # 分母限制為 100
print(y2) # 311/99
float(x) #轉換為浮點數
```

## 12.4 正則表達式 re

```
import re

# \d 0~9
# \D 非0~9
# \s 空白字符
# \S 非空白字符
# ^ 開頭
# $ 結尾
# re* 0個或多個
# re+ 1個或多個
# re? 0個或1個
# re{n} 恰好n次

s = 'a1bb2c2'
print(re.search(r'(?<=\d)[a-z](?=\d)', s).group())
# b
# 匹配前面和後面都要是數字的a~z字元

s = 'a1a a2b b3b'
print(re.search(r'(?<!a)\d(?!a)', s).group())
# 3
# 匹配前後都不是a的數字

s = 'aBc'
print(re.search(r'abc', s, re.I).group())
# aBc
# re.I 忽略大小寫
# re.S 任意字符 . 不受換行(\n)和空白限制

s = 'abbbbb'
print(re.search(r'ab+?', s).group())
# ab
# 非貪婪用法，匹配ab+但長度越短越好

s = 'abcde'
```

```
print(re.findall(r'[a-z]', s))
# ['a', 'b', 'c', 'd', 'e']

s = 'abbbba'
print(re.sub('(<?<=a).*(?=a)', '', s))
# aa
# 把.*的部分替換成空字串
```

## 12.5 Misc

```
# 轉為高精度整數比，(分子，分母)
x=0.2
x.as_integer_ratio() # (8106479329266893,
    9007199254740992)
x.is_integer() # 判斷是否為整數
x.__round__() # 四捨五入
int(eval(num.replace("/", "//"))) # parser string num
```

[illegible]

