

# Contents

<b>1</b>	<b>Basic</b>	<b>1</b>	10.4	PolygonArea	18
1.1	Default code	1	10.5	IsPointInPolygon	18
1.2	Misc	1	10.6	ConvexHull	18
1.3	Fast read & write	1	10.7	MinkowskiSum	19
1.4	Sort cmp	2	10.8	Polygon Shortest Distance	19
1.5	Discretization	2	10.9	ConvexHullTrick	19
1.6	Custom unordered_map	2	10.10	Polar Sort	19
1.7	__int128 read	2	10.11	PickTheorem	19
1.8	字典序 a 嚴格小於 b	2	10.12	最近點對	20
1.9	生成 n 位數的二進制組合	2	10.13	幾何中位數	20
1.10	Radom	2	10.14	矩陣掃描線	20
<b>2</b>	<b>對拍</b>	<b>2</b>	10.15	Polygon Circle intersection area	21
2.1	run.bat	2	10.16	Tangent line of two circles	21
2.2	run.sh	2	10.17	Circle intersection Point	21
<b>3</b>	<b>Flow &amp; Matching</b>	<b>2</b>	10.18	CircleCover	21
3.1	Dicnic	2	10.19	HalfPlaneIntersection	22
3.2	最大流最小花費	3	10.20	PolygonUnion	22
3.3	匈牙利匹配	3	10.21	PolygonCover	22
3.4	KM	4	<b>11</b>	<b>特殊題目</b>	<b>23</b>
<b>4</b>	<b>Graph</b>	<b>4</b>	11.1	包含子字串計數	23
4.1	Dijkstra	4	11.2	三維偏序	23
4.2	Bellman-Ford	4	<b>12</b>	<b>Python</b>	<b>23</b>
4.3	SPFA	4	12.1	Decimal	23
4.4	Floyd-Warshall	4	12.2	Fraction	23
4.5	歐拉路徑	5	12.3	正則表達式 re	23
4.6	BCC	5	12.4	Misc	24
4.7	SCC	5			
4.8	2SAT	6			
4.9	MaximalClique	6			
4.10	MaximumClique	6			
4.11	Minimum Mean Cycle	6			
4.12	Dominator Tree	7			
4.13	ManhattanMST	7			
<b>5</b>	<b>DP</b>	<b>8</b>			
5.1	數位 DP	8			
5.2	SOS DP	8			
<b>6</b>	<b>Math</b>	<b>8</b>			
6.1	Formulas	8			
6.2	l1l2l3mul	8			
6.3	Primes	8			
6.4	Coprime (互質 Pair)	8			
6.5	取樣定理	8			
6.6	Quick Pow	8			
6.7	Mat quick Pow	8			
6.8	Primes Table	9			
6.9	Phi 函數	9			
6.10	Factor Table	9			
6.11	卡塔蘭數	9			
6.12	Miller Rabin	9			
6.13	PollarRho	9			
6.14	PrimeFactorO(logn)	9			
6.15	O(1)mul	10			
6.16	Josephus Problem	10			
6.17	Harmonic Sum	10			
<b>7</b>	<b>Data Structure</b>	<b>10</b>			
7.1	BIT	10			
7.2	BIT 二維	10			
7.3	稀疏表 O(1) 區間最大最小值	10			
7.4	Segment Tree	11			
7.5	動態開點線段數	11			
7.6	動態開點線段數 2D	11			
7.7	持久化線段樹	12			
7.8	Time Segment Tree	12			
7.9	Treap	13			
7.10	PBDS	13			
<b>8</b>	<b>String</b>	<b>13</b>			
8.1	SA	13			
8.2	KMP	14			
8.3	Single Hash	15			
8.4	Double Hash	15			
8.5	Trie	15			
8.6	Z value	15			
8.7	MinRotation	15			
8.8	Manacher 馬拉車回文	15			
8.9	PalTree 回文樹	16			
8.10	DistinctSubsequence	16			
<b>9</b>	<b>Tree</b>	<b>16</b>			
9.1	LCA	16			
9.2	TreeHash	17			
9.3	輕重鍵剖分	17			
<b>10</b>	<b>Geometry</b>	<b>17</b>			
10.1	2D Definition	17			
10.2	Line Definition	18			
10.3	Basic	18			

# 1 Basic

## 1.1 Default code

```
#include<bits/stdc++.h>
#include<chrono> // for timing
#pragma GCC optimize("O3,unroll-loops")
#pragma target optimize("avx2,bmi,bmi2,lzcnt,popcnt")
#define IO ios_base::sync_with_stdio(0);cin.tie(0);cout
      .tie(0);
#define pii pair<int,int>
#define ft first
#define sd second
#define int long long
#define ld long double
#define PI acos(-1)
#define SZ(x) (int)x.size()
#define all(v) (v).begin(), (v).end()
#define _for(i,a,b) for(int i=(a);i<(b);++i)
using namespace std;
template<typename T>
ostream& operator<<(ostream& os,const vector<T>& vn){
    for(int i=0;i<vn.size();++i)os<<vn[i]<<" ";
    return os;
}
template<typename T>
ostream& operator<<(ostream& os,const set<T>& vn){
    for(typename set<T>::iterator it=vn.begin();it!=vn.
        end();++it)os<<*it<<" ";
    return os;
}
mt19937 mt(hash<string>()("Mashu_AC_Please")); //mt();
// mt19937 mt(chrono::steady_clock::now().
    time_since_epoch().count());
// g++ a.cpp -Wall -Wshadow -fsanitize=undefined -o a.
    exe
// ./a.exe
const int MXN=2e5+5;
const int INF=INT_MAX;
void sol() {}
signed main() {
    // auto start=chrono::high_resolution_clock::now();
    // #ifdef LOCAL
    // freopen("input.txt","r",stdin);
    // freopen("output.txt","w",stdout);
    // #endif
    IO
    int t=1;
    // cin>>t;
    while(t--) {sol();}
    // auto stop = chrono::high_resolution_clock::now()
    ;
    // auto duration = chrono::duration_cast<chrono::
    milliseconds>(stop - start);
    // cerr<<"Time:"<<duration.count()<<" ms\n";
}
```

## 1.2 Misc

```
iota(vec.begin(),vec.end(),1);// 產生1~size的整數列
stoi(s.begin(),s.end(),k);// 法1,字串轉成k進位int
string s;cin>>s;
int x=stoi(s,0,2); // 法2,2可以改其他進位
int bbb = bitset<10>(bb).to_ulong();//二進位轉十進位
__builtin_popcountll // 二進位有幾個1
__builtin_clzll // 左起第一個1前0的個數
__builtin_parityll // 1的個數的奇偶性
__builtin_mul_overflow(a,b,&res) // a*b是否溢位

// double 轉整數 請加 int b=round(a)
// 或是 int b=floor(a+0.5) (floor向下取整)
```

## 1.3 Fast read & write

```
inline int read() {
    char c = getchar(); int x = 0, f = 1;
```

```
while(c < '0' || c > '9') {if(c == '-') f = -1; c =
    getchar();}
while(c >= '0' && c <= '9') x = x * 10 + c - '0', c
    = getchar();
return x * f;
}
inline void write(int x){
    if(x<0) putchar('-'),x=-x;
    if(x>9) write(x/10);
    putchar(x%10+'0');
```

## 1.4 Sort cmp

```
struct cmp{inline bool operator()(const int a,const int
    b){return a<b;}};//common use
auto cmp=[](vector<int> a, vector<int> b) {return a[1]<
    b[1];};//for set use
set<vector<int>, decltype(cmp)> prepare, done;
```

## 1.5 Discretization

```
vector<int> vec;
sort(vec.begin(),vec.end());
vec.resize(unique(vec.begin(),vec.end())-vec.begin());
for(int i=0;i<n;++i){//+1是讓 index是1到N 可以不要
    arr[i]=lower_bound(vec.begin(),vec.end(),ll[i])-vec
        .begin()+1;
}
```

## 1.6 Custom unordered\_map

```
struct Type{
    int x;
    string y;
    bool operator==(const Type &other) const {
        return (x == other.x && y == other.y);
    }
};
struct hashes{
    size_t operator()(const Type &o) const {
        return ((hash<int>()(o.x)^(hash<string>()(o.y)
            <<1))>>1);
    }
};
//unordered_map<Type,int,hashes> map;
```

## 1.7 \_\_int128 read

```
// __int128_t p;
// lll n=qr(p);
#define lll __int128
template<class type_name> inline type_name qr(type_name
    sample)
{
    type_name ret=0,sgn=1;
    char cur=getchar();
    while(!isdigit(cur))
        sgn=(cur=='-'?-1:1),cur=getchar();
    while(isdigit(cur))
        ret=(ret<<1)+(ret<3)+cur-'0',cur=getchar();
    return sgn==-1?-ret:ret;
}
inline void print(__int128 x){
    if(x < 0){
        putchar('-');
        x = -x;
    }
    if(x > 9)
        print(x / 10);
    putchar(x % 10 + '0');
```

## 1.8 字典序 a 嚴格小於 b

```
template<class T> //字典序a嚴格小於b
bool lexicographicallySmaller(const vector<T> &a, const
vector<T> &b){
    int n=a.size();
    int m=b.size();
    int i;
    for(int i=0; i<n && i<m; ++i){
        if(a[i]<b[i]) return true;
        else if(b[i]<a[i]) return false;
    }
    return (i==n && i<m);
}
```

## 1.9 生成 n 位數的二進制組合

```
//產生1~n位數的二進位組合
vector<int> comb;
void genBinComb(int n){
    for(int i=1; i<=(1<n); ++i){
        string s = bitset<5>(i).to_string(); //bitset大
        //小要記得設
        comb.push_back(stoi(s));
    }
}
```

## 1.10 Radom

```
mt19937 gen(0x5EED);
int randint(int lb, int ub)
{ return uniform_int_distribution<int>(lb, ub)(gen); }
```

## 2 對拍

### 2.1 run.bat

```
@echo off
g++ ac.cpp -o ac.exe
g++ wa.cpp -o wa.exe
g++ gen1.cpp -o gen.exe

:loop
    echo %%x
    gen.exe > input
    ac.exe < input > ac
    wa.exe < input > wa
    fc ac wa
    if not errorlevel 1 goto loop
```

### 2.2 run.sh

```
for ((i=0;;i++))
do
    echo "$i"
    python3 gen.py > input
    ./ac < input > ac.out
    ./wa < input > wa.out
    diff ac.out wa.out || break
done
```

## 3 Flow & Matching

### 3.1 Dinic

```
// flow.init(n,s,t):有n個點(0~n-1)，起點s終點t
// flow.add_edge(u,v,f):建一條邊，從u點到v點流量為f
// flow.solve():回傳網路最大流答案
//時間複雜度:  $O(V^2 * E)$ 
struct Dinic{
    struct Edge{ int v,f,re; };
    int n,s,t,level[MXN];
    vector<Edge> E[MXN];
```

```
void init(int _n, int _s, int _t){
    n = _n; s = _s; t = _t;
    for (int i=0; i<n; i++) E[i].clear();
}
void add_edge(int u, int v, int f){
    E[u].push_back({v,f,(int)(E[v].size())});
    E[v].push_back({u,0,(int)(E[u].size()-1)});
}
bool BFS(){
    for (int i=0; i<n; i++) level[i] = -1;
    queue<int> que;
    que.push(s);
    level[s] = 0;
    while (!que.empty()){
        int u = que.front(); que.pop();
        for (auto it : E[u]){
            if (it.f > 0 && level[it.v] == -1){
                level[it.v] = level[u]+1;
                que.push(it.v);
            }
        }
    }
    return level[t] != -1;
}
int DFS(int u, int nf){
    if (u == t) return nf;
    int res = 0;
    for (auto &it : E[u]){
        if (it.f > 0 && level[it.v] == level[u]+1){
            int tf = DFS(it.v, min(nf,it.f));
            res += tf; nf -= tf; it.f -= tf;
            E[it.v][it.re].f += tf;
            if (nf == 0) return res;
        }
    }
    if (!res) level[u] = -1;
    return res;
}
int solve(int res=0){
    while ( BFS() )
        res += DFS(s,2147483647);
    return res;
} }
```

### 3.2 最大流最小花費

```
//最大流量上的最小花費
//最大流量優先，相同才是找最小花費，複雜度 $O(V^2 * E^2)$ 
// flow.init(n,s,t):有n個點(0~n-1)，起點s終點t
// flow.add_edge(u,v,f,c):建一條邊，從u點到v點流量為f，
// 每一單位流量的花費為c
// flow.solve():回傳一個pair(maxFlow,minCost)
// 限制：圖不能有負環
// 網路最大流的add_edge(u,v,f)可以無痛轉成最大流量上的
// 最小花費add_edge(u,v,1,f)即建立一條從u到v的邊流量為
// 1，單位流量花費為f
//  $O(V^2 * E^2)$ 
#define ll long long
struct zkwflow{
    static const int maxN=20000;
    struct Edge{ int v,f,re; ll w; };
    int n,s,t,ptr[maxN]; bool vis[maxN]; ll dis[maxN];
    vector<Edge> E[maxN];
    void init(int _n,int _s,int _t){
        n=_n,s=_s,t=_t;
        for(int i=0; i<n; i++) E[i].clear();
    }
    void add_edge(int u,int v,int f,ll w){
        E[u].push_back({v,f,(int)E[v].size(),w});
        E[v].push_back({u,0,(int)E[u].size()-1,-w});
    }
    bool SPFA() {
        fill_n(dis, n, LLONG_MAX);
        fill_n(vis, n, false);
        queue<int> q;
        q.push(s); dis[s]=0;
        while(!q.empty()) {
            int u = q.front(); q.pop();
            vis[u] = false;
            for(auto &it:E[u]){
                if(it.f>0 && dis[it.v]>dis[u]+it.w){
                    dis[it.v] = dis[u]+it.w;
```

```

        if(!vis[it.v]) {vis[it.v] = true; q
            .push(it.v);}
    }
}
if(dis[t]==LLONG_MAX) return false;
// 不管流量是多少，花費不能是正數時加上這行（最
// 小花費可行流）
// if(dis[t] >= 0) return false;
return true;
}
int DFS(int u, int nf) {
    if(u==t) return nf;
    int res = 0; vis[u] = true;
    for(int &i=ptr[u] ; i<(int)E[u].size() ; i++) {
        auto &it = E[u][i];
        if(it.f>0 && dis[it.v]==dis[u]+it.w && !vis
            [it.v]) {
            int tf = DFS(it.v, min(nf, it.f));
            res += tf;
            nf-=tf;
            it.f-=tf;
            E[it.v][it.re].f += tf;
            if(nf==0) { vis[u]=false; break; }
        }
    }
    return res;
}
pair<int,ll> solve(){
    int flow = 0; ll cost = 0;
    while (SPFA()){
        fill_n(ptr, n, 0);
        int f = DFS(s, INT_MAX);
        flow += f;
        cost += dis[t]*f;
    }
    return {flow, cost};
} // reset: do nothing
} flow;

```

### 3.3 匈牙利匹配

```

//匈牙利演算法-二分圖最大匹配
//記得每次使用需清空vis數組
//O(nm)
//其中Map為鄰接表(Map[u][v]為u和v是否有連接) S為紀錄這
//個點與誰匹配(S[i]為答案i和誰匹配)
const int M=505, N=505;
bool Map[M][N] = {0};
int S[N];
bool vis[N];
bool dfs(int u){
    for(int i=0;i<N;i++){
        if(Map[u][i]&&!vis[i]){ //有連通且未拜訪
            vis[i]=1; //紀錄是否走過
            if(S[i]==-1||dfs(S[i])){ //紀錄匹配
                S[i]=u;
                return true; //反轉匹配邊以及未匹配邊
                的狀態
            }
        }
    }
    return false;
}
//此二分圖為左邊M個點右邊N個點，跑匈牙利只要跑1~M就可以
//了，(S[右邊的點] -> 左邊的點)
memset(S,-1,sizeof(S));
int ans = 0;
for(int i=0;i<M;i++){
    memset(vis,0,sizeof(vis));
    if(dfs(i)) ans++;
    //跑匈牙利
}
cout<<ans<<"\n";
for(int i=0 ; i<N ; i++) {
    if(S[i]!=-1) cout<<"pair: "<<S[i]<<" "<<i<<"\n";
}

```

### 3.4 KM

```

//二分圖最大權完美匹配
//二分圖左邊的點都要匹配到右邊的點，且每條邊都有權重，
//求權重最大值，複雜度O(V^3)
// graph.init(n):二分圖左右各n個點
// graph.add_edge(u,v,w):建一條邊，從u點到v點權重為w
// graph.solve():回傳最大權重
struct KM{ // max weight, for min negate the weights
    int n, mx[MXN], my[MXN], pa[MXN];
    ll g[MXN][MXN], lx[MXN], ly[MXN], sy[MXN];
    bool vx[MXN], vy[MXN];
    void init(int _n) { // 1-based, N個節點
        n = _n;
        for(int i=1; i<=n; i++) fill(g[i], g[i]+n+1, 0)
            ;
    }
    void add_edge(int x, int y, ll w) {g[x][y] = w;} //
    // 左邊的集合節點x連邊右邊集合節點y權重為w
    void augment(int y) {
        for(int x, z; y; y = z)
            x=pa[y], z=mx[x], my[y]=x, mx[x]=y;
    }
    void bfs(int st) {
        for(int i=1; i<=n; ++i) sy[i]=INF, vx[i]=vy[i]
            =0;
        queue<int> q; q.push(st);
        for(;;) {
            while(q.size()) {
                int x=q.front(); q.pop(); vx[x]=1;
                for(int y=1; y<=n; ++y) if(!vy[y]){
                    ll t = lx[x]+ly[y]-g[x][y];
                    if(t==0){
                        pa[y]=x;
                        if(!my[y]){augment(y);return;}
                        vy[y]=1, q.push(my[y]);
                    }else if(sy[y]>t) pa[y]=x,sy[y]=t;
                }
            }
            ll cut = INF;
            for(int y=1; y<=n; ++y)
                if(!vy[y]&&cut>sy[y]) cut=sy[y];
            for(int j=1; j<=n; ++j){
                if(vx[j]) lx[j] -= cut;
                if(vy[j]) ly[j] += cut;
                else sy[j] -= cut;
            }
            for(int y=1; y<=n; ++y) if(!vy[y]&&sy[y]
                ]==0){
                if(!my[y]){augment(y);return;}
                vy[y]=1, q.push(my[y]);
            }
        }
    }
    ll solve(){ // 回傳值為完美匹配下的最大總權重
        fill(mx, mx+n+1, 0); fill(my, my+n+1, 0);
        fill(ly, ly+n+1, 0); fill(lx, lx+n+1, -INF);
        for(int x=1; x<=n; ++x) for(int y=1; y<=n; ++y)
            // 1-base
            lx[x] = max(lx[x], g[x][y]);
        for(int x=1; x<=n; ++x) bfs(x);
        ll ans = 0;
        for(int y=1; y<=n; ++y) ans += g[my[y]][y];
        return ans;
    }
} graph;

```

## 4 Graph

### 4.1 Dijkstra

```

int dis[N],vis[N];
void dijkstra(int s){//O(V^2+E)
    memset(dis,0x3f,sizeof(dis));
    memset(vis,0,sizeof(vis));
    dis[s] = 0;
    priority_queue<pii,vector<pii>,greater<pii>> pq;
    pq.push({0,s});
    while(!pq.empty()){
        auto [nowd,now] = pq.top();

```

```

pq.pop();
if(vis[now])
    continue;
vis[now] = 1;
for(auto [nxt,w]: graph[now]){
    if(dis[nxt] > dis[now] + w){
        dis[nxt] = dis[now] + w;
        pq.push({dis[nxt],nxt});
    }
}
}
}

```

## 4.2 Bellman-Ford

```
//總共m條邊，鬆弛n-1次->O(nm)
//在第n次做一次鬆弛，如果有點被鬆弛到，代表這張圖存在負環
for(int j = 0; j < n-1; j++){
    for(int i = 0; i < m; i++){ // 對於所有邊都嘗試鬆弛
        if(dis[ edge[i].to ] > dis[ edge[i].from ] +
            edge[i].weight){
            dis[ edge[i].to ] = dis[ edge[i].from ] +
                edge[i].weight;
        }
    }
}
```

### 4.3 SPFA

```
#define MXN 200005
struct SPFA{//O(kE) k:小常數
    int n;
    ll inq[MXN], len[MXN];
    vector<ll> dis;
    vector<pair<int, ll>> edge[MXN];
    void init(int _n){
        n = _n;
        dis.clear(); dis.resize(n, 1e18);
        for(int i = 0; i < n; i++){
            edge[i].clear();
            inq[i] = len[i] = 0;
        }
    }
    void addEdge(int u, int v, ll w){
        edge[u].push_back({v, w});
    }
    vector<ll> solve(int st = 0){
        deque<int> dq; //return {-1} if has negative cycle
        dq.push_back(st); //otherwise return dis from st
        inq[st] = 1; dis[st] = 0;
        while(!dq.empty()){
            int u = dq.front(); dq.pop_front();
            inq[u] = 0;
            for(auto [to, d] : edge[u]){
                if(dis[to] > d+dis[u]){
                    dis[to] = d+dis[u];
                    len[to] = len[u]+1;
                    if(len[to] > n) return {-1};
                    if(inq[to]) continue;
                    (!dq.empty()&&dis[dq.front()] > dis[to]?
                     dq.push_front(to) : dq.push_back(to));
                    inq[to] = 1;
                }
            }
        }
        return dis;
    }
} spfa;
```

## 4.4 Floyd-Warshall

```
for(int k=0;k<n;++k){//O(N^3)
    for(int i=0;i<n;++i){
        for(int j=0;j<n;++j){
            if(graph[i][k]!=MAX&&graph[k][j]!=MAX)//避免不連通圖
                graph[i][j]=min(graph[i][j],graph[i][k]+graph[k][j]);
        }
    }
}
```

```

    }
}

for(int k=0;k<n;++k){//判斷負環
    for(int i=0;i<n;++i){
        for(int j=0;j<n;++j){
            if(graph[i][k]!=MAX&&graph[k][j]!=MAX&&
                graph[k][k]<0){//避免不連通圖&&負環
                graph[i][j]=-MAX;
            }
        }
    }
}

// if(graph[a][b]==-MAX)
//     cout<<"-Infinity\n";
// else if(graph[a][b]==MAX)
//     cout<<"Impossible\n";
// else
//     cout<<graph[a][b]<<"\n";

```

## 4.5 歐拉路徑

```

const int MXN = 2e5+7;
struct ola{
    vector<pair<int, int>> edge[MXN];
    int ind[MXN], ru[MXN], use[MXN*3], es=0, n, go;
    void init(int _n) {
        n = _n;
        for(int i=0 ; i<n ; i++) edge[i].clear(), ind[i]=0;
    }
    void add_edge(int a, int b) {
        use[es] = 0;
        edge[a].push_back({b, es});
        edge[b].push_back({a, es++});
        ru[a]++; ru[b]++;
    }
    bool check() {
        int cnt = 0;
        for(int i=0 ; i<n ; i++) {
            if(ru[i]&1) go = i, cnt++;
        }
        if(cnt==0 || cnt==2) return true;
        return false;
    }
    vector<int> ans;
    void dfs(int x) {
        for(int i=ind[x] ; i<edge[x].size() ; i=ind[x]) {
            int u = edge[x][i].first, pos = edge[x][i].second;
            ind[x]++;
            if(!use[pos]) {
                use[pos]=1;
                dfs(u);
            }
        }
        ans.push_back(x);
    }
    vector<int> solve() {
        if(!check()) return {-1};
        ans.clear(); dfs(go);
        if(ans.size()!=es+1) return {-1};
        return ans;
    }
} euler;

```

#### 4.6 BCC

```
//無向圖上，不會產生割點的連通分量稱為點雙連通分量，
//base
#define PB push_back
#define REP(i, n) for(int i = 0; i < n; i++)
struct BccVertex {
    int n, nScc, step, dfn[MXN], low[MXN];
    vector<int> E[MXN], sccv[MXN];
    int top, stk[MXN];
    void init(int _n) {
        n = _n;
```

```

nScc = step = 0;
for (int i = 0; i < n; i++)
    E[i].clear();
}
void addEdge(int u, int v) {
    E[u].PB(v); E[v].PB(u);
}
void DFS(int u, int f) {
    dfn[u] = low[u] = step++;
    stk[top++] = u;
    for (auto v : E[u]) {
        if (v == f) continue;
        if (dfn[v] == -1) {
            DFS(v, u);
            low[u] = min(low[u], low[v]);
            if (low[v] >= dfn[u]) {
                int z;
                sccv[nScc].clear();
                do {
                    z = stk[--top];
                    sccv[nScc].PB(z);
                } while (z != v);
                sccv[nScc++].PB(u);
            }
        }
        else low[u] = min(low[u], dfn[v]);
    }
}
vector<vector<int>> solve() { //回傳每個點雙聯通分量
    vector<vector<int>> res;
    for (int i = 0; i < n; i++)
        dfn[i] = low[i] = -1;
    for (int i = 0; i < n; i++)
        if (dfn[i] == -1) {
            top = 0;
            DFS(i, i);
        }
    REP(i, nScc) res.PB(sccv[i]);
    return res;
}
} graph;

```

```

        if (!vst[i]) DFS(i);
        reverse(vec.begin(), vec.end());
        FZ(vst);
        for (auto v : vec)
            if (!vst[v]) {rDFS(v); nScc++;}
    }
} scc;

```

## 4.8 2SAT

[illegible]

#### 4.7 SCC

```
//在有向圖裡的任兩點u、v，皆存在至少一條u到v的路徑
//以及v到u的路徑
//fill zero 注意多筆測資要改fill
//注意要0base
#define PB push_back
#define FZ(x) memset(x, 0, sizeof(x))
const int MXN = 1e5;
struct Scc {
    int n, nScc, vst[MXN], bln[MXN]; //nScc 有幾個強連通分量
    vector<int> E[MXN], rE[MXN], vec;
    void init(int _n) {
        n = _n;
        for (int i = 0; i < MXN; i++)
            E[i].clear(), rE[i].clear();
    }
    void addEdge(int u, int v) {
        E[u].PB(v); rE[v].PB(u);
    }
    void DFS(int u) {
        vst[u] = 1;
        for (auto v : E[u])
            if (!vst[v]) DFS(v);
        vec.PB(u);
    }
    void rDFS(int u) {
        vst[u] = 1;
        bln[u] = nScc;
        for (auto v : rE[u])
            if (!vst[v]) rDFS(v);
    }
    void solve() {
        nScc = 0;
        vec.clear();
        FZ(vst);
        for (int i = 0; i < n; i++)
```

## 4.9 MaximalClique

```
//極大團
//對於一張圖選任意的點子集，如果不能在多選一個點使得選
//的點子集為更大的團
#define N 80
struct MaxClique{ // 0-base
    typedef bitset<N> Int;
    Int lnk[N] , v[N];
    int n;
    void init(int _n){
        n = _n;
        for(int i = 0 ; i < n ; i ++){
            lnk[i].reset(); v[i].reset();
        }
    }
    void addEdge(int a , int b)
    { v[a][b] = v[b][a] = 1; }
    int ans , stk[N], id[N] , di[N] , deg[N];
    Int cans;
    void dfs(int elem_num, Int candi, Int ex){
        if(candi.none() && ex.none()){
            cans.reset();
            for(int i = 0 ; i < elem_num ; i ++){
                cans[id[stk[i]]] = 1;
            }
            ans = elem_num; //cans=1 is in maximal clique
            return;
        }
        int pivot = (candi|ex)._Find_first();
        Int smaller_candi = candi & (~lnk[pivot]);
        while(smaller_candi.count()){
            int nxt = smaller_candi._Find_first();
            candi[nxt] = smaller_candi[nxt] = 0;
            ex[nxt] = 1;
            stk[elem_num] = nxt;
            dfs(elem_num+1, candi&lnk[nxt], ex&lnk[nxt]);
        }
    }
    int solve(){
        for(int i = 0 ; i < n ; i ++){

```



```

    id[i] = i; deg[i] = v[i].count();
}
sort(id , id + n , [&](int id1, int id2){
    return deg[id1] > deg[id2]; });
for(int i = 0 ; i < n ; i ++ ) di[id[i]] = i;
for(int i = 0 ; i < n ; i ++ )
    for(int j = 0 ; j < n ; j ++ )
        if(v[i][j]) lnk[di[i]][di[j]] = 1;
ans = 1; cans.reset(); cans[0] = 1;
dfs(0, Int(string(n, '1')), 0);
return ans;
} }solver;

```

#### 4.10 MaximumClique

```

//最大團:圖上最多可以選幾個點, 使選的彼此之間都有連邊
//最大獨立集:圖上最多可以選幾個點, 使選的彼此之間都沒有連邊
//最大獨立集通常會轉換為用補圖做最大團
//O(1.1888^n)
#define N 111
struct MaxClique{ // 0-base
    typedef bitset<N> Int;
    Int linkto[N] , v[N];
    int n;
    void init(int _n){
        n = _n;
        for(int i = 0 ; i < n ; i ++ ){
            linkto[i].reset(); v[i].reset();
        }
    }
    void addEdge(int a , int b)
    { v[a][b] = v[b][a] = 1; }
    int popcount(const Int& val)
    { return val.count(); }
    int lowbit(const Int& val)
    { return val._Find_first(); }
    int ans , stk[N];
    int id[N] , di[N] , deg[N];
    Int cans;
    void maxclique(int elem_num, Int candi){
        if(elem_num > ans){
            ans = elem_num; cans.reset();
            for(int i = 0 ; i < elem_num ; i ++ )
                cans[id[stk[i]]] = 1;
        }
        int potential = elem_num + popcount(candi);
        if(potential <= ans) return;
        int pivot = lowbit(candi);
        Int smaller_candi = candi & (~linkto[pivot]);
        while(smaller_candi.count() && potential > ans){
            int next = lowbit(smaller_candi);
            candi[next] = !candi[next];
            smaller_candi[next] = !smaller_candi[next];
            potential --;
            if(next == pivot || (smaller_candi & linkto[next]).count()){
                stk[elem_num] = next;
                maxclique(elem_num + 1, candi & linkto[next]);
            }
        }
    }
    int solve(){//回傳值為最大團的點數量
        for(int i = 0 ; i < n ; i ++ ){
            id[i] = i; deg[i] = v[i].count();
        }
        sort(id , id + n , [&](int id1, int id2){
            return deg[id1] > deg[id2]; });
        for(int i = 0 ; i < n ; i ++ ) di[id[i]] = i;
        for(int i = 0 ; i < n ; i ++ )
            for(int j = 0 ; j < n ; j ++ )
                if(v[i][j]) linkto[di[i]][di[j]] = 1;
        Int cand; cand.reset();
        for(int i = 0 ; i < n ; i ++ ) cand[i] = 1;
        ans = 1;
        cans.reset(); cans[0] = 1;
        maxclique(0, cand);
        return ans;
    } }solver;

```

#### 4.11 Minimum Mean Cycle

```

//給定一張有向圖, 邊上有權重, 要找到一個環其平均權重最小
/* minimum mean cycle O(VE) */
struct MMC{
#define E 101010
#define V 1021
#define inf 1e9
#define eps 1e-6
    struct Edge { int v,u; double c; };
    int n, m, prv[V][V], prve[V][V], vst[V];
    Edge e[E];
    vector<int> edgeID, cycle, rho;
    double d[V][V];
    void init( int _n )
    { n = _n; m = 0; }
    // WARNING: TYPE matters
    //建一條單向邊 (u, v) 權重為 w
    void addEdge( int vi , int ui , double ci )
    { e[ m ++ ] = { vi , ui , ci }; }
    void bellman_ford() {
        for(int i=0; i<n; i++) d[0][i]=0;
        for(int i=0; i<n; i++) {
            fill(d[i+1], d[i+1]+n, inf);
            for(int j=0; j<m; j++) {
                int v = e[j].v, u = e[j].u;
                if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
                    d[i+1][u] = d[i][v]+e[j].c;
                    prv[i+1][u] = v;
                    prve[i+1][u] = j;
                }
            }
        }
    }
    double solve(){//回傳值為最小平均權重 (小數)
        // returns inf if no cycle, mmc otherwise
        double mmc=inf;
        int st = -1;
        bellman_ford();
        for(int i=0; i<n; i++) {
            double avg=-inf;
            for(int k=0; k<n; k++) {
                if(d[n][i]<inf-eps) avg=max(avg, (d[n][i]-d[k][i])/(n-k));
                else avg=max(avg, inf);
            }
            if (avg < mmc) tie(mmc, st) = tie(avg, i);
        }
        fill(vst,0); edgeID.clear(); cycle.clear(); rho.clear();
        for (int i=n; !vst[st]; st=prv[i--][st]) {
            vst[st]++;
            edgeID.PB(prve[i][st]);
            rho.PB(st);
        }
        while (vst[st] != 2) {
            if(rho.empty()) return inf;
            int v = rho.back(); rho.pop_back();
            cycle.PB(v);
            vst[v]++;
        }
        reverse(ALL(edgeID));
        edgeID.resize(SZ(cycle));
        return mmc;
    } }mmc;

```

#### 4.12 Dominator Tree

```

// 給一張有向圖, 圖上有一個起點 s 可以走到所有點。
// 定義 "支配" 為從起點 s 出發, 所有能走到節點 x 的路徑
// 的最後一個必經點
// 最後 idom[i] 為點 i 的支配點
struct DominatorTree{ // O(n+m)
#define REP(i,s,e) for(int i=(s);i<=(e);i++)
#define REPD(i,s,e) for(int i=(s);i>=(e);i--)
    int n , s;
    vector< int > g[ MAXN ] , pred[ MAXN ];
    vector< int > cov[ MAXN ];
    int dfn[ MAXN ] , nfd[ MAXN ] , ts;
    int par[ MAXN ]; //idom[u] s到u的最後一個必經點
    int sdom[ MAXN ] , idom[ MAXN ];
    int mom[ MAXN ] , mn[ MAXN ];
    inline bool cmp( int u , int v )
    { return dfn[ u ] < dfn[ v ]; }
};

```

```

int eval( int u ){
    if( mom[ u ] == u ) return u;
    int res = eval( mom[ u ] );
    if(cmp( sdom[ mn[ mom[ u ] ] ] , sdom[ mn[ u ] ] ))
        mn[ u ] = mn[ mom[ u ] ];
    return mom[ u ] = res;
}
//節點數量，起點編號 1-base
void init( int _n , int _s ){
    ts = 0; n = _n; s = _s;
    REP( i , 1 , n ) g[ i ].clear(), pred[ i ].clear();
}
void addEdge( int u , int v ){
    g[ u ].push_back( v );
    pred[ v ].push_back( u );
}
void dfs( int u ){
    ts++;
    dfn[ u ] = ts;
    nfd[ ts ] = u;
    for( int v : g[ u ] ) if( dfn[ v ] == 0 ){
        par[ v ] = u;
        dfs( v );
    }
}
void build(){// 建立支配樹
    REP( i , 1 , n ){
        dfn[ i ] = nfd[ i ] = 0;
        cov[ i ].clear();
        mom[ i ] = mn[ i ] = sdom[ i ] = i;
    }
    dfs( s );
    REPD( i , n , 2 ){
        int u = nfd[ i ];
        if( u == 0 ) continue;
        for( int v : pred[ u ] ) if( dfn[ v ] ){
            eval( v );
            if( cmp( sdom[ mn[ v ] ] , sdom[ u ] ) )
                sdom[ u ] = sdom[ mn[ v ] ];
        }
        cov[ sdom[ u ] ].push_back( u );
        mom[ u ] = par[ u ];
        for( int w : cov[ par[ u ] ] ){
            eval( w );
            if( cmp( sdom[ mn[ w ] ] , par[ u ] ) )
                idom[ w ] = mn[ w ];
            else idom[ w ] = par[ u ];
        }
        cov[ par[ u ] ].clear();
    }
    REP( i , 2 , n ){
        int u = nfd[ i ];
        if( u == 0 ) continue;
        if( idom[ u ] != sdom[ u ] )
            idom[ u ] = idom[ idom[ u ] ];
    }
} domT;

```

#### 4.13 ManhattanMST

```

// 出來的是保證能做出MST的邊 邊是亂的 ~~魔術師~~
// 需要swap(u.first.first,u.first.second) + sort unique
#define ld long long
struct Pt{
    ld x,y;
    Pt(ld x=0,ld y=0):x(x),y(y){}
    Pt operator+(const Pt &a) const {return Pt(x+a.x, y+a.y); }
    Pt operator-(const Pt &a) const {return Pt(x-a.x, y-a.y); }
};
vector<pair<pair<int,int>,int>> ManhattanMST(vector<Pt> P) {
    vector<int> id(P.size());
    iota(id.begin(),id.end(), 0);
    vector<pair<pair<int,int>, int>> edg;
    for( int k = 0; k < 4; k++ ) {
        sort(id.begin(),id.end(), [&](int i, int j) {
            return (P[i] - P[j]).x < (P[j] - P[i]).y;
        });
        map<int, int> sweep;
        for( int i : id ) {

```

```

            auto it = sweep.lower_bound(-P[i].y);
            while( it != sweep.end() ) {
                int j = it->second;
                Pt d = P[i] - P[j];
                if( d.y > d.x ) {
                    break;
                }
                edg.push_back({{i, j},d.x + d.y});
                it = sweep.erase(it);
            }
            sweep[-P[i].y] = i;
        }
        for( Pt &p : P ) {
            if( k % 2 ) {
                p.x = -p.x;
            } else {
                swap(p.x, p.y);
            }
        }
    }
    return edg;
}

```

## 5 DP

### 5.1 數位 DP

```

// dp[位數][狀態]
// dp[pos][state]: 定義為目前位數在前導狀態為state的時
// 候的計數
// ex: 求數字沒有出現66的數量 L~r
// -> dp[pos][1] 可表示計算pos個位數在前導出現一個6的計
// 數 -> dp[3][1] 則計算 6XXX
// 模板的pos是反過來的，但不影響(只是用來dp記憶用)

// pos: 目前位數
// state: 前導狀態
// lead: 是否有前導0 (大部分題目不用但有些數字EX:00146
// 如果有影響時要考慮)
// limit: 是否窮舉有被num限制
vector<int> num;
int dp[20][state];
int dfs(int pos, int state, bool lead, bool limit) {
    if(pos==num.size()) {
        //有時要根據不同state回傳情況
        return 1;
    }
    if(limit==false && lead==false && dp[pos][state]
        !=-1) return dp[pos][state];
    int up = limit?num[pos]:9;
    int ans = 0;
    for(int i=0 ; i<=up ; i++) {
        //有時要考慮那些狀況要continue
        ans += dfs(pos+1, state|((check[i]==2), lead&&i
            ==0, limit&&i==num[pos]));
    }
    if(limit==false && lead==false) dp[pos][state] =
        ans;
    return ans;
}

```

### 5.2 SOS DP

```

for( int mask = 0; mask < (1 << N); mask++)
    F[mask] = A[mask];

for( int i = 0; i < N; i++)
    for( int mask = 0; mask < (1 << N); mask++)
        if( mask & (1 << i))
            F[mask] += F[mask ^ (1 << i)];

```

## 6 Math

### 6.1 Formulas



```
//五次方幂次和
a(n) = n^2*(n+1)^2*(2*n^2+2*n-1)/12
//四次方幂次和
a(n) = n*(n+1)*(2n+1)*(3n^3+3n-1)/30
```

## 6.2 l1laddmul

```
#define ll long long
#define l1l __int128
ll mypow(l1l n, l1l k, ll p){
    ll res=1;
    for(;k;k>>=1,n=n*n%p)if(k&1)res=res*n%p;
    return res;
}
ll mul(ll a, ll b, ll mod){
    l1l c=a%mod;
    return c*b%mod;
}
ll add(ll x, ll y, ll mod){
    l1l c=x+c+y;
    return c%mod;
}
```

## 6.3 Primes

```
mashu lucky prime : 91145149
1097774749, 1076767633, 100102021, 999997771
1001010013, 1000512343, 987654361, 999991231
999888733, 98789101, 987777733, 999991921, 1010101333
```

## 6.4 Coprime (互質 Pair)

```
// dp[i] 為認挑2個數字，最大公因數為i的組合數
const int mxn = 1e6+7;
int cnt[mxn], dp[mxn];
int sol(vector<int> vec) {
    for(int u: vec) cnt[u]++;
    for(int i=mxn-1; i>=1; i--) {
        int a=0, b=0;
        for(int j=i; j<mxn; j+=i) {
            a+=cnt[j]; b+=dp[j];
        }
        dp[i] = (a*(a-1)/2) - b;
    }
    return dp[1];
}
```

## 6.5 取樣定理

## 6.6 Quick Pow

```
// a^b
const int MOD = 1e9+7;
int qpow(int n, int k, int p) {
    int ret = 1;
    for(;k;k>>=1,n=n*n%p)if(k&1)ret=ret*n%p;
    return ret;
}
// a^(b^c) = a^(q*(p-1)+r) = a^r so let b^c mod p-1
bc=qpow(b,c,p-1);
ans=qpow(a,bc,p);
```

## 6.7 Mat quick Pow

```
struct mat{
    long long a[200][200],r,c; //resize
    mat(int _r,int _c){r=_r;c=_c;memset(a,0,sizeof(a));}
    void build(){for(int i=0;i<r;++i)a[i][i]=1;}
};
mat operator * (mat &x,mat &y){
    mat z(x.r,y.c);
    for(int i=0;i<x.r;++i)for(int j=0;j<y.c;++j)for(int k=0;k<y.c;++k)
        z.a[i][j]=(z.a[i][j]+x.a[i][k]*y.a[k][j]%MOD)%MOD;
    return z;
}
mat qpow(mat a,int k){
    mat r(a.r,a.r);r.build();while(k){if(k&1)r=r*a;a=a*a;k>>=1;}return r;
}
```

## 6.8 Primes Table

```
int np[MXN];
vector<int> vec;
void sol(){
    np[0]=np[1]=1;
    for(int i=2;i<MXN;++i){
        if(!np[i]){
            for(int j=i;j<MXN;j+=i){
                np[j]=1;
            }
            vec.push_back(i);
        }
    }
}
```

## 6.9 Phi 函數

```
// 計算小於n的數中與n互質的有幾個
// O(sqrtN)
// a^b mod c = a^(b%phi(c) + phi(c)) mod c : if b>= phi(c)
// = a^b mod c : if b< phi(c)
int phi(int n){
    int res = n, a=n;
    for(int i=2;i*i<=a;i++){
        if(a%i==0){
            res = res/i*(i-1);
            while(a%i==0) a/=i;
        }
    }
    if(a>1) res = res/a*(a-1);
    return res;
}
```

## 6.10 Factor Table

```
const int MXN = 1e7+7; //if>1e7 TLE
int np[MXN], fac[MXN], num[MXN];
// isprime, 最大質因數, 質因數數量
void table(){
    np[1]=1;
    for(int i=2;i<MXN;++i){
        if(np[i])continue;
        for(int j=i;j<MXN;j+=i){
            if(i!=j)np[j]=1;
            fac[j]=i;
            num[j]++;
        }
    }
}
//質因數分解
vector<int> res;
void div(int x){
    for(;x>1;x/=fac[x])res.push_back(fac[x]);
}
```

## 6.11 卡塔蘭數

```
// O(N), 要記得開Long Long 跟設定 MOD
cat[0]=1; cat[1]=1;
for(ll i=1 ; i<N ; i++) {
    cat[i] = cat[i-1]*(i*4-2)%MOD*qpow(i+1, MOD-2)%MOD;
}
```

## 6.12 Miller Rabin

```
// n < 4,759,123,141      3 : 2, 7, 61
// n < 1,122,004,669,633  4 : 2, 13, 23, 1662803
// n < 3,474,749,660,383  6 : pimes <= 13
// n < 2^64              7 :
// 2, 325, 9375, 28178, 450775, 9780504, 1795265022
// Make sure testing integer is in range [2, n-2] if
// you want to use magic.
LL magic[]={}
bool witness(LL a,LL n,LL u,int t){
    if(!a) return 0;
    LL x=mpow(a,u,n);
    for(int i=0;i<t;i++) {
        LL nx=mul(x,x,n);
        if(nx==1&&x!=1&&x!=n-1) return 1;
        x=nx;
    }
    return x!=1;
}
bool miller_rabin(LL n) {
    int s=(magic number size)
    // iterate s times of witness on n
    if(n<2) return 0;
    if(!(n&1)) return n == 2;
    ll u=n-1; int t=0;
    // n-1 = u*2^t
    while(!(u&1)) u>>=1, t++;
    while(s--){
        LL a=magic[s]%n;
        if(witness(a,n,u,t)) return 0;
    }
    return 1;
}
```

## 6.13 PollarRho

```
// does not work when n is prime O(n^(1/4))
ll f(ll x, ll mod){ return add(mul(x,x,mod),1,mod); }
ll pollard_rho(ll n) {
    if(!(n&1)) return 2;
    while(true){
        ll y=2, x=rand()%(n-1)+1, res=1;
        for(int sz=2; res==1; sz*=2) {
            for(int i=0; i<sz && res<=1; i++) {
                x = f(x, n);
                res = __gcd(abs(x-y), n);
            }
            y = x;
        }
        if (res!=0 && res!=n) return res;
    }
}
// 如果被卡隨機 用下面的
ll f(ll x,ll c,ll mod){return add(mul(x,x,mod),c,mod);}
ll pollard_rho(ll n){
    ll c=1,x=0,y=0,p=2,q,t=0;
    while(t++%128 or __gcd(p,n)==1){
        if(x==y)c++,y=f(x=2,c,n);
        if(q=mul(p,abs(x-y),n))p=q;
        x=f(x,c,n);y=f(y,c,n),c,n);
    }
    return __gcd(p,n);
}
```

## 6.14 PrimeFactorO(logn)

```
vector<ll> ret;
void fac(ll x){
    if(x<2)return;
    if (miller_rabin(x)) {
        ret.push_back(x);
        return;
    }
    ll f = pollard_rho(x);
    fac(f); fac(x/f);
}
```

## 6.15 O(1)mul

```
// < Long Long
ll mul(ll x,ll y,ll mod){
    ll ret=x*y-(ll)((long double)x/mod*y)*mod;
    // LL ret=x*y-(LL)((long double)x*y/mod+0.5)*mod;
    return ret<0?ret+mod:ret;
}
```

## 6.16 Josephus Problem

```
//base1 n people count k find lastone O(n)
int jo(int n, int k){return n>1?(jo(n-1,k)+k-1)%n+1:1;}
//base0 when k<n O(klogn)
int jo(int n, int k) {
    if (n == 1) return 0;
    if (k == 1) return n - 1;
    if (k > n) return (jo(n - 1, k) + k) % n;
    int f = jo(n - n / k, k) - n % k;
    return f + (f < 0 ? n : (f / (k - 1)));
}
//base1 when k=2 fast find mth
int jo2(int n, int m, int f=0){
    if(n == 1) return 1;
    int kill = (n + f) / 2;
    if(m <= kill) return 2 * m - f;
    return 2 * jo2(n - kill, m - kill, (n ^ f) & 1) -
        (1 ^ f);
}
```

## 6.17 Harmonic Sum

```
struct Harmonic{
    const double gamma = 0.5772156649;
    //求第N個調和級數
    double nthHarmonic(int n){
        double result = log(n)+gamma;
        return result;
    }
    //求項數n的Sn>k
    int findNearstN(int k){
        int n = exp(k-gamma)+0.5;
        return n;
    }
    // 16n
    // n/1 + n/2 + n/3 + ... + n/n
    //就是這東西
    [20,10,6,5,4,3,2,2,2,2,1,1,1,1,1,1,1,1]
    //這是N以下的全因數和
    int nthHarmonicSum9(int n){
        int inv2=qpow(2,MOD-2,MOD),ans=0;
        for(int i=1;i<=n;){
            int v = n/i; int j = n/v;
            int area=((j-i+1)%MOD)*((j+i)%MOD)%MOD*
                inv2%MOD; //梯形
            ans=(ans+v*area%MOD)%MOD;
            i=j+1;
        }
        return ans;
    }
};
```

## 7 Data Structure

### 7.1 BIT

```
//注意值域
const int N = 1e5+5;
int bit[N];
struct BIT {
    int n;
    void init(int _n){ n = _n;}
    int low(int x) {return x&-x;}
    void update(int x, int val) {
        while(x<n) bit[x]+=val, x+=low(x);
    }
    int query(int x) {
        int res = 0;
        while(x) res += bit[x], x-=low(x);
        return res;
    }
    int query(int l, int r) {return query(r) - query(l-1);}
```

### 7.2 BIT 二維

```
struct BIT {
    static const int mxn = 2005;
    int bit[mxn][mxn] = {0};
    int low(int x) {return x&-x;}
    void add(int x, int y, int val) {
        for(int i=x ; i<mxn ; i+=low(i)) for(int j=y ; j<mxn ; j+=low(j)) bit[i][j]+=val;
    }
    int query(int x, int y) {
        int ans = 0;
        for(int i=x ; i ; i-=low(i)) for(int j=y ; j ; j-=low(j)) ans+=bit[i][j];
        return ans;
    }
    int range_query(int a, int b, int x, int y) {
        return query(x, y) - query(x, b-1) - query(a-1, y) + query(a-1, b-1);
    }
} bit;
```

### 7.3 稀疏表 $O(1)$ 區間最大最小值

```
//st[i][j]表示[i,i+2^j-1]的最值,區間最大長度為Log2(n)
//i為1base
const int N = 5e4+5;
int stMax[N][20],stMin[N][20],a[N];
struct ST{
    int k;
    void build(int n,int a[]){
        k=log2(n);
        for(int i = 1; i <= n; i++) stMin[i][0] = stMax[i][0] = a[i];
        for(int j = 1; j <= k; j++){
            for(int i = 1; i + (1 << j) - 1 <= n; i++){
                stMax[i][j] = max(stMax[i][j-1], stMax[i + (1 << (j-1))][j-1]);
                stMin[i][j] = min(stMin[i][j-1], stMin[i + (1 << (j-1))][j-1]);
            }
        }
    }
    int queryMax(int l,int r){
        int j = log2(r-l+1);
        return max(stMax[l][j],stMax[r-(1<<j)+1][j]);
    }
    int queryMin(int l,int r){
        int j = log2(r-l+1);
        return min(stMin[l][j],stMin[r-(1<<j)+1][j]);
    }
}st;
```

## 7.4 Segment Tree

```
struct seg {
    #define left (index<<1)
    #define right (index<<1|1)
    static const int MXN = 200005;
    int val[MXN*4], tag[MXN*4];
    int a[MXN];
    void push(int index, int l, int r) {
        if(tag[index]!=0) {
            val[index]+=tag[index]*(r-l+1);
            if(l!=r) {
                tag[left] += tag[index];
                tag[right] += tag[index];
            }
            tag[index]=0;
        }
    }
    void pull(int index, int l, int r) {
        int mid = l+r>>1;
        push(left, l, mid);
        push(right, mid+1, r);
        val[index] = val[left]+val[right];
    }
    void build(int index, int l, int r) {
        if(l==r) {
            val[index] = a[l];
            return;
        }
        int mid = (l+r)>>1;
        build(left, l, mid);
        build(right, mid+1, r);
        pull(index, l, r);
    }
    void add(int index, int s, int e, int l, int r, int v) {
        if(e<l || r<s) return;
        if(l<=s && e<=r) {
            tag[index] += v;
            push(index, s, e);
            return;
        }
        int mid = (s+e)>>1;
        push(index, s, e);
        add(left, s, mid, l, r, v);
        add(right, mid+1, e, l, r, v);
        pull(index, s, e);
    }
    int query(int index, int s, int e, int l, int r) {
        if(e<l || r<s) return 0;
        if(l<=s && e<=r) {
            push(index, s, e);
            return val[index];
        }
        push(index, s, e);
        int mid = (s+e)>>1;
        return query(right, mid+1, e, l, r) + query(left, s, mid, l, r);
    }
} tree;
```

## 7.5 動態開點線段數

```
// tree.init(區間大小 0~n)
// tree.add(ql, qr, val) 區間加值
// tree.query(ql, qr) 區間總和查詢
struct seg {
    struct Node {
        int val, tag;
        Node *l, *r;
        Node(int v=0) : val(v), tag(0), l(nullptr), r(nullptr) {}
    };
    Node* root;
    int n;
    void init(int _n) {
        n = _n;
        root = new Node();
    }
    void push(Node* cur, int l, int r) {
```

```

    if(cur->tag) {
        cur->val += (r-l+1)*cur->tag;
        if(l!=r) {
            if(!cur->l) cur->l = new Node();
            if(!cur->r) cur->r = new Node();
            cur->l->tag += cur->tag;
            cur->r->tag += cur->tag;
        }
    }
    cur->tag = 0;
}

void pull(Node* node, int l, int r) {
    int mid = l+r>>1;
    push(node->l, l, mid);
    push(node->r, mid+1, r);
    node->val = node->l->val + node->r->val;
}

void add(Node* cur, int l, int r, int ql, int qr,
int val) {
    if(ql <= l && r <= qr) {
        cur->tag += val;
        push(cur, l, r);
        return;
    }
    if(!cur->l) cur->l = new Node();
    if(!cur->r) cur->r = new Node();

    int mid = (l + r) / 2;
    push(cur, l, r);
    if(ql<=mid) add(cur->l, l, mid, ql, qr, val);
    if(mid+1<=qr) add(cur->r, mid + 1, r, ql, qr,
        val);
    pull(cur, l, r);
}

int query(Node* cur, int l, int r, int ql, int qr)
{
    if(ql<=l && r<=qr) {
        push(cur, l, r);
        return cur->val;
    }
    if(!cur->l) cur->l = new Node();
    if(!cur->r) cur->r = new Node();
    int mid = l+r>>1;
    int ans = 0;
    push(cur, l, r);
    if(ql<=mid) ans+=query(cur->l, l, mid, ql, qr);
    if(mid+1<=qr) ans+=query(cur->r, mid+1, r, ql,
        qr);
    pull(cur, l, r);
    return ans;
}

int query(int ql, int qr) {
    return query(root, 0, n, ql, qr);
}

void add(int ql, int qr, int val) {
    add(root, 0, n, ql, qr, val);
}
} tree;

```

## 7.6 動態開點線段數 2D

```

struct segy {
    int n;
    struct Node {
        int val;
        Node *l, *r;
        Node(int v=0) : val(v), l(nullptr), r(nullptr)
        {}
    };
    segy(int _n=1e9) {n=_n; root=new Node();}
    Node* root;
    void init(int _n=1e9) {
        n = _n;
        root = new Node();
    }
    void pull(Node* node) {
        node->val = (node->l?node->l->val:0) + (node->r
            ?node->r->val:0);
    }
}

```

```

}
void add(Node* cur, int l, int r, int pos, int val)
{
    if(l==r) {
        cur->val += val;
        return;
    }
    int mid = l+r>>1;
    if(pos<=mid) {if(!cur->l) cur->l = new Node();
        add(cur->l, l, mid, pos, val);}
    else {if(!cur->r) cur->r = new Node();add(cur->
        r, mid + 1, r, pos, val);}
    pull(cur);
}

int query(Node* cur, int l, int r, int ql, int qr)
{
    if(ql<=l && r<=qr) {
        return cur->val;
    }
    int mid = l+r>>1;
    int ans = 0;
    if(ql<=mid) {if(!cur->l) cur->l = new Node();
        ans+=query(cur->l, l, mid, ql, qr);}
    if(mid+1<=qr) {if(!cur->r) cur->r = new Node();
        ans+=query(cur->r, mid+1, r, ql, qr);}
    pull(cur);
    return ans;
}

int query(int ql, int qr) {
    return query(root, 0, n, ql, qr);
}

void add(int pos, int val) {
    add(root, 0, n, pos, val);
}
}

};
struct segx {
    struct Node {
        segy * tree_y;
        Node *l, *r;
        Node(int m) : tree_y(new segy(m)), l(nullptr),
            r(nullptr) {}
    };
    Node* root;
    int n,m;
    segx(int _n=1e9, int _m=1e9) {n=_n; m=_m; root=new
        Node(m);}
    void init(int _n=1e9, int _m=1e9) {n=_n; m=_m; root
        =new Node(m);}
    void add(Node* cur, int l, int r, int qx, int qy,
        int val) {
        if(l!=r) {
            int mid = (l + r) / 2;
            if(qx<=mid) {if(!cur->l) cur->l = new Node(
                m); add(cur->l, l, mid, qx, qy, val);}
            else {if(!cur->r) cur->r = new Node(m); add
                (cur->r, mid + 1, r, qx, qy, val);}
        }
        cur->tree_y->add(qy, val);
    }
    int query(Node* cur, int lx, int rx, int qlx, int
        qrx, int qly, int qry) {
        if(qlx<=lx && rx<=qrx) {
            return cur->tree_y->query(qly, qry);
        }
        int mid = lx+rx>>1;
        int ans = 0;
        if(qlx<=mid) {if(!cur->l) cur->l = new Node(m);
            ans+=query(cur->l, lx, mid, qlx, qrx, qly,
                qry);}
        if(mid+1<=qrx) {if(!cur->r) cur->r = new Node(m)
            ); ans+=query(cur->r, mid+1, rx, qlx, qrx,
                qly, qry);}
        return ans;
    }
    int query(int qlx, int qly, int qrx, int qry) {
        return query(root, 0, n, qlx, qrx, qly, qry);
    }
    void add(int qx, int qy, int val) {
        add(root, 0, n, qx, qy, val);
    }
} tree;

```

## 7.7 持久化線段樹

```

struct seg {
    // 加值持久化線段樹
    struct Node {
        int val;
        Node *l, *r;
    };
    vector<Node*> version;
    void pull(Node* node) {
        node->val = node->l->val + node->r->val;
    }
    Node* build(int l, int r) {
        Node* node = new Node;
        if(l == r) {
            node->val = 0; // 初始值
            return node;
        }
        int mid = (l+r)/2;
        node->l = build(l, mid);
        node->r = build(mid+1, r);
        pull(node);
        return node;
    }
    Node* update(Node* cur, int l, int r, int pos, int v) {
        Node* node = new Node;
        if(l == r) {
            // 改成加值換這行
            // node->val = cur->val + v;
            node->val = v;
            return node;
        }
        int mid = (l+r)/2;
        if(pos <= mid) {
            node->l = update(cur->l, l, mid, pos, v);
            node->r = cur->r;
        } else {
            node->l = cur->l;
            node->r = update(cur->r, mid+1, r, pos, v);
        }
        pull(node);
        return node;
    }
    int query(Node* cur, int s, int e, int ql, int qr) {
        if(ql <= s && e <= qr) return cur->val;
        int ans = 0;
        int mid = (s+e)/2;
        if(ql <= mid) ans += query(cur->l, s, mid, ql, qr);
        if(mid+1 <= qr) ans += query(cur->r, mid+1, e, ql, qr);
        return ans;
    }
} tree;
// push 初始的樹
// tree.version.push_back(tree.build(1, n));

// update(舊版, 1, n, pos, v) return 新版
// 把pos值改成v

```

```

return ret;
}
node* update(node* pre, int s, int e, int pos, int v) {
    node *ret = new node();
    if(s == e) {ret->val = pre->val + v; return ret;}
    int mid = (s+e)>>1;
    if(pos <= mid) {
        ret->l = update(pre->l, s, mid, pos, v);
        ret->r = pre->r;
    } else {
        ret->r = update(pre->r, mid+1, e, pos, v);
        ret->l = pre->l;
    }
    ret->val = ret->l->val + ret->r->val;
    return ret;
}
void add(int pos, int v) {
    timing.push_back(update(timing.back(), 1, n, pos, v));
}
int que(node* pre, node* now, int l, int r, int k) {
    if(l == r) return r;
    int mid = (l+r)>>1;
    int diff = now->l->val - pre->l->val;
    // printf("now %d~%d diff %d\n", l, r, diff);
    if(diff >= k) return que(pre->l, now->l, l, mid, k);
    else return que(pre->r, now->r, mid+1, r, k-diff);
    return -1;
}
int query(int l, int r, int k) {
    l--;
    return que(timing[l], timing[r], 1, n, k);
}
int num[100005];
vector<int> sor;
map<int, int> mp;
signed main() {
    cin >> n >> q;
    timing.push_back(build(1, n));
    for(int i = 0; i < n; i++) {
        cin >> a; num[i] = a; sor.push_back(a);
    }
    // add: 1 1 1 2 1
    // num: 3 3 3 4 3
    // sor: 3 4
    sort(sor.begin(), sor.end());
    sor.erase(unique(sor.begin(), sor.end()), sor.end());
    for(int i = 0; i < n; i++) {
        int pos = lower_bound(sor.begin(), sor.end(), num[i]) - sor.begin() + 1;
        // printf("mp[%d] = %d\n", pos, num[i]);
        mp[pos] = num[i];
        num[i] = pos;
        add(num[i], 1);
    }
    while(q--) {
        int a, b, c; cin >> a >> b >> c;
        cout << mp[query(a, b, c)] << endl;
    }
}

```

## 7.8 Time Segment Tree

```

#include <bits/stdc++.h>
#define int long long int
using namespace std;
int n, q;
struct node {
    int val;
    node *l, *r;
    node(int v) {val = v; l = r = nullptr;}
    node() {val = 0; l = r = nullptr;}
};
vector<node*> timing;
node* build(int s, int e) {
    node *ret = new node();
    if(s == e) return ret;
    int mid = (s+e)>>1;
    ret->l = build(s, mid);
    ret->r = build(mid+1, e);
    ret->val = ret->l->val + ret->r->val;
}

```

## 7.9 Treap

```

struct Treap {
    int sz, val, pri, tag;
    Treap *l, *r;
    Treap(int _val) {
        val = _val; sz = 1;
        pri = rand(); l = r = NULL; tag = 0;
    }
};
int Size(Treap *a) {return a ? a->sz : 0;}
void pull(Treap *a) {
    a->sz = Size(a->l) + Size(a->r) + 1;
}
// val of a is always bigger than val of b
Treap* merge(Treap *a, Treap *b) {
    if(!a || !b) return a ? a : b;
    if(a->pri > b->pri) {
        a->r = merge(a->r, b);
    }
}

```

```

    pull(a);
    return a;
} else {
    b->l = merge( a , b->l );
    pull(b);
    return b;
}
}
// a<k, b>=k
void split(Treap *t, int k, Treap*&a, Treap*&b){
    if(!t) {a=b=NULL; return; }
    if(k <= t->val) {
        b = t;
        split(t->l, k, a, b->l);
        pull(b);
    }
    else {
        a = t;
        split(t->r, k, a->r, b);
        pull(a);
    }
}
Treap* add(Treap *t, int v) {
    Treap *val = new Treap(v);
    Treap *l = NULL, *r = NULL;
    split(t, v, l, r);
    return merge(merge(l, val), r);
}
Treap* del(Treap *t, int v) {
    Treap *l, *mid, *r, *temp;
    split(t, v, l, temp);
    split(temp, v+1, mid, r);
    return merge(l, r);
}

// base 1
int position(Treap *t, int p) {
    if(Size(t->l)+1==p) return t->val;
    if(Size(t->l)<p) return position(t->r, p-Size(t->l)-1);
    else return position(t->l, p);
}

//num of >= k
int query(Treap *t, int k) {
    if(!t) return 0;
    if(t->val==k) return Size(t->l)+1;
    if(t->val>k) return query(t->l, k);
    return Size(t->l)+1+query(t->r, k);
}

```

## 7.10 PBDS

```

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#define ordered_set tree<int, null_type, less<int>,
    rb_tree_tag, tree_order_statistics_node_update>
using namespace __gnu_pbds;
// ordered_set s;
// s.insert(1); s.erase(s.find(1));
// order_of_key(k) : Number of items strictly smaller
    than k .
// find_by_order(k) : K-th element in a set (counting
    from zero). (return iterator)

```

## 8 String

### 8.1 SA

```

#pragma GCC optimize("O3,unroll-loops")
#pragma target optimize("avx2,bmi,bmi2,lzcnt,popcnt")
#include<bits/stdc++.h>
#include<chrono>
#define mid (l + r) / 2
using namespace std;
const int N = 100010;
struct SA{
#define REP(i,n) for ( int i=0; i<int(n); i++ )
#define REP1(i,a,b) for ( int i=a; i<=int(b); i++ )

```

```

    bool _t[N*2];
    int _s[N*2], _sa[N*2], _c[N*2], x[N], _p[N], _q[N*2],
        hei[N], r[N];
    int operator [] (int i){ return _sa[i]; }
    void build(int *s, int n, int m){
        memcpy(_s, s, sizeof(int) * n);
        sais(_s, _sa, _p, _q, _t, _c, n, m);
        mkhei(n);
    }
    void mkhei(int n){
        REP(i,n) r[_sa[i]] = i;
        hei[0] = 0;
        REP(i,n) if(r[i]) {
            int ans = i>0 ? max(hei[r[i-1]] - 1, 0) : 0;
            while(_s[i+ans] == _s[_sa[r[i]-1]+ans]) ans++;
            hei[r[i]] = ans;
        }
    }
    void sais(int *s, int *sa, int *p, int *q, bool *t,
        int *c, int n, int z){
        bool uniq = t[n-1] = true, neq;
        int nn = 0, nmzx = -1, *nsa = sa + n, *ns = s + n,
            lst = -1;
#define MS0(x,n) memset((x),0,n*sizeof(*(x)))
#define MAGIC(XD) MS0(sa, n); \
        memcpy(x, c, sizeof(int) * z); \
        XD; \
        memcpy(x + 1, c, sizeof(int) * (z - 1)); \
        REP(i,n) if(sa[i] && !t[sa[i]-1]) sa[x[s[sa[i]
            ]-1]]++ = sa[i]-1; \
        memcpy(x, c, sizeof(int) * z); \
        for(int i = n - 1; i >= 0; i--) if(sa[i] && t[sa[i]
            ]-1]) sa[--x[s[sa[i]-1]]] = sa[i]-1;
        MS0(c, z);
        REP(i,n) uniq &= ++c[s[i]] < 2;
        REP(i,z-1) c[i+1] += c[i];
        if (uniq) { REP(i,n) sa[--c[s[i]]] = i; return; }
        for(int i = n - 2; i >= 0; i--) t[i] = (s[i]==s[i
            +1] ? t[i+1] : s[i]<s[i+1]);
        MAGIC(REP1(i,1,n-1) if(t[i] && !t[i-1]) sa[--x[s[i]
            ]]=p[q[i]=nn++]=i);
        REP(i, n) if (sa[i] && t[sa[i]] && !t[sa[i]-1]) {
            neq=lst<0||memcmp(s+sa[i],s+lst,(p[q[sa[i]]+1]-sa
                [i])*sizeof(int));
            ns[q[lst=sa[i]]]=nmzx+=neq;
        }
        sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmzx
            + 1);
        MAGIC(for(int i = nn - 1; i >= 0; i--) sa[--x[p[
            nsa[i]]]] = p[nsa[i]]);
    }
}sa;
int H[ N ], SA[ N ];
void suffix_array(int* ip, int len) {
    // should padding a zero in the back
    // ip is int array, len is array length
    // ip[0..n-1] != 0, and ip[len] = 0
    ip[len++] = 0;
    sa.build(ip, len, 128);
    for (int i=0; i<len; i++) {
        H[i] = sa.hei[i + 1];
        SA[i] = sa._sa[i + 1];
    }
    // resulting height, sa array \in [0,len)
}
bool check(string &s,string &t,int p){
    for(int i=0;i<t.size() && i+p<s.size();++i){
        if(t[i]<s[i+p])return 1;
        else if(t[i]>s[i+p]) return 0;
    }
    if(t.size()>s.size()-p) return 0;
    return 1;
}
//example for finding patterns in a string
string s,t;
int ip[N],len;
int main(){
    int n;
    cin>>s>>n;
    len = s.length();
    for(int i=0;i<len;++i) ip[i]=(int)s[i];
    ip[len] = 0;
    suffix_array(ip,len);

```



```

int l,r;
for(int i=0;i<n;++i){
    cin>>t;
    l = 0, r = s.size()-1;
    while(l!=r){
        if(check(s,t,SA[mid])) r=mid;
        else l = mid+1;
    }
    bool f=1;
    if(t.size()>s.size()-SA[l]){
        cout<<"NO\n",f=0;
        continue;
    }
    for(int j=0;j<t.size();++j){
        if(t[j]!=s[j+SA[l]]){
            cout<<"NO\n",f=0;
            break;;
        }
    }
    if(f) cout<<"YES\n";
}
}

```

## 8.2 KMP

```

// 回傳所有匹配成功的起始位置，s為文本，t為匹配字串
// nxt表示為匹配失敗時要退回的位置，也是t字串的相等前綴
// 後綴的最大長度
// *注意前綴後綴為長度最多為n-1的字子串
// nxt[j] = -1 if j=0
// 0 if 沒有相等的前綴後綴
// K k 為相等前綴後綴的最大長度
// 以下為例子
// j: 0 1 2 3 4 5 6
// t: a b a a b e
// nxt[j]: -1 0 0 1 1 2 0
// O(n+m)，n為s長，m為t長
const int MXN = 1e6+5;
int nxt[MXN];
vector<int> KMP(string s,string t){
    int slen = s.length(), tlen = t.length(), i=0,j=0,k
    =-1;
    nxt[0]=-1;
    while(j<tlen){//build nxt
        if(k==-1 || t[j]==t[k]) nxt[++j] = ++k;
        else k=nxt[k];
    }
    i=0,j=0;
    vector<int> ret;
    while(i<slen){// matching
        if(j==-1||s[i]==t[j]) i++,j++;
        else j=nxt[j];
        if(j==tlen){
            ret.push_back(i-tlen+1);//1-base
            j=nxt[j];
        }
    }
    return ret;
}

//另一版
//if t is the substring of s:
//if t in s:
bool cmp(string s, string t) {
    vector<int> front(t.size(), 0);
    for(int i=1, j=0 ; i<t.size() ; i++) {
        while(j>0 && t[i]!=t[j]) j = front[j-1];
        if(t[i]==t[j]) j++;
        front[i] = j;
    }
    int j=0, i=0;
    while(i<s.size()) {
        if(s[i]==t[j]) j++,i++;
        else {i += (j==0); j = (j<1?0:front[j-1]);}
        if(j>=t.size()) return true;
    }
    return false;
}

```

## 8.3 Single Hash

```

//字串雜湊前的idx是0-base，雜湊後為1-base
//H[R] - H[L-1] * p^(R-L+1)
//cmp的+modl是為了防止負數
//記得build完之後要buildPow
//小心遇到hash出負數要記得+modl
#define int long long
const int p = 75577, modl = 1e9 + 7, MXN = 1e6+5;
int Hash[MXN], qpow[MXN];
void build(const string& s) {
    Hash[0]=0;
    for(int i=1; i<=s.size(); i++)
        Hash[i] = (Hash[i-1] * p + s[i-1]) % modl;
}
void buildPow(){
    qpow[0]=1;
    for(int i=1;i<MXN;++i) qpow[i]=qpow[i-1]*p%modl;
}
bool cmp(int i, int j, int len) {
    return (Hash[i+len-1] - Hash[i-1] * qpow[len] %
        modl + modl) % modl ==
        (Hash[j+len-1] - Hash[j-1] * qpow[len] % modl +
        modl) % modl;
}
int get(int i, int j) {
    return (Hash[j]-Hash[i-1]*qpow[j-i+1]%modl+modl)%
        modl;
}

```

## 8.4 Double Hash

```

//字串雜湊前的idx是0-base，雜湊後為1-base
//即區間為 [0,n-1] -> [1,n]
//若要取得區間[L,R]的值則
//H[R] - H[L-1] * p^(R-L+1)
//cmp為比較從i開始長度為Len的字串和從j開始長度為Len的字
//串是否相同
//(h[i+Len-1] - h[i-1] * qpow(p, Len) % modl + modl)
#define int long long
#define x first
#define y second
const int P1 = 75577, P2 = 17, MOD = 1e9 + 7, MXN = 1e6
+5;
pair<int,int> Hash[MXN];
int qpow[2][MXN];
void build(const string& s){
    pair<int,int> val = make_pair(0,0);
    Hash[0]=val;
    for(int i=1; i<=s.size(); i++){
        val.x = (val.x * P1 + s[i-1]) % MOD;
        val.y = (val.y * P2 + s[i-1]) % MOD;
        Hash[i] = val;
    }
}
void buildPow(){
    qpow[0][0]=qpow[1][0]=1;
    for(int i=1;i<MXN;++i){
        qpow[0][i]=qpow[0][i-1]*P1%MOD;
        qpow[1][i]=qpow[1][i-1]*P2%MOD;
    }
}
bool cmp( int i, int j, int len ) {
    return ((Hash[i+len-1].x-Hash[i-1].x*qpow[0][len]%
        MOD+MOD)%MOD == (Hash[j+len-1].x-Hash[j-1].x*
        qpow[0][len]%MOD+MOD)%MOD)
        && ((Hash[i+len-1].y-Hash[i-1].y*qpow[1][len]%MOD+
        MOD)%MOD == (Hash[j+len-1].y-Hash[j-1].y*qpow
        [1][len]%MOD+MOD)%MOD);
}
pair<int, int> get(int i, int j) {
    return {(Hash[j].x-Hash[i-1].x*qpow[0][j-i+1]%MOD+
        MOD)%MOD, (Hash[j].y-Hash[i-1].y*qpow[1][j-i
        +1]%MOD+MOD)%MOD};
}

```

## 8.5 Trie

```
//cnt為記錄有多少個一樣的單詞且end的時候才有數字
const int MXN=1e6+5;//MXN取文本長
int trie[MXN][26], cnt[MXN], tot=0;//0 base
void update(string s){
    int p=0;//0 base
    for(int i=0; i<s.size(); ++i){
        int ch = s[i] - 'a';
        if(!trie[p][ch]) trie[p][ch] = ++tot;
        p = trie[p][ch];
    }
    cnt[p]++;
}
int query(string s){
    int p=0;
    for(int i=0; i<s.size(); ++i){
        int ch=s[i] - 'a';
        p = trie[p][ch];
        if(!p) return 0;
    }
    return cnt[p];
}
void visualizeTrie(int node = 0, int depth = 0) { //for debug
    for (int i = 0; i < 26; ++i) {
        if (trie[node][i]) {
            for (int j = 0; j < depth; ++j) cout << "
";
            cout << (char)('a' + i) << " (" << cnt[trie
                [node][i]] << ") \n";
            visualizeTrie(trie[node][i], depth + 1);
        }
    }
}
```

## 8.6 Z value

```
// O(n)
//z[i] = lcp(s[1...], s[i...])
//1base
int z[MXN];
void Z_value(const string& s) {
    int i, j, left, right, len = s.size();
    left=right=0; z[0]=len;
    for(i=1; i<len; ++i) {
        j=max(min(z[i-left], right-i), 0);
        for(; i+j<len&&s[i+j]==s[j]; ++j);
        z[i]=j;
        if(i+z[i]>right) {
            right=i+z[i];
            left=i;
        }
    }
}
```

## 8.7 MinRotation

```
//rotate(begin(s), begin(s)+minRotation(s), end(s))
//For example, rotations of acab are acab, caba, abac,
//and baca.
//find Lexicographically minimal rotation of a string
int minRotation(string s) {
    int a = 0, N = s.size(); s += s;
    for(int b=0; b<N; b++) for(int k=0; k<N; k++) {
        if(a+k == b || s[a+k] < s[b+k])
            {b += max(0, k-1); break;}
        if(s[a+k] > s[b+k]) {a = b; break;}
    } return a;
}
```

## 8.8 Manacher 馬拉車回文

```
// O(N)求以每個字元為中心的最長回文半徑
// 頭尾以及每個字元間都加入一個
// 沒出現過的字元，這邊以 '@' 為例
// s為傳入的字串，Len為字串長度
// z為儲存以每個字元為中心的回文半徑+1(有包含 '@' 要小心)
// ex: s = "abaac" -> "@a@b@a@a@c@"
// z = [12141232121]
```

```
const int MXN = 1e6+5;
int z[2*MXN];
char s[2*MXN];
void z_value_pal(char *s, int len, int *z){
    len=(len<1)+1;
    for(int i=len-1; i>=0; i--){
        s[i]=i&1?s[i>1]:'@';
        z[0]=1;
        for(int i=1, l=0, r=0; i<len; i++){
            z[i]=i<r?min(z[l+1-i], r-i):1;
            while(i-z[i]>=0&&i+z[i]<len&&s[i-z[i]]==s[i+z[i]])
                ++z[i];
            if(i+z[i]>r) l=i, r=i+z[i];
        }
    }
    // cin>>s;
    // z_value_pal(s, strlen(s), z);
    // int mx=-1, mxi=0;
    // for(int i=0; i<strlen(s); ++i)
    //     if(mx<z[i]) mx = z[i], mxi = i;
    // mx--;
    // for(int i=mx-mx; i<=mxi+mx; ++i)
    //     if(s[i]!='@') cout<<s[i];
}
```

## 8.9 PalTree 回文樹

```
// Len[s]是對應的回文長度
// num[s]是有幾個回文後綴
// cnt[s]是這個回文字串在整個字串中的出現次數
// fail[s]是他長度次長的回文後綴，aba的fail是a
const int MXN = 1000010;
struct PalT{
    int nxt[MXN][26], fail[MXN], len[MXN];
    int tot, lst, n, state[MXN], cnt[MXN], num[MXN];
    int diff[MXN], sfail[MXN], fac[MXN], dp[MXN];
    char s[MXN]={-1};
    int newNode(int l, int f){
        len[tot]=l, fail[tot]=f, cnt[tot]=num[tot]=0;
        memset(nxt[tot], 0, sizeof(nxt[tot]));
        diff[tot]=(l>0?1-len[f]:0);
        sfail[tot]=(l>0&&diff[tot]==diff[f]?sfail[f]:f);
        return tot++;
    }
    int getfail(int x){
        while(s[n-len[x]-1]!=s[n]) x=fail[x];
        return x;
    }
    int getmin(int v){
        dp[v]=fac[n-len[sfail[v]]-diff[v]];
        if(diff[v]==diff[fail[v]])
            dp[v]=min(dp[v], dp[fail[v]]);
        return dp[v]+1;
    }
    int push(){
        int c=s[n] - 'a', np=getfail(lst);
        if(!(lst=nxt[np][c])){
            lst=newNode(len[np]+2, nxt[getfail(fail[np])
                ][c]);
            nxt[np][c]=lst; num[lst]=num[fail[lst]]+1;
        }
        fac[n]=n;
        for(int v=lst; len[v]>0; v=sfail[v])
            fac[n]=min(fac[n], getmin(v));
        return ++cnt[lst], lst;
    }
    void init(const char *_s){
        tot=lst=n=0;
        newNode(0, 1), newNode(-1, 1);
        for(; _s[n];) s[n+1]=_s[n], ++n, state[n-1]=push();
        for(int i=tot-1; i>1; i--) cnt[fail[i]]+=cnt[i];
    }
} palt;
// state 數組
// state[i] 代表第 i 個字元為結尾的最長回文編號(編號
// 是甚麼不重要)
// S = "abacaaba"
// 以第 2(0-base) 個字元為結尾的最長回文是 aba
```

```
// 以第 7(0-base) 個字元為結尾的最長回文是 aba
// 兩個最長回文都相同，因此 state[2] 會等於 state[7]

// Len 數組
// 求出某個 state 的長度
//
// S = "aababa"
//
// (0-base)
// len[state[1]] = 2 ( "aa" )
// len[state[3]] = 3 ( "aba" )
// len[state[5]] = 5 ( "ababa" )

// num 數組
// 某個state的回文有幾個回文後綴
//
// 假設某個 state 代表的回文為 = "ababa" 為例
// state 代表的回文的 num = 3
// -> ababa -> aba -> a

// cnt 數組
// 某個 state 的回文在整個字串中出現次數
//
// S = "aababaa"
// state[3] 代表的回文為 "aba" 在整個字串中出現 2
// 次
// 因此 cnt[state[3]] = 2

// fail數組
// 每個 state 的次長回文後綴的 state 編號
//
// S = "ababa"
// len[fail[4]] = 3 (fail[4] = "aba" )
// len[fail[2]] = 1 (fail[2] = "a" )
// len[fail[0]] = 0 (fail[0] = "" 空字串)
// 0 所代表的 state 是空字串
```

## 8.10 DistinctSubsequence

```
//預設為小寫字母
//return the number of distinct non-empty subsequences
//of sting
#define int long long
int mod = 1e9 + 7;
vector<int> cnt(26);
int distinct_subsequences(string s) {
    for (char c : s)
        cnt[c - 'a'] = accumulate(begin(cnt), end(cnt), 1LL,
            ) % mod;
    return accumulate(begin(cnt), end(cnt), 0LL) % mod;
}
```

## 9 Tree

### 9.1 LCA

```
//先建edge[MXN]
//跑dfs，再跑makeanc
//之後才可以呼叫lca
// 0-base
const int MXN=1e5;
const int logN=__lg(MXN);
int tin[MXN],tout[MXN],anc[MXN][logN+1];
vector<int> edge[MXN];
int ti=0;
void dfs(int x,int f){
    anc[x][0]=f;
    tin[x]=ti++;
    for(int u:edge[x]){
        if(u==f)continue;
        dfs(u,x);
    }
    tout[x]=ti++;
}
// x is y's anc
inline bool isanc(int x,int y){
```

```
    return tin[x]<=tin[y] && tout[x]>=tout[y];
}
int lca(int x,int y){
    if(isanc(x,y))return x;
    if(isanc(y,x))return y;
    for(int i=logN;i>=0;--i){
        if(!isanc(anc[y][i],x)){
            y=anc[y][i];
        }
    }
    return anc[y][0];
}
void makeanc(int n){
    for(int i=1;i<=logN;++i){
        for(int j=0;j<n;++j){
            anc[j][i] = anc[anc[j][i-1]][i-1];
        }
    }
}
```

### 9.2 TreeHash

```
// 1. dfs 先做子樹
// 2. 葉節點的hash值為1
// 3. 對於節點x，其hash值為紀錄x的所有子樹的hash值(紀錄
// 到temp)，然後由小排到大(排除子樹的隨機問題)
// 4. n表示節點x有幾個子樹，p和MOD通常為一個很大的質
// 數，由此算出x的hash值
// 5. 樹根的hash值即為整顆樹的hash值，若兩顆樹的hash值
// 相同，則兩棵樹就是同構
const int MXN = 200005;
int subtree_sz[MXN];
int hash_[MXN];
int base = 44560482149;
int MOD = 274876858367;
int dfs(int x, int fa, vector<int>* edge){
    vector<int> temp;
    subtree_sz[x] = 1;
    for(int child : edge[x]){
        if(child==fa) continue;
        temp.push_back(dfs(child, x, edge));
        subtree_sz[x] += subtree_sz[child];
    }
    sort(temp.begin(), temp.end());
    int ret = subtree_sz[x];
    for(int v : temp){
        ret = (((ret * base + v + ret) % MOD + ret) %
            MOD + v) % MOD ;
    }
    hash_[x] = ret;
    return ret;
}
```

### 9.3 輕重鏈剖分

```
const int MXN = 2e5+7;
int top[MXN], son[MXN], dfn[MXN], rnk[MXN], dep[MXN],
    father[MXN];
vector<int> edge[MXN];
int dfs1(int v, int fa, int d) {
    int maxsz = -1, maxu, total = 1;
    dep[v] = d;
    father[v] = fa;
    for(int u: edge[v]) {
        if(fa == u) continue;
        int temp = dfs1(u, v, d+1);
        total += temp;
        if(temp>maxsz) {
            maxsz = temp;
            maxu = u;
        }
    }
    if(maxsz== -1) son[v] = -1;
    else son[v] = maxu;
    return total;
}
```

```

int times = 1;
void dfs2(int v, int fa) {
    rnk[times] = v;
    dfn[v] = times++;
    top[v] = (fa == -1 || son[fa] != v ? v : top[fa]);
    if(son[v] != -1) dfs2(son[v], v);
    for(int u: edge[v]) {
        if(fa == u || u == son[v]) continue;
        dfs2(u, v);
    }
}
//rnk: 剖分後的編號 (rnk[時間] = 原點)
//dfn: 剖分後的編號 (dfn[原點] = 時間)
//top: 剖分的頭頭
//son: 剖分的重兒子

```

## 10 Geometry

### 10.1 2D Definition

```

#define ld long double
const ld eps=1e-10;
int dcmp(ld x){if(fabs(x)<eps) return 0;else return x<0?-1:1;}
struct Pt{
    ld x,y;
    Pt(ld x=0,ld y=0):x(x),y(y){}
    Pt operator+(const Pt &a) const {
        return Pt(x+a.x, y+a.y); }
    Pt operator-(const Pt &a) const {
        return Pt(x-a.x, y-a.y); }
    Pt operator*(const ld &a) const {
        return Pt(x*a, y*a); }
    Pt operator/(const ld &a) const {
        return Pt(x/a, y/a); }
    ld operator*(const Pt &a) const { //dot
        return x*a.x + y*a.y; }
    ld operator^(const Pt &a) const { //cross
        return x*a.y - y*a.x; }
    bool operator<(const Pt &a) const {
        return x < a.x || (x == a.x && y < a.y); }
    //return dcmp(x-a.x) < 0 || (dcmp(x-a.x) == 0
    // && dcmp(y-a.y) < 0); }
    bool operator>(const Pt &a) const {
        return x > a.x || (x == a.x && y > a.y); }
    //return dcmp(x-a.x) > 0 || (dcmp(x-a.x) == 0
    // && dcmp(y-a.y) > 0); }
    bool operator==(const Pt &a) const {
        return dcmp(x-a.x) == 0 && dcmp(y-a.y) == 0; }
    // return x == other.x && y == other.y;
    bool operator!=(const Pt &a) const {
        return !(*this == a); }
};
typedef Pt Vec;
ld Dot(Vec a,Vec b){return a.x*b.x+a.y*b.y;}
ld Cross(Vec a,Vec b){return a.x*b.y-a.y*b.x;}
ld Length(Vec a){return sqrt(Dot(a,a));}
int Sgn(double x){return (x > -eps) - (x < eps); } //
return 0: x==0, 1: x>0, -1: x<0
ld Angle(Vec a,Vec b){return acos(Dot(a,b)/Length(a)/
Length(b));} //弧度
ld Degree(Vec a,Vec b){return Angle(a,b)*180/acos(-1);}
//角度
ld Ori(Pt a,Pt b,Pt c){return Cross(b-a,c-a);} //1.(a,b)
x(a,c)的面積 2. a在bc左側>0 3. a在bc右側<0 4. a在
bc上==0
Vec Rotate(Vec a,ld rad){return Vec(a.x*cos(rad)-a.y*
sin(rad),a.x*sin(rad)+a.y*cos(rad));} //逆時針旋轉,
rad為弧度
Vec Normal(Vec a){ld L=Length(a);return Vec(-a.y/L,a.x/
L);} //單位法向量,確保a不是零向量
Vec Unit(Vec x) { return x / Length(x); } //單位向量
Vec Perp( const Vec v ){ return { v.y , -v.x };} //垂直
向量
bool argcmp(const Pt &a, const Pt &b) { // 極角cmp: arg
(a) < arg(b)
int f = (Pt{a.y, -a.x} > Pt{b.y, -b.x}) ? 1 : -1 * (a != Pt
{});
int g = (Pt{b.y, -b.x} > Pt{a.y, -a.x}) ? 1 : -1 * (b != Pt
{});

```

```

return f == g ? (a ^ b) > 0 : f < g;
}
struct Circle {
    Pt o; ld r;
    Circle(Pt _o=Pt(0, 0), ld _r=0):o(_o), r(_r) {}
};

```

### 10.2 Line Definition

```

struct Line {
    Pt a, b, v; // start, end, end-start
    ld ang;
    Line(Pt _a=Pt(0, 0), Pt _b=Pt(0, 0)):a(_a),b(_b) { v
        = b-a; ang = atan2(v.y, v.x); }
    bool operator<(const Line &L) const {
        return ang < L.ang;
    }
};
int PtSide(Pt p, Line L) { //return 1:左側 0:線上 -1:右
側
return Sgn(Ori(L.a, L.b, p));
}
bool PtOnSeg(Pt p, Line L) { //點是否在線段上
return Sgn(Ori(L.a, L.b, p)) == 0 && Sgn((p - L.a)
* (p - L.b)) <= 0;
}
Pt Proj(Pt p, Line l) { //點到線段的投影點
Pt dir = Unit(l.b - l.a);
return l.a + dir * (dir * (p - l.a));
}
bool isInter(Line l, Line m) { //判斷兩線段是否相交
if (PtOnSeg(m.a, l) || PtOnSeg(m.b, l) ||
PtOnSeg(l.a, m) || PtOnSeg(l.b, m))
return true;
return PtSide(m.a, l) * PtSide(m.b, l) < 0 &&
PtSide(l.a, m) * PtSide(l.b, m) < 0;
}
Pt LineInter(Line l, Line m) { //兩線段交點
double s = Ori(m.a, m.b, l.a), t = Ori(m.a, m.b, l.
b);
return (l.b * s - l.a * t) / (s - t);
}

```

### 10.3 Basic

```

//確保兩直線P+tv和Q+tw有唯一交點且Cross(v,w)非零
Pt getLineIntersect(Line a, Line b) {
    Pt p1 = a.a, p2 = a.b, q1 = b.a, q2 = b.b;
    ld f1 = (p2-p1)^(q1-p1), f2 = (p2-p1)^(p1-q2), f;
    if(dcmp(f=f1+f2) == 0)
        return dcmp(f1)?Pt(NAN,NAN):Pt(INFINITY,INFINITY);
    return q1*(f2/f) + q2*(f1/f);
}
//點到直線距離
double distanceToLine(Pt p,Pt a,Pt b){
    Vec v1=b-a,v2=p-a;
    return fabs(Cross(v1,v2)/Length(v1));
}
//點到線段距離
double distanceToSegment(Pt p,Pt a,Pt b){
    if(a==b) return Length(p-a);
    Vec v1=b-a,v2=p-a,v3=p-b;
    if(dcmp(Dot(v1,v2))<0) return Length(v2);
    else if(dcmp(Dot(v1,v3))>0) return Length(v3);
    else return fabs(Cross(v1,v2)/Length(v1));
}
//點到直線投影
Pt GetLineProjection(Pt p,Pt a,Pt b){
    Vec v=b-a;
    return a+v*(Dot(v,p-a)/Dot(v,v));
}
//點p於直線ab的對稱點
Pt getSymmetryPoint(Pt p,Pt a,Pt b){
    Pt q=getLineProjection(p,a,b);
    return q*2-p;
}

```

```

}

//判斷線段相交(剛好交一點), 若兩線段共線->c1=c2=0
bool isSegmentProperIntersection(Pt a1,Pt a2,Pt b1,Pt
b2){
    double c1=Cross(a2-a1,b1-a1),c2=Cross(a2-a1,b2-a1),
    c3=Cross(b2-b1,a1-b1),c4=Cross(b2-b1,a2-b1);
    return dcmp(c1)*dcmp(c2)<0&& dcmp(c3)*dcmp(c4)<0;
}

//判斷線段相交(只要有交點即可)
bool isSegmentNotProperIntersection(Pt a1,Pt a2,Pt b1,
Pt b2){
    return max(a1.x,a2.x)>=min(b1.x,b2.x)&&max(b1.x,b2.
x)>=min(a1.x,a2.x)&&max(a1.y,a2.y)>=min(b1.y,b2
.y)&&max(b1.y,b2.y)>=min(a1.y,a2.y)
&&dcmp(Cross(a1-b1,a2-b1))*dcmp(Cross(a1-b2,a2-b2))
<=0&&dcmp(Cross(b1-a1,b2-a1))*dcmp(Cross(b1-a2,
b2-a2))<=0;
}

//點是否在線段上
bool isOnSegment(Pt p,Pt a1,Pt a2){
    return dcmp(Cross(a1-p,a2-p))==0&&dcmp(Dot(a1-p,a2-
p))<=0;
}

```

## 10.4 PolygonArea

```

//須注意Long Long 及 加上絕對值
double polygonArea(Pt* p,int n){
    double area=0;
    for(int i=1;i<n-1;++i){
        area+=Cross(p[i]-p[0],p[i+1]-p[0]);
    }
    return area/2;
}

```

## 10.5 IsPointInPolygon

```

//判斷點是否在多邊形內部
int isPointInPolygon(Pt p,Pt* poly,int n){
    int wn=0;
    for(int i=0;i<n;++i){
        if(isOnSegment(p,poly[i],poly[(i+1)%n])) return
-1;//在邊界上
        int k=dcmp(Cross(poly[(i+1)%n]-poly[i],p-poly[i
]));
        int d1=dcmp(poly[i].y-p.y);
        int d2=dcmp(poly[(i+1)%n].y-p.y);
        if(k>0&&d1<=0&&d2>0) wn++;
        if(k<0&&d2<=0&&d1>0) wn--;
    }
    if(wn!=0) return 1;//內部
    return 0;//外部
}

```

## 10.6 ConvexHull

```

//若要求高精度用dcmp比較
//若是搞int點要記得定義改int
//輸入不能有重複點, 注意h的點未排序!
//若需保留共線點, 把while裡的Ori判斷式改成<=0
void hull(vector<Pt> &dots) { // n=1 => ans = {}
    sort(dots.begin(), dots.end());
    vector<Pt> ans(1, dots[0]);
    for (int ct = 0; ct < 2; ++ct, reverse(all(dots)))
        for (int i = 1, t = SZ(ans); i < SZ(dots); ans.
            push_back(dots[i++]))
            while (SZ(ans) > t && Ori(ans[SZ(ans) - 2], ans.
                back(), dots[i]) <= 0)
                ans.pop_back();
    ans.pop_back(), ans.swap(dots);
}

```

## 10.7 MinkowskiSum

```

//定義: 給兩點集合A,B, Minkowski Sum是A+B={a+b|a∈A,b∈B}
的凸包
//視覺化為一個凸包 A 繞著凸包 B 轉一圈, O(N)
//兩個凸多邊形的 Minkowski sum, 也會是凸多邊形
//P 和 Q 組成的 Minkowski sum 最多有 |P|+|Q| 個點
//在凸包 A 和 B 上的邊也會在 Minkowski sum 上出現
//傳入的點集合不用逆時針排序
vector<Pt> Minkowski(vector<Pt> A, vector<Pt> B) { // |
A|, |B|>=3
    hull(A), hull(B);
    vector<Pt> C(1, A[0] + B[0]), s1, s2;
    for (int i = 0; i < SZ(A); ++i)
        s1.push_back(A[(i + 1) % SZ(A)] - A[i]);
    for (int i = 0; i < SZ(B); ++i)
        s2.push_back(B[(i + 1) % SZ(B)] - B[i]);
    for (int i = 0, j = 0; i < SZ(A) || j < SZ(B);)
        if (j >= SZ(B) || (i < SZ(A) && Cross(s1[i], s2[j])
            >= 0))
            C.push_back(B[j % SZ(B)] + A[i++]);
        else
            C.push_back(A[i % SZ(A)] + B[j++]);
    return hull(C), C;
}

```

## 10.8 Polygon Shortest Distance

```

//給兩多邊形, 求最短距離
double PolyDist(vector<Pt> A, vector<Pt> B) {
    for (auto &p : B) p = {-p.x, -p.y};
    auto C = Minkowski(A, B); // assert SZ(C) > 0
    if (isPointInPolygon(Pt{}, C.data(), C.size()))
        return 0;
    double ans = distanceToSegment(Pt{}, C.back(), C[0])
;
    for (int i = 0; i + 1 < SZ(C); ++i) {
        ans = min(ans, distanceToSegment(Pt{}, C[i], C[i
+ 1]));
    }
    return ans;
}

```

## 10.9 ConvexHullTrick

```

struct Convex {
    int n;
    vector<Pt> A, V, L, U;
    //init, pass convex hull points
    Convex(const vector<Pt> &_A) : A(_A), n(_A.size())
    { // n >= 3
        auto it = max_element(all(A));
        L.assign(A.begin(), it + 1);
        U.assign(it, A.end()), U.push_back(A[0]);
        for (int i = 0; i < n; ++i) {
            V.push_back(A[(i + 1) % n] - A[i]);
        }
    }
    PtSide(Pt p, Line L) {
        return dcmp((L.b - L.a)^(p - L.a));
    }
    int inside(Pt p, const vector<Pt> &h, auto f) {
        auto it = lower_bound(all(h), p, f);
        if (it == h.end()) return 0;
        if (it == h.begin()) return p == *it;
        return 1 - dcmp((p - *prev(it))^( *it - *prev(
it)));
    }
    // 1. whether a given point is inside the Convex
    Hull
    // ret 0: out, 1: on, 2: in
    int inside(Pt p) {
        return min(inside(p, L, less{}), inside(p, U,
            greater{}));
    }
    static bool cmp(Pt a, Pt b) { return dcmp(a ^ b) >
        0; }
}

```



```
// 2. Find tangent points of a given vector
// ret the idx of far/closer tangent point
int tangent(Pt v, bool close = true) {
    assert(v != Pt{});
    auto l = V.begin(), r = V.begin() + L.size() - 1;
    if (v < Pt{}) l = r, r = V.end();
    if (close) return (lower_bound(l, r, v, cmp) - V.begin()) % n;
    return (upper_bound(l, r, v, cmp) - V.begin()) % n;
}

// 3. Find 2 tang pts on CH of a given outside point
// return index of tangent points
// return {-1, -1} if inside CH
array<int, 2> tangent2(Pt p) {
    array<int, 2> t{-1, -1};
    if (inside(p) == 2) return t;
    if (auto it = lower_bound(all(L), p); it != L.end() and p == *it) {
        int s = it - L.begin();
        return {(s + 1) % n, (s - 1 + n) % n};
    }
    if (auto it = lower_bound(all(U), p, greater{}); it != U.end() and p == *it) {
        int s = it - U.begin() + L.size() - 1;
        return {(s + 1) % n, (s - 1 + n) % n};
    }
    for (int i = 0; i != t[0]; i = tangent((A[t[0] = i] - p), 0));
    for (int i = 0; i != t[1]; i = tangent((p - A[t[1] = i]), 1));
    return t;
}

int find(int l, int r, Line L) {
    if (r < l) r += n;
    int s = PtSide(A[l % n], L);
    return *ranges::partition_point(views::iota(l, r), [&](int m) {
        return PtSide(A[m % n], L) == s;
    }) - l;
};

// 4. Find intersection point of a given line
// intersection is on edge (i, next(i))
vector<int> intersect(Line L) {
    int l = tangent(L.a - L.b), r = tangent(L.b - L.a);
    if (PtSide(A[l], L) == 0) return {l};
    if (PtSide(A[r], L) == 0) return {r};
    if (PtSide(A[l], L) * PtSide(A[r], L) > 0) return {};
    return {find(l, r, L) % n, find(r, l, L) % n};
}
};
```

## 10.10 Polar Sort

```
//極角排序，從270度開始逆時針排序
bool cmp(const Pt& lhs, const Pt& rhs) {
    if (Cross((lhs < Pt{}), (rhs < Pt{})))
        return (lhs < Pt()) < (rhs < Pt());
    return Cross(lhs, rhs) > 0;
}

/* 若要以p[i]為原點排序->計算v=p[j]-p[i]
for(int j=0;j<n;++j){
    if(i!=j){
        Vector v = p[j]-p[i];
        node[nodeSz++] = {v,j};
    }
}
sort(node,node+nodeSz,cmp);
*/
```

## 10.11 PickTheorm

```
int area,in,on;//area:多邊形面積 in:內部格點數 on:邊界格點數
```

```
void PickTheorm(Pt* p,int n){
    area=polygonArea(p,n);
    for(int i=0;i<n;++i){
        on+=__gcd(abs((int)p[i].x-(int)p[(i+1)%n].x),
            abs((int)p[i].y-(int)p[(i+1)%n].y));
    }
    in=abs(area)+1-on/2;
}
```

## 10.12 最近點對

```
//最近點對距離注意若整數要define double Long Long
double closestEuclideanDistance(Pt* p,int n){
    sort(p,p+n);
    set<Pt> s={{p[0].y,p[0].x}};
    int j = 0;
    Pt t;
    double dd=LLONG_MAX,d;
    for(int i=1;i<n;++i){
        d = sqrt(dd);
        while(j<i && p[j].x < p[i].x-d){
            s.erase({p[j].y,p[j++].x});
        }
        auto l = s.lower_bound({p[i].y-d,p[i].x-d});
        auto u = s.upper_bound({p[i].y+d,p[i].x+d});
        for(auto it=l;it!=u;it++){
            t = {it->y,it->x};
            dd =min(dd, Dot(p[i]-t,p[i]-t));
        }
        s.emplace(p[i].y,p[i].x);
    }
    return dd;
}
```

## 10.13 幾何中位數

```
//回傳為到每個頂點距離和最小的點
Pt weiszfeld(const Pt *p,int n){
    double nn=n;
    Pt cur = p[0], next;
    for(int i=1;i<n;++i)
        cur.x+=p[i].x, cur.y+=p[i].y;
    cur.x/=nn, cur.y/=nn;
    double w,numerX,numerY,denomin;
    while(1){
        numerX=numerY=denomin=0;
        bool update=0;
        double d;
        for(int i=0;i<n;++i){
            d=length(cur-p[i]);
            if(d>eps){
                w = 1.0/d;
                numerX+=w*p[i].x;
                numerY+=w*p[i].y;
                denomin+=w;
                update=1;
            }else{
                next = p[i];
                break;
            }
        }
        if(update){
            next.x = numerX/denomin;
            next.y = numerY/denomin;
        }
        if(length(cur-next)<eps) break;
        cur = next;
    }
    return next;
}
```

## 10.14 矩陣掃描線

```
#include <bits/stdc++.h>
#define int long long int
using namespace std;
int n, st[1000005<<2], lazy[1000005<<2], old[1000005<<2];
```



```

vector<tuple<int, int, int, int>> v;
vector<int> sor;
void pull(int index, int l, int r) {
    if(lazy[index]) st[index] = old[index];
    else if(l==r) st[index] = 0;
    else st[index] = st[index<<1|1]+st[index<<1];
    // printf("pull %lld~%lld, %lld\n", l, r, st[index]);
    return;
}
void insert(int index, int s, int e, int l, int r, int k) {
    //printf("insert: range %lld~%lld, query %lld~%lld\n", s, e, l, r);
    if(l<=s && e<=r) {
        lazy[index] +=k;
        pull(index, s, e);
        return;
    }
    int mid = (s+e)/2;
    if(l<=mid) insert(index<<1, s, mid, l, r, k);
    if(mid<r) insert(index<<1|1, mid+1, e, l, r, k);
    pull(index, s, e);
}
void input(int index, int l, int r) {
    if(l==r) {
        old[index] = sor[l]-sor[l-1];
        return;
    }
    int mid = (l+r)/2;
    input(index<<1, l, mid);
    input(index<<1|1, mid+1, r);
    old[index] = old[index<<1] + old[index<<1|1];
    //cout<<L<<" to "<<r<<" is "<<old[index]<<endl;
    return;
}
// int diff=1000005;
signed main(){
    cin >> n;
    int l, r, d, u;
    for (int i = 0; i < n; i++){
        cin >> l >> d >> r >> u;
        // l+=diff;
        // d+=diff;
        // r+=diff;
        // u+=diff;
        sor.push_back(d);
        sor.push_back(u);
        v.push_back({l, d, u, 1});
        v.push_back({r, d, u, -1});
    }
    set<int> temp(sor.begin(), sor.end());
    sor = vector<int>(temp.begin(), temp.end());
    sort(sor.begin(), sor.end());
    for(int i=0; i<v.size(); i++) {
        auto [a, b, c, k] = v[i];
        v[i] = make_tuple(a, (int)(lower_bound(sor.begin(), sor.end(), b)-sor.begin()), (int)(lower_bound(sor.begin(), sor.end(), c)-sor.begin()), k);
    }
    input(1, 1, sor.size()-1);
    // cout<<"get: ";
    // for(int i: sor) cout<<i<<" "; cout<<endl;
    sort(v.begin(), v.end());
    int pre=0;
    int ans=0;
    for(auto [pos, a, b, k]: v) {
        if(pre!=pos) {
            ans+=(pos-pre)*st[1];
            pre = pos;
        }
        insert(1, 1, sor.size()-1, a+1, b, k);
        // printf("now act: pos %lld, %lld~%lld, act: %lld\n", pos, a+1, b, k);
        // printf("now ans: %lld\n", st[1]);
    }
    cout<<ans<<endl;
}

```

## 10.15 Polygon Circle intersection area

```

//給兩圓，回傳兩圓的交點，若相交回傳的兩點一樣的話要用
//dcmp判斷
vector<Pt> interCircle(Circle c1, Circle c2) {
    Pt o1 = c1.o, o2 = c2.o;
    ld r1 = c1.r, r2 = c2.r;
    if( Length( o1 - o2 ) > r1 + r2 ) return {};
    if( Length( o1 - o2 ) < max(r1, r2) - min(r1, r2) )
        return {};
    ld d2 = ( o1 - o2 ) * ( o1 - o2 );
    ld d = sqrt(d2);
    if( d > r1 + r2 ) return {};
    Pt u = (o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d));
    ld A = sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2-d));
    ;
    Pt v = Pt( o1.y-o2.y , -o1.x + o2.x ) * A / (2*d);
    return {u+v, u-v};
}

```

## 10.16 Tangent line of two circles

```

//給兩圓，求兩圓的外切線或內切線
vector<Line> go( const Circle& c1 , const Circle& c2 ,
    int sign1 ){
    // sign1 = 1 for outer tang, -1 for inner tang
    vector<Line> ret;
    double d_sq = Dot( c1.o - c2.o , c1.o - c2.o );
    if( d_sq < eps ) return ret;
    double d = sqrt( d_sq );
    Pt v = ( c2.o - c1.o ) / d;
    double c = ( c1.r - sign1 * c2.r ) / d;
    if( c * c > 1 ) return ret;
    double h = sqrt( max( 0.0 , 1.0 - c * c ) );
    for( int sign2 = 1 ; sign2 >= -1 ; sign2 -= 2 ){
        Pt n = { v.x * c - sign2 * h * v.y ,
            v.y * c + sign2 * h * v.x };
        Pt p1 = c1.o + n * c1.r;
        Pt p2 = c2.o + n * ( c2.r * sign1 );
        if( fabs( p1.x - p2.x ) < eps and
            fabs( p1.y - p2.y ) < eps )
            p2 = p1 + Perp( c2.o - c1.o );
        ret.push_back( { p1 , p2 } );
    }
    return ret;
}

```

## 10.17 Circle intersection Point

```

//傳入多邊形和圓形，回傳多邊形和圓形的交集面積
ld PCIntersect(vector<Pt> v, Circle cir) {
    for(int i = 0; i < (int)v.size(); ++i) v[i] = v[i]
        - cir.o;
    ld ans = 0, r = cir.r;
    int n = v.size();
    for(int i = 0; i < n; ++i) {
        Pt pa = v[i], pb = v[(i+1)%n];
        if(Length(pa) < Length(pb)) swap(pa, pb);
        if(dcmp(Length(pb)) == 0) continue;
        ld s, h, theta;
        ld a = Length(pb), b = Length(pa), c = Length(pb-pa);
        ld cosB = (pb*(pb-pa))/a/c, B = acos(cosB);
        if(cosB > 1) B = 0;
        else if(cosB < -1) B = PI;
        ld cosC = (pa*pb)/a/b, C = acos(cosC);
        if(cosC > 1) C = 0;
        else if(cosC < -1) C = PI;
        if(a > r) {
            s = (C/2)*r*r;
            h = a*b*sin(C)/c;
            if(h < r && B < PI/2) s -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
        }
        else if(b > r) {
            theta = PI - B - asin(sin(B)/r*a);
            s = 0.5*a*r*sin(theta) + (C-theta)/2*r*r;
        }
    }
}

```

```

    else s = 0.5*sin(C)*a*b;
    ans += abs(s)*dcmp(v[i]^v[(i+1)%n]);
}
return abs(ans);
}

```

## 10.18 CircleCover

```

#define N 100
#define D long double
struct CircleCover{ //O(N^2LogN)
    int C; Circle c[ N ]; //填入C(圓數量),c(圓陣列,0base)
    bool g[ N ][ N ], overlap[ N ][ N ];
    // Area[i] : area covered by "at least" i circles
    D Area[ N ];
    void init( int _C ){ C = _C; } //總共 _C 個圓
    bool CCinter( Circle& a , Circle& b , Pt& p1 , Pt& p2 )
    {
        Pt o1 = a.o , o2 = b.o;
        D r1 = a.r , r2 = b.r;
        if( Length( o1 - o2 ) > r1 + r2 ) return {};
        if( Length( o1 - o2 ) < max(r1, r2) - min(r1, r2) )
            return {};
        D d2 = ( o1 - o2 ) * ( o1 - o2 );
        D d = sqrt(d2);
        if( d > r1 + r2 ) return false;
        Pt u=(o1+o2)*0.5 + (o1-o2)*((r2*r2-r1*r1)/(2*d2));
        D A=sqrt((r1+r2+d)*(r1-r2+d)*(r1+r2-d)*(-r1+r2+d));
        Pt v=Pt( o1.y-o2.y , -o1.x + o2.x ) * A / (2*d2);
        p1 = u + v; p2 = u - v;
        return true;
    }
    struct Teve {
        Pt p; D ang; int add;
        Teve() {}
        Teve(Pt _a, D _b, int _c):p(_a), ang(_b), add(_c){}
        bool operator<(const Teve &a)const
        {return ang < a.ang;}
    }eve[ N * 2 ];
    // strict: x = 0, otherwise x = -1
    bool disjuct( Circle& a, Circle &b, int x )
    {return dcmp( Length( a.o - b.o ) - a.r - b.r ) > x;}
    bool contain( Circle& a, Circle &b, int x )
    {return dcmp( a.r - b.r - Length( a.o - b.o ) ) > x;}
    bool contain(int i, int j){
        /* c[j] is non-strictly in c[i]. */
        return (dcmp(c[i].r - c[j].r) > 0 ||
            (dcmp(c[i].r - c[j].r) == 0 && i < j) ) &&
            contain(c[i], c[j], -1);
    }
    void solve(){
        for( int i = 0 ; i <= C + 1 ; i ++ )
            Area[ i ] = 0;
        for( int i = 0 ; i < C ; i ++ )
            for( int j = 0 ; j < C ; j ++ )
                overlap[i][j] = contain(i, j);
        for( int i = 0 ; i < C ; i ++ )
            for( int j = 0 ; j < C ; j ++ )
                g[i][j] = !(overlap[i][j] || overlap[j][i] ||
                    disjuct(c[i], c[j], -1));
        for( int i = 0 ; i < C ; i ++ ){
            int E = 0, cnt = 1;
            for( int j = 0 ; j < C ; j ++ )
                if( j != i && overlap[j][i] )
                    cnt ++;
            for( int j = 0 ; j < C ; j ++ )
                if( i != j && g[i][j] ){
                    Pt aa, bb;
                    CCinter(c[i], c[j], aa, bb);
                    D A=atan2(aa.y - c[i].o.y, aa.x - c[i].o.x);
                    D B=atan2(bb.y - c[i].o.y, bb.x - c[i].o.x);
                    eve[E ++] = Teve(bb, B, 1);
                    eve[E ++] = Teve(aa, A, -1);
                    if(B > A) cnt ++;
                }
            if( E == 0 ) Area[ cnt ] += PI * c[i].r * c[i].r;
            else{
                sort( eve , eve + E );
                eve[E] = eve[0];
                for( int j = 0 ; j < E ; j ++ ){
                    cnt += eve[j].add;

```

```

                Area[cnt] += (eve[j].p ^ eve[j + 1].p) * 0.5;
                D theta = eve[j + 1].ang - eve[j].ang;
                if (theta < 0) theta += 2.0 * PI;
                Area[cnt] +=
                    (theta - sin(theta)) * c[i].r*c[i].r * 0.5;
            }
        }
    }
}

```

## 10.19 HalfPlaneIntersection

```

//O(NLgN)
// for point or line solution, change > to >=
bool onleft(Line L, Pt p) {
    return dcmp(L.v^(p-L.a)) > 0;
} // segment should add Counterclockwise
// assume that Lines intersect
// 傳入每條方程式的兩點方程式
// 回傳形成的凸多邊形頂點
// (半平面為像量 ab 的逆時針方向)
//注意題目輸入的點要是逆時針排序
vector<Pt> HPI(vector<Line>& L) {
    sort(L.begin(), L.end()); // sort by angle
    for(auto l:L){
        cerr<<l.a.x<<" "<<l.a.y<<" "<<l.b.x<<" "<<l.b.y<<"
            "<<l.ang<<"\n";
    }
    int n = L.size(), fir, las;
    Pt *p = new Pt[n];
    Line *q = new Line[n];
    q[fir=las=0] = L[0];
    for(int i = 1 ; i < n ; i++) {
        while(fir < las && !onleft(L[i], p[las-1])) las--;
        while(fir < las && !onleft(L[i], p[fir])) fir++;
        q[++las] = L[i];
        if(dcmp(q[las].v^q[las-1].v) == 0) {
            las--;
            if(onleft(q[las], L[i].a)) q[las] = L[i];
        }
        if(fir < las) p[las-1] = getLineIntersect(q[las-1],
            q[las]);
    }
    while(fir < las && !onleft(q[fir], p[las-1])) las--;
    if(las-fir <= 1) return {};
    p[las] = getLineIntersect(q[las], q[fir]);
    int m = 0;
    vector<Pt> ans(las-fir+1);
    for(int i = fir ; i <= las ; i++) ans[m++] = p[i];
    return ans;
}

```

## 10.20 PolygonUnion

```

//O(N^2LgN)
//傳入二維vector，每個vector代表一個多邊形，每個多邊形
//的點必須按照順時針或逆時針順序
//回傳聯集多邊形的面積
ld tri(Pt o, Pt a, Pt b){ return (a-o) ^ (b-o);}
double polyUnion(vector<vector<Pt>> py){ //py[0~n-1]
    must be filled
    int n = py.size();
    int i,j,ii,jj,ta,tb,r,d; double z,w,s,sum=0,tc,td,
        area;
    vector<pair<double,int>> c;
    for(i=0;i<n;i++){
        area=py[i][py[i].size()-1]^py[i][0];
        for(int j=0;j<py[i].size()-1;j++) area+=py[i][j]^py
            [i][j+1];
        if((area/=2)<0) reverse(py[i].begin(),py[i].end());
        py[i].push_back(py[i][0]);
    }
    for(i=0;i<n;i++){
        for(ii=0;ii+1<py[i].size();ii++){
            c.clear();
            c.emplace_back(0.0,0); c.emplace_back(1.0,0);
            for(j=0;j<n;j++){
                if(i==j) continue;
                for(jj=0;jj+1<py[j].size();jj++){
                    ta=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj]))

```

```

        tb=dcmp(tri(py[i][ii],py[i][ii+1],py[j][jj
            +1]));
        if(ta==0 && tb==0){
            if((py[j][jj+1]-py[j][jj])*(py[i][ii+1]-py[
                i][ii])>0&&j<i){
                c.emplace_back(segP(py[j][jj],py[i][ii],
                    py[i][ii+1]),1);
                c.emplace_back(segP(py[j][jj+1],py[i][ii
                    ],py[i][ii+1]),-1);
            }
        }else if(ta>=0 && tb<0){
            tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
            td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
            c.emplace_back(tc/(tc-td),1);
        }else if(ta<0 && tb>=0){
            tc=tri(py[j][jj],py[j][jj+1],py[i][ii]);
            td=tri(py[j][jj],py[j][jj+1],py[i][ii+1]);
            c.emplace_back(tc/(tc-td),-1);
        } } }
        sort(c.begin(),c.end());
        z=min(max(c[0].first,0.0),1.0); d=c[0].second; s
            =0;
        for(j=1;j<c.size();j++){
            w=min(max(c[j].first,0.0),1.0);
            if(!d) s+=w-z;
            d+=c[j].second; z=w;
        }
        sum+=(py[i][ii]^py[i][ii+1])*s;
    } }
    return sum/2;
}

```

## 10.21 PolygonCover

```

//傳入二維vector，每個vector代表一個多邊形，每個多邊形
//的點必須按照順時針或逆時針順序
// return Area[i] : area covered by "at Least" i
// polygon
//O(N^2LgN)
vector<double> PolyUnion(const vector<vector<Pt>> &P) {
    auto reorder = [&](vector<Pt> &v) { //排序成逆時針
        且最左下角的點在最前面
        rotate(v.begin(), min_element(all(v)), v.end())
        ;
        if (Ori(v[0], v[1], v[2]) < 0) {
            reverse(all(v));
        }
    };
    for(auto &i: py){ reorder(i);}
    const int n = P.size();
    vector<double> Area(n + 1);
    vector<Line> Ls;
    for (int i = 0; i < n; i++)
        for (int j = 0; j < P[i].size(); j++)
            Ls.push_back({P[i][j], P[i][(j + 1) % P[i].
                size()]});
    auto cmp = [&](Line &l, Line &r) {
        Pt u = l.b - l.a, v = r.b - r.a;
        if (argcmp(u, v)) return true;
        if (argcmp(v, u)) return false;
        return PtSide(l.a, r) < 0;
    };
    sort(all(Ls), cmp);
    for (int l = 0, r = 0; l < Ls.size(); l = r) {
        while (r < Ls.size() and !cmp(Ls[l], Ls[r])) r
            ++;
        Line L = Ls[l];
        vector<pair<Pt, int>> event;
        for (auto &ls : Ls) {
            Pt c = ls.a, d = ls.b;
            if (Sgn((L.a - L.b) ^ (c - d)) != 0) {
                int s1 = PtSide(c, L) == 1;
                int s2 = PtSide(d, L) == 1;
                if (s1 ^ s2) event.emplace_back(
                    LineInter(L, {c, d}), s1 ? 1 : -1);
            } else if (PtSide(c, L) == 0 and Sgn((L.a -
                L.b) ^ (c - d)) > 0) {
                event.emplace_back(c, 2);
                event.emplace_back(d, -2);
            }
        }
    }
}

```

```

    }
    sort(all(event), [&](auto i, auto j) {
        return (L.a - i.ft) * (L.a - L.b) < (L.a -
            j.ft) * (L.a - L.b);
    });
    int cov = 0, tag = 0;
    Pt lst{0, 0};
    for (auto [p, s] : event) {
        if (cov >= tag) {
            Area[cov] += lst ^ p;
            Area[cov - tag] -= lst ^ p;
        }
        if (abs(s) == 1) cov += s;
        else tag += s / 2;
        lst = p;
    }
    for (int i = n - 1; i >= 0; i--) Area[i] += Area[i
        + 1];
    for (int i = 1; i <= n; i++) Area[i] /= 2;
    return Area;
}

```

## 11 特殊題目

### 11.1 包含子字串計數

```

// * 給一個字串s
// * 求長度為Len且有包含s的字串有幾種
// * 呼叫solve(s, Len)
const int len = 1005;
int aut[len][26];
int dp[len][len];
const int mod = 1e9+7;
void prefix(string &s, vector<int> &pi) {
    for(int i=1, j=0 ; i<s.size() ; i++) {
        while(j>0 && s[i]!=s[j]) j = pi[j-1];
        if(s[i]==s[j]) j++;
        pi[i] = j;
    }
}
void automata(string &s, vector<int> &pi) {
    for(int i=0 ; i<s.size() ; i++) {
        for(int c=0 ; c<26 ; c++) {
            if(i>0 && c+'A' != s[i]) aut[i][c] = aut[pi
                [i-1]][c];
            else aut[i][c] = i + (c + 'A'==s[i]);
        }
    }
}
int quai(int x, int n) {
    if(n==0) return 1;
    int mid = quai(x,n/2);
    mid = mid*mid%mod;
    if(n&1) return mid*x%mod;
    return mid;
}
int solve(string s, int len) {
    vector<int> pi(s.size(), 0);
    prefix(s, pi);
    automata(s, pi);
    int n = s.size(), ans = quai(26, len);
    dp[0][0] = 1;
    for(int i=0 ; i<len ; i++) {
        for(int j=0 ; j<n ; j++) {
            for(int c=0 ; c<26 ; c++) {
                dp[i+1][aut[j][c]] += dp[i][j];
                dp[i+1][aut[j][c]] %= mod;
            }
        }
    }
    for(int i=0 ; i<n ; i++) ans = (ans - dp[len][i] +
        mod)%mod;
    return ans;
}

```

### 11.2 三維偏序

```
// vec
```

```
// {{a, b, c},
// {a, b, c},
// ...
// {a, b, c}}
// 貼上 BIT 模板
// 三維偏序
// a <= a, b <= b, c <= c
map<vector<int>, int> cnt;
int cdq(vector<vector<int>> &vec, int l, int r) {
    if(l==r) return 0;
    int mid = l+r>>1;
    int ans = cdq(vec, l, mid)+cdq(vec, mid+1, r);
    vector<vector<int>> temp;
    for(int i=l, j=mid+1; i<=mid || j<=r; ) {
        while(i<=mid && (j>r || vec[i][1] <= vec[j][1])) {
            bit.add(vec[i][2], cnt[vec[i]]); temp.
            push_back(vec[i++]);
        }
        if(j<=r) {
            temp.push_back(vec[j]);
            ans += bit.query(vec[j][2]);
        }
    }
    for(int i=l; i<=mid; i++) bit.add(vec[i][2], -cnt[vec[i]]);
    for(int i=l; i<=r; i++) vec[i] = temp[i-1];
    return ans;
}
int solve(vector<vector<int>> &vec) {
    bit.init(2e5+5);
    for(vector<int> v: vec) cnt[v]++;
    sort(vec.begin(), vec.end());
    vec.erase(unique(vec.begin(), vec.end()), vec.end());
    return cdq(vec, 0, vec.size()-1);
}
```

```
# re+ 1個或多個
# re? 0個或1個
# re{n} 恰好n次

s = 'a1bb2c2'
print(re.search(r'(?<=\\d)[a-z](?=\\d)', s).group())
# b
# 匹配前面和後面都要是數字的a~z字元

s = 'a1a a2b b3b'
print(re.search(r'(?<!a)\\d(?:!a)', s).group())
# 3
# 匹配前後都不是a的數字

s = 'aBc'
print(re.search(r'abc', s, re.I).group())
# aBc
# re.I 忽略大小寫
# re.S 任意字元 . 不受換行(\\n)和空白限制

s = 'abbbbbb'
print(re.search(r'ab+?', s).group())
# ab
# 非貪婪用法，匹配ab+但長度越短越好

s = 'abcde'
print(re.findall(r'[a-z]', s))
# ['a', 'b', 'c', 'd', 'e']

s = 'abbbba'
print(re.sub(r'(?<=a).*(?=a)', '', s))
# aa
# 把.*的部分替換成空字串
```

## 12 Python

### 12.1 Decimal

```
from decimal import Decimal, getcontext, ROUND_FLOOR
getcontext().prec = 250 # set precision (MAX_PREC)
getcontext().Emax = 250 # set exponent limit (MAX_EMAX)
getcontext().rounding = ROUND_FLOOR # set round floor
itwo,two,N = Decimal(0.5),Decimal(2),200
pi = angle(Decimal(-1))
```

### 12.4 Misc

```
# 轉為高精度整數比，(分子，分母)
x=0.2
x.as_integer_ratio() # (8106479329266893,
9007199254740992)
x.is_integer() # 判斷是否為整數
x.__round__() # 四捨五入
int(eval(num.replace("/", "//"))) # parser string num
```

### 12.2 Fraction

```
from fractions import Fraction
import math
"""專門用來表示和操作有理數，可以進行算"""
frac1 = Fraction(1) # 1/1
frac2 = Fraction(1, 3) # 1/3
frac3 = Fraction(0.5) # 1/2
frac4 = Fraction('22/7') # 22/7
frac5 = Fraction(8, 16) # 自動約分為 1/2
frac9 = Fraction(22, 7)
frac9.numerator # 22
frac9.denominator # 7
x = Fraction(math.pi)
y2 = x.limit_denominator(100) # 分母限制為 100
print(y2) # 311/99
float(x) #轉換為浮點數
```

### 12.3 正則表達式 re

```
import re

# \\d 0~9
# \\D 非0~9
# \\s 空白字元
# \\S 非空白字元
# ^ 開頭
# $ 結尾
# re* 0個或多個
```

A full-page sheet of white graph paper featuring a uniform grid of thin black lines. The grid consists of small squares, with thicker vertical and horizontal lines dividing the page into larger sections. There are no margins, text, or other markings on the paper.

A full-page sheet of white graph paper featuring a uniform grid of thin black lines. The grid consists of small squares, typical of standard graph paper used for mathematics or engineering. There are no margins, text, or other markings on the page.